

ASM+C & C+ASM
ABI std. compliancy
Code optimization

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Cross compiler

- A cross compiler is a compiler capable of creating executable code for a platform other than the one on which the compiler is running
- For example, a compiler that runs on a Windows 10 PC but generates a code that runs on ARM SoC is a cross compiler
- Cross compiling is a typical step for embedded application written in C language
- When integrating ASM functions, a special care is needed
 - To correctly identify the functions arguments
 - To use the proper resources to return a result
- The knowledge of the ABI standard is fundamental in this context.

Branch to main.c from startup.s

- Startup.s

; Have to import the main function address

; Executes a branch to main

```
IMPORT __main
LDR    R0, =__main
BX     R0
```

- Main.c

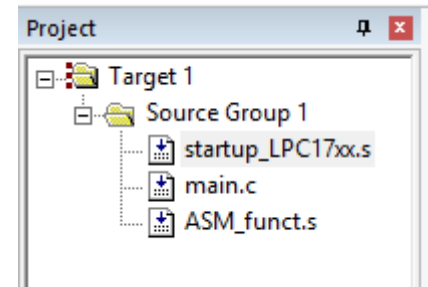
```
int main(){
    while(1) ;
}
```

```
Reset_Handler  PROC
EXPORT Reset_Handler [WEAK]
IMPORT __main
LDR    R0, =__main
BX     R0
ENDP
```

ABI_C+ASM project (I)

(in the example folder - startup_LPC17xx.s)

- In our setup we find the following source codes:
 - startup_LPC17xx.s
 - main.c
 - ASM_funct.s
- The startup_LPC17xx.s file includes the branch to main in the reset handler



```
Reset_Handler    PROC
                 EXPORT Reset_Handler             [WEAK]
                 IMPORT __main
                 LDR     R0, =__main
                 BX      R0
                 ENDP
```

ABI_C+ASM project (II)

(in the example folder - main.c)

- The main.c file is a C source code
 - It invokes a function called ASM_func with 6 parameters
 - After executing the called function, it enters in an endless loop.

```
extern int ASM_func(int, int, int, int, int, int);
```

```
int main(void) {
```

```
    int i=0xFFFFFFFF, j=2, k=3, l=4, m=5, n=6;
```

```
    volatile int r=0;
```

```
    r = ASM_func(i, j, k, l, m, n);
```

```
    while(1);
```

```
}
```

ABI_C+ASM project (III)

(in the example folder - ASM_func.s)

- Inline ASM

`__ASM("SVC #0x10");`

Per chiamare assembler in c

- External ASM function invoked by a C function

`r = ASM_func(i, j, k, l, m, n);`

Where parameters are stored?

ASM_func

How to return results?

```
AREA asm_functions, CODE, READONLY
EXPORT ASM_func
; save current SP for a faster access
; to parameters in the stack
MOV r12, sp
; save volatile registers
STMFD sp!, {r4-r8, r10-r11, lr}
; extract argument 4 and 5 into R4 and R5
LDR r4, [r12]
LDR r5, [r12, #4]
; setup a value for R0 to return
MOV r0, r5
; restore volatile registers
LDMFD sp!, {r4-r8, r10-r11, pc}
END
```

Per chiamare una procedure

Parameters are in R0-R3 (a1-a4)

