CS 241: Systems Programming Lecture 18. System Calls I

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What is an operating system?

An **Operating System (OS)** is software that manages all of the resources of a computer. It acts an interface between software and hardware.

Operating system tasks

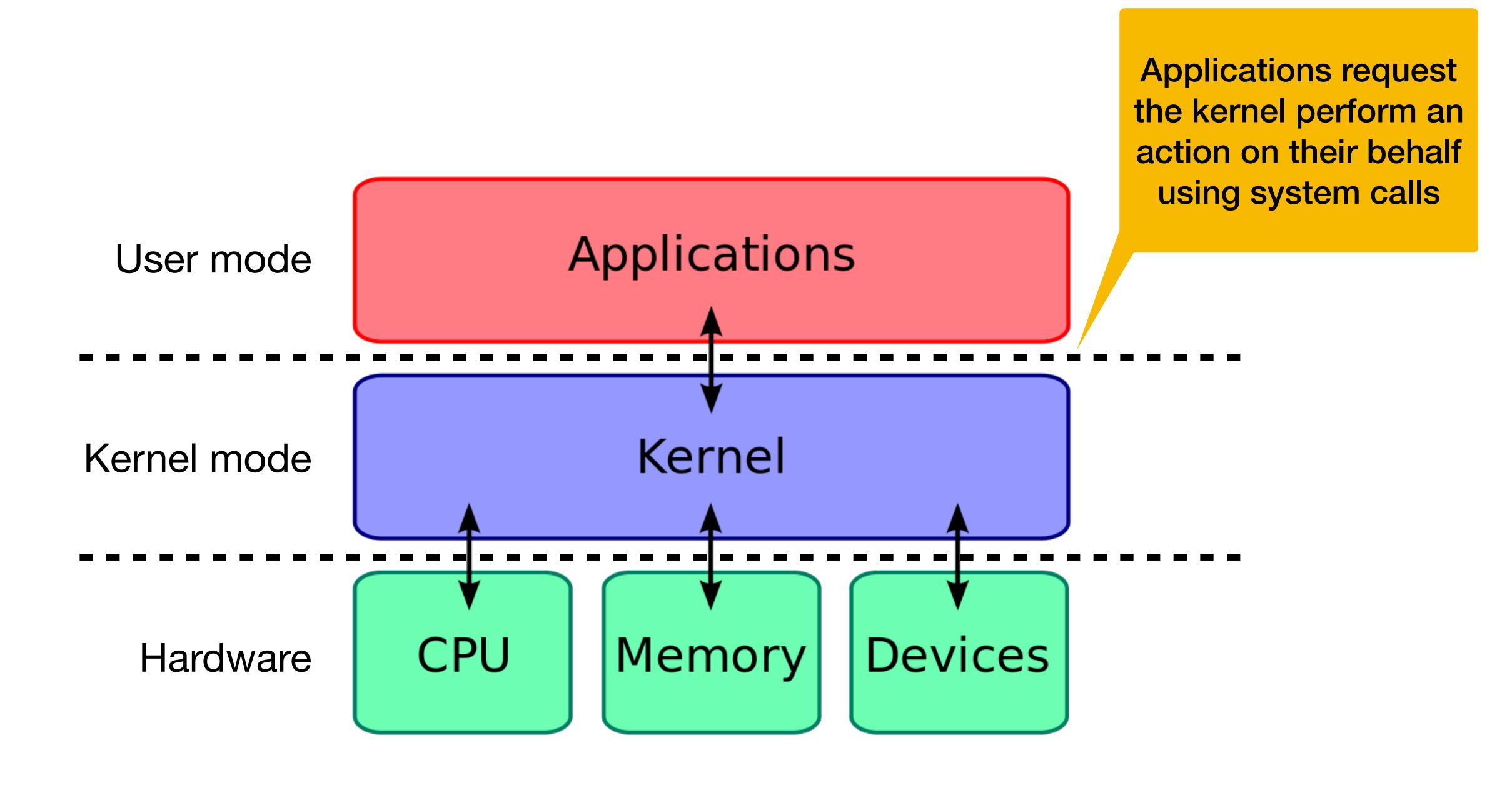
Managing the resources of a computer

CPU, memory, network, etc.

Coordinate the running of all other programs

OS can be considered as a set of programs

kernel – name given to the core OS "program"



https://en.wikipedia.org

Do we need an operating system?

A. Yes

B. No

C. I don't know/I'm not sure

System calls

Programs talk to the OS via system calls

- Set of functions to request access to resources of the machine
- System calls vary by operating system and computer architecture

Types of system calls

- Input/output (may be terminal, network, or file I/O)
- File system manipulation (e.g., creating/deleting files/directories)
- Process control (e.g., process creation/termination)
- Resource allocation (e.g., memory)
- Device management (e.g., talking to USB devices)
- Inter-process communication (e.g., pipes and sockets)

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Most basic UNIX system call: exit

Programs (normally) end by calling exit() or returning from main()... which calls exit()

The exit system call takes an exit status as its only parameter

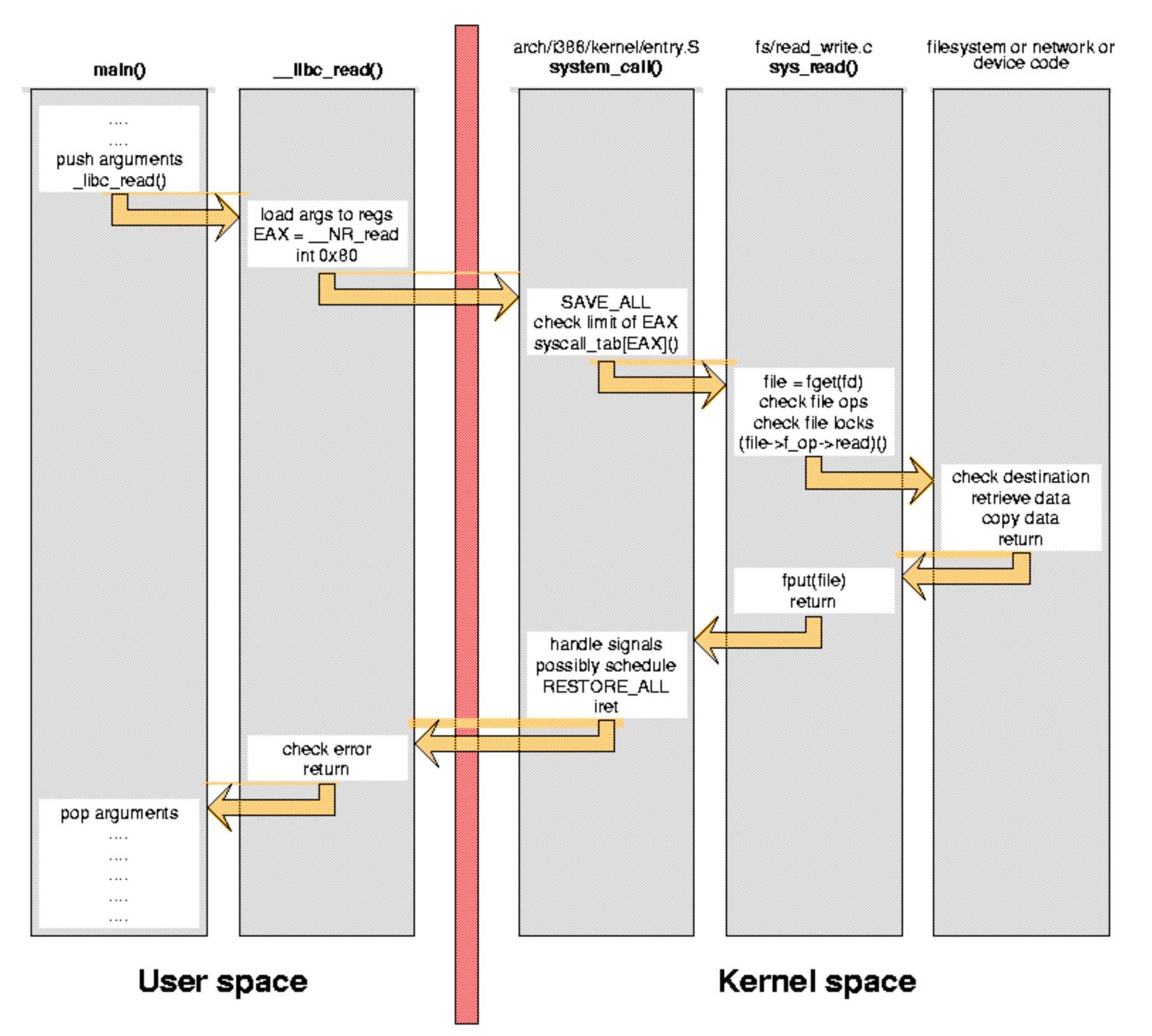
When the kernel receives an exit system call from a program, it

- cleans up all of the resources associated with the program
- notifies the program that created the exiting program (the parent) that a child has exited

System calls as API

System calls are an example of an application programming interface (API)

- Each system call is assigned a small integer (the system call number)
- System calls are performed by setting up the arguments (often in registers) and using a dedicated "system call" or "interrupt" instruction
- The kernel's system call handler calls an appropriate function based on the system call number
- Data (and success/failure) is returned to the application



http://www.linux.it/~rubini/docs/ksys/

System calls and libc

C standard library

- Some functions make no system calls (e.g., strcpy (3))
- Some functions "wrap" a single system call (e.g., open (2))
- Some functions have complex behavior and might make a variable number of system calls (e.g., malloc(3))

We're going to focus on the libc wrappers for the system calls

These live in section 2 of the manual: open(2), _exit(2), fork(2)

System calls and Rust

OS vendors make changes to their system calls over time

Different computer architectures use different system call numbers

To deal with this, the system call interface lives in libc:

- To make a system call, applications call functions in libc
- libc places the system call number and arguments in the correct registers and traps into the kernel

In Rust, we have two options

- 1. Use higher-level functionality provided by the standard library
- 2. Call functions in libc

Unsafe Rust

Rust lets you be unsafe

Rust has an unsafe keyword that lets you perform unsafe operations

- Call functions marked as unsafe (including everything in the libc)
- Dereference raw pointers (we'll talk about these shortly)
- Modify a mutable global variable
- Implement an unsafe trait (we'll talk about traits in a few lectures)
- Access fields of a C-style union

To make system calls, we'll need unsafe for the first two

The purpose of unsafe

The compiler (and language) is conservatively correct

It ensures that the programs (that don't use unsafe) are memory safe

It rejects programs that are safe due to its inability to prove them safe

unsafe provides a way to bypass those limitations

unsafe limits the scope of where memory errors can occur to precisely those regions of the code marked unsafe

Unsafe functions/methods

Functions and methods can be marked as unsafe by using the unsafe keyword

```
Unsafe functions can only be called from within an unsafe block (or
unsafe function)
unsafe fn does_unsafe_things() -> i32 { 0 }

fn main() {
    let x = unsafe {
        does_unsafe_things()
    };
    println!("{x}");
}
```

Functions in other languages

Rust can call functions in other languages (usually C functions)

All such external functions are unsafe and can only be called from unsafe blocks

Why did the Rust designers require that functions written in other languages be called from an unsafe block?

A. Select A when you have an answer

Raw pointers

We've seen pointers in Rust

- Shared references (e.g., &i32)
- Mutable references (e.g., &mut i32)
- Boxes
- Slices

Rust has two additional pointer types

- Constant pointer (e.g., *const i32)
- Mutable pointer (e.g., *mut i32)

Pointers

Pointers are like their reference counterparts but without some restrictions

References must always point to valid, aligned objects of the appropriate type

Additionally, mutable references may not be aliased

Pointers may be invalid (including null) or point to a misaligned object

Mutable pointers may alias

Alignment

Alignment of a value refers to its address in memory

An **aligned** value is one whose address is a multiple of its size in bytes (at least for primitive types like i32 or usize, structs are aligned at the alignment of their largest member)

A misaligned value is one whose address is not a multiple of its size (or its largest member)

Rust (and most programming languages) require values be aligned

This restriction comes from hardware which often doesn't support misaligned memory reads/writes or performs them more slowly

Creating a pointer from a reference

```
fn pointer_stuff(ptr: *const i32) { }
fn main() {
    let x = 10;
    // Cast the reference to a pointer
    let p = &x as *const i32;
    pointer_stuff(p);
    // Implicit conversion
    pointer_stuff(&x);
```

Creating a mutable pointer

```
fn pointer_stuff(ptr: *const i32) { }
fn main() {
    let mut x = 10;
    // Cast the mutable reference to a mutable pointer
    let p = \&mut x as *mut i32;
    // Implicit conversion from *mut i32 to *const i32
    pointer_stuff(p);
    // Implicit conversion from &mut i32 to *const i32
    pointer_stuff(&mut x);
```

The reason to create a pointer is to read or write the memory it points to

```
fn main() {
    let mut x = 10;
    let mut y = 20;
    let x_ptr = &mut x as *mut _;
    let y_ptr = &mut y as *mut _;
   unsafe {
        let tmp = *x_ptr; // Read
        *x_ptr = *y_ptr; // Read + write
        *y_ptr = tmp; // Write
  // after
```

The _ causes the compiler to use type inference to determine the type, in this case: *mut i32

What will x and y be at the "after" line in this code?

	X	Y
a	10	20
b	10	10
С	20	20
d	20	10

Reading or writing values

The reason one wants to create a pointer is to read or write the memory it points to

```
fn main() {
    let mut x = 10;
    let mut y = 20;
    let x_ptr = &mut x as *mut _;
    let y_ptr = &mut y as *mut _;
    println!("Before: x={x} y={y}");
    unsafe {
        let tmp = *x_ptr; // Read
        *x_ptr = *y_ptr; // Read + write
        *y_ptr = tmp; // Write
    println!("After: x={x} y={y}");
```

The _ causes the compiler to use type inference to determine the type, in this case: *mut i32

```
Output:
Before: x=10 y=20
After: x=20 y=10
```

Pointer from a slice (or Vec or String)

```
let v: Vec<char> = vec!['\vec'; 1000];
let p: *const char = v.as_ptr();

let s = String::from("Pointers!");
let p: *const u8 = s.as_ptr();
```

