Programming Abstractions

Week 8-1: MiniScheme C

What can MiniScheme do at this point?

MiniScheme B has constant numbers

MiniScheme B has pre-bound symbols that are in the init-env

Recall

```
(parse input) — Parses the input, at this point either a number or a variable, and returns a (lit-exp num) or (var-exp sym)
```

(eval-exp tree e) — Evaluates the parse tree in the environment e, returning a value

What does (parse 15) return?

- A. 15
- B. '(lit-exp 15)
- C. It's an error

What does (parse 'z) return?

C. It's an error

What does (eval-exp (var-exp 'z) environment) do?

- A. Returns what z is bound to in environment
- B. It's an error
- C. It looks up with z is bound to, returning the result or causing an error if z is not bound
- D. Something else

What does (eval-exp (lit-exp 108) environment) do?

- A. Returns what 108 is bound to in environment
- B. It's an error
- C. It looks up with 108 is bound to, returning the result or causing an error if 108 is not bound
- D. Returns 108
- E. Something else

Homeworks 6 and 7

Multiple steps, each adding parts to the MiniScheme interpreter

For each new type of expression

- Add a new data type
 - ift-exp
 - let-exp
 - etc.
- Add constructors, recognizers and accessors
- Modify parse to produce those
- Modify eval-exp to interpret them

```
EXP → number
      symbol
      (if EXP EXP EXP)
      (let(LET-BINDINGS)EXP)
      (letrec (LET-BINDINGS) EXP)
        lambda ( PARAMS ) EXP )
      ( set! symbol EXP )
      (begin EXP*)
      (EXP EXP^*)
LET-BINDINGS → LET-BINDING*
LET-BINDING → [ symbol EXP ]
PARAMS → symbol*
```

Let's add arithmetic and some list proceduresMiniScheme C

Let's add +, -, *, /, car, cdr, cons, etc.

Students find this to be the hardest part of the project

- It's the first complex part
- It contains some things that make more sense later, once we add lambda expressions

Many ways to call procedures

```
(+ 2 3)
((lambda (x y) (+ x y) 2 3)
(let ([f +]) (f 2 3)
```

The parser can't identify primitive procedures like + because symbols like f may be bound to primitive procedures

It can't tell because the parser does not have access to the environment

All that the parser can do is recognize a procedure application and parse

- the procedure; and
- the arguments

Enter lists

So far, the input to MiniScheme A and B has just been a number or a symbol

If the input is a list, then the kind of expression it represents depends on the first element

- If the first element is 'lambda, it's a lambda expression
- If the first element is 'let, it's a let expression
- ► If the first element is 'if, it's an if-then-else expression
- etc.

Applications don't have keywords, so any nonempty list for which the first element is not one of our supported keywords is an application

Procedure applications

MiniScheme C

An app-exp is a new data type that stores

- The parse tree for a procedure
- A list of parse trees for the arguments

Procedures to implement

- (app-exp proc args)
- (app-exp? exp)
- (app-exp-proc exp)
- (app-exp-args exp)

Recursive implementation Parsing

Expressions are recursive: $EXP \rightarrow (EXP EXP^*)$

When parsing an application expression, you want to parse the sub expressions using parse

How should you parse the arguments?

```
Consider input that looks like ((lambda (x y) x) 2 3) or (f 4 5 6)
```

The procedure part can be parsed with (parse (first input))

How should you parse the arguments?

Evaluating an app-exp

Evaluate the procedure part

Evaluate each of the arguments

If the procedure part evaluates to a primitive procedure, call a procedure you'll write that will perform the operation on the arguments

► E.g., if the primitive procedure is *, then you'll want to call * on the arguments

The tricky part is what does it mean to evaluate the procedure part?

Evaluating the procedure part of an app-exp

Consider the input '(+ 2 3 4)

The procedure part is '+ which will be parsed as '(var-exp +)

Variable reference expressions are evaluated by looking the symbol up in the current environment

Therefore, we need our initial environment to contain a binding for the symbol '+ (and all the other primitive procedures we want to support)

prim-proc data type

We can create a new data type prim-proc

```
  (prim-proc symbol)
```

- (prim-proc? value)
- (prim-proc-symbol value)

Adding primitives to our initial environment

```
(define primitive-operators
 '(+ - * /))
(define prim-env
  (env primitive-operators
       (map prim-proc primitive-operators)
       empty-env))
(define init-env
  (env '(x y) '(23 42) prim-env))
```

Evaluating an app-exp

Recall: app-exp stores the parse tree for the procedure and a list of parse trees for the arguments

We need to evaluate all of those; add something like the following to eval-exp

Applying a procedure

The apply-proc procedure takes an evaluated procedure and a list of evaluated arguments

It can look at the procedure and determine if it's a primitive procedure

- If so, it will call apply-primitive-op
- If not, it's an error for now; later, we'll add code to deal with non-primitive procedure (i.e., normal lambdas)

Applying primitive operations

(apply-primitive-op op args)

apply-primitive-op takes a symbol (such as '+ or '*) and a list of arguments

You probably want something like

```
(define (apply-primitive-op op args)
  (cond [(eq? op '+) (apply + args)]
       [(eq? op '*) (apply * args)]
       ...
       [else (error ...)]))
```

What is returned by (parse '(* 2 3))?

```
A. '((prim-proc *) 2 3)
B. '((prim-proc *) (lit-exp 2) (lit-exp 3))
C. '(app-exp (prim-proc *) ((lit-exp 2) (lit-exp 3)))
D. '(var-exp * (lit-exp 2) (lit-exp 3))
E. '(app-exp (var-exp *) ((lit-exp 2) (lit-exp 3)))
```

When evaluating an app-exp, the procedure and each of the arguments are evaluated. For example, when evaluating the result of (parse '(- 20 5)), there will be three recursive calls to eval-exp, the first of which is evaluating (var-exp '-).

What is the result of evaluating (var-exp '-)?

- B. '(app-exp -)
- C. '(prim-proc -)
- D. It's an error because requires arguments

What is the result of (eval-exp (parse '(* 4 5)) init-env)?

```
A. 20
```

- B. '(app-exp (var-exp *) ((lit-exp 4) (lit-exp 5)))
- C. '(prim-proc * 4 5)
- D. '(prim-proc (var-exp *) (lit-exp 4) (lit-exp 5))
- E. '(app-exp (prim-proc *) 4 5)

Why go to all that trouble?

In a later version of MiniScheme, we'll implement lambda

We'll deal with this by adding a line to apply-proc that will apply closures

Adding other primitive procedures

In addition (pardon the pun) to +, -, *, and /, you'll add several other primitive procedures

- add1
- ▶ sub1
- negate
- ► list
- cons
- car
- cdr

And you'll add a new variable null bound to the empty list

What does (car (list 3 5 2)) parse to?

What does (car (list 3 5 2)) parse to?

Adding additional primitive procedures

- 1. Add the procedure name to primitive-operators
- 2. Add a corresponding line to the cond in apply-primitive-op

```
E.g.,
[(eq? op 'car) (car (first args))]
[(eq? op 'list) args]
```

What can MiniScheme C do?

Numbers

Pre-defined variables

Procedure calls to built-in procedures