

Problem 1

The goal is to implement Newtons method for a generic function $F:\mathbb{R}^2\longrightarrow\mathbb{R}^2$. I will use a symbolic computation library to compute the Jacobian J_F , which means the function needs to consist of sympy-compatible primitives. Except for this restriction, the functions can be generic. I have attempted to abstract away the sympycode, because it is just boilerplate that pollutes the interesting part.

LISTING I: Helper functions for symbolic manipulation

```
1 # file: src/alg.py
2 import numpy as np
3 from numpy import linalg as LA
4 import sympy as sp
5 import scipy.sparse as scsp
8
   def symbolic jac(py fn):
9
        # Computes a symbolic jacobian matrix
10
        # for a function f : R^2 -> R^2.
11
        # compute the entries of the vector by
       # evaluating the function for sp-symbols
       x, y = sp.symbols('x y')
       f1, f2 = py_fn(x, y)
15
16
       F = sp.Matrix([f1, f2])
17
18
       return F.jacobian([x, y])
19
20
21 def callable fn(symbolic):
       # Create a function that substitutes
22
23
       # for the symbolic values.
       x, y = sp.symbols('x y')
24
25
       return sp.lambdify(
26
            [x, y], symbolic, 'numpy'
```

Armed with some auxiliary functions to handle the symbolic computation, we implement the iteration using the Newton method equation

$$\mathbf{x} = \mathbf{x} - \mathbf{J}_{\mathrm{F}}^{-1} \mathbf{F}(\mathbf{x})$$

LISTING 2: Newton's method

```
# file: src/alg.py
2 MAX ITER = 100
   def solve (F, \times0, tol=1E-6):
5
        x, y = x0
6
7
        J = symbolic jac(F)
8
9
        # singular jacobian means trouble
10
        Jfn = callable fn(J)
11
        assert LA.det(Jfn(x, y)) != 0
12
13
        # function-version of the Jacobian
14
        Ji = callable fn(J.inv())
15
16
        def step(f, Ji f, x):
17
            # computes the next iteration using the
18
            # Newton method equation.
19
            # r is the previous step
20
            return x - Ji f(*x).dot(f(*x))
21
22
        for _ in range(MAX_ITER):
23
            px, py = x, y
24
            x, y = step(F, Ji, (x, y))
25
            yield x, y
26
27
            # check the tolerance criteria
28
            if LA.norm(F(x, y)) < tol:
29
                break
30
            if LA.norm((x - px, y - py)) < tol:
31
                break
32
33
   def last(it):
34
        # run an iterator to the end
35
        x = None
        for x in it: pass
37
        return x
    # interactive session
   >>> from src.alg import solve, last
   >>> def F(x, y):
           return x**2 + y**2 - 2, x - y
   >>> solve(F, (-1,0))
   <qenerator...>
   >>> last(solve(F, (-1,0)))
    (-1.00000000013107, -1.00000000013107)
   >>> last(solve(F, (1,0)))
    (1.00000000013107, 1.00000000013107)
```

The interactive session shows how the function can be used, (it may not be so obvious since it is implemented as a generator-function, so we can collect the error *from outside*; single resposibility principle and so on) and that it is correct at least for two points in different basins of attraction for the equation

$$\mathbf{F}(x,y) = \begin{pmatrix} x^2 + y^2 - 2 \\ x - y \end{pmatrix},\,$$

which has its true roots in (-1, -1) and (1, 1).

QUADRATIC CONVERGENCE

We want to verify that Newtons method converges quadratically, also in the multivariable case. To see this, we want to evaluate the limit

$$\mu = \lim_{n \to \infty} \frac{\left\| \mathbf{x}_{n+1} - \mathbf{x}_n \right\|_2}{\left\| \mathbf{x}_n - \mathbf{x}_{n-1} \right\|_2^2}.$$

Obviously, we don't have an infinite number of terms of $\{x_n\}$. The best we can do is approximate μ by the ratios of our finite sequence.

