## Homework #1

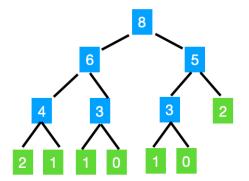
1. When larger values of N are tested with my written PAD function, the program takes a long time, often several minutes, to run the function. This is because my PAD function uses binary recursion, calling PAD(n-2) and PAD(n-3) within every function call. Each call of PAD branches into two more calls, so the algorithm grows exponentially. Therefore, increasing N severely increases the runtime of the PAD algorithm. An iterative solution would be favored for space and efficiency, in which previous values are stored to calculate the next value.

## Test cases:

PAD(0) = 1 PAD(1) = 1 PAD(2) = 1 PAD(3) = 2 PAD(4) = 2 PAD(5) = 3 PAD(6) = 4 PAD(7) = 5 PAD(8) = 7 PAD(9) = 9 PAD(10) = 12

2. The relationship between the output of SUMS(n) and PAD(n) is: SUMS(n) = PAD(n) - 1. This is because when constructing a full binary tree representing the recursive calls of PAD, we will have a full binary tree with N leaf nodes and 2N-1 total nodes. The number of leaf nodes, N, is the result of PAD(n), and the total number of nodes, 2N-1, represents the total number of calls to PAD. However, the leaf nodes (N total) are PAD calls that do not perform addition since they simply return 1, while each non-leaf node performs one addition. Therefore, (2N-1)-N = N-1, so SUMS(n) = PAD(n)-1.

For example, in the binary tree below, constructed when PAD is called on n=8, the function calls create a binary tree with 13 total nodes and 7 leaf nodes, so the SUMS(8)=6 and PAD(8)=7, which fits with the observation that SUMS(n)=PAD(n)-1 since 6=((7\*2)-1)-7.



Test cases:

```
SUMS(0) = 0

SUMS(1) = 0

SUMS(2) = 0

SUMS(3) = 1

SUMS(4) = 1

SUMS(5) = 2

SUMS(6) = 3

SUMS(7) = 4

SUMS(8) = 6

SUMS(9) = 8

SUMS(10) = 11
```

## 3. Test cases: