

CTA

$$A := \alpha \mid A \rightarrow A \mid N$$

$$t =_v x^A \mid \lambda x^A t \mid tt \mid \text{cond}_x(t, t) \mid 0 \mid \text{st}$$

$$\frac{}{\Gamma, x:A \vdash x:A} (\text{Ax}) \quad \frac{\Gamma \vdash t:A \rightarrow B \quad \Gamma \vdash u:A}{\Gamma \vdash tu:B} (\text{App}) \quad \frac{\Gamma, x:A \vdash t:B}{\Gamma \vdash \lambda x^A t:A \rightarrow B} (\text{Abs})$$

$$\frac{\Gamma \vdash f:A \quad \Gamma, x:N \vdash g:A}{\Gamma \vdash \text{cond}_x(f, g):N \rightarrow A} (\text{cond}) \quad \frac{}{\Gamma \vdash 0:N} (\text{Zero}) \quad \frac{\Gamma \vdash t:N}{\Gamma \vdash \text{st}:N} (\text{Succ})$$

$$\frac{\Gamma \vdash t:A \rightarrow B}{\Gamma, x:A \vdash tx:B} (\eta)$$

\vdash is the greatest fixed pt that satisfies $(D_1) \sim (D_7)$

$$(D_1) ((\Gamma, x:A), x, A) \in \vdash$$

$$(D_2) (\Gamma, t, A \rightarrow B), (\Gamma, u, A) \in \vdash \Rightarrow (\Gamma, tu, B) \in \vdash$$

$$(D_3) ((\Gamma, x:A), t, B) \in \vdash \Rightarrow (\Gamma, \lambda x^A t, A \rightarrow B) \in \vdash$$

$$(D_4) (\Gamma, 0, N) \in \vdash$$

$$(D_5) (\Gamma, t, N) \in \vdash \Rightarrow (\Gamma, \text{st}, N) \in \vdash$$

$$(D_6) (\Gamma, t, A), ((\Gamma, x:N), u, A) \in \vdash \Rightarrow (\Gamma, \text{cond}_x(t, u), N \rightarrow A) \in \vdash$$

$$(D_7) (\Gamma, t, A \rightarrow B) \in \vdash \Rightarrow ((\Gamma, x:A), tx, B) \in \vdash$$

$$\Gamma \vdash t:A \stackrel{\text{def}}{\iff} (\Gamma, t, A) \in \vdash$$

$$\frac{\frac{\frac{\frac{\frac{\frac{\frac{x:N \vdash \text{sum}(x):N \rightarrow N}{x:N \vdash \text{sum}(x):N \rightarrow N}}{x:N, y:N \vdash \text{sum}(x)(y):N}}{x:N, y:N \vdash S(\text{sum}(x)(y)):N}}{x:N \vdash \text{cond}_y(x, S(\text{sum}(x)(y))):N \rightarrow N}}{\vdash \lambda x. \text{cond}_y(x, S(\text{sum}(x)(y))):N \rightarrow N \rightarrow N}}{\vdash \text{sum}:N \rightarrow N \rightarrow N}$$

$$\text{sum} = \lambda x. \text{cond}_y(x, S(\text{sum}(x)(y)))$$

global trace condition : Any infinite path contains a progressing trace

Reduction

(1-step full reduction)

$$(\lambda x. t) u \xrightarrow{1} t[u/x]$$

$$\text{cond}_x(t, u)(0) \xrightarrow{1} t$$

$$\text{cond}_x(t, u)(\varsigma v) \xrightarrow{1} u[v/x]$$

& closed under (finite?) substitution

(1-step safe reduction) $t \xrightarrow{1} u$ with "safe" condition
(not inside cond)

Possibly infinitary many times reduction $\xrightarrow{\omega}$ (full reduction)

$\xrightarrow{\omega}$ (safe reduction)

Finitely many times reduction $\xrightarrow{*}, \xrightarrow{*}$

Lem (Substitution) $\Gamma, x:A \vdash t:B \quad \& \quad \Gamma \vdash u:A \quad \Rightarrow \quad \Gamma \vdash t[u/x]:B$

Prop Subject reduction

Def closed t^N is total $\stackrel{\text{def}}{\iff} t \xrightarrow{*} n \in N$
closed t^N is a value $\stackrel{\text{def}}{\iff} t \in N$

closed $t^{A \rightarrow B}$ is total $\stackrel{\text{def}}{\iff} \forall a:A : \text{closed value}. \quad (ta)^B : \text{total}$

closed $t^{A \rightarrow B}$ is a value $\stackrel{\text{def}}{\iff} t \text{ is Total}$

$t[\vec{x}]$ is total $\stackrel{\text{def}}{\iff} \forall \vec{a}: \vec{A} : \text{closed value}. \quad t[\vec{a}]: \text{total}$
 $\vec{x}: \vec{A}$
 $\vec{x} = \text{FV}(t)$

Lem ① $t^A : \text{closed} \& t \rightarrow u : \text{total} \Rightarrow t : \text{total}$

② $\vec{x} = FV(t) \& \vec{z} : \vec{\beta} \& \left\{ \begin{array}{l} \vec{u} : \vec{\beta} \\ \vec{a} : \vec{A} \end{array} \right\} \text{closed value}, t^{\vec{A} \rightarrow N} [\vec{u}] \vec{a} : \text{total} \Rightarrow t : \text{total}$

∴ ① Ind on A

$$\bullet A = N \quad \frac{t^N}{\text{closed.}} \xrightarrow{\text{SR}} \frac{u^N}{\substack{\text{closed} \\ \vec{u}}} \xrightarrow{\text{total}} n \quad \therefore t^N : \text{total}$$

• $A = A_1 \rightarrow A_2$ Take $a : A_1 : \text{closed total}$

$$\frac{(ta)^{A_2}}{\text{closed}} \xrightarrow{\text{SR}} \frac{(u^{A_1 \rightarrow A_2} a)^{A_2}}{\substack{\text{total by } u : \text{total}}} \quad \therefore ta : \text{total by I.H.}$$

Then $t : \text{total.}$

② Assume that $\vec{x} = FV(t) \& \vec{z} : \vec{\beta} \& \left\{ \begin{array}{l} \vec{u} : \vec{\beta} \\ \vec{a} : \vec{A} \end{array} \right\} \text{closed value}, t^{\vec{A} \rightarrow N} [\vec{u}] \vec{a} : \text{total}$

Ind on $|\vec{A}|$

• $|\vec{A}| = 0 : \text{By def}$

• $\vec{A} = A_0 \vec{A}' : \text{Take } \vec{u} : \text{closed values, } \vec{a}' : \text{closed values, } a_0^{A_0} : \text{closed value}$

By asmp. $(t[\vec{u}] a_0 \vec{a}')^N : \text{total}$

By I.H. $(t[\vec{u}] a_0)^{\vec{A}' \rightarrow N} : \text{total}$

$\therefore t[\vec{u}]^{A_0 \rightarrow \vec{A}' \rightarrow A} : \text{total by def}$

By def $t[\vec{z}] : \text{total}$

Thm t : not total $\Rightarrow \Gamma \not\vdash t : A$ in CT λ

\therefore Assume that t : not total & $\vec{z} : \vec{D} \vdash t : \vec{A} \rightarrow N$ has a CT λ -proof Π

By Lem②, $\exists \vec{\alpha}, \vec{d}$: closed values s.t. $t[\vec{d}] \vec{\alpha}$: not total

We inductively construct $(e_i, \vec{d}_i, \vec{\alpha}_i)$ for each $i \in N$

s.t. $\begin{cases} e_i : \text{node of } \Pi \text{ where } \Pi|_{e_i} = (\vec{z}_i : \vec{D}_i \vdash t_i : \vec{A}_i \rightarrow N) \text{ & } e_{i+1} \text{ is a child of } e_i \\ t_i : \text{not total} \\ \vec{d}_i, \vec{\alpha}_i : \text{closed values s.t. } (t_i[\vec{d}_i] \vec{\alpha}_i)^N : \text{not total} \text{ & } \left(\begin{array}{c} \vec{d}_i \\ \vec{\alpha}_i \end{array} \right) \xrightarrow{\vec{d}_{i+1}} \left(\begin{array}{c} \vec{d}_{i+1} \\ \vec{\alpha}_{i+1} \end{array} \right) \end{cases}$

• $(e_0, \vec{d}_0, \vec{\alpha}_0) \stackrel{\text{def}}{=} (\varepsilon, \vec{d}, \vec{\alpha})$ where. $\Pi|_{e_0} = (\vec{z} : \vec{D} \vdash t : \vec{A} \rightarrow N)$

connection from i -th to $(i+1)$ -th

• Assume that we already have $(e_i, \vec{d}_i, \vec{\alpha}_i)$

Case (Ax): $\Pi|_{e_i} = \Gamma, x : A \vdash \frac{x : A}{t_i}$ not the case (x : total)

Case (Zero): $\Pi|_{e_i} = \Gamma \vdash \frac{0 : N}{t_i}$ not the case (0 : total)

Case (Succ): $\Pi|_{e_{i+1}} \stackrel{\text{def}}{=} \frac{\Gamma \vdash t_{i+1} : N}{\Gamma \vdash \frac{St_{i+1} : N}{t_i}}$ By IH $(St_{i+1})[\vec{d}_i] : \text{not total} \text{ & } \vec{\alpha}_i = \emptyset$

$\therefore t_{i+1}[\vec{d}_i] : \text{not total}$
Take $\vec{d}_{i+1} \stackrel{\text{def}}{=} \vec{d}_i$ & $\vec{\alpha}_{i+1} \stackrel{\text{def}}{=} \emptyset$ & $\left(\begin{array}{c} \vec{d}_i \\ \emptyset \end{array} \right) \xrightarrow{\vec{d}_{i+1}} \left(\begin{array}{c} \vec{d}_{i+1} \\ \emptyset \end{array} \right)$

Case (Abs): $\Pi|_{e_{i+1}} \stackrel{\text{def}}{=} \frac{\Gamma, z : A \vdash t_{i+1} : \vec{A} \rightarrow N}{\Pi|_{e_i} = \Gamma \vdash \frac{\lambda z. t : A \rightarrow \vec{A} \rightarrow N}{t_i}}$ By IH $(\lambda z. t)[\vec{d}_i] \frac{\vec{\alpha}_i}{\vec{\alpha}_i} : \text{not total}$

$$(\lambda z. t)[\vec{d}_i] \vec{\alpha}_i = (\lambda z. t[\vec{d}_i, z]) \vec{\alpha}_i \rightarrow \underline{t[\vec{d}_i, a] \vec{\alpha}_i}$$

Take $\vec{d}_{i+1} \stackrel{\text{def}}{=} \vec{d}_i, a$ & $\vec{\alpha}_{i+1} \stackrel{\text{def}}{=} \vec{\alpha}_i$ & $\left(\begin{array}{c} \vec{d}_i \\ \vec{\alpha}_i \end{array} \right) \xrightarrow{\vec{d}_{i+1}} \left(\begin{array}{c} \vec{d}_{i+1} \\ \vec{\alpha}_{i+1} \end{array} \right)$ not total by & Lem①

Case (App): $\Pi|_{e_i} = \frac{\Gamma \vdash t : C \rightarrow \vec{A} \rightarrow N}{\Gamma \vdash \frac{tu : \vec{A} \rightarrow N}{t_i}}$ By IH. $(tu)[\vec{d}_i] \vec{\alpha}_i : \text{not total}$

Subcase: $u[\vec{d}_i] : \text{not total}$ (let $C = \vec{C} \rightarrow N$. There exist $\vec{C} : \text{closed total}$ s.t. $u[\vec{d}_i] \vec{C} : \text{not total}$)

Take $\Pi|_{e_{i+1}} \stackrel{\text{def}}{=} (\Gamma \vdash u : C)$, $t_{i+1} \stackrel{\text{def}}{=} u$, $\vec{d}_{i+1} \stackrel{\text{def}}{=} \vec{d}_i$, $\vec{\alpha}_{i+1} \stackrel{\text{def}}{=} \vec{C}$ & $\left(\begin{array}{c} \vec{d}_i \\ \vec{\alpha}_i \end{array} \right) \xrightarrow{\vec{d}_{i+1}} \left(\begin{array}{c} \vec{d}_{i+1} \\ \vec{C} \end{array} \right)$

Subcase: $u[\vec{d}_i] : \text{total}$.

Take $\Pi|_{e_{i+1}} \stackrel{\text{def}}{=} (\Gamma \vdash t : A \rightarrow \vec{A} \rightarrow N)$, $t_{i+1} \stackrel{\text{def}}{=} t$, $\vec{d}_{i+1} \stackrel{\text{def}}{=} \vec{d}_i$, $\vec{\alpha}_{i+1} \stackrel{\text{def}}{=} u[\vec{d}_i], \vec{\alpha}_i$ & $\left(\begin{array}{c} \vec{d}_i \\ \vec{\alpha}_i \end{array} \right) \xrightarrow{\vec{d}_{i+1}} \left(\begin{array}{c} \vec{d}_{i+1} \\ u[\vec{d}_i], \vec{\alpha}_i \end{array} \right)$

$$\text{Case (1)} : \overline{\Gamma \vdash t : N \rightarrow A} \quad \text{By IH. } (tx)[\vec{d}d] \vec{a}_i : \text{not total}$$

$$\overline{\Pi e_i = \frac{\Gamma \vdash t : N \vdash \underline{tx} : A}{(\vec{x}: \text{fresh})}} \quad \text{if } \vec{d}_i$$