LISTING 3: Computing the sequence of ratios

```
approx = list(solve(F, (1000, 0), tol=1E-15))
2
3
   def pairs(L):
4
       yield from zip(L[1:], L)
5
6
   def diffs(L):
7
       for (xn, yn), (xm, ym) in pairs(L):
8
            yield (xn - xm, yn - ym)
9
10
   def norms (vs):
11
            for v in vs: yield LA.norm(v)
12
   norms_of_diffs = list(norms(diffs(approx)))
13
14
   for p, q in pairs(norms of diffs):
15
16
       print(p/q**2)
```

Which produces the following values. Note that we need to use a very low tolerance, otherwise we will not see anything resembling convergence at all.

Table 1: The sequence $\{\mu_n\}$ of ratios of error.

Iteration	μ_n	
1	0.001 414	
2	0.002 828	
3	0.005 656	
4	0.011 309	
5	0.022 596	
6	0.045 009	
7	0.088 582	
8	0.166701	
9	0.272763	
10	0.341 980	
11	0.353 357	
12	0.353 553	
13	0.354 073	

It seems like the sequence settles on $\mu \approx 0.35$, which indicates quadratic convergence.

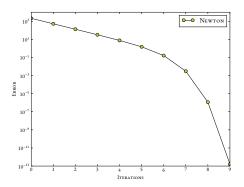


FIGURE I: Convergence of Newton's method with the norm $\|\mathbf{F}(x,y)\|_2$

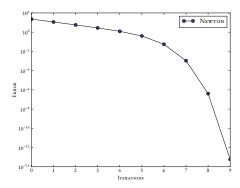


Figure 2: Convergence of Newton's method with the norm $\|\mathbf{x}_{k+1} - \mathbf{x}_k\|_2$

The figures indicate that the convergence is at least superlinear. It is not easy to read from an image exactly how fast the convergence is, but with the estimated μ , quadratic convergence seems likely.

Convergence along diagonals

We want to see what happens as the method converges for two initial guesses that converges to different solutions.

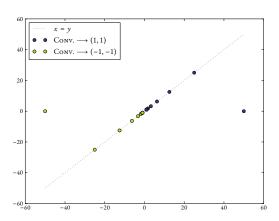


FIGURE 3: Convergence of Newton's method along the line x = y, starting at $\mathbf{x} = (-50, 0)$ and $\mathbf{x} = (50, 0)$.

As we can see, the solutions instantly "jump" to the x = y diagonal, and converges along it.

OPTIONAL PROBLEM

Given

$$f(z) = z^3 - 1$$

as a function $f: \mathbb{C} \longrightarrow \mathbb{C}$, we want to inspect the basins of attraction for Newton's mehod. We can view f as a function $f: \mathbb{R}^2 \longrightarrow \mathbb{R}^2$ instead, and use the solver we already have.

For every point in the complex plane, we are interested in which of the three roots of f Newton's method converges to.

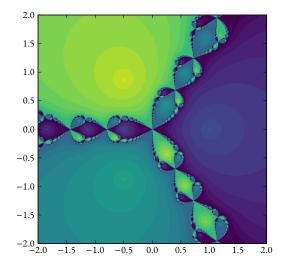


FIGURE 4: Basins of attraction for Newton's method applied on $f(z) = z^3 - 1$.

FIGURE 4 shows how the convergence forms three "basins" in a beautiful fractal pattern.

Problem 2

We want to consider the linear system

$$A\mathbf{u} = \mathbf{f}$$

where

$$A = (L + (\Delta x)^2 k^2 I)$$

is a matrix in $\mathbb{R}^{n^2 \times n^2}$, and $\Delta x = 1/n$. Notice that A is an operator that operates on vectors in \mathbb{R}^{n^2} , corresponding to an $n \times n$ lattice in a domain Ω :

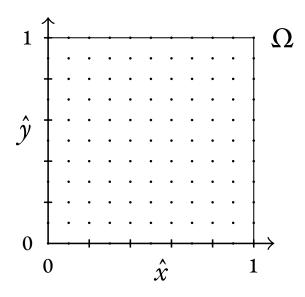


FIGURE 5: Lattice points in Ω .

For any function f(x, y) defined on Ω , we can let

$$\mathbf{f} = \begin{pmatrix} f_1 & f_2 & \cdots & f_l & \cdots & f_{n^2} \end{pmatrix}; \quad f_l = f(x_i, y_i)\,,$$

where

$$\begin{cases} x_i = i \cdot \Delta x \\ y_j = j \cdot \Delta x \\ l = (j-1) n + i \end{cases}$$

This corresponds exactly to "sampling" f at the lattice points in Ω : If we consider a concrete function

$$f(x,y) \coloneqq \exp\left(-50\left(\left(x - \frac{1}{2}\right) + \left(y - \frac{1}{2}\right)\right)\right)$$

defined on Ω , the sampled data looks something like FIGURE 6 for a relatively low choice of n.

