

Embedded Systems - Real-Time Scheduling

part 1

Timing constraints: origin and characterization

**Basic concepts of *real-time*
Real-Time Systems requirements**

“Real-Time” Definitions

There is a large variety of definitions related with **Real-time**, real-time **systems**, real-time **services** and **real-time functionality**.

They all express the **dependency of a computing system** with respect to **time** as imposed by **a given external process, either physical or logical**.

Real-Time Computing

The **results of the computations** must be

- Logically correct
- Produced in time

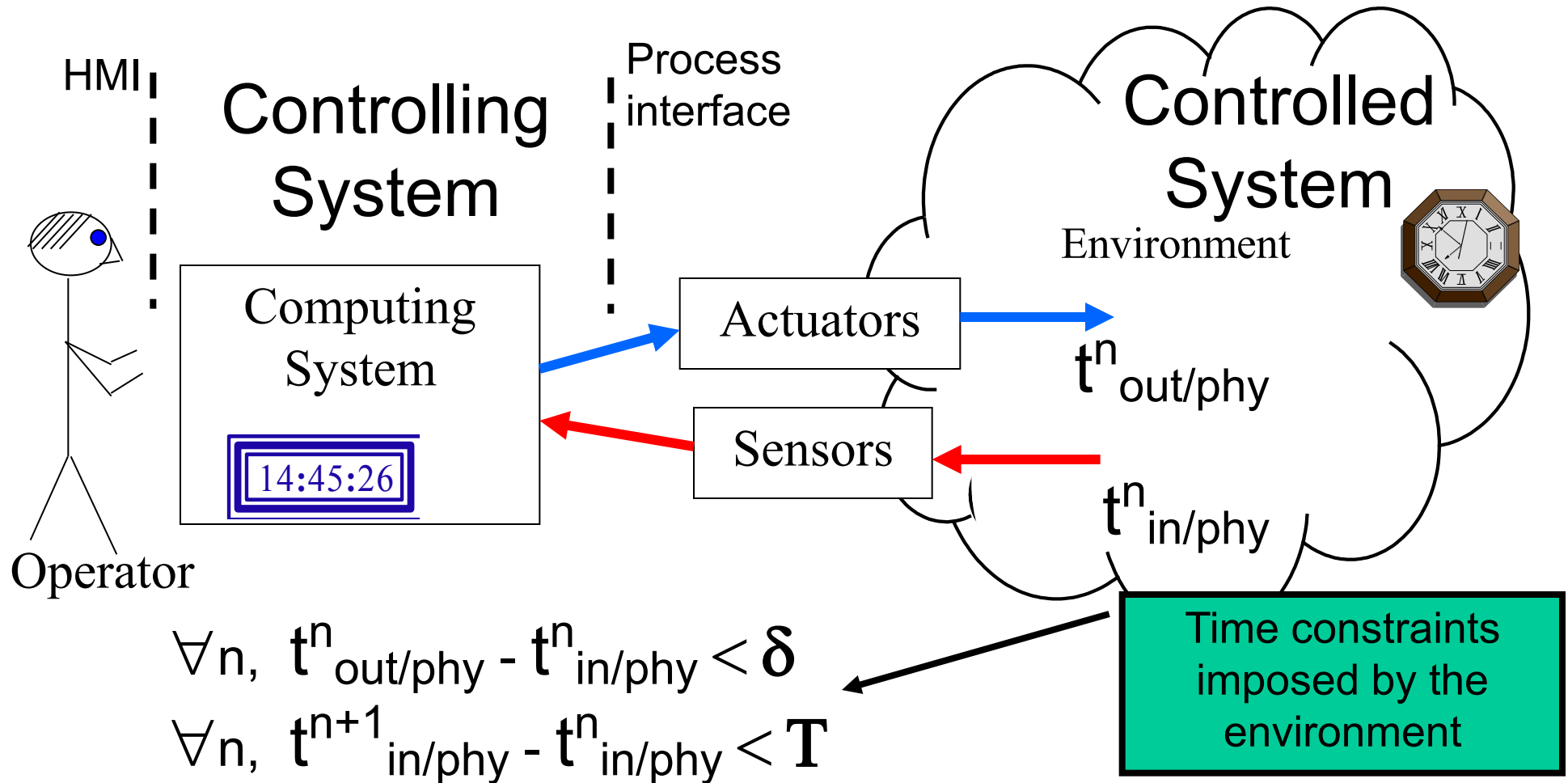
(Stankovic, 1988)

Timeliness



**Logical
correctness**

Real-Time System



Real-Time System

- Generally, when a computing system is able to **keep the pace** of a given process (physical or logical) and, if needed, **influence it in a desired way**, then it is a **real-time system**.
- All living beings are *real-time systems* with respect to their natural habitats, which establish their real-time.
- Similarly, when building intelligent machines (programmable) to **interact with a given process**, we need to use programming techniques and SW infrastructures that allow having confidence on their capacity to actuate in time.

Purpose of studying Real-Time Systems

- The main **objective** of a Real-Time Systems (RTS) discipline is to develop **appropriate techniques** for the
 - **Design / synthesis,**
 - **Analysis and**
 - **Verification**of such systems that allow obtaining **guarantees** that their **temporal behavior is adequate to the dynamics** of the process they must interact with.

Special features of Real-Time Systems

Given their **dependence with time**, some of the most important aspects that characterize a real-time system are:

- **Execution time of its computations**
- **Response time to events**
- **Regularity of generating periodic events**

Special features of Real-Time Systems

These aspects are influenced by several internal parameters:

- Execution time
 - Code structure (language, conditionals, cycles)
 - Operating system or kernel (system calls)
 - DMA, caches, instructions pipeline
- Response time and regularity
 - Multi-tasking
 - Access to shared resources (buses, disks, comm ports,...)
 - HW Interrupts

Requirements of Real-Time Systems

Typical **requirements** imposed on real-time systems:

- **Functional**

- **Temporal**

- **Dependability**

- **...**

Non-Functional

Functional requirements

Example

Acquiring environment data

- Sample process variables (**real-time entities**) either continuous or discrete

Direct Digital Control (DDC)

- Realize low-level feedback control (sensor/actuators loops)

Human-Machine Interface (HMI)

- Provide information on the system state, carry out logging, perform configuration

Functional requirements

Acquiring environment data

The **process real-time entities**, upon acquisition, are internally accessible to the computing system through **local images** (internal variables).