$$\text{Take } t_{i+1} \stackrel{\text{def}}{=} t \quad \& \quad \vec{d}_{i+1} \stackrel{\text{def}}{=} \vec{d} \quad \& \quad \vec{a}_{i+1} \stackrel{\text{def}}{=} d \vec{a}_i \quad \& \quad \left(\begin{array}{c|c} \vec{d} & \vec{d} \\ \hline d & d \\ \vec{a}_i & \vec{a}_i \end{array} \right) \xrightarrow{\vec{d}d} \left(\begin{array}{c|c} \vec{d} & \vec{d} \\ \hline d & d \\ \vec{a}_i & \vec{a}_i \end{array} \right) \xrightarrow{\vec{d}d} \dots$$

$$\text{Case (cond)} : \frac{\Gamma \vdash t : A \quad \Gamma, z : N \vdash u : A}{\Gamma \vdash \text{cond } z. (t, u) : N \rightarrow A} \quad \text{By IH. cond } z. (t, u)[\vec{d}_i] \vec{a} : \text{not total}$$

Subcase ($n=0$) :

$$\text{cond } z. (t[\vec{d}_i], u[\vec{d}_i, \vec{z}]) \vec{a} \rightarrow \underbrace{t[\vec{d}_i] \vec{a}}_{\text{not total by } \text{Lem ①}}$$

$$\text{Take } \overline{\Pi e_m = (\Gamma \vdash t : A)}$$

$$t_{i+1} \stackrel{\text{def}}{=} t$$

$$\vec{d}_{i+1} \stackrel{\text{def}}{=} \vec{d}_i$$

$$\vec{a}_{i+1} \stackrel{\text{def}}{=} \vec{a}$$

$$\& \quad \left(\begin{array}{c|c} \vec{d}_i & \vec{d}_i \\ \hline 0 & 0 \\ \vec{a} & \vec{a} \end{array} \right) \xrightarrow{\vec{d}d} \left(\begin{array}{c|c} \vec{d}_i & \vec{d}_i \\ \hline 0 & 0 \\ \vec{a} & \vec{a} \end{array} \right)$$

Subcase ($n = S_{n'}$) :

$$\text{cond } z. (t[\vec{d}_i], u[\vec{d}_i, \vec{z}]) (S_{n'}) \vec{a} \rightarrow \underbrace{u[\vec{d}_i, n'] \vec{a}}_{\text{not total by } \text{Lem ①}}$$

$$\text{Take } \overline{\Pi e_m = (\Gamma, z : N \vdash u : A)}$$

$$t_{i+1} \stackrel{\text{def}}{=} u$$

$$\vec{d}_{i+1} \stackrel{\text{def}}{=} \vec{d}_i, n'$$

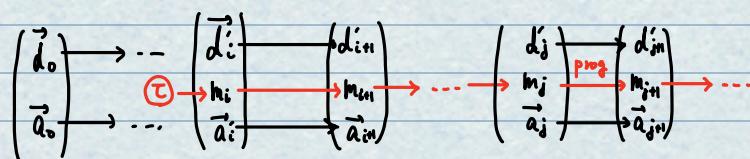
$$\vec{a}_{i+1} \stackrel{\text{def}}{=} \vec{a}$$

$$\& \quad \left(\begin{array}{c|c} \vec{d}_i & \vec{d}_i \\ \hline S_{n'} & n' \\ \vec{a} & \vec{a} \end{array} \right) \xrightarrow{\vec{d}d} \left(\begin{array}{c|c} \vec{d}_i & \vec{d}_i \\ \hline S_{n'} & n' \\ \vec{a} & \vec{a} \end{array} \right)$$

$\textcircled{*}$ $(e_i)_{i \in \omega}$ is an inf path in Π

By regularity, $(e_i)_{i \in \omega}$ contains progressing trace $(\tau_i)_{i \in \omega}$

We have



$$m_{j+1} = m_j \text{ if } j : \text{not prog.pt.}$$

$$m_{j+1} = m_j - 1 \text{ if } j : \text{not prog.pt.} \therefore (m_j)_{j \in \omega} \text{ is infinitely decreasing seq. of } \geq 0 \text{ contradiction!} //$$

