CSCI 210: Computer Architecture
Lecture 8: Computer Representation of MIPS Instructions

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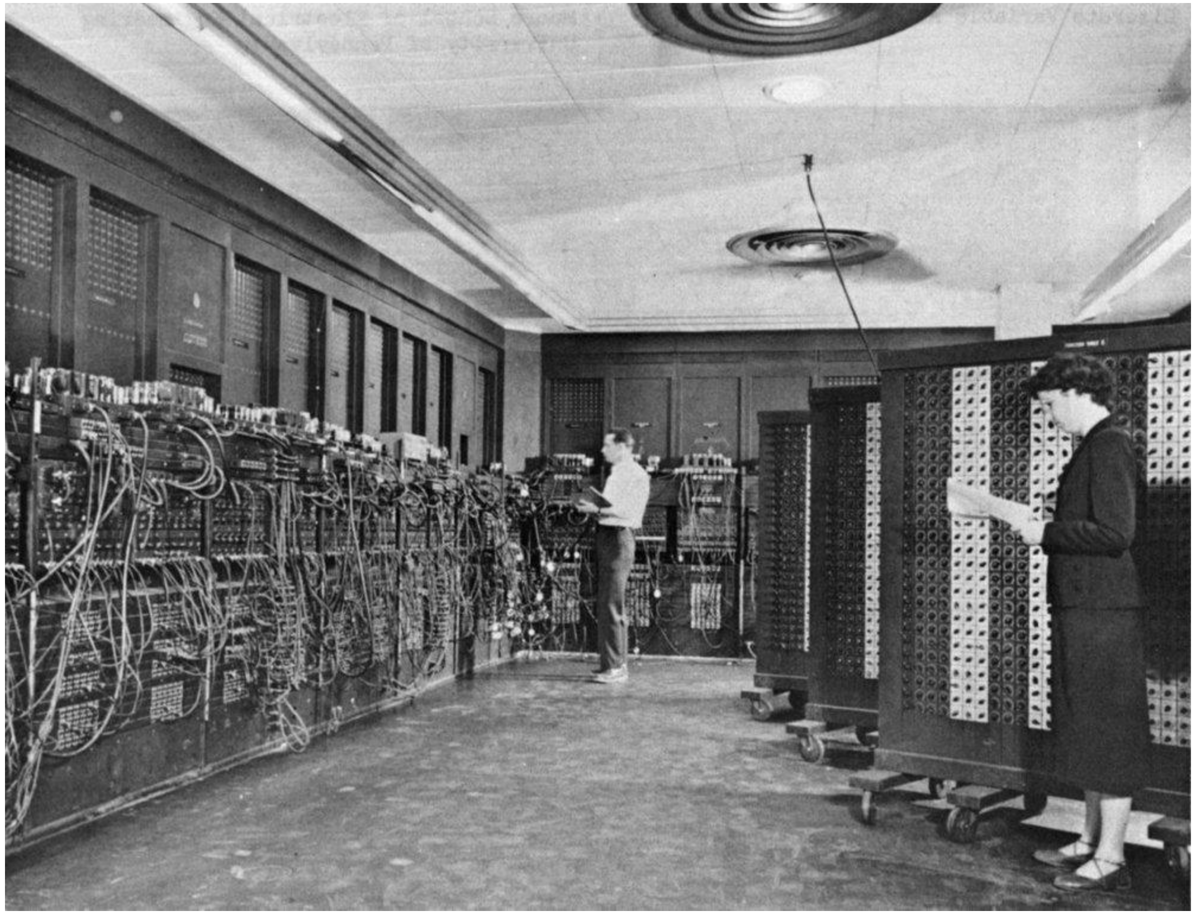
Announcements