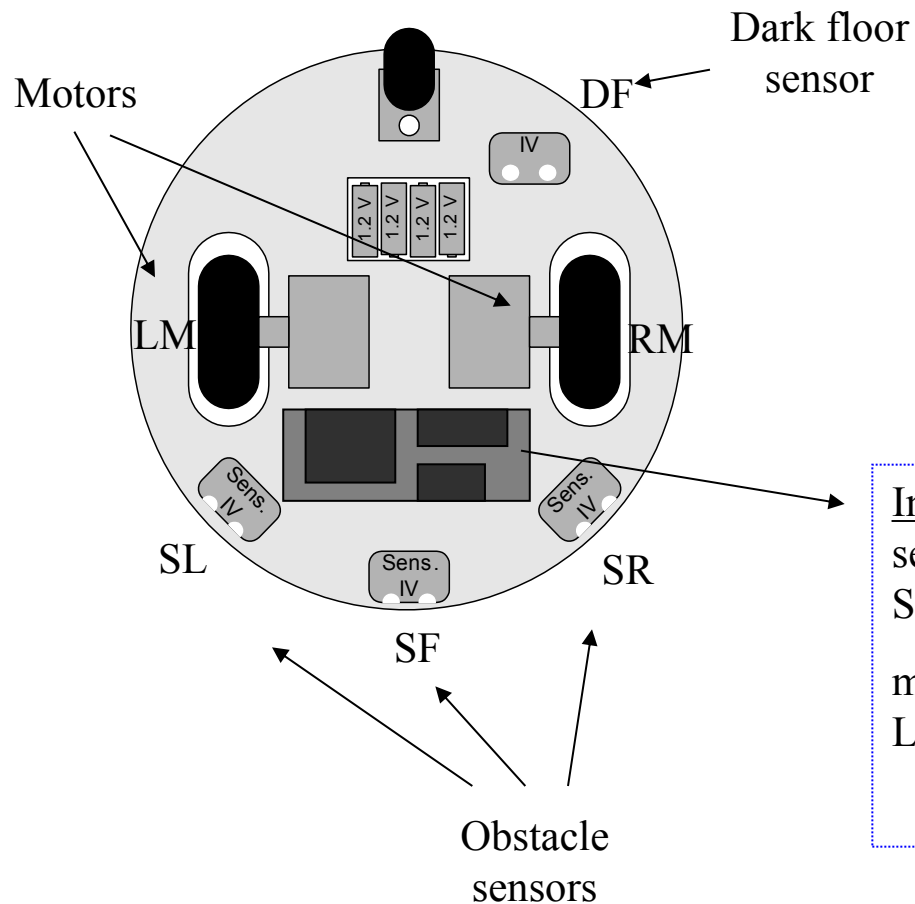
Each **local image** of a real-time entity has a **limited temporal validity** due to the process dynamics.

The **set of local images** of the real-time entities forms the **real-time database**.

The **real-time database** needs to be **updated** whenever there is a **change** in the state/value of a real-time entity.

Functional requirements

Example: A simple mobile robot



RT entities:

sensors:

SL, SF, SR, DF

Motors speeds:

LM, RM

Internal images:

sensors:

SL', SF', SR', DF'

motors speeds:

LM', RM'

RT database

Temporal requirements

Arise from the **dynamics of the process** to be controlled/monitored

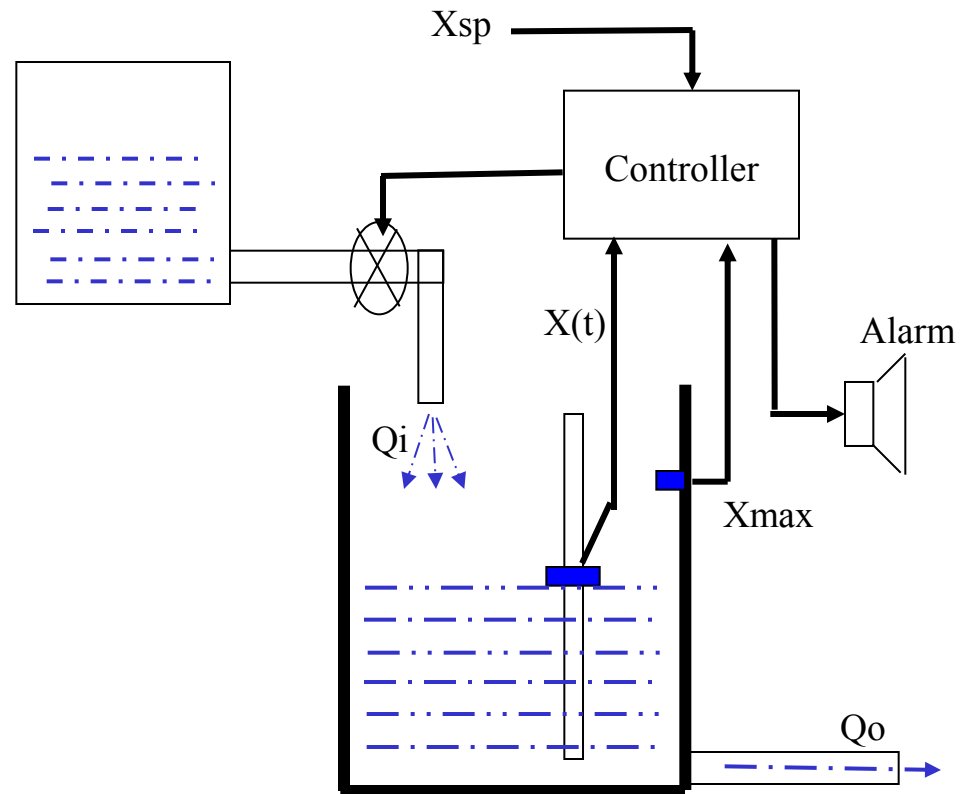
Sets **constraints**:

- To the **observation delays** of the system state
- To the **computing delays** of the control/actuation values
- To the **delay variations** of the previous (*jitter*)

In some cases, such constraints must be met in
all instances (including the worst-case)
and **not only on average terms**.