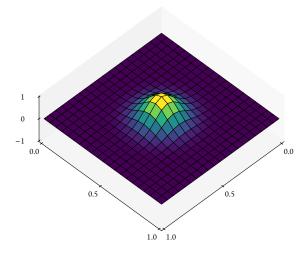


FIGURE 6: f(x, y) sampled across a lattice in Ω .

This sampling can be done succinctly in python, using a generator function to produce the lattice:

LISTING 4: Program to sample functions over lattices

sample (F, n) produces exactly the vector \mathbf{f} , for a given function F, sampled over a lattice with n points along each axis.

NOTE THAT i have programmed it in a way that does not use indeces directly *on purpose*, because we usually work with 1-indexed arrays in mathematics, and 0-indexed arrays in Python, so it is difficult to reason back and forth because you have to mentally translate the indexing, or write functions to translate for you. Either way it becomes a mess.

To solve the system $A\mathbf{u} = \mathbf{f}$, we want to use fix-point iteration:

$$A\mathbf{u} = \mathbf{f}$$

$$(M - N)\mathbf{u} = \mathbf{f}$$

$$M\mathbf{u} = N\mathbf{u} + \mathbf{f}$$

$$\mathbf{u} = M^{-1}(N\mathbf{u} + \mathbf{f})$$

in this equation, $M^{-1}N$ and $M^{-1}f$ are constants. It is sensible (and more efficient) to compute them once upfront and give them names. This gives the more orderly iteration

$$\mathbf{u} \leftarrow C\mathbf{u} + \mathbf{g}$$
,

which converges by the Banach fixed-point theorem if $\mathbf{x} \mapsto C\mathbf{u} + \mathbf{g}$ is a contraction, which is the case as long as

$$0 < \rho(C) := \max_{\lambda \text{ eig. of } C} \lambda < 1.$$

In python, this can be implemented as follows

LISTING 5: *N-dimensional solver*.

```
# src/alg.py
   def spectral radius(M):
3
       return np.max(np.abs(LA.eigvals(M)))
4
5
   def solve nd fpi(M, N, f):
6
        \# solves the linear system (M - N)u = b by
7
        # fix-point iteration u = inv(M)(Nu + b).
8
9
       Mi = LA.inv(M)
10
       C = Mi.dot(N)
       q = Mi.dot(f)
```

provided we already have a choice of M and N. Several ways to choose these matrices are possible, and we want to be able to choose.

Listing 6: Argument-"parser" and choice of M, N.

```
# src/alg.py
   def jacobi mat(A):
       # Jacobi method
        M = np.diag(A.diagonal())
        N = M - A
7
        return M, N
   def gs mat(A):
10
        # Gauss-Seidel
11
        M = np.tril(A)
12.
        N = M - A
13
        return M, N
14
15
   def sor_mat(A, omega):
        # successive over-relaxation
16
17
        D = np.diag(A.diagonal())
18
        L = np.tril(A, k=-1)
19
        M = D + omega*L
        N = M - A
20
21
        return M, N
22
23
   def choose matrices(A, method="jacobi", omega=1.0):
24
        # pick a method based on the parameter
25
        # i have included some redundant parameters
26
        # so it is possible to write "shorthand"
27
        # upper-case also works.
28
        M, N = {
29
            # Jacobi method
            "jacobi":
30
                              jacobi mat,
31
                              jacobi mat,
32
            # Gauss-Seidel method
            "gauss-seidel": gs_mat,
33
34
                              gs_mat,
35
            # Successive over-relaxation
36
            "sor": lambda A: sor mat(A, omega),
37
        } [method.lower()](A)
```

```
39
        return M, N
40
41
42
   def solve nd(A, f, method="jacobi", omega=1.0):
        \# solve Ax = f.
43
44
        # returns an iterator over tuples (u, r),
        # where u is successively better solutions,
45
        # and r is the residual vector.
46
47
        # omega is unused unless "sor" is specified
48
49
        M, N = choose matrices(A,
50
                                method=method,
51
                                omega=omega)
52
53
        for u in solve nd fpi(M, N, f):
54
            # compute the residual
55
            r = f - A.dot(u)
56
            yield u, r
```

Now, given some matrix M, and some vector \mathbf{b} , we can solve systems of the form $M\mathbf{u} = \mathbf{b}$. Notice that, again, the system is implemented in such a way that it produces successively better and better approximations. It also computes a residual vector which it gives us along with each approximation as a tuple (\mathbf{u}, \mathbf{r}) . This makes the API a little clunky if all we want to do is compute the solution of some system, but it makes it easy to work with the data as a sequence. If all we want is the solution, we need the left-hand element of the last tuple.

Testing the solver for a trivial problem

We want to make sure our code is correct for a simple problem, just so that we can have *some* confidence that it actually works. Solving a simple system, which we know has a solution, such as

$$\begin{cases} 3x - y = 1 \\ 2x + 2y = 0 \end{cases}$$

would give us a good hint about possible errors we might have made. (By insertion it is easy to verify that x = 1/4

and y = -1/4 is a solution)

```
1 # interactive session
2 >>> from src.alg import solve nd, last
3 >>> import numpy as np
5
   >>> A = np.array([[3, -1],
                     [2, 2]])
   >>> b = np.array([1, 0])
10
   >>> u, v = last(solve nd(A, b, method="Jacobi"))
   >>> u # the left-hand element of the last tuple
11
12
   array([ 0.25, -0.25])
13
14
   >>> u, v = last(solve nd(A, b, method="GS"))
15 >>> u
   array([ 0.25, -0.25])
16
17
18 >>> u, v = last(solve nd(A, b, method="SOR"))
19 >>> 11
20 array([ 0.25, -0.25])
```

As we can see, all the methods give the correct solution.

SOLVING THE ORIGINAL PROBLEM

We can find a solution to our original problem, fixing n = 10 and k = 1/100, by setting up

Here, src.alg.A is an algorithm provided for us that generates A. f is the vector obtained when sampling F := f(x, y) across the lattice defined by n = 10. The sample-function is the same one we developed earlier.

Given these structures, we are interested in comparing the performance of each of the three methods when solving the system $A\mathbf{u} = \mathbf{f}$.

i. Comparing the convergence

We are interested in looking at the relative residual $\|r_n\|_2 / \|r_0\|_2$ for each iteration n of the algorithms. These can be computed with the following Python-program:

```
def right(it):
1
        # "unzips" an iterator of tuples,
2
3
        # producing only the right-hand elements.
4
        for , x in it: yield x
5
6
   def relative residual(appr):
        # we need a list, because
8
        # we can't peek iterators.
9
        assert type(appr) == list
10
        # compute the norm of the first element.
11
        a = LA.norm(appr[0][1])
12
        for r in right(iter(appr)):
13
            yield LA.norm(r) / a
```

FIGURE 7 shows the relative residuals for each of the methods, including three different values for ω in the case of successive over-relaxation.

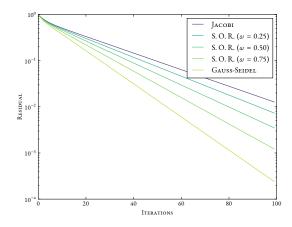


Figure 7: The residuals of each method for a given iteration.

ii. Spectral radius and convergence

Considering FIGURE 7, it is quite clear that after a certain number of iterations, the convergence is linear. In hind-sight, this is not so surprising. Consider \mathbf{u} as a linear combination of an eigen-basis $\beta = \{v_i\}$ given by the eigenvectors of C. Then the operation of C on

$$\mathbf{u} = \sum_{i} a_{i} \cdot v_{i}$$

is to scale each component of \mathbf{u} by the corresponding eigenvalue:

$$C\mathbf{u} = \sum_{i} \lambda_{i} \cdot a_{i} \cdot v_{i}$$

$$\implies C^{n}\mathbf{u} = \sum_{i} \lambda_{i}^{n} \cdot a_{i} \cdot v_{i}.$$

Since $0 < \rho(C) < 1$, each component converges by itself, so the sequence converges in its entirety, but *only as fast* as the slowest component converges! Naturally, the slowest converging component lies on the span of the eigenvector of C corresponding to $\lambda = \rho(C)$. The components in this subspace converges linearly, with $\mu = \rho(C)$.