- Problem Set 2 due Friday

- Lab 1 available now

CS History: ENIAC



U.S. Army photo of ENIAC

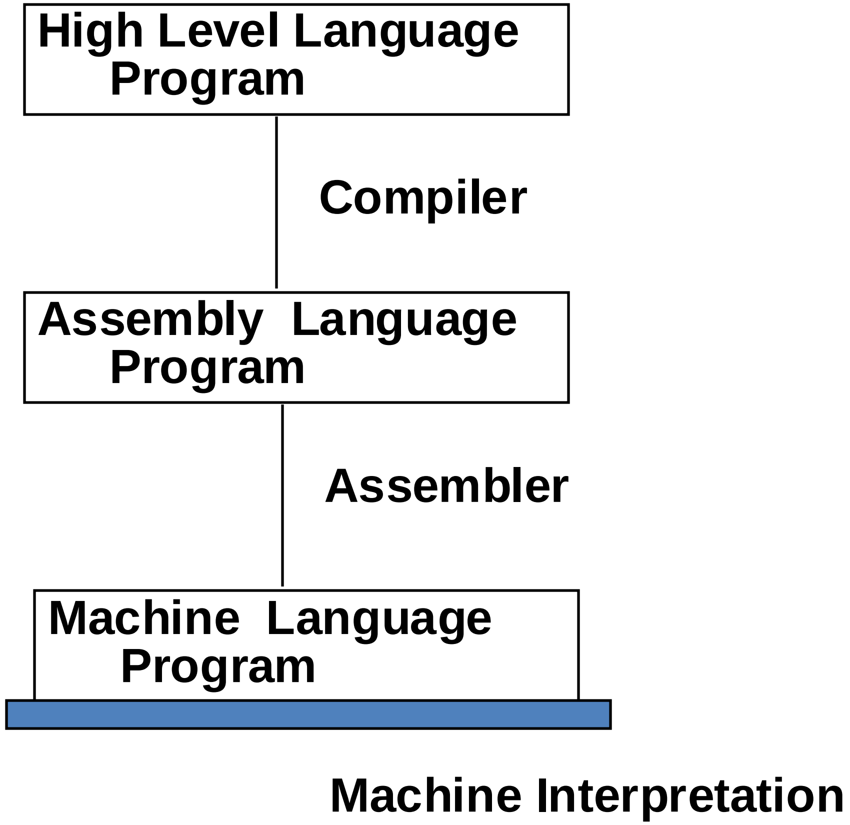
- Electronic Numerical Integrator And Computer
- First programmable, electronic, general-purpose computer
- Created by the US Army in 1945
- Designed to compute ballistic tables during WWII
- Originally didn't have storage
- Decimal, not binary!

CS History: ENIAC

- Programmers were Kay McNulty, Jean Bartok, Betty Snyder, Marlyn Meltzer, Fran Bilas, and Ruth Lichterman.
- Selected from a group of 200 women employed hand calculating equations for the army
- Programmed by connecting components with cables and setting switches
- Kay McNulty developed the use of subroutines
- Betty Snyder and Jean Bartok went on to help develop the first commercial computers

U.S. Army photo

How to Speak Computer




```
temp = v[k];
v[k] = v[k+1];
v[k+1] = temp;
```

```
lw  $15,   0($2)
lw  $16,   4($2)
sw  $16,   0($2)
sw  $15,   4($2)
```

```
10001100011000100000000000000000
10001100111100100000000000000100
10101100111100100000000000000000
10101100011000100000000000000100
```

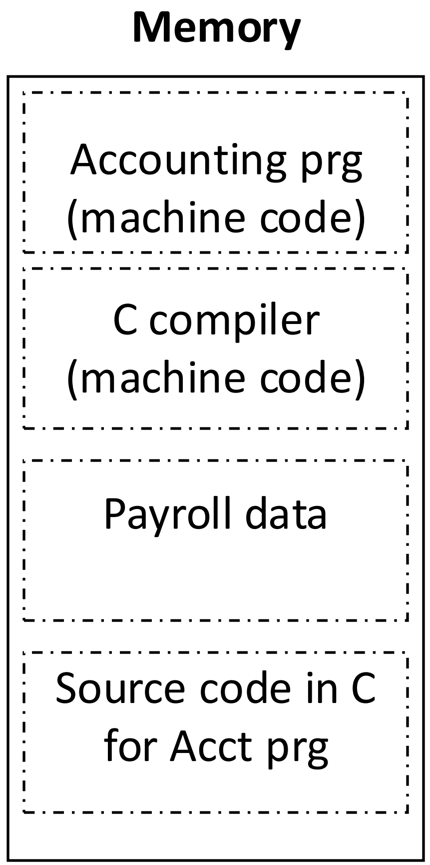
Two Key Principles of Machine Design

1. Instructions are represented as numbers and, as such, are indistinguishable from data

2. Programs are stored in alterable memory (that can be read or written to) just like data

Stored-program concept

- Programs can be shipped as files of binary numbers – binary compatibility
- Computers can inherit ready-made software provided they are compatible with an existing ISA and OS – leads industry to align around a small number of ISAs



What happens if someone writes new machine code in the memory where your program is stored, overwriting your program?

A. The program will crash.

B. The old instructions will run.

C. The new instructions will run.

D. None of the above

Recall: Instruction Set Architecture

- Definition of how to access the hardware from software
- Supported instructions, registers, etc . . .

