Julia: From multicore to GPU parallelism

presented by Steffen Haug

WHAT IS JULIA?

(Briefly)

1. High level programming language



- 1. High level programming language
- 2. Dynamically typed



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- 2. Dynamically typed
- 3. For HPC



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- 2. Dynamically typed
- 3. For HPC
- 4. Novel compilation infrastructure (Central on CPU and GPU)

Untyped AST specialized at runtime



- 1. High level programming language
- 2. Dynamically typed
- 3. For HPC
- 4. Novel compilation infrastructure (Central on CPU and GPU)

Untyped AST specialized at runtime

"Looks like Python - runs like Fortran"



```
function square!(xs)
  for i in eachindex(xs)
    xs[i] ^= 2
  end
end

V = rand(10000);
@code_native syntax=:intel square(V)
```

Function that squares (mutates) a vector.

```
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  for i in eachindex(xs)
    xs[i] ^= 2
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```



```
square:
 mov rax, gword ptr [rdi + 8]
 test rax. rax
 je .LBB0_7
 mov rcx, gword ptr [rdi]
 nov edx. 1
 cmp rax, 16
 jb .LBB0_5
 mov rsi. rax
 and rsi, -16
 lea rdx. [rsi + 1]
 xor edi. edi
LBB0_3:
 vmovups xmm0, xmmword ptr [rcx + 8*rdi]
 vmovups xmm1, xmmword ptr [rcx + 8*rdi + 32]
 vmovups xmm2, xmmword ptr [rcx + 8*rdi + 64]
 vmovups xmm3, xmmword ptr [rcx + 8*rdi + 96]
 vinsertf128 ymm8, ymm8, xmmword ptr [rcx + 8*rdi + 16], 1
 vinsertf128 vmm1, vmm1, xmmword ptr [rcx + 8*rdi + 48], 1
 vinsertf128 ymm2, ymm2, xmmword ptr [rcx + 8*rdi + 80], 1
 vinsertf128 ymm3, ymm3, xmmword ptr [rcx + 8*rdi + 112], 1
 vmulpd vmm0, vmm0, vmm0
 vmulpd ymm1, ymm1, ymm1
 vmulpd vmm2, vmm2, vmm2
 vmulpd ymm3, ymm3, ymm3
 vextractf128 xmmword ptr [rcx + 8*rdi + 16], ymm0, 1
 vmovupd xmmword ptr [rcx + 8*rdil. xmm8
 vextractf128 xmmword ptr [rcx + 8*rdi + 48], ymm1, 1
 vmovupd xmmword ptr [rcx + 8*rdi + 32], xmm1
 vextractf128 xmmword ptr [rcx + 8*rdi + 88], vmm2, 1
 vmovupd xmmword ptr [rcx + 8*rdi + 64], xmm2
 vextractf128 xmmword ptr [rcx + 8*rdi + 112], vmm3, 1
 vmovupd xmmword ptr [rcx + 8*rdi + 96], xmm3
 add rdi, 16
 cmp rsi. rdi
 ine IRRR 3
 cmp rax, rsi
 ie .LBB0 7
IRRA 5-
 dec rdx
LBB0 6:
 vmovsd xmm0, qword ptr [rcx + 8*rdx]
 vmulsd xmm0, xmm0, xmm0
 vmovsd qword ptr [rcx + 8*rdx], xmm0
 inc rdx
 cmp rdx, rax
 jb .LBB0_6
LBB0_7:
 vzeroupper
 ret
```



Remove some

```
.LBB0 3:
 vmovuns xmm0, xmmword ntr [rcx + 8*rdi]
 vmovups xmm1, xmmword ptr [rcx + 8*rdi + 32]
 vmovups xmm2, xmmword ptr [rcx + 8*rdi + 64]
 vmovuns xmm3, xmmword ntr [rcx + 8*rdi + 96]
 vinsertf128 vmm0, vmm0, xmmword ptr [rcx + 8*rdi + 16], 1
 vinsertf128 ymm1, ymm1, xmmword ptr [rcx + 8*rdi + 48], 1
 vinsertf128 vmm2, vmm2, xmmword ptr [rcx + 8*rdi + 80], 1
 vinsertf128 ymm3, ymm3, xmmword ptr [rcx + 8*rdi + 112], 1
 vmulpd ymm0, ymm0, ymm0
 vmulpd vmm1, vmm1, vmm1
 vmulpd ymm2, ymm2, ymm2
 vmulpd vmm3, vmm3, vmm3
 vextractf128 xmmword ptr [rcx + 8*rdi + 16], ymm0, 1
 vmovupd xmmword ptr [rcx + 8*rdi], xmm0
 vextractf128 xmmword ptr [rcx + 8*rdi + 48], ymm1, 1
 vmovupd xmmword ptr [rcx + 8*rdi + 32], xmm1
 vextractf128 xmmword ptr [rcx + 8*rdi + 80], ymm2, 1
 vmovund xmmword ntr [rcx + 8*rdi + 64], xmm2
 vextractf128 xmmword ptr [rcx + 8*rdi + 112], vmm3, 1
 vmovupd xmmword ptr [rcx + 8*rdi + 96], xmm3
 add rdi. 16
 cmp rsi, rdi
 ine .LBB0 3
```

Julia has unrolled the loop and used SIMD-instructions!