Max 4

Stacked parameters

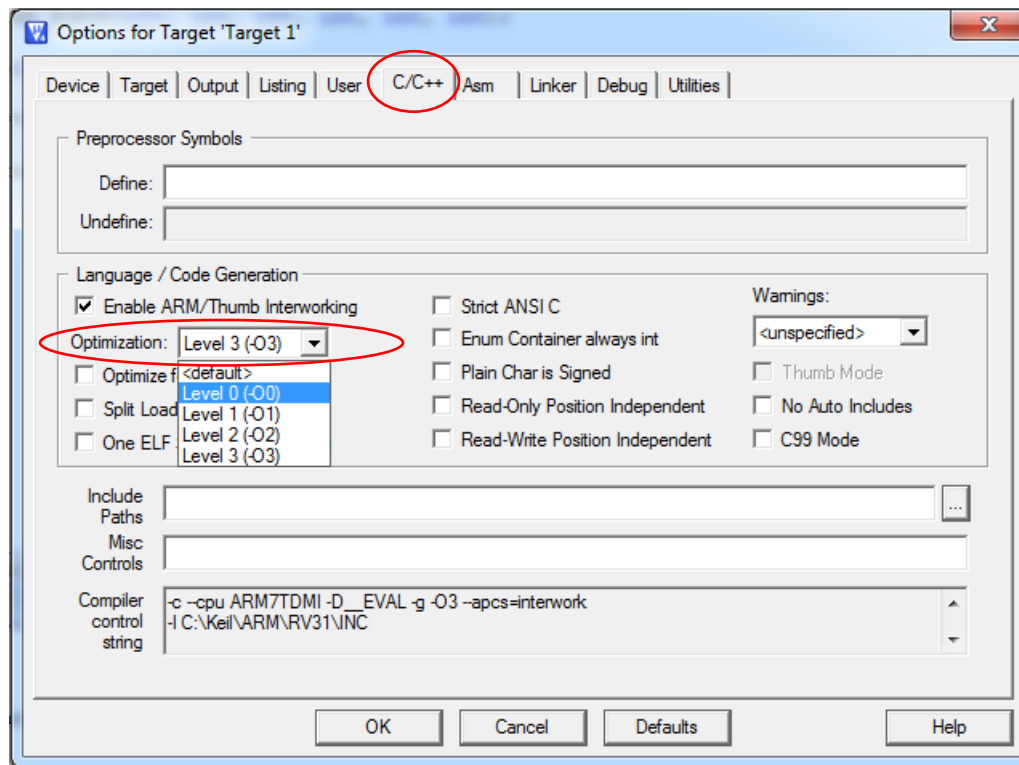
Se ce ne sono + di 4 ne usate lo stack

C and ASM helpful directives

- **EXPORT**: makes visible a function outside the file defining it
- **IMPORT**: makes visible a function from other files
- **extern**: Permits to import a variable from other file (where it is defined)

COMPILER OPTIMIZATION

- As we are now working in C language, the compiler can be asked to optimize the resulting executable file (i.e., the actual machine code).



Compiler optimization and the volatile attribute

- Higher optimization levels can reveal problems in some programs that are not apparent at lower optimization levels
- This happens when, for example, missing the volatile qualifiers
- The declaration of a variable as volatile tells the compiler that the variable can be modified at any time externally to the implementation
 - by the operating system,
 - by another thread of execution such as an interrupt routine or signal handler,
 - by hardware.

Example on volatile variables

Nonvolatile version of buffer loop

```
int buffer_full;
int read_stream(void)
{
    int count = 0;
    while (!buffer_full)
    {
        count++;
    }
    return count;
}
```

Volatile version of buffer loop

```
volatile int buffer_full;
int read_stream(void)
{
    int count = 0;
    while (!buffer_full)
    {
        count++;
    }
    return count;
}
```

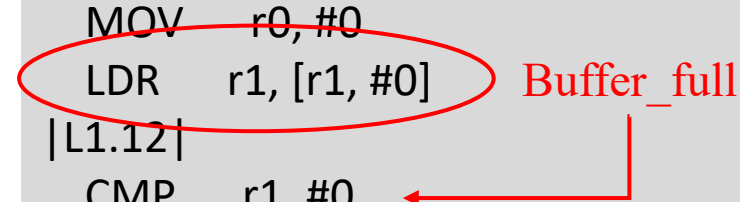
Example on volatile variables

Nonvolatile version of buffer loop

```
int buffer_full;

i read_stream PROC
{
    LDR    r1, |L1.28|
    MOV    r0, #0
    LDR    r1, [r1, #0]
    |L1.12|
    CMP    r1, #0
    ADDEQ  r0, r0, #1
    BEQ    |L1.12|    ; infinite loop
}
BX    lr
ENDP

|L1.28|
DCD    ||.data||
AREA ||.data||, DATA, ALIGN=2
buffer_full
DCD    0x00000000
```