Cor $\vdash t : A$ in CT- $\lambda \Rightarrow t : \text{total}$

Cor $\vdash t : N \Rightarrow \exists n \in \mathbb{N} (t \xrightarrow{*} n)$

Unique N.F

Def t, u : closed terms of type A. We define $t \sim_A u$

$$t \sim_n u \stackrel{\text{def}}{\iff} \forall t', u' \in \text{NF} \left(t \xrightarrow{*} t' \& u \xrightarrow{*} u' \Rightarrow t' = u' \right)$$

$$t \sim_{A \rightarrow B} u \stackrel{\text{def}}{\iff} \forall a_1^A, a_2^B : \text{closed} \left(a_1 \sim_A a_2 \Rightarrow ta_1 \sim_B ua_2 \right)$$

$$t[\vec{x}] \sim_A u[\vec{x}] \stackrel{\text{def}}{\iff} \forall \vec{a}_1, \vec{a}_2 : \text{closed} \left(\vec{a}_1 \sim_A \vec{a}_2 \Rightarrow t[\vec{a}_1] \sim_A u[\vec{a}_2] \right)$$

Lem $t \sim_N t \Rightarrow t$ has a unique n.f.

\therefore trivial,

Lem $t \sim_n u \Rightarrow \exists n \in \mathbb{N} \left(\begin{array}{l} t \xrightarrow{*} n \& u \xrightarrow{*} n \& \\ \& \forall t' \in \text{NF} (t \xrightarrow{*} t' \Rightarrow t' = n) \\ \& \forall u' \in \text{NF} (u \xrightarrow{*} u' \Rightarrow u' = n) \end{array} \right)$

\therefore By the previous Cor, $t \xrightarrow{*} m$ & $u \xrightarrow{*} n$. By $t \sim_n u$, $m = n$

| If $t \xrightarrow{*} t' \in \text{NF}$. By $t \sim_n u$, $t' = n$

| If $u \xrightarrow{*} u' \in \text{NF}$. By $t \sim_n u$, $u' = n$ //

Lem $\forall \vec{u}, \vec{u}', \vec{a}, \vec{a}' : \text{closed}$

$$\left(\vec{u} \sim \vec{u}' \& \vec{a} \sim \vec{a}' \Rightarrow t[\vec{u}] \vec{a} \sim_n t'[\vec{u}'] \vec{a}' \right) \Rightarrow t[\vec{x}] \sim_{A \rightarrow N} t'[\vec{x}]$$

\therefore Assume $\forall \vec{u}, \vec{u}', \vec{a}, \vec{a}' : \text{closed} \left(\vec{u} \sim \vec{u}' \& \vec{a} \sim \vec{a}' \Rightarrow t[\vec{u}] \vec{a} \sim_n t'[\vec{u}'] \vec{a}' \right)$.

| Take $\vec{u}, \vec{u}', \vec{a}, \vec{a}'$ s.t. $\vec{u} \sim \vec{u}'$ & $\vec{a} \sim \vec{a}'$. Then $t[\vec{u}] \vec{a} \sim_n t'[\vec{u}'] \vec{a}'$.

$$\therefore t[\vec{u}] \sim_{A \rightarrow N} t'[\vec{u}']$$

$$\therefore t[\vec{x}] \sim_{A \rightarrow N} t'[\vec{x}] \quad \checkmark$$

Thm $t \not\vdash t \Rightarrow \Gamma \not\vdash t : A \text{ in CT}\lambda$

∴ Assume that $t \not\vdash t$ & $\vec{z} : \vec{D} \vdash t : \vec{A} \rightarrow N$ has a $\text{CT}\lambda$ -proof Π (Show contradiction)

Then, by Lem, $t[\vec{d}] \vec{a} \not\vdash_N t'[\vec{d}'] \vec{a}'$ for some $\vec{a} \sim_{\vec{A}} \vec{a}'$ & $\vec{d} \sim_{\vec{D}} \vec{d}'$

We inductively construct $(e_i, \vec{d}_i, \vec{a}_i, \vec{d}'_i, \vec{a}'_i)$ for each $i \in \omega$

where $\Pi e_i = \vec{x}_i : \vec{D}_i \vdash t_i : \vec{A}_i \rightarrow N$

$$\& \vec{d}_i \sim_{\vec{D}_i} \vec{d}'_i \quad \& \vec{a}_i \sim_{\vec{A}_i} \vec{a}'_i$$

$$\& t_i[\vec{d}_i] \vec{a}_i \not\vdash t_i[\vec{d}'_i] \vec{a}'_i$$

& $(\vec{d}_i \xrightarrow{\vec{a}_i} \vec{d}_{i+1}), (\vec{d}'_i \xrightarrow{\vec{a}'_i} \vec{d}'_{i+1})$ are trace compatible

- $(e_0, \vec{d}_0, \vec{a}_0, \vec{d}'_0, \vec{a}'_0) \stackrel{\text{def}}{=} (\varepsilon, \vec{d}, \vec{a}, \vec{d}', \vec{a}')$ where $\Pi \varepsilon = (\vec{z} : \vec{D} \vdash t : \vec{A} \rightarrow N)$

