

14. (Aside) Johnson-Lindenstrauss

Friday, February 9, 2024

3:16 PM

Note: these notes use different notation

Fact If $X_i \sim N(0, 1)$ are independent, then $Z = \sum_{i=1}^M X_i^2 \sim \chi_M^2$ ("chi-squared")

and if $Z \sim \chi_M^2$ then $\mathbb{E}[Z] = M$ and $\forall \varepsilon \in (0, 1/2)$

$$\mathbb{P}[|Z - M| \geq \varepsilon \cdot M] \leq 2 \cdot \exp(-\frac{M}{4}(\varepsilon^2 - \varepsilon^3))$$

could also use Bernstein's ineq. for sub-exponential r.v.

Lemma Fix $x \in \mathbb{R}^N$, let $A \in \mathbb{R}^{M \times N}$ be random w/ $A_{ij} \sim N(0, 1)$ iid

let $y = Ax$, then w.p. $\geq 1 - 2 \exp(-\frac{M}{4}(\varepsilon^2 - \varepsilon^3))$,

$$(1 - \varepsilon) \|x\|^2 \leq \|\frac{1}{\sqrt{M}} y\|^2 \leq (1 + \varepsilon) \|x\|^2$$

proof wlog let $\|x\| = 1$.

- y_i is also a Gaussian, as it is a weighted sum of Gaussians

(Recall: if $X \sim \mathcal{L}_1$ and $Z \sim \mathcal{L}_2$ then $X + Y \sim \mathcal{L}_1 * \mathcal{L}_2$ convolution)

The sum of normal distributions (i.e. multimodal) isn't Gaussian, it's a mixture model

- so it's completely characterized by its mean + variance

$$\mathbb{E}[y_i] = \sum_j x_j \mathbb{E}[a_{ij}] = 0$$

$$\text{Var}[y_i] = \mathbb{E}[y_i^2] = \mathbb{E}[(\sum_{j=1}^N a_{ij} x_j)^2] = \mathbb{E}[\sum_{j=1}^N \sum_{j'=1}^N a_{ij} a_{ij'} x_j x_{j'}]$$

$$\text{and } y_i \perp y_j \text{ (i} \neq \text{j)} = \|x\|^2 + 0 = 1$$

$$\Rightarrow \|y\|^2 \sim \chi_M^2$$

$$\text{and } \left| \|y\|^2 - M \right| \geq \varepsilon \cdot M \Leftrightarrow \left| \frac{1}{M} \|y\|^2 - 1 \right| \geq \varepsilon \stackrel{= \|x\|^2}{=} \text{so use Fact above.}$$

Thm (Johnson-Lindenstrauss 1984) [one of many variants]

Let $X = \{x_1, \dots, x_k\} \subseteq \mathbb{R}^N$ and $\varepsilon \in (0, 1/2)$. If $M \geq \frac{16}{\varepsilon^2} \log(k)$

then \exists a Lipschitz continuous map $f: \mathbb{R}^N \rightarrow \mathbb{R}^M$ s.t. $\forall x, y \in X$

$$(1 - \varepsilon) \|x - y\|^2 \leq \|f(x) - f(y)\|^2 \leq (1 + \varepsilon) \|x - y\|^2 \quad (*)$$

proof Draw $A \in \mathbb{R}^{M \times N}$ as before, think of $f(x) = \frac{1}{\sqrt{M}} \cdot Ax$

For a fixed x, y then (via linearity) $(*)$ holds w.p. $\geq 1 - 2 \cdot \exp(-\frac{M}{4}(\varepsilon^2 - \varepsilon^3))$

and we have $\binom{k}{2} = \frac{k(k-1)}{2}$ pairs, so

$$\begin{aligned} \mathbb{P}[\text{any pair } x, y \text{ fails } (*)] &\leq \sum_{\text{all pairs}} \mathbb{P}[\text{fixed pair } x, y \text{ fails } (*)] \\ &\leq \frac{k(k-1)}{2} \cdot 2 \cdot \exp(-\frac{M}{4}(\varepsilon^2 - \varepsilon^3)) \end{aligned}$$

so choose $M = \frac{16}{\varepsilon^2} \log(k) \dots$

$$= k(k-1) \exp(-\frac{16}{4\varepsilon^2}(\varepsilon^2 - \varepsilon^3) \log(k))$$

Facts: $\tilde{a} \sim N(\tilde{\mu}, \Sigma)$

$$\Rightarrow C \cdot \tilde{a} \sim N(C \cdot \tilde{\mu}, C \Sigma C^T)$$

14a. Johnson-Lindenstrauss

Monday, February 12, 2024

11:04 AM

$$\begin{aligned} &= k \cdot (k-1) \cdot k^{(-4(1-\frac{1}{2}\epsilon))} \quad \text{and } \epsilon < \frac{1}{2} \Rightarrow 1 - \frac{1}{2}\epsilon > -\frac{1}{2} \\ &< k \cdot (k-1) \cdot k^{-2} = \frac{k-1}{k} < 1 \quad \text{or } -4(1-\frac{1}{2}\epsilon) < -2 \end{aligned}$$

