

# **Programming Language Translation**

#### Practical 5: week beginning 13 August 2018

Hand in this prac sheet before the start of your next practical, correctly packaged in a transparent folder with your solutions and the "cover sheet". Unpackaged and late submissions will not be accepted. Since the practical is done on a group basis, please hand in one copy of the cover sheet for each member of the group. These will be returned to you in due course, signed by the marker. Lastly, please resist the temptation to carve up the practical, with each group member only doing one task. The group experience is best when you discuss each task together.

### **Objectives**

In this practical you are to

• develop a recursive descent parser and associated *ad hoc* scanner "from scratch" that will analyse a set of C or C++ declaration statements.

Copies of this handout and the cover sheet are available on the RUConnected course page.

#### **Outcomes**

When you have completed this practical you should understand:

- the inner workings of an ad hoc scanner;
- the inner workings of a recursive descent parser;
- how to test that a scanner and parser behave correctly;
- (hopefully) C and C++ declarations a little better than before.

#### To hand in (30 marks)

This week you are required to hand in, besides the cover sheet:

- A listing of the final version of your source program, and some listings of input and output files;
- Electronic copies of the sources of your program.

Please submit your prac in the "hand-in" box inside the laboratory at the start of the next practical session.

WARNING. This exercise really requires you to do some real thinking and planning. Please do not just sit at a computer and hack away as most of you are wont to do. Sit in your groups and discuss your ideas with one another and with the demonstrators. If you don't do this you will probably find that the whole exercise turns into a nightmare, and I don't want that to happen.

Keep the cover sheet and your solutions until the end of the semester. Check carefully that your mark has been entered into the Departmental Records.

A rule not stated there, but which should be obvious, is that you are not allowed to hand in another student's or group's work as your own. Attempts to do this will result in (at best) a mark of zero and (at worst) severe disciplinary action and the loss of your DP. You are allowed -- even encouraged -- to work and study with other students, but if you do this you are asked to acknowledge that you have done so on *all* cover sheets and

with suitable comments typed into *all* listings. You are expected to be familiar with the University Policy on Plagiarism.

#### Task 0 Creating a working directory and unpacking the prac kit

Unpack the prac kit PRAC5.ZIP. In it you will find the skeleton of a system adapted for intermediate testing of a scanner (and to which you will later add a parser), and some simple test data files -- but you really need to learn to develop your own test data.

#### Task 1 Get to grips with the problem

It is generally acknowledged, even by experts, that the syntax of declarations in C and C++ can be quite difficult to understand. This is especially true for programmers who have learned Pascal, Modula-2 or C# before turning to a study of C or C++. Simple declarations like

```
int x, list[100];
```

present few difficulties (x is a scalar integer, list is an array of 100 integers). However, in developing more complicated examples like

it is easy to confuse the placement of the various brackets, parentheses and asterisks, perhaps even writing syntactically correct declarations that do not mean what the author intended. By the time one is into writing (or reading) declarations like

```
bool (*(*f())[])();
int (*(*g[50])())[15];
```

there may be little consolation to be gained from learning that C was designed so that the syntax of declarations (defining occurrences) should mirror the syntax for access to the corresponding quantities in expressions (applied occurrences).

A grammar that describes many of the forms that such declarations can take might be expressed in Cocol as follows (please note however, that this grammar does not allow ALL C/C++ declarations):

```
COMPILER Cdecls

/* Describe a subset of the forms that C declarations can assume */

CHARACTERS

digit = "0123456789" .

letter = "ABCDEFGHIJKLMNOPQRSTUVWXYZabcdefghijklmnopqrstuvwxyz_" .

IGNORE CHR(0) .. CHR(31)

TOKENS

number = digit { digit } .

ident = letter { letter | digit } .
```

```
PRODUCTIONS

Cdecls = { DecList } EOF .

DecList = Type OneDecl { "," OneDecl } ";" .

Type = "int" | "void" | "bool" | "char" .

OneDecl = "*" OneDecl | Direct .

Direct = (ident | "(" OneDecl ")" ) [ Suffix ] .

Suffix = Array { Array } | Params .

Params = "(" [ OneParam { "," OneParam } ] ")" .

OneParam = Type [ OneDecl ] .

Array = "[" [ number ] "]" .

END Cdecls.
```

Tempting as it might be simply to use Coco/R to produce a program that will analyse declarations, this week we should like you to produce such a recognizer more directly, by developing a program in the spirit of the one you will find in Chapter 8 of the notes.