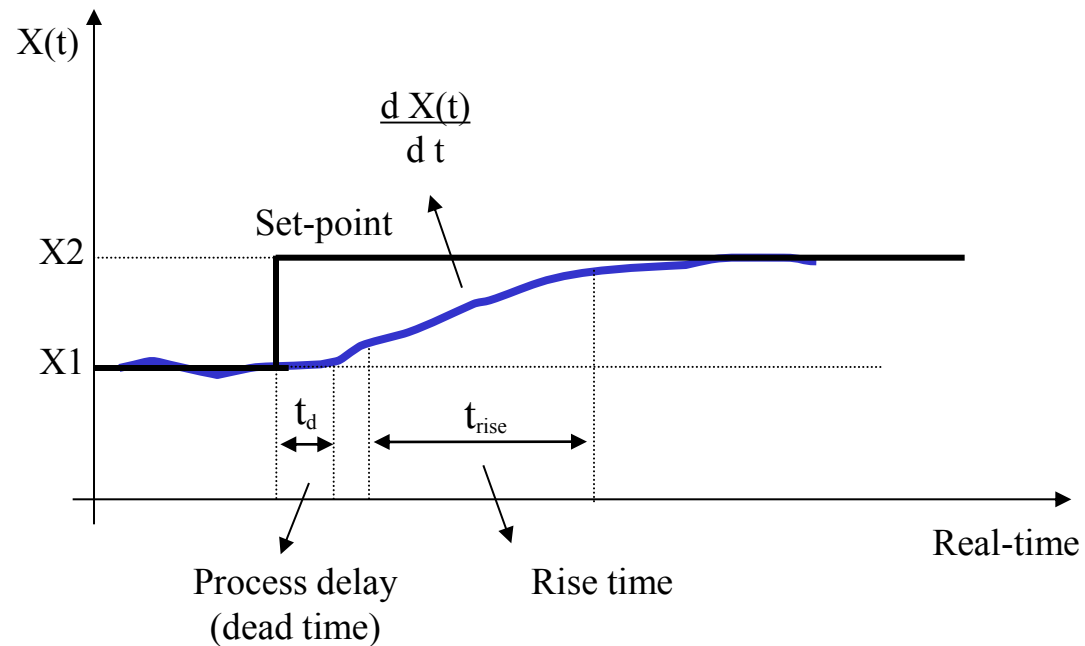
Temporal requirements

Control of the level of a liquid in a tank



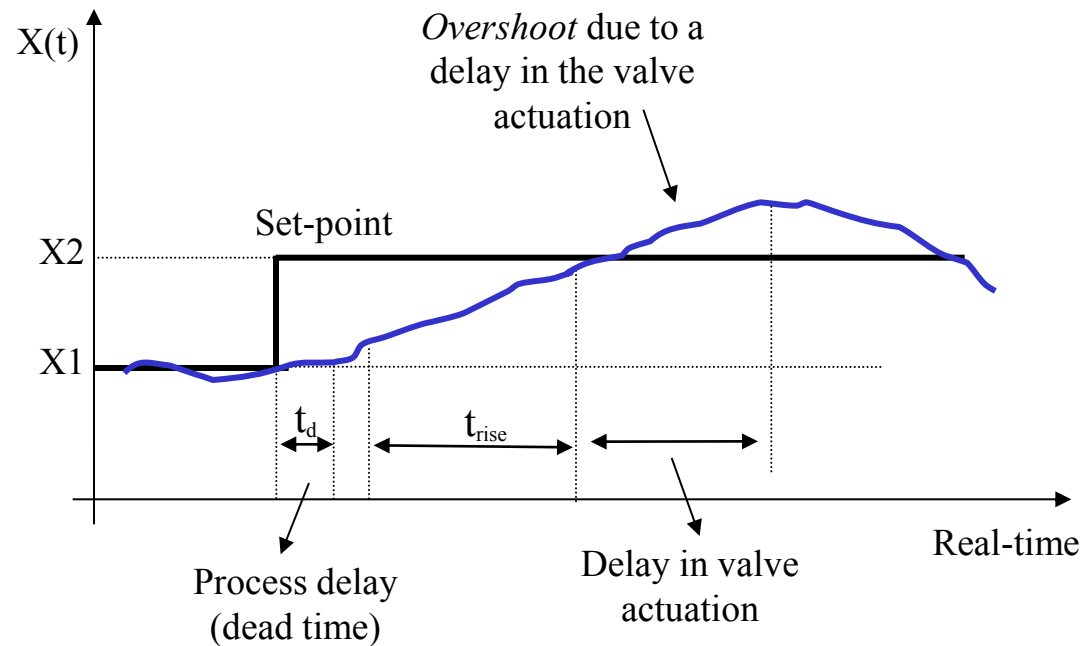
Temporal requirements

Control of the level of a liquid in a tank



Temporal requirements

Delayed actuation –control degradation



Temporal requirements

(Control systems)

Sampling period – T_s ($\sim < 1/10 t_{\text{rise}}$ - *quasi-continuous control*)