So $\mathbb{P}[(*) \text{ holds } \forall \text{ pairs}] > 0$

via the probabilistic method this means such a map f must exist. \square

Discussion

- In practice, for a specific A that we draw, we want it to be very likely, so make M larger in that case
- $O(1/\epsilon^2)$ dependence grows quickly so it's best for low accuracy ... but analysis is tight up to a $\log(1/\epsilon)$ factor
 $\log(k)$ dependence is "correct"
- Independent of original dimension N !] Big deal!
- Works if A_{ij} Sub-Gaussian
- Faster "JL-inspired" transforms exist... ask me about them?
- See David Woodruff's 2014 monograph for an example of a chaining argument to use classical JL to apply to a whole subspace.
- Preserving $\|\Phi(x-y)\|^2 = (1 \pm \epsilon)\|x-y\|^2 \quad \forall x, y \in V$
also preserves inner products. Assume $0 \in V$ so $\|\Phi x\|^2 \approx \|x\|^2 \quad \forall x \in V$
Then the polarization identity (for \mathbb{R} scalars) is
$$\langle x, y \rangle = \frac{1}{2}(\|x\|^2 + \|y\|^2 - \|x-y\|^2)$$

so $\langle \Phi x, \Phi y \rangle = \frac{1}{2}(\|\Phi x\|^2 + \|\Phi y\|^2 - \|\Phi(x-y)\|^2)$
$$\approx \frac{1}{2}(\|x\|^2 + \|y\|^2 - \|x-y\|^2) = \langle x, y \rangle.$$

14b. Johnson-Lindenstrauss + chaining

Sunday, February 18, 2024

8:30 PM

$$a = (1 \pm \epsilon)b \text{ shorthand for } (1 - \epsilon)b \leq a \leq (1 + \epsilon)b$$

Recap: before we talked about preserving distance $\|\Phi x - \Phi y\|^2 = (1 \pm \epsilon) \|x - y\|^2$
and if this is for k points (ie. $k \cdot (k-1)/2$ pairs), can do it w/ random matrix w/ $m \geq O(\frac{1}{\epsilon^2} \log(k/\delta))$

To simplify, look only at $\|\Phi x\|^2 = (1 \pm \epsilon) \|x\|^2$ for k points
(to go from k to $\tilde{k} = O(k^2)$ pairs, it's in the log term so not a big deal)

What if we want $\|\Phi x\|^2 = (1 \pm \epsilon) \|x\|^2$ for all x in a subspace V
In particular, $A = V \cap S^{n-1}$ unit sphere
 $\nwarrow d$ -dim subspace of \mathbb{R}^n an ∞ # of points!

Idea #1: Cover A w/ an $\epsilon/2$ -net, let this be " B ". Then find a $\epsilon/2$ J.L matrix for $|B|$ points (ie. $N_2(\epsilon/2, A)$ points). Use triangle-ineq.
That works but it gives a large number ...
 $N_2(\epsilon/2, A) \leq (1 + 2/\epsilon)^d$... nasty ϵ -dependence.

Idea #2 Clever chaining argument (cf. Woodruff book) $1 \pm \epsilon \|\Phi b\|^2 = (1 \pm \epsilon) \|b\|^2$
Let B be a $\epsilon/2$ net for A (and $B \subseteq A$, so $\forall b \in B, \|b\| \leq 1$) actually assume
Fix $a \in A$, goal is $\|\Phi a\|^2 = (1 \pm \epsilon) \|a\|^2$. Note $\|a\| = 1$ $\langle \Phi b, \Phi b' \rangle = (1 \pm \epsilon) \langle b, b' \rangle$
Write $a = b^0 + a^0$ w/ $b^0 \in B$ being closest pt. in B , so $\|a^0\| \leq 1/2$ since B is a $1/2$ -net which is similar

Then $B \subseteq A \subseteq V$ (subspace)

so $a^0 \in V$ also. So $\frac{a^0}{\|a^0\|} \in A = V \cap S^{n-1}$

Write $\frac{a^0}{\|a^0\|} = b^1 + a^1$ w/ $b^1 \in B, \|a^1\| \leq 1/2$, etc.

So, continuing,

$$a = b^0 + a^0 = b^0 + \|a^0\| (b^1 + a^1)$$

$$= b^0 + \|a^0\| (b^1 + \|a^1\| (b^2 + a^2))$$

$$= \sum_{k=0}^{\infty} c^k \text{ w/ } c^0 = b^0, c^1 = \|a^0\| \cdot b^1, c^2 = \|a^0\| \cdot \|a^1\| \cdot b^2 \dots$$

So

so $\|c^k\| \leq 2^{-k}$ Idea: we re-use the same b^i sometimes, just w/ different coefficients!
being a bit "loose" w/ details

$$\|\Phi a\|^2 = \|\Phi \sum_{k=0}^{\infty} c^k\|^2 = \sum_{0 \leq i < j < \infty} \|\Phi c^i\|^2 + 2 \langle \Phi c^i, \Phi c^j \rangle$$

$$= \left(\sum_{0 \leq i < j < \infty} \underbrace{\|\Phi c^i\|^2 + 2 \langle \Phi c^i, \Phi c^j \rangle}_{= \|\sum_{i=j}^{\infty} c^k\|^2} \right) \pm 2 \epsilon \left(\sum_{\substack{i \leq j \\ 2^{-i} \leq 2^{-j} \\ 0 \leq i}} \underbrace{\|c^i\| \cdot \|c^j\|}_{\leq 2^{-i}} \right)$$

$$= \|a\|^2 + O(\epsilon)$$