- Assume that we already have $(e_i, \vec{d}_i, \vec{a}_i, \vec{d}'_i, \vec{a}'_i)$

Case (Ax): $\Pi e_i = \Gamma, x : A \vdash \frac{x : A}{t_i}$ not the case (by $x[\vec{d}_i] \sim x[\vec{d}'_i]$)

Case (Zero): $\Pi e_i = \Gamma \vdash 0 : N$ not the case (by $0[\vec{d}_i] = 0 \sim 0 = 0[\vec{d}'_i]$)

Case (Succ): $\Pi e_{i+1} \stackrel{\text{def}}{=} \frac{\Gamma \vdash t_{i+1} : N}{\Pi e_i \vdash \frac{St_{i+1} : N}{t_i}}$

By IH $(St_{i+1})[\vec{d}_i] \not\vdash_N (St_{i+1})[\vec{d}'_i] \quad \& \quad \vec{a}_i = \vec{a}'_i = \emptyset$

∴ $\exists u, u' \in \text{NF} (u \neq u' \& S(t_{i+1}[\vec{d}_i]) \xrightarrow{*} u \& S(t_{i+1}[\vec{d}'_i]) \xrightarrow{*} u')$

∴ $u = Su \& u' = Su' \& u \neq u' \& t_{i+1}[\vec{d}_i] \xrightarrow{*} u \& t_{i+1}[\vec{d}'_i] \xrightarrow{*} u'$ for some $u, u' \in \text{NF}$

∴ $t_{i+1}[\vec{d}_i] \not\vdash_N t_{i+1}[\vec{d}'_i]$

Take $\vec{d}_{i+1} \stackrel{\text{def}}{=} \vec{d}_i \& \vec{d}'_{i+1} \stackrel{\text{def}}{=} \vec{d}'_i \& \vec{a}_{i+1} = \vec{a}'_{i+1} = \emptyset \& (\vec{d}_i \xrightarrow{\vec{a}_i} \vec{d}_{i+1}) \xrightarrow{id} (\vec{d}'_i \xrightarrow{\vec{a}'_i} \vec{d}'_{i+1})$

$$\text{Case (Abs)} : \quad \Pi e_{i+1} = \frac{\Gamma, x:A \vdash t_{i+1} : \vec{A} \rightarrow N}{\Pi e_i = \Gamma \vdash \lambda x. t_{i+1} : A \rightarrow \vec{A} \rightarrow N}$$

$$\text{By IH, } (\lambda x. t_i[\vec{d}_i, x]) \xrightarrow[\vec{a}_i]{\vec{a}} \not\vdash_N (\lambda x. t_i[\vec{d}'_i, x]) \xrightarrow[\vec{a}'_i]{\vec{a}'}$$

$$\therefore t_i[\vec{d}_i, a] \not\vdash_N t_i[\vec{d}'_i, a'] \quad \left(\begin{array}{l} \text{since } (\lambda x \dots) a \vec{e} \not\vdash_N (\lambda x \dots) a' \vec{e}' \\ \downarrow \quad \downarrow \\ t_i[\vec{l}, \vec{a}] \not\vdash_N t_i[\vec{l}', \vec{a}'] \end{array} \right)$$

$\therefore \neq$

Take $\vec{d}_{i+1} \stackrel{\text{def}}{=} \vec{d}_i, a$ & $\vec{d}'_{i+1} \stackrel{\text{def}}{=} \vec{d}'_i, a'$
 $\vec{a}_{i+1} \stackrel{\text{def}}{=} \vec{a}$ & $\vec{a}'_{i+1} \stackrel{\text{def}}{=} \vec{a}'$

Then

$$\text{Case (App)} : \quad \frac{\Gamma \vdash t : A \rightarrow \vec{A} \rightarrow N \quad \Gamma \vdash u : A}{\Pi e_i = \Gamma \vdash tu : \vec{A} \rightarrow N}$$

$$\text{By IH, } (t[\vec{d}_i])(u[\vec{d}_i]) \vec{a}_i \not\vdash_N (t[\vec{d}'_i])(u[\vec{d}'_i]) \vec{a}'_i$$

$$\text{Subcase 1: } u[\vec{d}_i] \not\vdash_{\vec{B} \rightarrow N} u[\vec{d}'_i].$$

$$u[\vec{d}_i] \vec{b} \not\vdash_N u[\vec{d}'_i] \vec{b}' \quad \& \quad \vec{b} \underset{\exists}{\sim} \vec{b}' \quad \text{for some } \vec{b}, \vec{b}'$$

Take $\Pi e_{i+1} \stackrel{\text{def}}{=} \Gamma \vdash u : A$
 $\vec{d}_{i+1} \stackrel{\text{def}}{=} \vec{d}_i$ & $\vec{d}'_{i+1} \stackrel{\text{def}}{=} \vec{d}'_i$ & $\vec{a}_{i+1} \stackrel{\text{def}}{=} \vec{b}$ & $\vec{a}'_{i+1} \stackrel{\text{def}}{=} \vec{b}'$

Then

$$\text{Subcase 2: } u[\vec{d}_i] \sim_A u[\vec{d}'_i].$$

Take $\Pi e_{i+1} \stackrel{\text{def}}{=} \Gamma \vdash t : A \rightarrow \vec{A} \rightarrow N$

$$\vec{d}_{i+1} \stackrel{\text{def}}{=} \vec{d}_i, \quad \vec{d}'_{i+1} \stackrel{\text{def}}{=} \vec{d}'_i \quad \& \quad \vec{a}_{i+1} \stackrel{\text{def}}{=} u[\vec{d}_i], \vec{a}_i \quad \& \quad \vec{a}'_{i+1} \stackrel{\text{def}}{=} u[\vec{d}'_i], \vec{a}'_i$$

Then

$$\text{Case } (\gamma) : \quad \begin{array}{c} \text{Te}_{i+1} \stackrel{\text{def}}{=} \frac{\Gamma \vdash t_{i+1} : A \rightarrow B}{\Gamma, x:A \vdash t_{i+1} x:B} \\ \text{Te}_i \stackrel{\text{def}}{=} \end{array}$$

By IH, $\vec{d}_i = \vec{c}, c$ & $\vec{d}'_i = \vec{c}', c'$ & $t_{i+1}[\vec{c}] c \vec{a}_i \not\sim_N t_{i+1}[\vec{c}'] c' \vec{a}'_i$

Take $\vec{d}_{i+1} \stackrel{\text{def}}{=} \vec{c}$ & $\vec{d}'_{i+1} \stackrel{\text{def}}{=} \vec{c}'$ & $\vec{a}_{i+1} \stackrel{\text{def}}{=} c, \vec{a}_i$ & $\vec{a}'_{i+1} = c, \vec{a}'_i$

Then
$$\left(\begin{array}{c|cc} \vec{c} & \xrightarrow{\text{id}} & \vec{c} \\ \vec{c} & \xrightarrow{\text{id}} & \vec{c} \\ \hline \vec{a}_i & \xrightarrow{\text{id}} & \vec{a}_i \end{array} \right) \quad \& \quad \left(\begin{array}{c|cc} \vec{c}' & \xrightarrow{\text{id}} & \vec{c}' \\ \vec{c}' & \xrightarrow{\text{id}} & \vec{c}' \\ \hline \vec{a}'_i & \xrightarrow{\text{id}} & \vec{a}'_i \end{array} \right)$$

Case (cond) :

$$\text{Te}_i = \frac{\Gamma \vdash t : A \quad \Gamma, z:N \vdash u : A}{\Gamma \vdash \text{cond}_z(t, u) : N \rightarrow A}$$

By IH, $\text{cond}_z(t[\vec{d}_i], u[\vec{d}_i, z]) \not\sim_{\vec{a}_i} \text{cond}_z(t[\vec{d}'_i], u[\vec{d}'_i, z])$

By $b \sim_n b'$, $\exists n \in N (b \xrightarrow{*} n \xleftarrow{*} b')$

Case $n=0$: $\text{cond}_z(t[\vec{d}_i], u[\vec{d}_i, z]) \not\sim_{\vec{a}_i} \text{cond}_z(t[\vec{d}'_i], u[\vec{d}'_i, z])$

$$\begin{array}{ccc} \text{cond}_z(t[\vec{d}_i], u[\vec{d}_i, z]) \not\sim_{\vec{a}_i} & & \text{cond}_z(t[\vec{d}'_i], u[\vec{d}'_i, z]) \not\sim_{\vec{a}'_i} \\ \downarrow * & & \downarrow + \\ \text{cond}_z(t[\vec{d}_i], u[\vec{d}_i, z]) \circ b & & \text{cond}_z(t[\vec{d}'_i], u[\vec{d}'_i, z]) \circ b' \\ \downarrow & & \downarrow \\ t[\vec{d}_i] b & \not\sim & t[\vec{d}'_i] b' \end{array}$$

