

### Strategy for the TSP Approximation Solution:

- The strategy first employed randomness to select a starting node in the initial tour
- Then, a greedy approach was used by sorting the neighbors of the current vertex and selecting the nearest unvisited one.
- Lastly, the strategy returned to a random approach by using simulated annealing to make small improvements to the tour generated by the greedy approach.

### Walk Through and Time Complexity Analysis of the TSP Approximation:

- The main method begins by parsing input from files
  - Information is parsed into variables vertices (the number of vertices), edges (the number of edges), u (current node), v (next node), and weight (the edge weight between u and v)
- Three dictionaries are instantiated and are called graph (key is the index of u in vertex map, the value is another dictionary where that key is the index of v in vertex map and the value is the weight), vertex\_map (key is the vertex, value is the index of that vertex), and rev\_map (a reverse of vertex\_map)
- There is a for loop which runs for every line of input to fill in the three dictionaries. Its runtime can be simplified to be  $O(n^2)$  because of the formula for number of edges in a complete graph.
- There is a while loop which runs for 45 seconds using a timer.
- Within this while loop, it first makes a random, greedy tour. This brings us to the nearest\_neighbor\_tour() function
- The nearest\_neighbor\_tour() function picks a random starting vertex and then uses a greedy approach of sorting the unvisited neighbors of the current vertex by edge weight and going to the nearest one first. This function has a runtime of  $O(n^2)$  because it has a for loop which loops through all of the vertex's neighbors ( $n-1$ ) and it sorts those neighbors within that loop ( $n$ ).
- After that tour has been made, it is passed as a parameter to the adjust function.
- The adjust function begins by setting the variable sample\_fraction based on the number of vertices. Then, a Boolean variable called improved is set to true and then a while loop occurs which continues looping while improved is true.
  - Within that while loop, a list of tuples is initialized using nested for loops, which take  $O(n^2)$  time
  - Next to get a random sample of those pairs, random.sample() is called, which takes  $O(n^2)$
  - Then, there is a for loop which iterates on the sampled pairs, so that runs  $n^2$  times. Within that for loop, the tour is copied which will take  $O(n)$ . so this for loop is  $O(n^3)$ .
- Next, in the main, the tour weight is calculated using a function that has a for loop that runs  $n-1$  times. So,  $O(n)$ .
- Based on this, the time complexity class is  **$O(n^3)$** , but it is within two loops, one affected by a timer and the other affected by whether the code improves.