



# CLOUD COMPUTING CONCEPTS

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## SNAPSHOTS

Lecture C

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CONSISTENT CUTS

# CUTS

- **Cut** = time frontier at each process and at each channel
- Events at the process/channel that happen before the cut are “in the cut”
  - And happening after the cut are “out of the cut”

# CONSISTENT CUTS

**Consistent Cut:** a cut that obeys causality

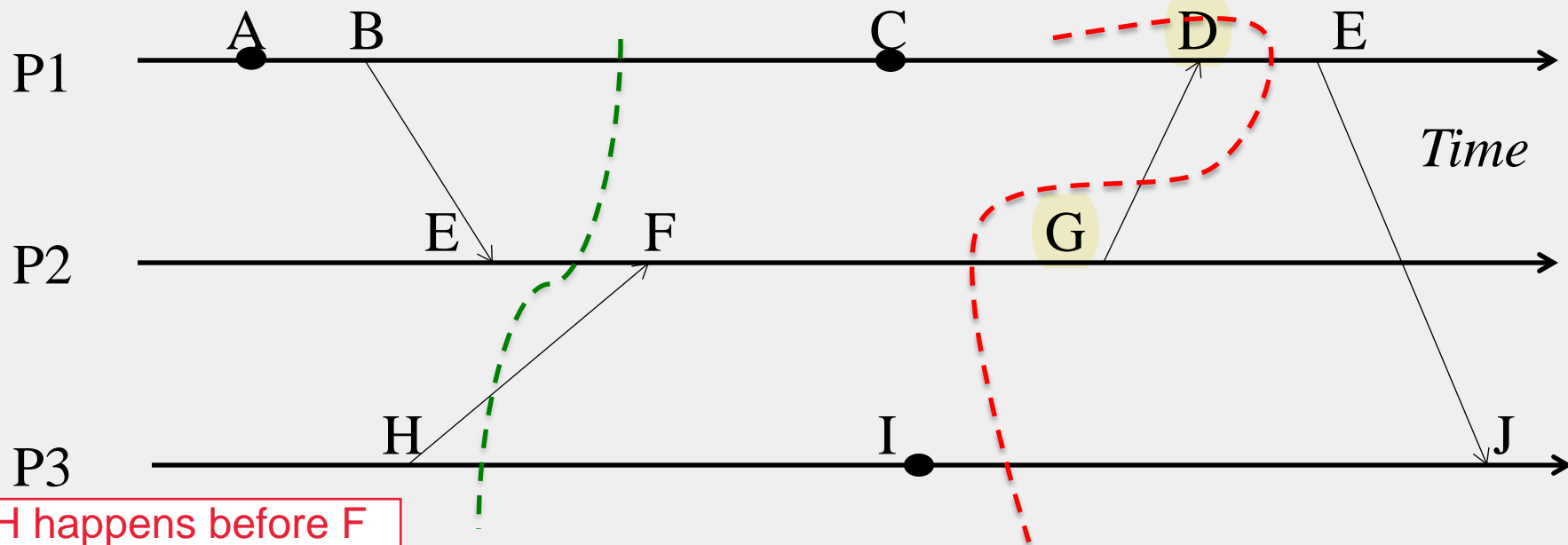
- A cut  $C$  is a consistent cut if and only if:
  - for (each pair of events  $e, f$  in the system)
    - Such that event  $e$  is in the cut  $C$ , and if  $f \rightarrow e$  ( $f$  happens-before  $e$ )
      - Then: Event  $f$  is also in the cut  $C$