Take $\text{Te}_{i+1} \stackrel{\text{def}}{=} \Gamma \vdash t : A$

$\vec{d}_{i+1} \stackrel{\text{def}}{=} \vec{d}_i$ & $\vec{d}'_{i+1} \stackrel{\text{def}}{=} \vec{d}'_i$ & $\vec{a}_{i+1} \stackrel{\text{def}}{=} b$ & $\vec{a}'_{i+1} = b'$

$$\left(\begin{array}{c|cc} \vec{d}_i & \xrightarrow{\text{id}} & \vec{d}_i \\ b & \xrightarrow{\text{id}} & b \\ \hline \vec{a}_i & \xrightarrow{\text{id}} & \vec{a}_i \end{array} \right) \quad \& \quad \left(\begin{array}{c|cc} \vec{d}'_i & \xrightarrow{\text{id}} & \vec{d}'_i \\ b' & \xrightarrow{\text{id}} & b' \\ \hline \vec{a}'_i & \xrightarrow{\text{id}} & \vec{a}'_i \end{array} \right)$$

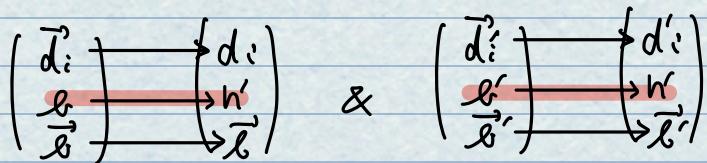
Case $n=S^m$: $\text{cond}_z(t[\vec{d}_i], u[\vec{d}_i, z]) \not\sim_{\vec{a}_i} \text{cond}_z(t[\vec{d}'_i], u[\vec{d}'_i, z])$

$$\begin{array}{ccc} \text{cond}_z(t[\vec{d}_i], u[\vec{d}_i, z]) \not\sim_{\vec{a}_i} & & \text{cond}_z(t[\vec{d}'_i], u[\vec{d}'_i, z]) \not\sim_{\vec{a}'_i} \\ \downarrow * & & \downarrow + \\ \text{cond}_z(t[\vec{d}_i], u[\vec{d}_i, z](S^m)b) & & \text{cond}_z(t[\vec{d}'_i], u[\vec{d}'_i, z](S^m)b') \\ \downarrow & & \downarrow \\ u[\vec{d}_i, n] b & \not\sim & u[\vec{d}'_i, n] b' \end{array}$$

Take $\text{Te}_{i+1} \stackrel{\text{def}}{=} \Gamma, z:N \vdash u:A$

$$\vec{d}_{i+1} \stackrel{\text{def}}{=} \vec{d}_i, n' \quad \& \quad \vec{d}'_{i+1} \stackrel{\text{def}}{=} \vec{d}'_i, n' \quad \& \quad \vec{a}_{i+1} \stackrel{\text{def}}{=} \vec{a}_i \quad \& \quad \vec{a}'_{i+1} \stackrel{\text{def}}{=} \vec{a}'_i$$

Then



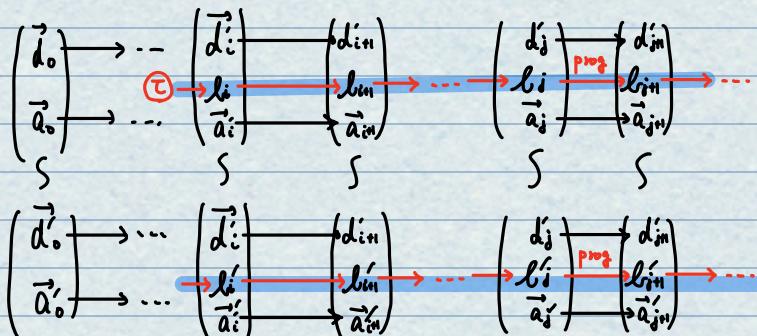
④ For traces of type N , $b \sim_N b'$ has a unique n.f. $\in N$

$$\begin{array}{c} b \rightsquigarrow n \\ \text{ns} \\ b' \rightsquigarrow n' \\ \text{progpt} \\ \downarrow \\ S_n \end{array}$$

⑤ $(e_i)_{i \in \omega}$ is an inf path in Π

By regularity, $(e_i)_{i \in \omega}$ contains progressing trace $(\tau_i)_{i \in \omega}$ of type N

We have



Define $m_j \stackrel{\text{def}}{=} \text{n.f. of } b_j : N$

$$m_{j+1} = m_j \text{ if } j: \text{not prog.pt.}$$

$$m_{j+1} = m_j - 1 \text{ if } j: \text{not prog.pt.}$$

$\therefore (m_j)_{j \in \omega}$ is infinitely decreasing seq. of ≥ 0 contradiction! //