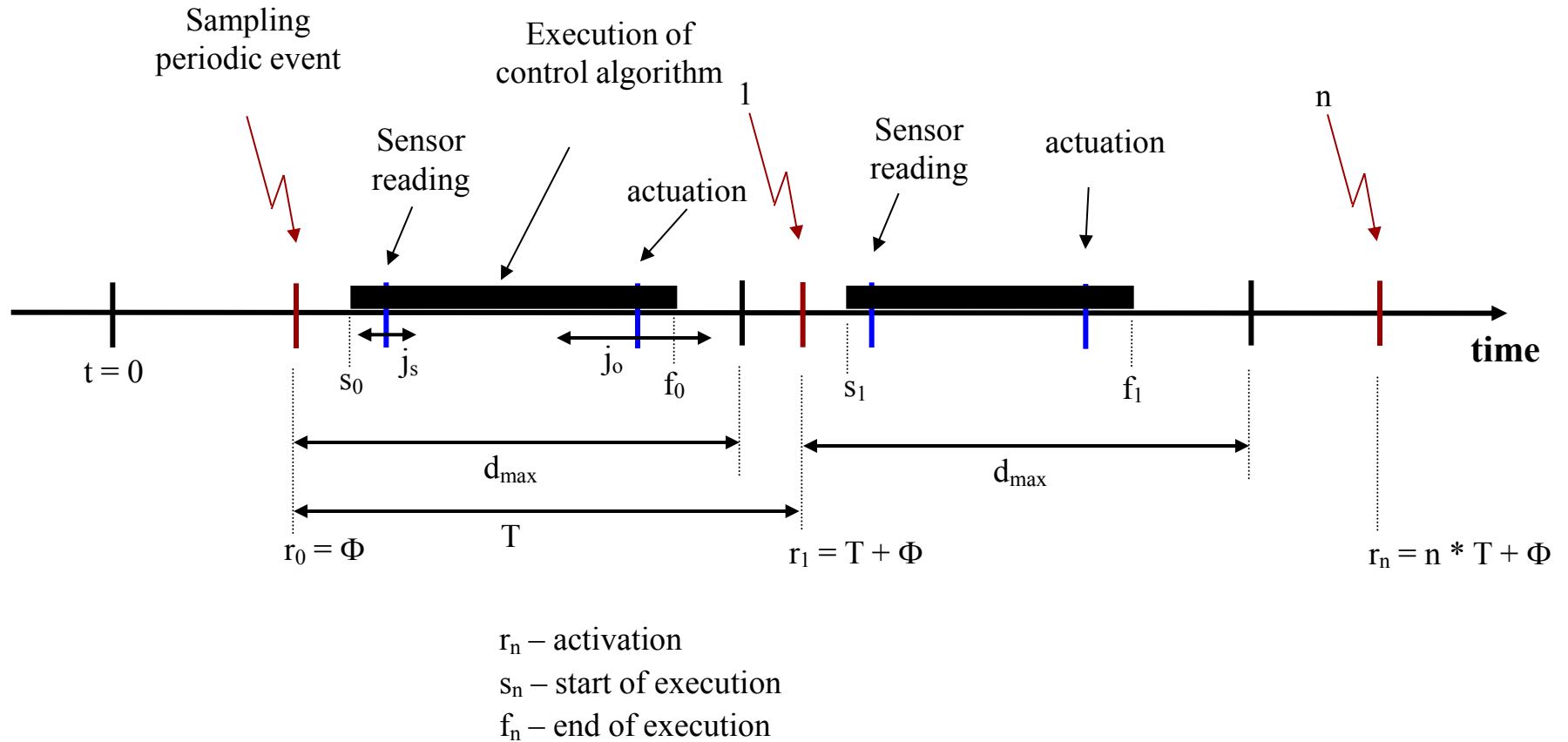
Maximum delay in the valve response – d_{max} ($< T_s$)
(control delay – can be compensated in many cases with a model)

Variations in the level reading (*jitter*) – $j_{s,\text{max}}$ ($\ll d_{\text{max}}$)

Variations in the valve actuation delay (*jitter*) – $j_{o,\text{max}}$ ($\ll d_{\text{max}}$)
(harder to compensate – degradation of the quality of the control)

Maximum delay in alarm signaling – $d_{\text{al,max}}$

Temporal requirements

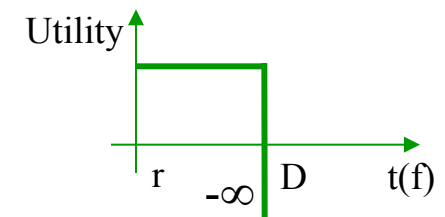
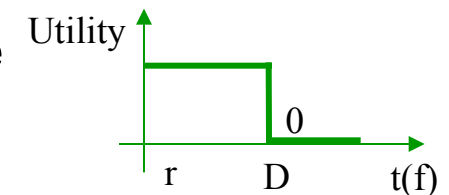
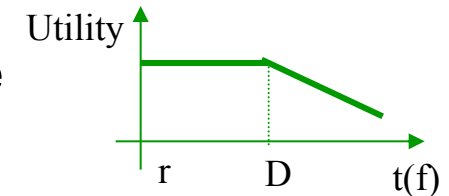


Temporal requirements

Classification of real-time constraints:

(according to the utility of the result to the application)

- **Soft** – Time constraint according to which the associated value keeps **some utility** to the application even after a deadline D , despite a possible degradation of the quality of service (videophone)
- **Firm** – Time constraint according to which the associated value loses **any utility** to the application after a deadline D (pre-paid services validation)
- **Hard** – Time constraint that, when not met, may generate a **catastrophic failure** (critical control)



Temporal requirements

Classification of Real-Time Systems:

(according to the type of time constraints they have to cope with)

- **Soft Real-Time** – There are *firm* or *soft* time constraints, only (e.g., simulators, multimedia systems)
- **Hard Real-Time** – There is at least one *hard* time constraint. These are **safety-critical** systems (e.g. transportation systems, weapons control, process control of dangerous materials...)

Dependability Requirements

Real-Time Systems are frequently used in **critical applications** either in terms of **safety** and **economy** (e.g., power plants, traffic control (air, rail, road), process industry...).

This leads to a specific requirement for:

High reliability – In *hard real-time* systems one typically finds **ultra high reliability** requirements ($\lambda < 10^{-9}$ failures/hour – cannot be verified experimentally! 1 failure in 114155 years!).

Dependability Requirements

Important aspects concerning **critical safety**:

- **Stable interfaces** between the critical and non-critical subsystems in order to prevent **error propagation**.
- **Well defined worst-case scenarios**. The system must possess the **adequate resources** to withstand a worst-case operational scenario without need for probabilistic arguments, i.e., it should provide **service guarantees** even in such cases.
- **Architecture** composed by **autonomous subsystems**, whose properties can be verified independently from the other subsystems (composability).

Wrapping up

- Notion of **real-time** and **real-time system**
- Purpose of studying RTS – obtain **guarantees** of **adequate temporal behavior**
- Aspects to consider: **execution time**, **response time**, **regularity** of generating/handling periodic events
- RTS requirements: **functional**, **temporal** and **dependability**
- Notion of **real-time database**
- **Soft**, **firm** and **hard** time constraints
- **Hard real time** vs **soft real time**
- Importance of the **worst-case scenario**

Embedded Systems - Real-Time Scheduling

part 2

Computing models

**Models of tasks with explicit time constraints
Logic and temporal control (by events -ET and by time -TT)**

Computing models

Transformational model

- A program starts and finishes, transforming input data in output results.

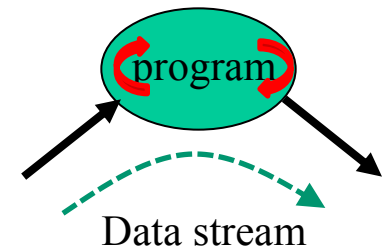
Reactive model

- A program executes indefinitely a sequence of interactions, for example operating over a continuous input data stream.