# EXAMPLE

G == send event

D == receive event

inconsistent == D in the cut G not. BUT G happens before D



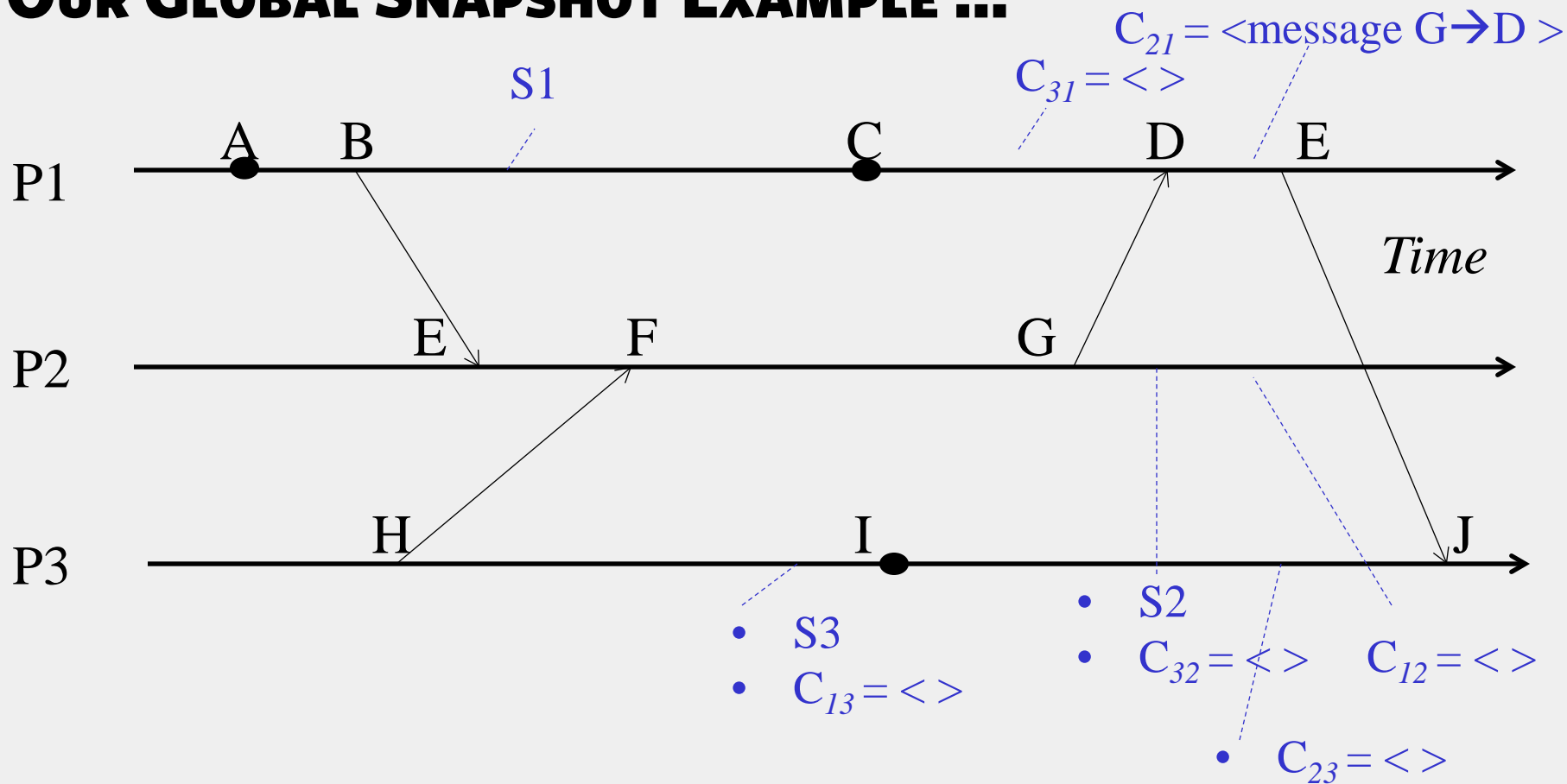
H happens before F  
== a cut thru them  
still consistent cut