The essence of this program is that it will eventually have a main method that will

- use a command line parameter to retrieve the file name of a data file;
- from this file name derive an output file name with a different extension;
- open these two files;
- initialize the "character handler";
- initialize the "scanner";
- start the "parser" by calling the routine that is to parse the goal symbol;
- close the output file and report that the system parsed correctly.

In this practical you are to develop such a scanner and parser, which you should try in easy stages. So for Task 1, study the grammar above and the skeleton program from the kit (Declarations.cs) as shown below. In particular, note how the character handler section has been programmed.

```
// This is a skeleton program for developing a parser for C declarations
// P.D. Terry, Rhodes University, 2015
using Library;
using System;
using System. Text;
class Token {
 public int kind;
 public string val;
 public Token(int kind, string val) {
   this.kind = kind;
   this.val = val;
 } // constructor
} // Token
class Declarations {
 static InFile input;
```

```
static OutFile output;
   static string NewFileName(string oldFileName, string ext) {
   // Creates new file name by changing extension of oldFileName to ext
     int i = oldFileName.LastIndexOf('.');
     if (i < 0) return oldFileName + ext; else return oldFileName.Substring(0,
i) + ext;
   } // NewFileName
   static void ReportError(string errorMessage) {
   // Displays errorMessage on standard output and on reflected output
     Console.WriteLine(errorMessage);
     output.WriteLine(errorMessage);
   } // ReportError
   static void Abort(string errorMessage) {
   // Abandons parsing after issuing error message
    ReportError(errorMessage);
     output.Close();
     System.Environment.Exit(1);
   } // Abort
   const int
    noSym
                = 0,
     EOFSym
               = 1;
     // and others like this
   const char EOF = ' \setminus 0';
   static bool atEndOfFile = false;
   // Declaring ch as a global variable is done for expediency -
   // global variables are not always a good thing
   static char ch; // look ahead character for scanner
   static void GetChar() {
   // Obtains next character ch from input, or CHR(0) if EOF reached
   // Reflect ch to output
     if (atEndOfFile) ch = EOF;
     else {
      ch = input.ReadChar();
      atEndOfFile = ch == EOF;
      if (!atEndOfFile) output.Write(ch);
   } // GetChar
   // Declaring sym as a global variable is done for expediency
   // global variables are not always a good thing
   static Token sym;
```

```
static void GetSym() {
 // Scans for next sym from input
   while (ch > EOF && ch <= ' ') GetChar();
   StringBuilder symLex = new StringBuilder();
   int symKind = noSym;
   // *********over to you!
   sym = new Token(symKind, symLex.ToString());
 } // GetSym
/* ++++ Commented out for the moment
 static void Accept(int wantedSym, string errorMessage) {
 // Checks that lookahead token is wantedSym
   if (sym.kind == wantedSym) GetSym(); else Abort(errorMessage);
 } // Accept
 static void Accept(IntSet allowedSet, string errorMessage) {
 // Checks that lookahead token is in allowedSet
   if (allowedSet.Contains(sym.kind)) GetSym(); else Abort(errorMessage);
  } // Accept
 static void CDecls() {}
+++++ */
 public static void Main(string[] args) {
   // Open input and output files from command line arguments
   if (args.Length == 0) {
     Console.WriteLine("Usage: Declarations FileName");
     System.Environment.Exit(1);
   input = new InFile(args[0]);
   output = new OutFile(newFileName(args[0], ".out"));
                                            // Lookahead character
   GetChar();
// To test the scanner we can use a loop like the following:
   do {
     GetSym();
                                           // Lookahead symbol
     OutFile.StdOut.Write(sym.kind, 3);
     OutFile.StdOut.WriteLine(" " + sym.val); // See what we got
   } while (sym.kind != EOFSym);
/* After the scanner is debugged we shall substitute this code:
   GetSym();
                                  // Lookahead symbol
   CDecls();
                                  // Start to parse from the goal symbol
   // if we get back here everything must have been satisfactory
   Console.WriteLine("Parsed correctly");
* /
```

```
output.Close();
} // Main
} // Declarations
```