```
void square(double *xs, int N) {
 for (int i = 0; i < N; i++) {
   xs[i] *= xs[i];
.L1:
  vmovupd vmm0.YMMWORD PTR [rdi+rsi*1]
  vmovupd ymm1,YMMWORD PTR [rdi+rsi*1+0x20]
  vmovupd vmm2.YMMWORD PTR [rdi+rsi*1+0x40]
  vmovupd ymm3,YMMWORD PTR [rdi+rsi*1+0x60]
  vmulpd ymm0,ymm0,ymm0
  vmulpd vmm1.vmm1.vmm1
  vmulpd ymm2,ymm2,ymm2
  vmulpd vmm3, vmm3, vmm3
  vmovupd YMMWORD PTR [rdi+rsi*1],ymm0
  vmovupd YMMWORD PTR [rdi+rsi*1+0x20].vmm1
  vmovupd YMMWORD PTR [rdi+rsi*1+0x40],ymm2
  vmovupd YMMWORD PTR [rdi+rsi*1+0x60],ymm3
        rsi.0xffffffffffff80
  sub
        rdx,rsi
  cmp
  ine
        .L1
```

This is what Clang does.

(Slightly better due to Julia-arrays being strided)

Julia uses LLVM



Julia can do the same optimizations we expect in C/C++

Julia uses LLVM



Julia can do the same optimizations we expect in C/C++

 $... because \ it \ is \ literally \ the \ same \ optimizer!$

But there are some <u>huge</u> caveats

1. There is some *slight* compilation overhead at startup

However: As problem size grows, relative compilation time decreases.

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- 2. Julia is garbage-collected
- 3. We have no explicit control over memory (Allocation, copying, ...)

Accidental heap operations can really hurt!

But there are some <u>huge</u> caveats

- 1. There is some *slight* compilation overhead at startup
- 2. Julia is garbage-collected
- 3. We have no explicit control over memory (Allocation, copying, ...)
- 4. The whole thing breaks if Julia can't figure out the types

There are many ways to shoot yourself in the foot! \cdots

```
Δx = 0.1

function riemann!(X, x)

N = Length(x)

X[0] = 0.0

for i in 2:N

X[i] = X[i - 1] + Δx * X[i]

end

v = rand(100000)

V = zeros(100000)

Qtime riemann!(V, v)
```

```
const \( \Delta \times = 0.1 \)

function riemann!(X, x)

N = length(x)

X[0] = 0.0

for i in 2:N

X[i] = X[i - 1] + \( \Delta \times X[i] \)

end

end

v = rand(100000)

V = zeros(100000)

@time riemann!(V, v)
```

Two seemingly quite similar programs... $\,$

```
Δx = 0.1

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    for i in 2:N
        X[i] = X[i - 1] + \( \Delta \times \times X[i] \)
    end
end

v = rand(100000)
v = zeros(100000)
0time riemann!(V, v)
```

0.000175 seconds

One is $\sim 60 \times$ faster than the other!

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    end
end

v = rand(100000)
v = zeros(100000)
@time riemann!(V, V)
```

0.000175 seconds

0.010550 seconds (499.49 k allocations)

And 500K allocations? 🤯

```
Δx = 0.1

function riemann!(X, x)

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end

v = rand(100000)

V = zeros(100000)

Qtime riemann!(V, v)
```

```
const \( \Delta x = 0.1 \)
function riemann!(X, \( x \))
    N = length(x)
    X[0] = 0.0
    for i in 2:N
         X[i] = X[i - 1] + \( \Delta x \) X[i]
    end
end

v = rand(100000)
v = zeros(1000000)
@time riemann!(V, \( v \))
```

0.000175 seconds

And 500K allocations? 🔯

 Δx could be reassigned \implies Julia can't assume its type!

```
Δx = 0.1

function riemann!(X, x)

N = length(x)

X[0] = 0.0

for i in 2:N

X[i] = X[i - 1] + Δx * X[i]

end

end

v = rand(100000)

V = zeros(100000)

@time riemann!(V, v)
```

```
const \( \Delta x = 0.1 \)
function riemann!(X, X)
    N = length(X)
    X[e] = 0.0
    for in 2:N
        X[i] = X[i - 1] + \( \Delta x * X[i] \)
    end
end

v = rand(100000)
v = zeros(1000000)
@time riemann!(V, v)
```

0.000175 seconds

And 500K allocations? 🔯

Leads to runtime type checks and dynamic dispatching...

```
Δx = 0.1

function riemann!(X, x)

N = length(x)

X[0] = 0.0

for i in 2:N

X[i] = X[i - 1] + Δx * X[i]

end

end

v = rand(100000)

V = zeros(100000)

@time riemann!(V, v)
```

0.000175 seconds

And 500K allocations? 🔯

 \ldots and forces intermediate calculations onto the heap!

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    for i in 2:N
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end
v = rand(100000)
V = zeros(100000)
@time riemann!(V, v)
```

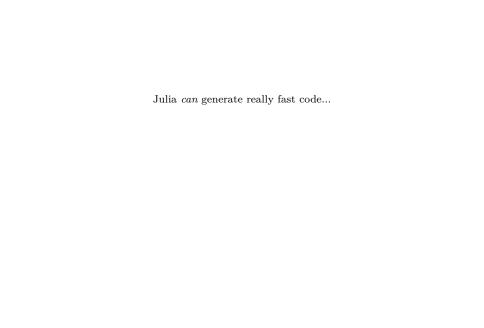
0.000175 seconds

And 500K allocations?



...and forces intermediate calculations onto the heap!







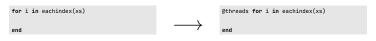
Julia can generate really fast code...

But the *lack of explicit control* requires a lot of attention.

Even more so when we parallelize our programs.

Multicore

using .Threads: @threads



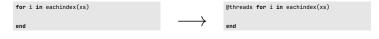
Similar API to OpenMP: Just annotate loops with a macro.



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The normal caveats apply

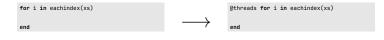
1. Spawning threads involve some overhead



Similar API to OpenMP: Just annotate loops with a macro.

The normal caveats apply

- 1. Spawning threads involve some overhead
- 2. Memory is shared (Race conditions, deadlocks)



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The normal caveats apply

- 1. Spawning threads involve some overhead
- 2. Memory is shared (Race conditions, deadlocks)
- 3. Possible speedup severely limited by core count



Similar API to OpenMP: Just annotate loops with a macro.

THE NORMAL CAVEATS APPLY

- 1. Spawning threads involve some overhead
- 2. Memory is shared (Race conditions, deadlocks)
- 3. Possible speedup severely limited by core count
- 4. Domain decomposition sometimes necessary to get speedup on simple problems

GPU PARALLELISM



LEVEL 1: "VECTORIZED" ARRAY API

(Vectorized in the sense that you only operate on arrays, Numpy style)

```
# In-place operations.
v = CUDA.rand(10000000)
v *= 10.0
v ^= 2
v •= sqrt.(v)

# Parallel scan.
w = CUDA.zeros(10000000)
accumulate!(max, w, v)
```

v, w of type CuArray that manages a device-side buffer.

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v •= sqrt.(v)

# Parallel scan.
w = CUDA.zeros(10000000)
accumulate!(max, w, v)
```

Block- and gridsizes automatically determined.

```
# FFT (10M-point FFT takes 200ms on 980)
using CUDA.CUFFT
\( \omega = \text{fft}(v) \)

# Kernels get fused
\( \omega = \omega \cdot \cdot
```

Bindings to CUDA libraries like cuFFT and cuBLAS (Also cuSPARSE, cuSOLVER, ...)

```
# FFT (10M-point FFT takes 200ms on 980)
using CUDA.CUFFT
\omega = fft(v)

# Kernels get fused
w •= w.^2 .+ v

# Check out what happens on the GPU
@device_code_sass w •= w.^2 .+ v
```

GPU code generated on the fly using LLVMs PTX/SASS backend... (Same sort of infrastructure as specializing Julia functions)

```
# FFT (10M-point FFT takes 200ms on 980)
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ω = fft(v)

# Kernels get fused
w •= w.^2 .+ v

# Check out what happens on the GPU
@device_code_sass w •= w.^2 .+ v
```

Enables kernel fusion and global optimization (!)

```
# FFT (10M-point FFT takes 200ms on 980)
using CUDA.CUFFT

\omega = fft(v)

# Kernels get fused
w \cdot = w.^2 .+ v

# Check out what happens on the GPU
@device_code_sass w \cdot = w.^2 .+ v
```

Like with CPU code, you can inspect the generated GPU code. (SASS/PTX code is too long/ugly to put on a slide)