Real-time model

- **Reactive model** in which the output stream has to be synchronized with the input data stream, thus imposing time constraints on the program execution.



Determinism versus Predictability

(within reactive systems)

Determinism

- Feeding a program with the **same sequence of data inputs** originates the **same sequence of data outputs**.
(logical property)

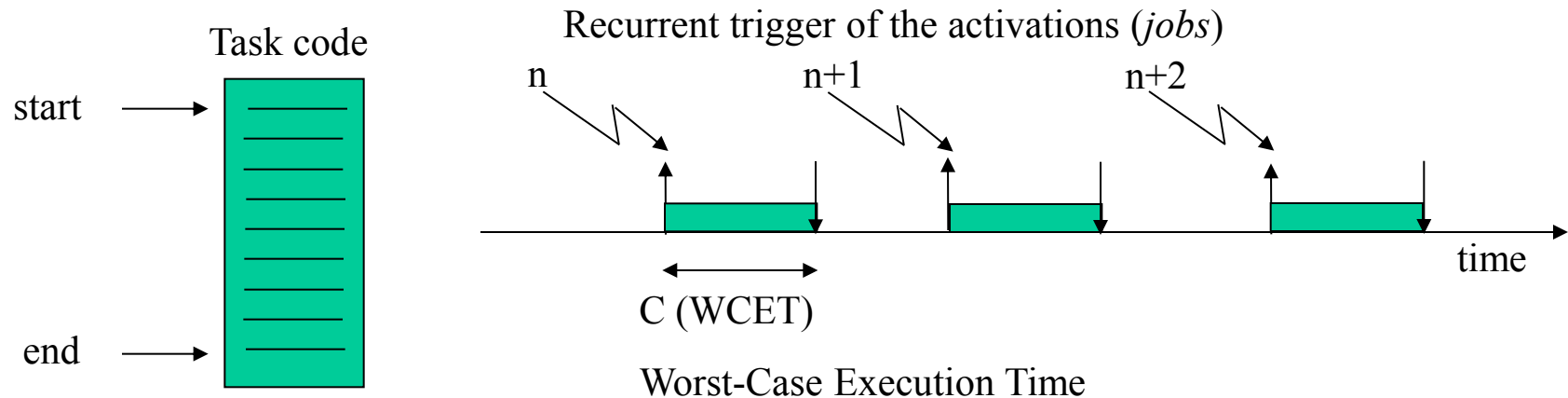
Predictability

- feeding a program with the **same input sequence** will generate the **same output sequence** with a known **delay** or within a know **time window**.
(temporal determinism – logical and temporal property)

Real-Time model

Definition of **task** (recurrent)

Infinite sequence of activations (instances or jobs), each executing a set of operations (a **given function**).

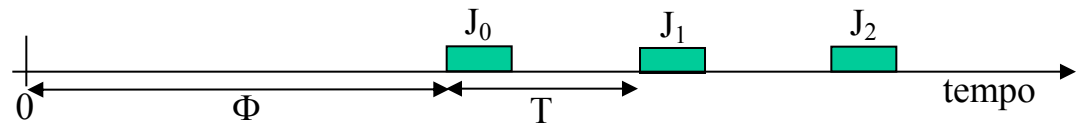


Real-Time model

Concerning recurrency of arrivals, a task can be

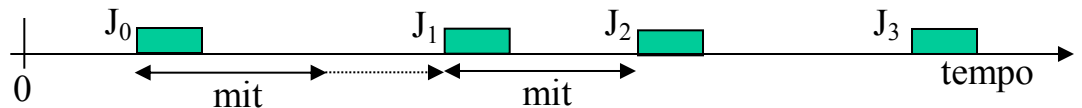
- **periodic**

Instance n activated at $a_n = n * T + \Phi$



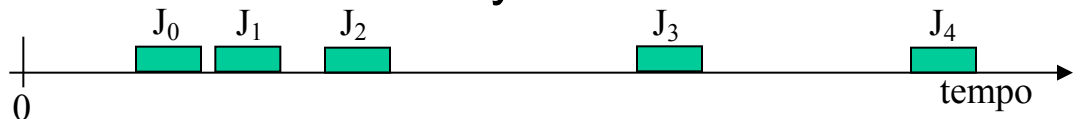
- **sporadic**

Minimal time between consecutive activations (minimum inter-arrival time – *mit*)



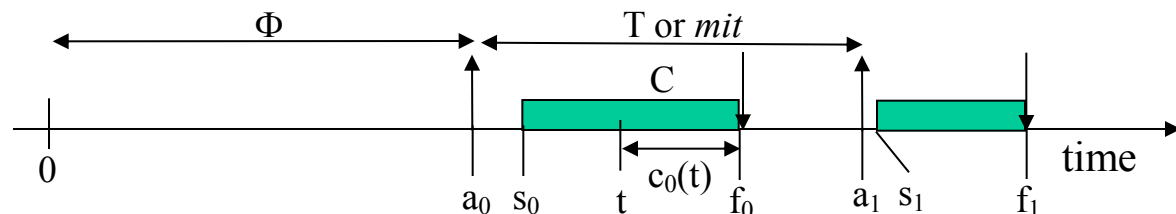
- **aperiodic**

Can only be characterized stochastically



Real-Time model

Characterization of tasks



C – Worst-case execution time (WCET)

T – period (periodic)

Φ –relative phase/offset = instant of 1st activation (periodic)

mit – *minimum inter-arrival time* (sporadic)

a_n – activation instant of the n th instance

s_n – start of execution of the n th instance

f_n – finishing time of the n th instance

$c_n(t)$ – maximum remaining execution time at instant t of the n th instance

Real-Time model

Tasks requirements:

- **Temporal** – time constraints on the **finishing instants** or for generating certain **output events**.
- **Precedence** – establish a certain **order of execution** among tasks.
- **Resource usage** – need for using **shared resources** (e.g. communication ports, buffers in shared memory, global variables, system peripherals...). May imply the use of **atomic operations** (which execution sequence cannot be interrupted)