#### Task 2 First steps towards a scanner [10 marks]

Next, develop the scanner by completing the getSym method, whose goal in life is to recognize tokens. Tokens for this application could be defined by an enumeration of

```
noSym, intSym, charSym, boolSym, voidSym, numSym, identSym,
lparenSym, rparenSym, lbrackSym, rbrackSym, pointerSym,
commaSym, semicolonSym, EOFSym
```

The scanner can (indeed, must) be developed on the pretext that an initial character ch has been read. When called, it must (if necessary) read past any "white space" in the input file until it comes to a character that can form part (or all) of a token. It must then read as many characters as are required to identify a token, and assign the corresponding value from the enumeration to the kind field of an object called, say, sym -- and then read the next character ch (remember that the parsers we are discussing always look one position ahead in the source).

Test the scanner with a program derived from the skeleton, which should be able to scan the data file and simply tell you what tokens it can find, using the simple loop in the main method as supplied. At this stage do not construct the parser, or attempt to deal with comments. A simple data file for testing can be found in the files SAMPLEO.CPP, a longer one in SAMPLE1.CPP, one that simply has a list of all tokens in TEST.TXT.

You can compile your program by giving the command

NB: This task must be handed in before 5 pm on Thursday 16<sup>th</sup> August. Please submit a hardcopy of your GetSym() and auxiliary routines.

## Task 3 Handling comments [5 marks]

Next, refine the scanner so that it can deal with (that is, safely ignore) C++ style comments (of both sorts) in found the set of declarations. Suitable data files for testing in are to be the files SAMPLE2.CPP and SAMPLE3.CPP.

You cannot possibly expect to start on Task 4 until such time as the scanner is working properly, so test it thoroughly, please!

#### Task 4 At last, a parser! [15 marks]

This task is to develop the associated parser as a set of routines, one for each of the non-terminals suggested in the grammar above. These methods should, where necessary, simply call on the GetSym scanner routine to deliver the next token from the input. As discussed in Chapter 8, the system hinges on the premise that each time a parsing routine is called (including the initial call to the goal routine) there will already be a token waiting in the variable sym, and whenever a parsing routine returns, it will have obtained the follower token in readiness for the caller to continue parsing (see discussion on page 102). It is to make communication between all these routines easy that we declare the lookahead character ch and the lookahead token sym to be fields "global" to the Declarations class.

Of course, anyone can write a recognizer for input that is correct. The clever bit is to be able to spot incorrect input, and to react by reporting an appropriate error message. For the purposes of this exercise it will be sufficient first to develop a simple routine on the lines of the accept routine that you see on page 105, that simply issues a stern error message, closes the output file, and then abandons parsing altogether.

Something to think about: If you have been following the lectures, you will know that associated with each nonterminal A is a set FIRST(A) of the terminals that can appear first in any string derived from A. So that's why we learned to use the IntSet class in practical 1! Library routines especially developed for this course, are available on RUConnected (Library.cs).

#### A note on testing

To test your parser you might like to make use of the data files supplied. One of these (SAMPLE1.CPP) has a number of correct declarations. Another (SAMPLE4.CPP) has a number of incorrect declarations. Your parser should, of course, be able to parse SAMPLE1.CPP easily, and you should be able to "compile" this same file by issuing a command like BCC SAMPLE1.CPP just to verify this. Using BCC to "compile" SAMPLE4.CPP should be fun, but parsing SAMPLE4.CPP with your system will be a little more frustrating unless you added syntax error recovery, as the parser will simply stop as soon as it finds the first error. You might like to create a number of "one-liner" data files to make this testing stage easier. Feel free to experiment! But, above all, do test your program out.

#### Task 5 Food for thought

Due to lack of time available, you have not been asked to incorporate syntax error recovery techniques into your parser. However, since error recovery is a vital part of creating a usable parser, and the topic is examinable, you might like to investigate how this could be achieved. The model solution will cover this as well.