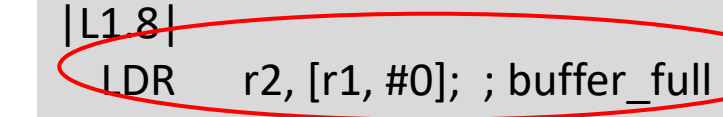


Volatile version of buffer loop

```
volatile int buffer_full;

i read_stream PROC
{
    LDR    r1, |L1.28|
    MOV    r0, #0
    |L1.8|
    LDR    r2, [r1, #0]; ; buffer_full
    CMP    r2, #0
    ADDEQ  r0, r0, #1
    BEQ    |L1.8|
}
BX    lr
ENDP

|L1.28|
DCD    ||.data||
AREA ||.data||, DATA, ALIGN=2
buffer_full
DCD    0x00000000
```



ABI standard for ARM

ABI for the ARM Architecture (Base Standard)



Application Binary Interface for the ARM[®] Architecture The Base Standard

Document number:

ARM IHI 0036B, current through ABI release 2.09

Date of Issue:

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Abstract

This document describes the structure of the Application Binary Interface (ABI) for the ARM architecture, and links to the documents that define the base standard for the ABI for the ARM Architecture. The base standard governs inter-operation between independently generated binary files and sets standards common to ARM-based execution environments.

Keywords

ABI for the ARM architecture, ABI base standard, embedded ABI

| Register | Synonym | Special | Role in the procedure call standard |
|----------|---------|----------------|---|
| r15 | | PC | The Program Counter. |
| r14 | | LR | The Link Register. |
| r13 | | SP | The Stack Pointer. |
| r12 | | IP | The Intra-Procedure-call scratch register. |
| r11 | v8 | | Variable-register 8. |
| r10 | v7 | | Variable-register 7. |
| r9 | | v6 SB TR | Platform register. The meaning of this register is defined by the platform standard. |
| r8 | v5 | | Variable-register 5. |
| r7 | v4 | | Variable register 4. |
| r6 | v3 | | Variable register 3. |
| r5 | v2 | | Variable register 2. |
| r4 | v1 | | Variable register 1. |
| r3 | a4 | | Argument / scratch register 4. |
| r2 | a3 | | Argument / scratch register 3. |
| r1 | a2 | | Argument / result / scratch register 2. |
| r0 | a1 | | Argument / result / scratch register 1. |

Can be freely used to hold
local variables

If there are more than 4
formal arguments, they
have to be saved in the
stack

Table 2, Core registers and AAPCS usage

Passing arguments

- The first four registers r0-r3 (a1-a4) are used to pass argument values into a subroutine and to return a result value in r0-r1 from a function.
 - A subroutine must preserve the contents of the registers r4-r8, r10, r11 and SP
- The base standard provides for passing arguments in core registers (r0-r3) and on the stack.
 - For subroutines that take a small number of parameters, only registers are used, greatly reducing the overhead of a call.

STACK management

- The stack implementation is *full-descending*, with the current extent of the stack held in the register SP (r13).
- The stack will, in general, have both a *base* and a *limit* though in practice an application may not be able to determine the value of either.

Disassembly

```
4: int main(void){
5:
⇒ 0x00000180 B50E      PUSH      {r1-r3,lr}
6:      int i=0xFFFFFFFF, j=2, k=3, l=4, m=5, n=6;
0x00000182 F04F34FF  MOV      r4,#0xFFFFFFFF
0x00000186 2502      MOVS     r5,#0x02
0x00000188 2603      MOVS     r6,#0x03
0x0000018A 2704      MOVS     r7,#0x04
0x0000018C F04F0805  MOV      r8,#0x05
0x00000190 F04F0906  MOV      r9,#0x06
7:      volatile int r=0;
8:
0x00000194 2000      MOVS     r0,#0x00
0x00000196 9002      STR      r0,[sp,#0x08]
9:      r = ASM_func(i, j, k, l, m, n);
10:
0x00000198 463B      MOV      r3,r7
0x0000019A 4632      MOV      r2,r6
0x0000019C 4629      MOV      r1,r5
0x0000019E 4620      MOV      r0,r4
0x000001A0 E9CD8900  STRD     r8,r9,[sp,#0]
0x000001A4 F000F83E  BL.W     ASM_func (0x00000224)
0x000001A8 9002      STR      r0,[sp,#0x08]
11:      while(1);
0x000001AA BF00      NOP
0x000001AC E7FE      B        0x000001AC
```

