

## Amortized Analysis

### CLRS (Chap. 17)

Bound  $\sum a_i$  by  $\sum t_i$  where  $a_i$  is amortized cost and  $t_i$  is real cost.

Costs only considered in sequences.

Prove bounds on average running-time *of* the worst case, **not** the average case.

### Aggregate Analysis

- All operations considered to be same running time.
- Multipop stack example: can only be  $O(n)$  since can only pop what you push.

### Banker / Accounting

- Different operations can have different running times
- Difference of  $a_i - t_i$  is \*credit\*, that can help 'pay' for future operations
- "Total credit must be non-negative at all times"
  - Negative credit is like undercharging earlier operations with the promise of paying them later
- Stack example
  - Pushing uses one dollar \*immediately\*
  - Popping has no charge, we use credit (we take the dollar off the plate at the top of the stack)
  - "charging push operation a little bit more, we can charge pop nothing"
  - there is always a positive number of credit since pushing adds one dollar.

### Physicist / Potential

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### Dynamic Tables

### Okasaki