LEVEL 2: JULIA ON THE GPU

```
function cusquare!(xs)
    i = (blockIdx().x - 1) * blockDim().x + threadIdx().x
    if i < length(xs)
        xs[i] *= xs[i]
    end
    return nothing
end

Nt = 1024
Nb = ceil(Int, length(v) / Nt)
@cuda threads=Nt blocks=Nb cusquare!(v)</pre>
```

```
function cusquare!(xs)
    i = (blockIdx().x - 1) * blockDim().x + threadIdx().x
    if i < length(xs)
        xs[i] *= xs[i]
    end
    return nothing
end

Nt = 1024
Nb = ceil(Int, length(v) / Nt)
@cuda threads=Nt blocks=Nb cusquare!(v)</pre>
```

The @cuda macro transforms typed Julia IR into GPU code!



```
function cusquare!(xs)
    i = (blockIdx().x - 1) * blockDim().x + threadIdx().x
    if i < length(xs)
        xs[i] *= xs[i]
    end
    return nothing
end

Nt = 1024
Nb = ceil(Int, length(v) / Nt)
@cuda threads=Nt blocks=Nb cusquare!(v)</pre>
```

Basically, almost any sensible Julia can just work on the GPU

```
function cusquare!(xs)
    i = (blockIdx().x - 1) * blockDim().x + threadIdx().x
    if i < length(xs)
        xs[i] *= xs[i]
    end
    return nothing
end

Nt = 1024
Nb = ceil(Int, length(v) / Nt)
@cuda threads=Nt blocks=Nb cusquare!(v)</pre>
```

However, there is no escape hatch: The code must type check, and can't allocate

PITFALLS

PITFALLS

CuArray controls all memory transfer

PITFALLS

Accidentally reading back to the CPU is really painful!

Future Outlook

LEVEL 3: MULTI-GPU / MULTI-NODE

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CUDA.jl composes with DistributedArrays.jl

DistributedArrays.jl: Domain decomposition and distribution algorithms

CUDA.jl: Manages device-side buffers and host-device communication

Both understands Julias abstract concept of an "Array", so they work together. *

https://www.sciencedirect.com/science/article/pii/S0965997818310123, (Section 6.3.2. Distributed GPU Arrays).

^{*} Tim Besard et al. "Rapid software prototyping for heterogeneous and distributed platforms". In: Advances in Engineering Software (2019). URL:

Julia makes bold claims

Julia makes bold claims: "Looks like Python – runs like C/C++/Fortran"

Julia makes bold claims: "Only dynamic language to achieve petascale"

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It is a truth with a big asterisk.

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If there is one thing that is true, it is that you can't escape

LOCATION, LOCATION

Julia makes bold claims: "Only dynamic language to achieve petascale"

If there is one thing that is true, it is that you can't escape

LOCATION, LOCATION

Julia has real potential, but it is not a free lunch! (Especially on the GPU)

Julia makes bold claims: "Only dynamic language to achieve petascale"

If there is one thing that is true, it is that you can't escape

LOCATION, LOCATION

I spend a lot of time "fighting the abstraction"

Julia makes bold claims: "Only dynamic language to achieve petascale"

If there is one thing that is true, it is that you can't escape

LOCATION, LOCATION

Allocation, garbage-collection, copies, CUDA device $\leftrightarrow \! host$ transfers, etc. are all implicit!

Julia makes bold claims: "Only dynamic language to achieve petascale"

If there is one thing that is true, it is that you can't escape

LOCATION, LOCATION

Allocation, garbage-collection, copies, CUDA device $\leftrightarrow \! host$ transfers, etc. are all implicit!

As Julia takes away your control, you have to be careful.

Julia makes bold claims: "Only dynamic language to achieve petascale"

If there is one thing that is true, it is that you can't escape

LOCATION, LOCATION

Allocation, garbage-collection, copies, CUDA device $\leftrightarrow \! host$ transfers, etc. are all implicit!

The abstraction leaks!

What niche does Julia even fill?

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If you don't understand the compilation process, Julia is basically unusable.

What niche does Julia even fill?

 ${\it Julia\ wont't\ replace\ lower-level\ languages}.$

That being said...

I quite like Julia

What niche does Julia even fill?

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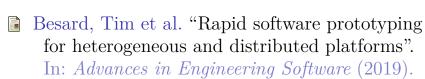
Julia makes HPC a $\mathit{first\text{-}class}$ $\mathit{citizen}$ in the "Matlab workflow"

What niche does Julia even fill?

Julia makes HPC a first-class citizen in the "Matlab workflow"

- Interactive REPL/Notebook-based workflow
- 2. Good for fucking around with data
- 3. Sensible package manager
- 4. "Math stuff" is easily accessible
- 5. Visualization is made easy





URL: https://www.sciencedirect.com/science/

article/pii/S0965997818310123.