Real-Time model

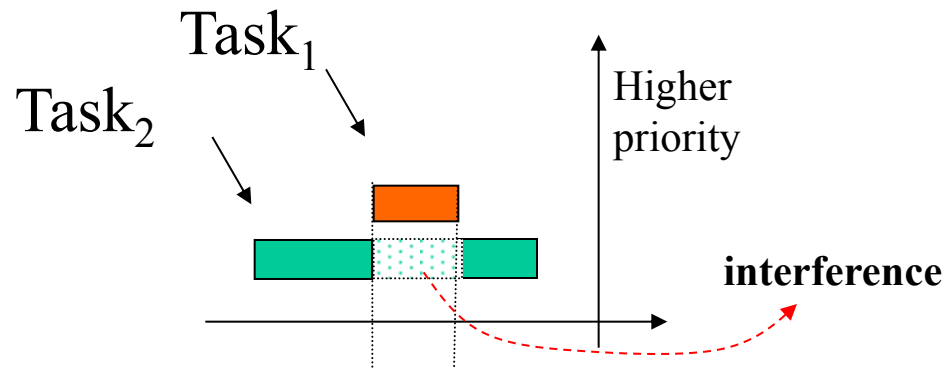
Preemption

- When a task can be **temporarily suspended** for the execution of another one with **higher priority**, it is **preemptable**.
- When a system makes use of the preemptability of tasks is it called **preemptive**.
- A set of tasks is said to admit **full preemption** when all its tasks can be preempted in any point of their execution (independent tasks)

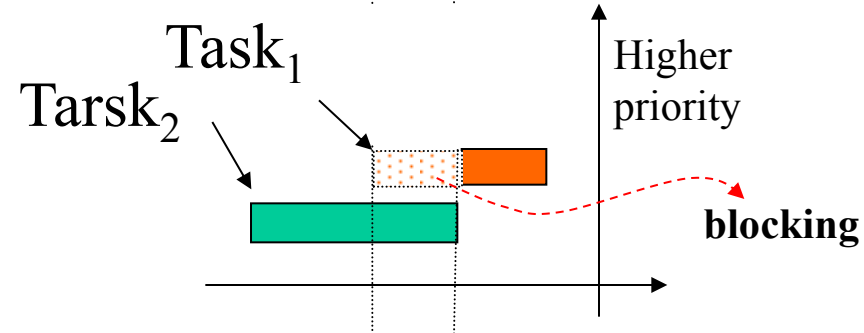
Note: the access to **shared resources** (thus tasks with **dependencies**) may impose restrictions on the level of preemptability of a task.

Real-Time model

- **Execution with preemption**



- **Execution without preemption**



Real-Time model

The **temporal requirements** can be of diverse types:

- **Deadline** – Upper bound to the finish instant of a task.
- **Window** – Lower and upper bounds to the finish instant of a task.
- **Synchronism** – Upper bound to the interval between two output events (there are types, too).
- **Distance** – Upper bound to the interval between the finishing time of two consecutive jobs

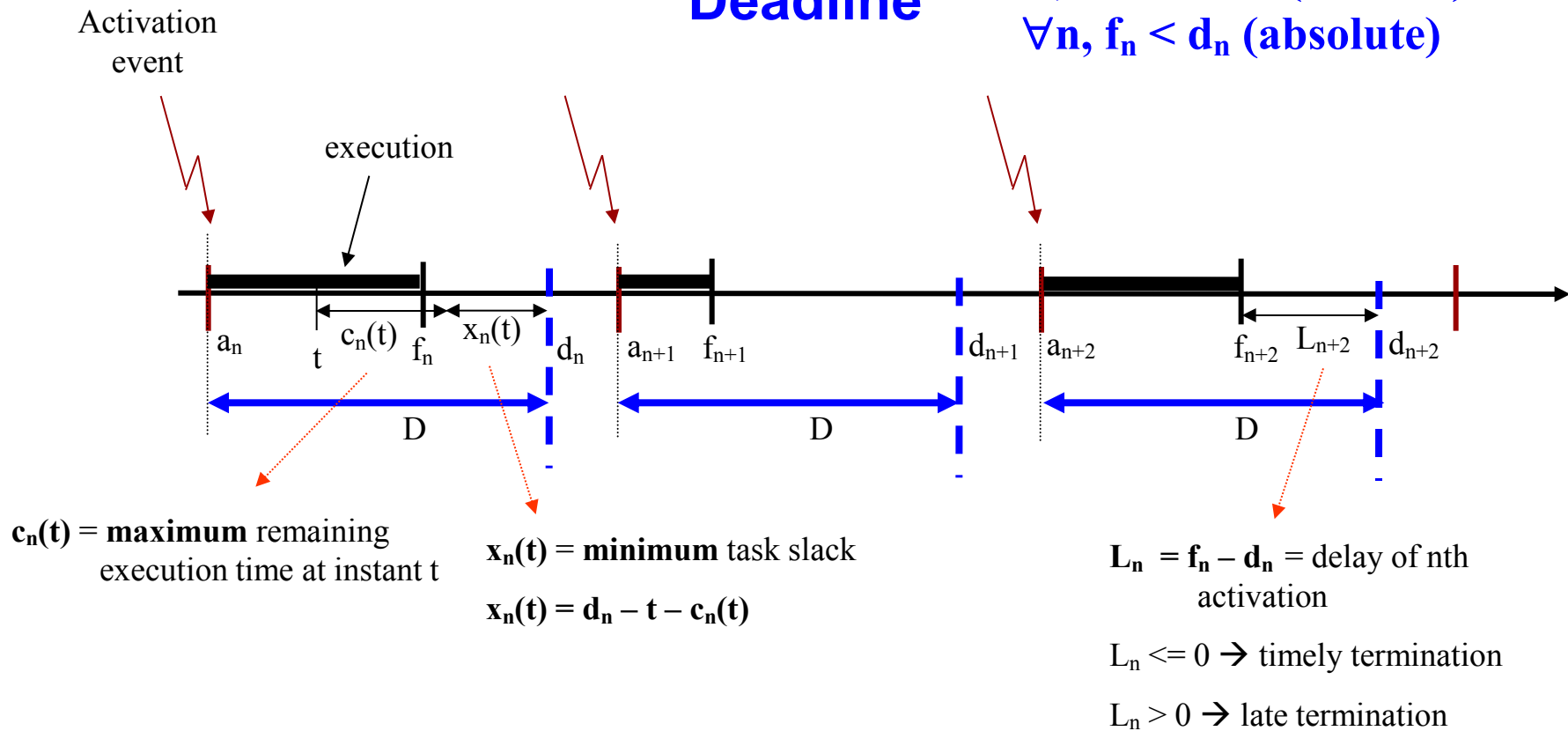
The ***deadline*** type is the most common!