Key ISA decisions

destination operand *operation*

y = x + b

source operands

- operations
  - how many?
  - which ones
- operands
  - how many?
  - location
  - types
- instruction format
  - size
  - how many formats?

add r1, r2, r5

how does the computer know what
0001 0100 1101 1111
means?

RISC versus CISC (Historically)

- Complex Instruction Set Computing
  - Larger instruction set
  - More complicated instructions built into hardware
  - Variable number of clock cycles per instruction

- Reduced Instruction Set Computing
  - Small, highly optimized set of instructions
  - Memory accesses are specific instructions
  - One instruction per clock cycle (only the very first RISCs!)

A = A*B

RISC

```
lw    $t0, 0(A)
lw    $t1, 0(B)
mul   $s1, $t0, $t1
sw    $s1, 0(A)
```

CISC

```
mul   B, A
```

Which of these is faster?

RISC

```
lw    $t0, 0(A)
lw    $t1, 0(B)
mul   $s1, $t0, $t1
sw    $s1, 0(A)
```

CISC

```
mul   B, A
```

RISC vs CISC

RISC

- More work for compiler/assembly programmer
- More RAM used to store instructions
- Less complex hardware

CISC

- Less work for compiler/assembly programmer
- Fewer instructions to store
- More complex hardware

So . . . Which System “Won”?

- Most processors are RISC
- BUT the x86 (Intel) is CISC
- x86 breaks down CISC assembly into multiple, RISC-like, machine language instructions
- Distinction between RISC and CISC is less clear
  - Some RISC instruction sets have more instructions than some CISC sets

The computer figures out what format an instruction is from

A. Codes embedded in the instruction itself.

B. A special register that is loaded with the instruction.

C. It tries each format and sees which one forms a valid instruction.

D. None of the above

Instruction Formats What does each bit mean?

- Having many different instruction formats...
  - complicates decoding
  - uses more instruction bits (to specify the format)