Using Python, we can easily compute the spectral radiuses of the iteration matrices, essentially producing an upper bound for the rate of convergance for each algorithm.

Table 2: The spectral radiuses of the iteration matrices for each method

Algorithm	Spectral Radius	
Jacobi	0.9594	
Gauss-Seidel	0.9206	
S. O. R. $(\omega = 0.25)$	0.9539	
S. O. R. $(\omega = 0.50)$	0.9465	
S. O. R. ($\omega = 0.75$)	0.9362	

TABLE 3 reveals that they all have similar rates of convergence, but the successive over-relaxation method seems to get better as ω increases. This motivates a more systematic search for ω .

By searching by brute force in I=[0,100], it is revealed that $\omega \approx 1.6$ gives $\rho \approx 0.74$. This is a promising result, so we search in that general area (I=[1,2]). We find that $\omega \approx 1.5628$ gives $\rho \approx 0.7198$, which is a much faster converging approximation than any of the originals.

iii. Comparing the relative time

Now that we know (at least we think we know) the optimal choice of ω , it makes sense to benchmark all three algorithms. The solver i presented earlier does not return early based on a tolerance-criterium, because letting them all run for exactly the same number of iterations makes nicer plots and easier (to read) code. Now that we want the algorithms to compete, we don't want to naïvely run them all for the same amount of iterations. (That would be a pretty dumb benchmark) We can use the norm of the residual vector, $\|\mathbf{f} - A\mathbf{u}\|_2$ as a measure of error, and return whenever this is close to zero. Since we want to really pressure the algorithms, it makes sense to use a very large value for MAX_ITER (i used 5000) and use a low tolerance (i used $1 \cdot 10^{-10}$)

To benchmark the algorithms, i will make them solve the same problem in a long loop. This is simple, because we already have the code to setup the problem we looked at previously, but it is not a very sophisticated way to benchmark programs.

TABLE 3: The time it takes to run 1000 iterations.

Algorithm	Time (µs)	# Iterations
Jacobi	13718915	552
Gauss-Seidel	9938928	276
S. O. R. ($\omega = 1.5628$)	5126427	77

Sure enough, the Gauss-Seidel algorithm runs faster than the Jacobi method, and the successive over-relaxation algorithm with optimal ω is way faster than both of them. Since this is exactly what we expected, the benchmark is probably not too far off.

iv. Relation between k, Δx and rate of convergence

It is clear from the previous discussion that what ultimately determines the rate of convergence is the spectral radius of $C = M^{-1}N$. The wave number k is fixed for a given problem, so we can not choose it freely.

If we consider that A is deconposed into M and N, in such a way that M is lower triangular, we may notice that M, regardless of our choice of algorithm, has $4-(k\Delta x)^2$ all along the diagonal. This means that M^{-1} has spectral radius $\rho=1/\left(4-(k\Delta x)^2\right)$. Regardless of choice of algorithm.

We know that $0 < \Delta x \le 1$, since n is a positive number. It is therefore technically possible for ρ to blow up given values k > 1, which makes it possible for $(k \Delta x)^2$ to take the value 4. I will ignore this, because it is not the case in our problem. We see that the absolute lowest we can make $\rho(M^{-1})$ is 1/4, regardless the choice of algorithm, by choosing a really small Δx .

On the other hand, considering the Jacobi decomposition, i found that by making Δx really small, $\rho(N)$ ap-

proached 4. This means that if we just let $\Delta x \longrightarrow 0$, The operation of N on the space is to stretch something by a factor 4, and the operation of M^{-1} is to scale back down with a factor 1/4, essentially taking us no-where.

Recall that Δx is just how dense our lattice is; this result says intuitively that if we want a very precise solution, the iteration converges very slowly.

Some more calculations reveals why the successive over-relaxation algorithm converges so much faster compared to the Jacobi method: We essentially "move" a part of N into M, and it turns out that for a suitable choice of the relaxation coefficient ω , this *decreases* the spectral radius of N way more than it increases the spectral radius of M^{-1} . Since the spectral radius ultimately places an upper bound on the rate of convergence for an iterative procedure, this means that successive over-relaxation methods converge quickly.