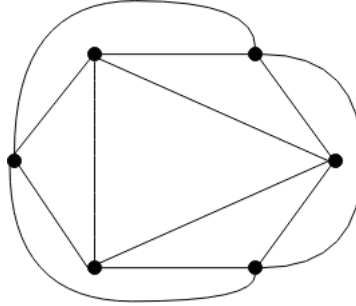


MATH 222 Assignment 2

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1. Consider a tree T with some vertex v of degree 4.
 Since a tree cannot have any cycles, every path from v taking a unique edge incident to v must eventually terminate in at least one leaf.
 Therefore T must have at least 4 leaves.
 Now let's consider if T has another vertex u of degree 3.
 Since T is a tree, there must be a single $v-u$ path in T .
 And since there must be a single $u-v$ path and there can be no cycles, exactly one of the edges incident to v must yield a $v-u$ path and exactly one of the edges incident to u must yield a $u-v$ path.
 So there are exactly 3 paths from v taking a unique edge incident to v that terminate in at least one leaf, and similarly there are exactly 2 paths from u taking a unique edge incident to u that terminate in at least one leaf.
 $\therefore T$ contains at least 5 leaves.
2. For a tree $T = (V, E)$, there must have a single $u-v$ path in $T \forall u, v \in V$
 Therefore
- 3.



Example of a 4-regular plane graph with 6 vertices.

Since $|V| - |E| + R = 2 \iff 6 - 12 + 8 = 2$ holds, we can verify that this example holds.

4. We can use the theorem that $|E| \leq 3|V| - 6$ for any planar graph G with $|V| \geq 3$.
 Assume that both G and \overline{G} are planar.
 Therefore for G :

$$|E| \leq 3|V| - 6$$

$$|E| \leq 3|11| - 6$$

$$|E| \leq 27$$

And therefore for \overline{G} , since the maximum number of edges in \overline{G} is the difference between the number of edges in K_1 and the maximum number of edges in G :

$$\begin{aligned} |E| &\leq 3|V| - 6 \\ \frac{11(11-1)}{2} - 27 &\leq 3|11| - 6 \\ 28 &\leq 27 \end{aligned}$$

Which is a contradiction.

\therefore either G or \overline{G} must be non-planar.

5. First we prove that if $\deg(v) < k$ and G is k -colorable then $G - v$ is k -colorable.

We know that k must be larger than the maximum degree vertex in G .

Since G is k -colorable using Brooks' theorem, we know that the maximum degree of a vertex in G must be $k - 1$.

Since the maximum degree of $G - v$ is less or equal the maximum degree of G because some edges in G may have been removed, therefore the maximum degree of $G - v$ must be $k - 1$.

\therefore if G is k -colorable, $G - v$ must be k -colorable.

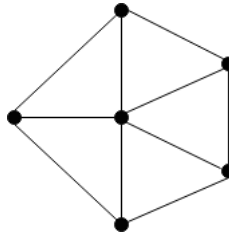
Now we prove that if $G - v$ is k -colorable then G is k -colorable.

Using Brooks's theorem the maximum degree of $G - v$ is $k - 1$.

And since $v < k$, G will have a maximum degree of k since each of the vertices in $G - v$ will have at most one more edge in G .

Since G has a maximum degree of k , using Brooks' theorem again we can easily see that G will be k -colorable.

6. a) As a counter-example, consider C_5 in which $\chi(C_5) = 3$ yet contains no subgraph isomorphic to K_3 .
b) Consider the Wheel graph W_6 with $\chi(G) = 4$.



While W_6 has a subgraph homeomorphic to K_4 , it has no subgraph isomorphic to K_4 .

\therefore the statement is false.

- c) Since a $\chi(G) = 2$, G must be bipartite.

Since for G to be planar, using kuratowskis' theorem, G must have no subgraph that is homeomorphic to K_5 or $K_{3,3}$, but K_5 is not bipartite, so we only consider $K_{3,3}$.

Now consider if $G = K_{6,6}$

We know a bipartite graph $K_{m,n}$ has a Hamiltonian cycle iff $m = n$ so G has a Hamiltonian cycle.

We know that a graph with at least 2 vertices has an Eulerian circuit iff it is connected and every vertex is of even degree, so clearly G has a Eulerian circuit.

Clearly $K_{3,3}$ is an isomorphic subgraph of $K_{6,6}$ so $K_{3,3}$ is a homeomorphic subgraph to $K_{6,6}$.

$\therefore G = K_{6,6}$ is not planar.

\therefore the statement is false