Real-Time model

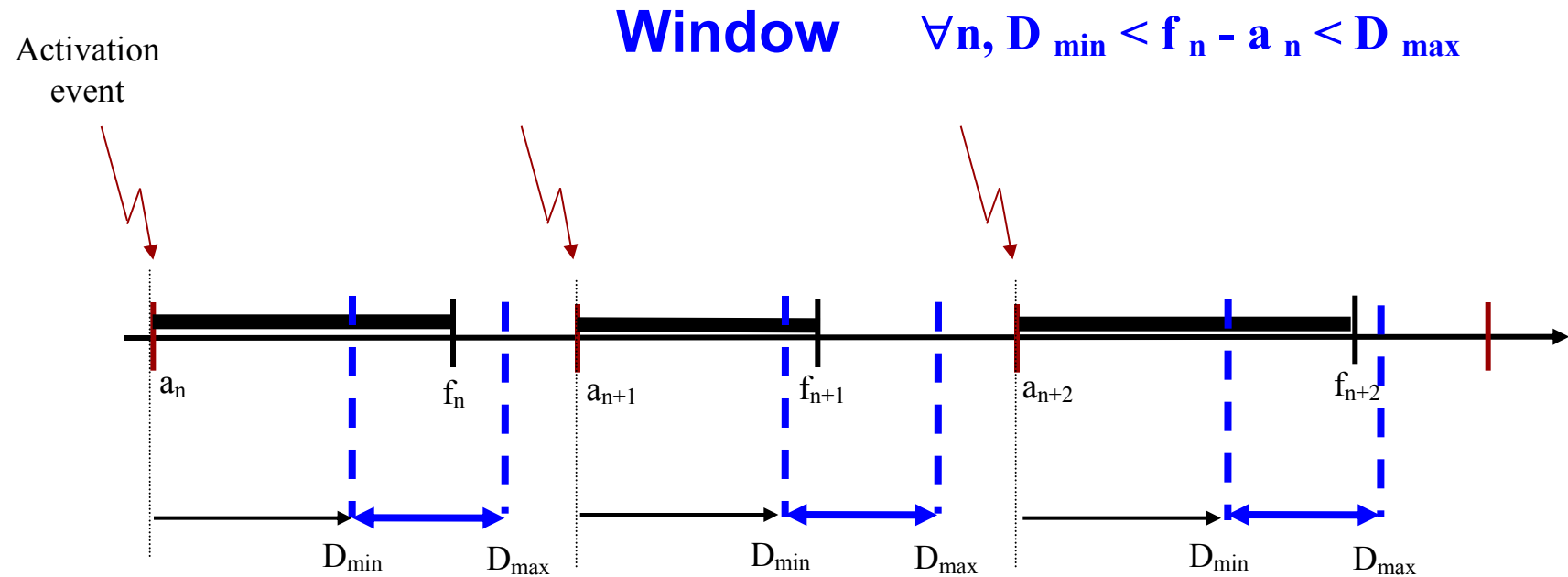
Deadline

$$\forall n, f_n - a_n < D \text{ (relative)}$$

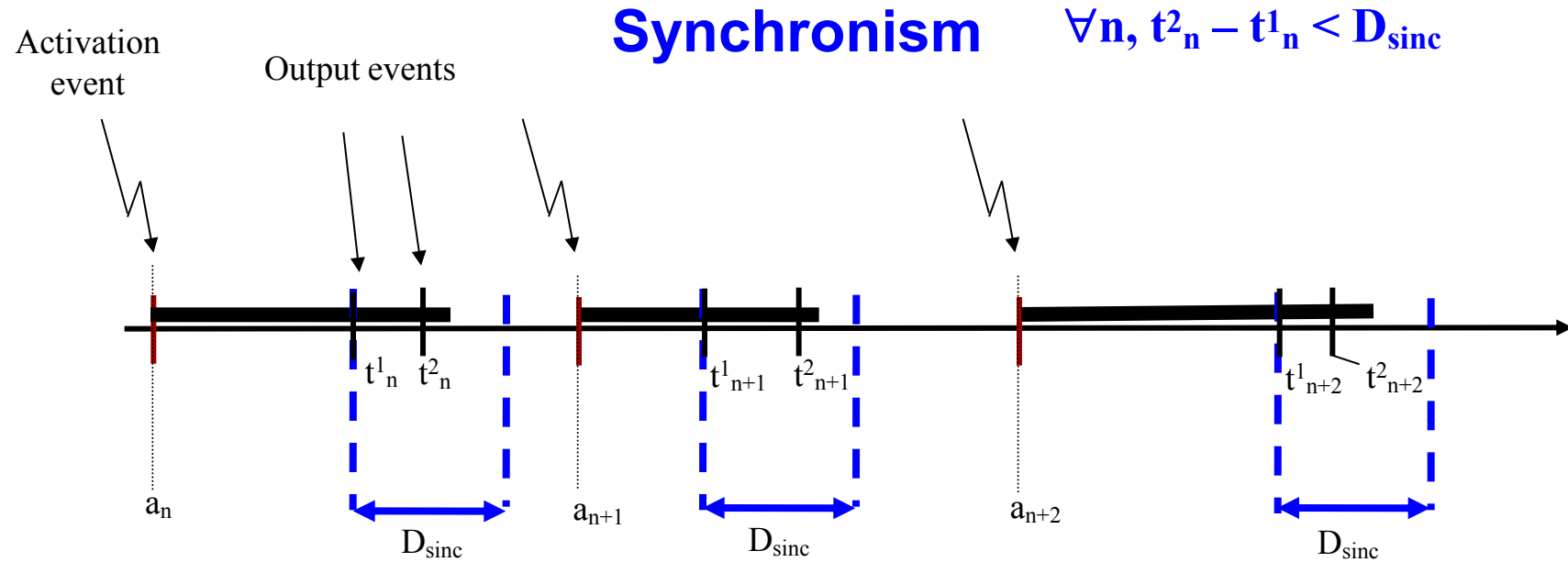
$$\forall n, f_n < d_n \text{ (absolute)}$$