Consistent Cut

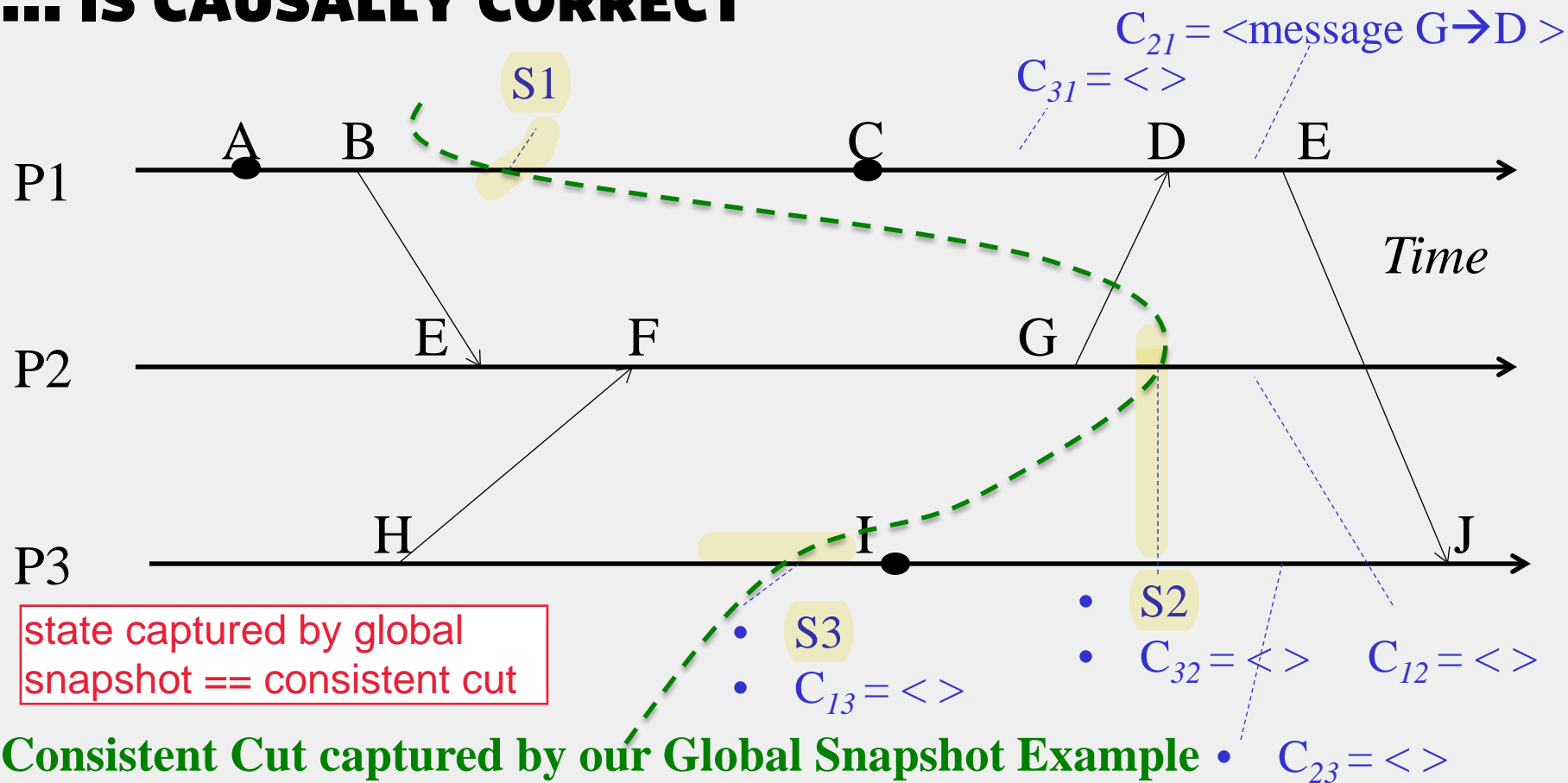
Inconsistent Cut

G → D, but only D is in cut

# OUR GLOBAL SNAPSHOT EXAMPLE ...



# ... IS CAUSALLY CORRECT



# IN FACT...

- Any run of the Chandy-Lamport Global Snapshot algorithm creates a consistent cut

# CHANDY-LAMPORT GLOBAL SNAPSHOT ALGORITHM CREATES A CONSISTENT CUT



Let's quickly look at the proof

- Let  $e_i$  and  $e_j$  be events occurring at  $P_i$  and  $P_j$ , respectively, such that
  - $e_i \rightarrow e_j$  ( $e_i$  happens before  $e_j$ )
- The snapshot algorithm ensures that
  - if  $e_j$  is in the cut then  $e_i$  is also in the cut.
- That is: if  $e_j \rightarrow \langle P_j \text{ records its state} \rangle$ , then
  - It must be true that  $e_i \rightarrow \langle P_i \text{ records its state} \rangle$ .



# CHANDY-LAMPORT GLOBAL SNAPSHOT ALGORITHM CREATES A CONSISTENT CUT

- If  $e_j \rightarrow \langle P_j \text{ records its state} \rangle$ , then it must be true that  $e_i \rightarrow \langle P_i \text{ records its state} \rangle$ .
  - By contradiction, suppose  $e_j \rightarrow \langle P_j \text{ records its state} \rangle$  and  $\langle P_i \text{ records its state} \rangle \rightarrow e_i$
  - Consider the path of app messages (through other processes) that go from  $e_i \rightarrow e_j$
  - Due to FIFO ordering, markers on each link in above path will precede regular app messages
  - Thus, since  $\langle P_i \text{ records its state} \rangle \rightarrow e_i$ , it must be true that  $P_j$  received a marker before  $e_j$
  - Thus  $e_j$  is not in the cut  $\Rightarrow$  contradiction

# NEXT

- What is the Chandy-Lamport algorithm used for?