Variable allocation in register and initialization

Volatile variable allocation in the stack

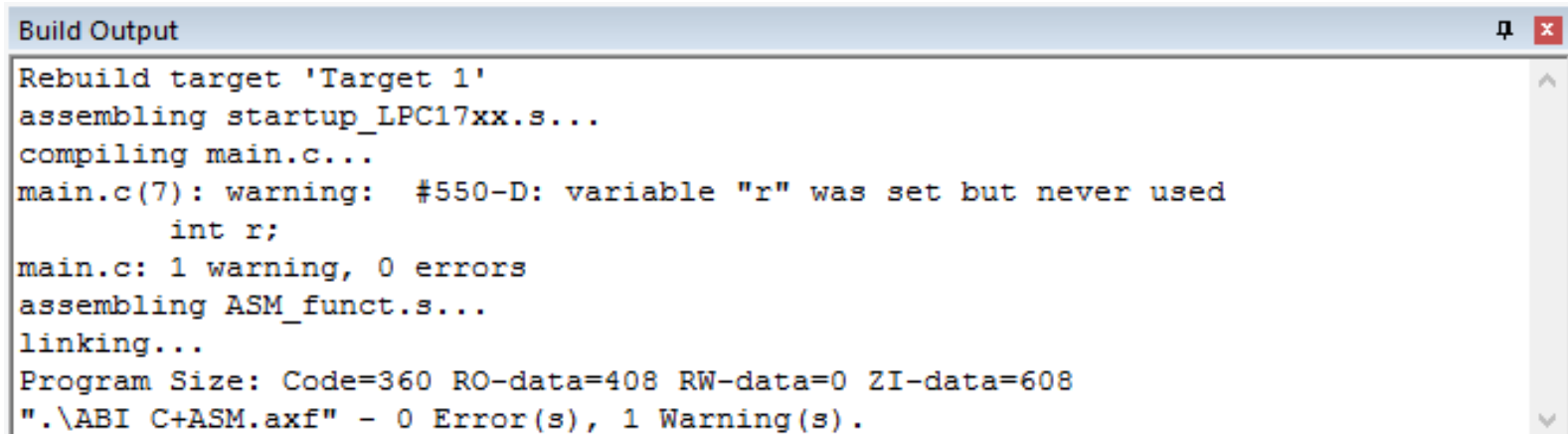
Parameters i, j, k, l setup in registers r0-r3

Parameters m, n setup in stack

Returned value is written at his address

What happens if don't declare r as volatile?

- A warning is signaled about missing usefulness of variable r
- An optimization of the machine code is implemented by the compiler.



```
Build Output
Rebuild target 'Target 1'
assembling startup_LPC17xx.s...
compiling main.c...
main.c(7): warning: #550-D: variable "r" was set but never used
    int r;
main.c: 1 warning, 0 errors
assembling ASM_funct.s...
linking...
Program Size: Code=360 RO-data=408 RW-data=0 ZI-data=608
".\ABI_C+ASM.axf" - 0 Error(s), 1 Warning(s).
```

What happens if don't declare r as volatile?

Disassembly

```
4: int main(void){
5:
→ 0x00000180 B51C      PUSH      {r2-r4,lr}
6:          int i=0xFFFFFFFF, j=2, k=3, l=4, m=5, n=6;
7:          int r;
8:
0x00000182 F04F34FF  MOV      r4,#0xFFFFFFFF
0x00000186 2502      MOVS     r5,#0x02
0x00000188 2603      MOVS     r6,#0x03
0x0000018A 2704      MOVS     r7,#0x04
0x0000018C F04F0805  MOV      r8,#0x05
0x00000190 F04F0906  MOV      r9,#0x06
9:          r = ASM_func(i, j, k, l, m, n);
10:
0x00000194 463B      MOV      r3,r7
0x00000196 4632      MOV      r2,r6
0x00000198 4629      MOV      r1,r5
0x0000019A 4620      MOV      r0,r4
0x0000019C E9CD8900  STRD     r8,r9,[sp,#0]
0x000001A0 F000F83C  BL.W     ASM_func (0x0000021C)
11:          while(1);
0x000001A4 BF00      NOP
0x000001A6 E7FE      B        0x000001A6
```

Variable r is not allocated and the ASM function is considered as "void" function