Strings hold their characters UTF-8 encoded in a Vec<u8>

Gives a pointer to the first element of the slice

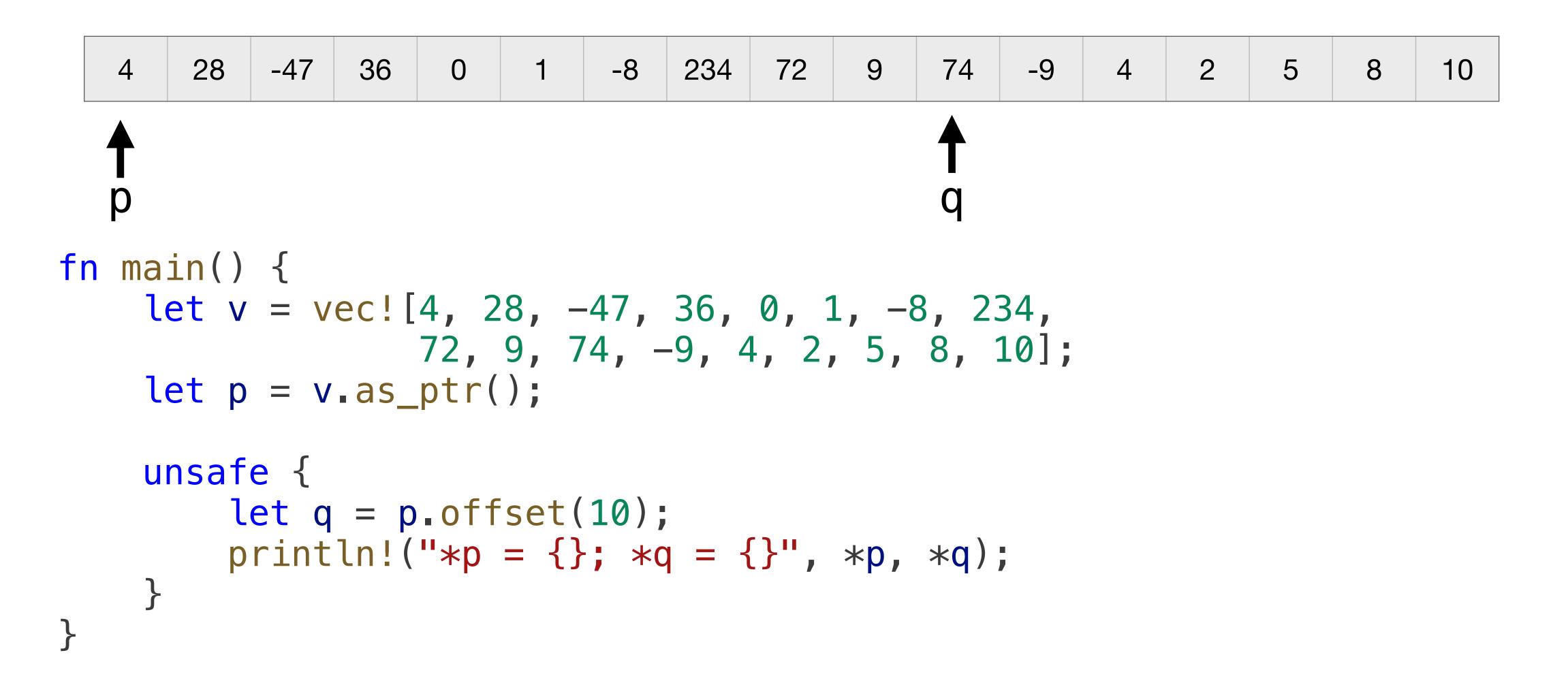
Use .as_mut_ptr() to get a *mut _ rather than *const _

Is the .as_ptr() method necessary or can we just cast the reference?

```
let s = String::from("Pointers!");
let p = &s as *const u8;
```

- A. .as_ptr() is necessary
- B. Casting the reference also works
- C. .as_ptr() is necessary for a String but casting would work for a Vec

Creating pointers from other pointers



Null pointers

```
Use std::ptr::null() and std::ptr::null_mut() to create
*const _ or *mut _

Use .is_null() to test if a pointer is null

let ptr: *mut i32 = std::ptr::null_mut();
println!("{}", ptr.is_null());
```

libc crate

The libc crate exposes libc functions/types/constants in Rust

System call wrapper in libc	Rust declaration of the function	Notes
<pre>void _exit(int status)</pre>	<pre>fn _exit(status: c_int) -> !</pre>	Exit (doesn't return)
<pre>pid_t getpid(void)</pre>	<pre>fn getpid() -> pid_t</pre>	Get the process ID
<pre>ssize_t read(int fildes, void *buf,</pre>	<pre>fn read(fd: c_int, buf: *mut c_void,</pre>	Read data from a file
<pre>int rename(const char *old,</pre>	<pre>fn rename(oldname: *const c_char,</pre>	Renames files

4

Types of arguments

Arguments to syscalls fall into a few basic types

C type	Libc crate's equivalent	Normal Rust equivalent (on many platforms)	Notes
int	c_int	i32	Normal integer
size_t/ssize_t	size_t/ssize_t	usize/isize	Represents the size of things ssize_t is used to return -1 as an error
char *	*const c_char *mut c_char	&CStr CString *const i8/*mut i8	0-byte terminated string
void *	*const c_void *mut c_void		A pointer to anything
Pointers to structs	Pointers to structs with the same name		References can be converted to pointers

C strings

C considers a string to be a sequence of nonzero (often signed) bytes followed by a byte with value 0

Here's "Hello eg"



Rust considers a string to be a sequence of u8 of UTF-8 encoded characters and an associated length

$$len = 17$$

There's no difference between -16 and 240 other than interpretation. Both have binary value 11110000

The kernel uses C strings

The kernel, being written in C, uses C strings

More importantly, the system call interface uses C strings

Converting between a C string and a Rust string isn't difficult but can be subtle

- Who owns the data?
- Is the string from the kernel valid UTF-8?