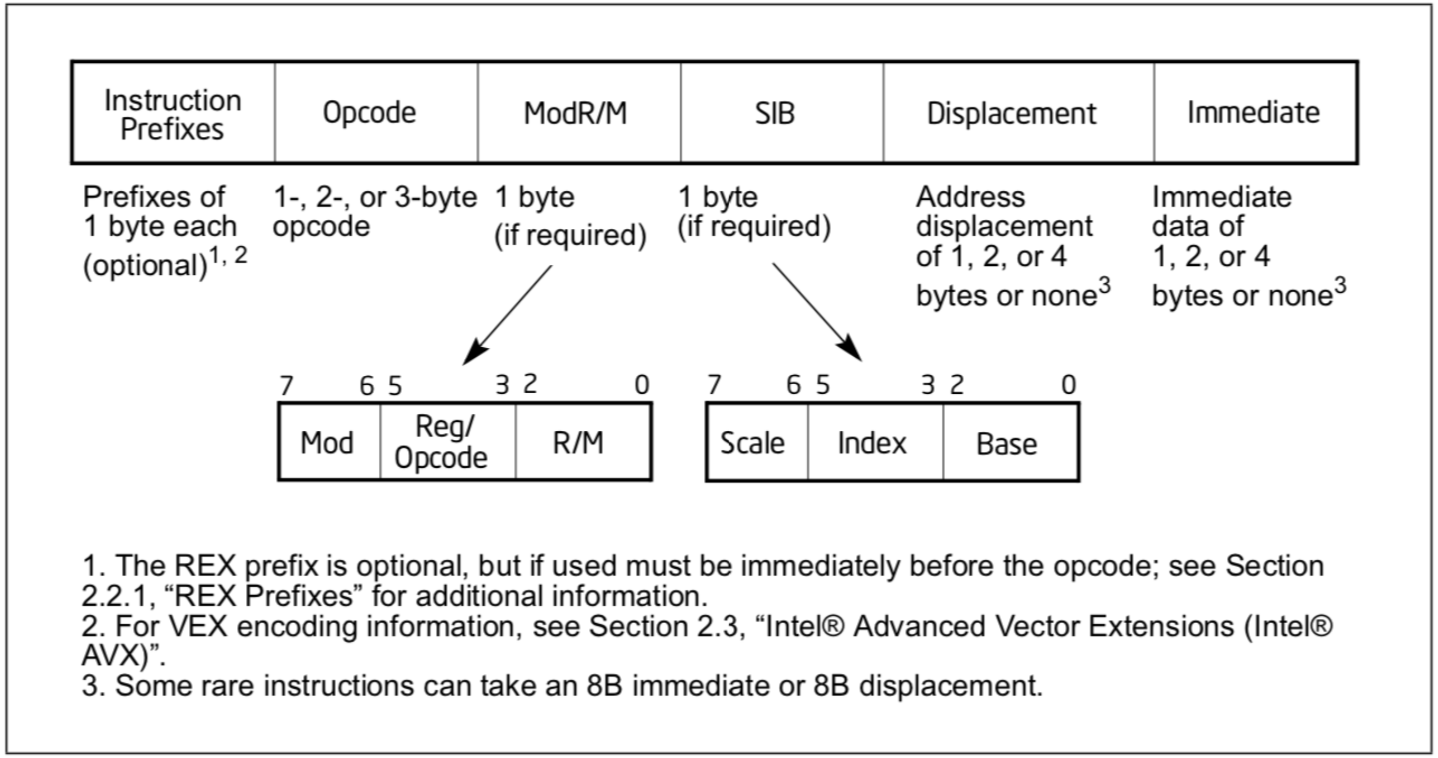


Figure 2-1. Intel 64 and IA-32 Architectures Instruction Format

Representing Instructions

- MIPS instructions
  - Encoded as 32-bit instruction words
  - Small number of formats encoding operation code (opcode), register numbers, …
  - Regularity!

	6 bits	5 bits	5 bits	5 bits	5 bits	6 bits
R-type	**opcode**	**rs**	**rt**	**rd**	**sa**	**funct**
I-type	**opcode**	**rs**	**rt**	**immediate**		
J-type	**opcode**	**target**				

Your architecture supports 16 instructions and 16 registers (r0–r15). You have fixed width instructions which are 16 bits. How many register operands can you specify (explicitly) in an add instruction?

A. ≤ 1 operand
B. ≤ 2 operands
C. ≤ 3 operands
D. ≤ 4 operands
E. None of the above

Hint: Remember you need to specify which instruction it is, and all the registers

MIPS Instruction Formats

	6 bits	5 bits	5 bits	5 bits	5 bits	6 bits
R-type	**opcode**	**rs**	**rt**	**rd**	**sa**	**funct**
I-type	**opcode**	**rs**	**rt**	**immediate**		
J-type	**opcode**	**target**				

Which row contains correct examples of instructions with the given types?

	R-type	I-type
A	addi	sw
B	addi	sub
C	add	sw
D	add	sub
E	None of the above	

MIPS Instruction Fields

- MIPS fields are given names to make them easier to refer to

op	rs	rt	rd	shamt	funct

op	6-bits	opcode that specifies the operation
rs	5-bits	register file address of the first source operand
rt	5-bits	register file address of the second source operand
rd	5-bits	register file address of the result’s destination
shamt	5-bits	shift amount (for shift instructions)
funct	6-bits	function code augmenting the opcode

MIPS Arithmetic Instructions Format

t0 = s1 – s2

```
sub $t0, $s1, $s2
```

0	17	18	8	0	0x22
opcode	**rs**	**rt**	**rd**	**sa**	**funct**

R-format Example

op	rs	rt	rd	shamt	funct
6 bits	5 bits	5 bits	5 bits	5 bits	6 bits

```
add $t0, $s1, $s2
```

CORE INSTRUCTION SET

NAME, MNEMONIC		FOR-MAT	OPERATION (in Verilog)		OPCODE / FUNCT (Hex)
Add	add	R	R[rd] = R[rs] + R[rt]	(1)	$0 / 20_{hex}$
Add Immediate	addi	I	R[rt] = R[rs] + SignExtImm	(1,2)	8_{hex}
Add Imm. Unsigned	addiu	I	R[rt] = R[rs] + SignExtImm	(2)	9_{hex}
Add Unsigned	addu	R	R[rd] = R[rs] + R[rt]		$0 / 21_{hex}$

NAME	NUMBER	USE
$zero	0	The Constant Value 0
$at	1	Assembler Temporary
$v0-$v1	2-3	Values for Function Results and Expression Evaluation
$a0-$a3	4-7	Arguments
$t0-$t7	8-15	Temporaries
$s0-$s7	16-23	Saved Temporaries
$t8-$t9	24-25	Temporaries
$k0-$k1	26-27	Reserved for OS Kernel
$gp	28	Global Pointer
$sp	29	Stack Pointer
$fp	30	Frame Pointer
$ra	31	Return Address

Reading

- Next lecture: Bitwise Operations
  - Section 2.7
- Problem Set 2 due Friday
- Lab 1 due Monday