Architetture dei Sistemi di Elaborazione

Delivery date: October 25<sup>th</sup> 2023

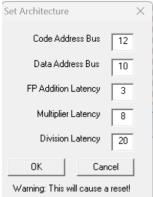
Laboratory 2

Expected delivery of lab\_02.zip must include:
- program 1.s and program 2.s

- This file, filled with information and possibly compiled in a pdf format.

Please, configure the winMIPS64 simulator with the *Base Configuration* provided in the following (*in italics not user controllable configuration*):

- Code address bus: 12
- Data address bus: 10
- Pipelined FP arithmetic unit (latency): 3 stages
- Pipelined multiplier unit (latency): 8 stages
- divider unit (latency): not pipelined unit, 20 clock cycles
- Forwarding is enabled
- Branch prediction is disabled
- Branch delay slot is disabled
- Integer ALU: 1 clock cycle
- Data memory: 1 clock cycle
- Branch delay slot: 1 clock cycle.



1) Write an assembly program (**program\_1.s**) for the *winMIPS64* architecture described before able to implement the following piece of code described at high-level:

```
for (i = 0; i < 64; i++) { v5[i] = ((v1[i]*v2[i]) + v3[i]) + v4[i]; v6[i] = v5[i]/(v4[i]+v1[i]); v7[i] = v6[i]*(v2[i]+v3[i]); }
```

Assume that the vectors v1[], v2[], v3[], and v4[] are allocated previously in memory and contain 64 double precision **floating point** values; assume also that v1[] and v4[] do not contain 0 values. Additionally, the vectors v5[], v6[], v7[] are empty vectors also allocated in memory.

<u>Calculate</u> the data memory footprint of your program:

Data	Number of Bytes
V1	512
V2	512
V3	512
V4	512
V5	512
V6	512
V7	512
Total	3584

Are there any issues? Yes, where and why? No? Do you need to change something?

Your answer: Con il data address bus impostato a 10 in memoria c'è spazio per 2^10 byte, cioè 1024 byte. Per poter eseguire correttamente questo programma ho almeno bisogno dei 3584 byte come mostrato nella tabella di cui sopra. Per poter fare ciò sono obbligato ad aumentare il data address bus fino a 12, così da avere 4096 byte.

**ATTENTION**: winMIPS64 has a limitation due to the underlying software. There is a limitation in the string length when declaring a vector. Split the vectors elements in multiple lines (it also increases the readability).

a. Compute the CPU performance equation (CPU time) of the previous program following the next directions, assume a clock frequency of 1MHz:

CPU time = 
$$(\sum_{i=1}^{n} CPI_i \times IC_i) \times Clock$$
 cycle period

- Count manually, the number of the different instructions  $(IC_i)$  executed in the program
- Assume that the  $CPI_i$  for every type of instructions equals the number of clock cycles in the instruction EXE stage, for example:
  - integer instructions CPI = 1
  - LD/SD instructions CPI = 1
  - FP MUL instructions CPI = 8
  - FP DIV instructions CPI = 20
  - ...
- b.Compute by hand again the CPU performance equation assuming that you can improve the FP Multiplier or the FP Divider by speeding up by 2 only one of the units at a time:
  - Pipelined FP multiplier unit (latency): 8 → 4 stages
     Or
  - FP Divider unit (latency): not pipelined unit,  $20 \rightarrow 10$  clock cycles

Table 1: CPU time by hand

	CPU Time	CPU Time	CPU Time	
	initial (a)	(b – MUL speed up)	(b – DIV speed up)	
program_1.S	3.650 ms	3.138 ms	3.010 ms	

c. Using the simulator calculate again the CPU time and complete the following table:

Table 2: CPU time using the simulator

CPU Time		CPU Time	CPU Time	
	initial (a)	(b – MUL speed up)	(b – DIV speed up)	
program_1.S	3.398 ms	2.886 ms	2.758 ms	

Are there any difference? If yes, where and why? If Not, provide some comments in the following:

Your answer: Sì, è normale che ci siano delle differenze in quanto nella prima tabella il calcolo viene fatto in modo approssimativo contando solo le Execute mentre nella seconda si prendono i tempi calcolati dal simulatore. In ogni caso, il calcolo a mano, nonostante abbia l'obiettivo di dare solo un'idea del tempo totale, rivela come considerando solo le execute si possa avere un'ottima approssimazione del dato reale; infatti la differenza tra i primi tempi e i secondi è di soli 252ms. Tra l'altro, il fatto che il tempo di differenza sia sempre lo stesso suggerisce che lo scarto non sia dovuto né dalla moltiplicazione né dalla divisione

d. Using the simulator and the *Base Configuration*, <u>disable the Forwarding option</u> and compute how many clock cycles the program takes to execute.

Table 3: forwarding disabled

	Number	of	IPC	(Instructions	Per
	clock cycles		Clock)		
program_1.S	4743		0.2165		

Enable one at a time the **optimization features** that were initially disabled and collect statistics to fill the following table (fill all required data in the table before exporting this file to pdf format to be delivered).

Table 4: Program performance for different processor configurations

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	Program	Forwarding		Branch	Target	Delay Slot		Forwarding +	
				Buffer				Branch	Target
								Buffer	
		IPC	CC	IPC	CC	IPC	CC	IPC	CC
	Program_1.S	0.3022	3398	0.2209	4648	0.2346	81	0.3076	3339

2) Using the WinMIPS64 simulator, validate experimentally the Amdahl's law, defined as follows:

$$speedup_{overall} = \frac{execution time_{old}}{execution time_{new}} = \frac{1}{(1 - fraction_{enhanced}) + \frac{fraction_{enhanced}}{speedup_{enhanced}}}$$

- a. Using the program developed before: program 1.s
- b. Modify the processor architectural parameters related with multicycle instructions (Menu→Configure→Architecture) in the following way:
  - 1) Configuration 1
    - Starting from the *Base Configuration*, change only the FP addition latency to 6
  - 2) Configuration 2
    - Starting from the *Base Configuration*, change only the Multiplier latency to 4
  - 3) Configuration 3
    - Starting from the *Base Configuration*, change only the division latency to 10

Compute by hand (using the Amdahl's Law) and using the simulator the speed-up for any one of the previous processor configurations. Compare the obtained results and complete the following table.

Table 5: program 1.s speed-up computed by hand and by simulation

Proc. Config.	Base config. [c.c.]	Config. 1	Config. 2	Config. 3
Speed-up comp.				
By hand	<u>3650</u>	<u>0.826</u>	<u>1.163</u>	<u>1.213</u>
By simulation	<u>3398</u>	0.898	<u>1.177</u>	1.232

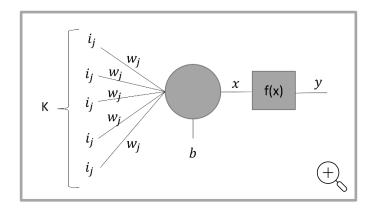
3) Write an assembly program (**program\_2.s**) for the winMIPS64 architecture able to compute the output (y) of a **neural computation** (see the Fig. below):

$$X = \sum_{j=0}^{K-1} i_j * w_j + b$$
$$V = f(X)$$

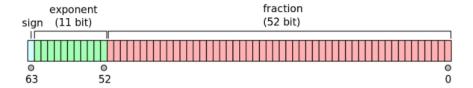
where, to prevent the propagation of NaN (Not a Number), the activation function f is defined as:

$$f(x) = \begin{cases} 0, & \text{if the exponent part of } x \text{ is equal to } 0x7ff \\ x, & \text{otherwise} \end{cases}$$

Assume the vectors i and w respectively store the inputs entering the neuron and the weights of the connections. They contain K=30 double precision **floating point** elements. Assume that b is a double precision **floating point** constant and is equal to 0xab, and y is a double precision **floating point** value stored in memory. Compute y.



Below is reported the encoding of IEEE 754 double-precision binary floating-point format:



Given the Base Configuration, run your program and extract the following information.

	Number of clock cycles	IPC (Instructions Per Clock)	CPI (Clock per Instructions)
program_2.S	400	0.47	2.128