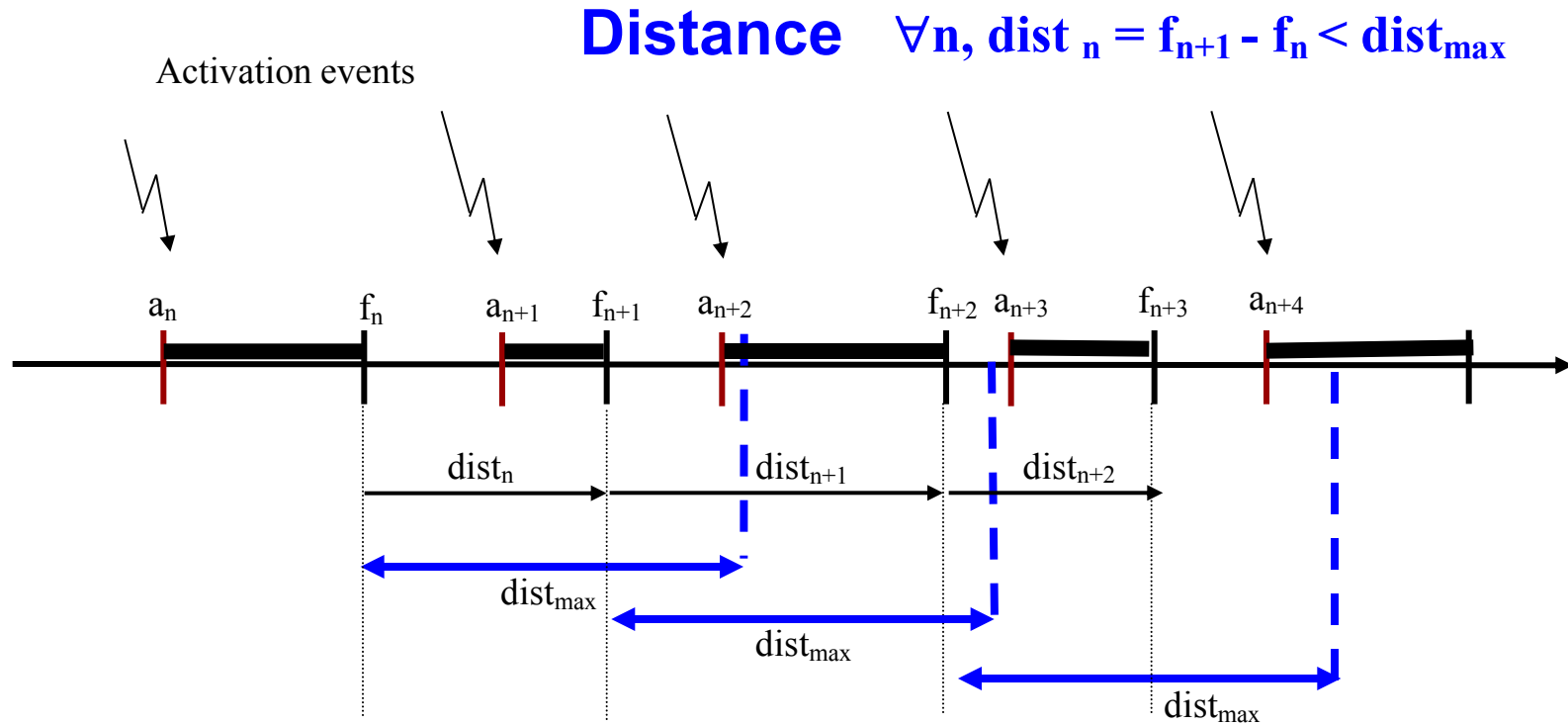
Real-Time model



Real-Time model



Real-Time model



Real-Time model

Examples of tasks characterization:

- **Periodic:** $\tau_i = \tau_i (C_i, \Phi_i, T_i, D_i)$

$$\tau_1 = \tau_1 (2, 5, 10, 10) \quad \tau_2 = \tau_2 (3, 10, 20, 20)$$

- **Sporadic:** Similar to periodic but with mit_i in the place of T_i and Φ_i is not typically used (could mean a minimal interval until the first activation).

$$\tau_i = \tau_i (C_i, mit_i, D_i)$$

$$\tau_1 = \tau_1 (2, 5, 5) \quad \tau_2 = \tau_2 (3, 10, 7)$$

Logic and temporal control

Logic control

Program control flow, i.e., effective sequence of operations to be executed – **fundamental to determine C (WCET)**

Temporal control

Control of the **execution instants** of the program operations (e.g., activations, verification of time constraints,...)

Temporal control

Triggering tasks (or functions)

By time (*time-triggered*)

The execution of tasks (or functions) is triggered by a control signal based on the progression of time (e.g., a periodic interrupt).

By events (*event-triggered*)

The execution of tasks (or functions) is triggered by an asynchronous control signal generated upon a change in the system state (e.g., by an external interrupt).

Temporal control

Activating tasks by Time

time-triggered (TT) systems

- Typically used in automatic control (sampling of continuous variables).
- There is a common time base (allows establishing offset relationships or relative phasing)
- CPU utilization is constant even when there are no variations in the system state.
- **Well defined worst-case situation**

Temporal control

Activating tasks by Events

event-triggered (ET) systems

- Typically used in monitoring sporadic conditions in the system state (e.g., alarms or asynchronous service requests).
- CPU utilization (and not only CPU!) is variable according to the frequency of event occurrence.
 - **Poorly defined worst-case situation**
either *probabilistic arguments* are used
or the *maximum event rate must be bounded*

Temporal control

Example: Consider the following task sets and determine the worst-case response time of each task

- TT $\{\tau_i = \tau_i (C_i=1, \Phi_i=i, T_i=5, D_i=T_i \ i=1..5)\}$
- ET $\{\tau_i = \tau_i (C_i=1, (\Phi_i=0), \text{mit}_i=5, D_i=\text{mit}_i \ i=1..5)\}$

Moreover, determine the average and maximum **CPU utilization**

**** utilization = $\sum_{i=1..N}(\text{execution time}/\text{activation period})$ ****

for both cases, considering that on average the ET tasks are activated once every 100 time units.

Wrapping up

- Computing models (**real-time model**)
- Real-time tasks: periodic, sporadic and aperiodic
- Temporal constraints: **deadline**, window, synchronism and distance
- **Logic control** and **temporal control**
- **Event-triggered** versus **time-triggered** tasks