## Programming Abstractions

Lecture 5: Variations on let

#### Announcements

Office hours 14:00–15:00 tomorrow (half hour later than usual!)

Homework 1 due Friday

# What values does this code return? (define (foo x)

- A. 10
- B. 11
- C. 12

- D. Some other value
- E. Error

#### 

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- B. 11
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- D. Some other value
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## A common problem

When writing programs, it's not uncommon to define some local variables in terms of other local variables

Example: Return the elements of a list of numbers that are at least as large as the first element (the head) of the list, in reverse order

This doesn't work; we can't use head in the definition of bigger

#### The issue

The issue is the scope of the binding for head: just the body of the let

One (bad) work around would be to use multiple lets

## Sequential let

## Returning to our example

#### A more realistic example

Write a procedure (split-by pred lst) that splits lst into two lists, the first contains all of the elements that match pred, the second contains all the elements that do not match pred

```
(split-by even? (range 10)) => '((0 2 4 6 8) (1 3 5 7 9))
(split-by (\lambda (x) (< x 3)) (range 5)) => '((0 1 2) (3 4))
(define (split-by pred 1st)
```

## Another problem: recursion

Often, we're going to want to define a recursive procedure but we can't do that with let or let\*

We can't use fact in the definition of fact

#### Recursive let

```
(letrec ([id1 s-exp1] [id2 s-exp2]...) body)
```

All of the s-exps can refer to all of the ids

This is used to make recursive procedures

## Recursive let drawback (subtle)

The values of the identifiers we're binding can't be used in the bindings

Invalid (the value of x is used to define y)

Valid (the *value* of x isn't used to define y, it's only used when y is called)