&CStr and CString

```
To pass a Rust string to the kernel, use std::ffi::CString
let cstr = CString::new(some_str)?;
let ptr = cstr.as_ptr();
To convert a C string into a Rust string, use std::ffi::CStr
let normal_str = CStr::from_ptr(ptr).to_str()?.to_string();
CString::new() will return an Err(err) if some_str contains a 0 byte
■to_str() will return an Err(err) if the CStr points to non UTF-8 data
```

Example: Home directories

```
struct passwd *getpwnam(const char *name);
struct passwd {
   char *pw_name;
                       /* username */
   char *pw_passwd;
                     /* user password */
   uid_t pw_uid;
                     /* user ID */
                 /* group ID */
        pw_gid;
   gid_t
                        /* user information */
   char
         *pw_gecos;
   char *pw_dir;
                   /* home directory */
        *pw_shell;
                   /* shell program */
   char
```

Declarations in Rust's libc crate

```
pub unsafe extern "C" fn getpwnam(name: *const c_char)
-> *mut passwd;
#[repr(C)]
pub struct passwd {
    pub pw_name: *mut c_char,
    pub pw_passwd: *mut c_char,
    pub pw_uid: uid t,
    pub pw_gid: gid_t,
    pub pw_gecos: *mut c_char,
    pub pw_dir: *mut c_char,
    pub pw_shell: *mut c_char,
```

Handling errors

When a system call fails

- the C wrapper returns -1 (or NULL, in the case of returning a pointer)
- the per-thread global variable errno is set to an integer specifying the reason

Rust's std::io::last_os_error() reads errno and constructs a std::io::Error which we can use with a Result

```
use std::ffi::{CStr, CString};
type Result<T> = std::result::Result<T, Box<dyn std::error::Error>>;
fn get_home_dir_for_user(user: &str) -> Result<String> {
    let user = CString::new(user)?;
    unsafe {
        let pwd: *const libc::passwd = libc::getpwnam(user.as_ptr());
        if pwd.is_null() {
            return Err(std::io::Error::last_os_error().into());
        if (*pwd).pw_dir.is_null() {
            return Err("No home directory found".into());
        let home_dir = CStr::from_ptr((*pwd).pw_dir).to_str()?.to_string();
        Ok(home_dir)
```

Why do we use system calls instead of making a function call directly to the function in the kernel that will handle our system call request?

Discuss with your group and select A on your clickers when you have a reason (or multiple reasons)

Input/output system calls

Open a file: open(2)

#include <fcntl.h>

... indicates 0 or more additional arguments. In this case, open() takes exactly 2 or 3 arguments

```
int open(char const *path, int oflag, ...);
```

```
    O_RDONLY open for reading only
```

O WRONLY open for writing only

O RDWR open for reading and writing

O_APPEND append on each write

► O_TRUNC truncate size to 0

O_CREAT create file if it does not exist

► O_EXCL error if O_CREAT and the file exists

► O_NONBLOCK do not block on open or for data to become available

Last arg is the "int mode" -- see chmod(2) and umask(2)

Returns file descriptor on success, -1 on error

Bitwise OR the flags together,
e.g.,
O_WRONLY | O_CREAT

File descriptors

Integer index into OS file table for this process

3 are automatically created for you

- STDIN_FILENO 0 standard input
- STDOUT_FILENO 1 standard output
- STDERR_FILENO 2 standard error

These are what are used in shell redirection

\$./a.out 2> errors.txt

Read data: read(2)

```
#include <unistd.h>
```

```
ssize_t read(int fildes, void *buf, size_t nbyte);
```

- Attempts to read nbytes from filedes storing data in buf
- Returns the number of bytes read
- ► Upon **EOF**, returns 0
- ► Upon error, returns –1 and sets errno

Write data: write(2)

```
#include <unistd.h>
```

```
ssize_t write(int fildes, void const *buf, size_t nbyte);
```

- Attempts to write nbyte of data to the object referred to by filedes from the buffer buf
- Upon success, returns number of bytes are written
- On error, returns -1 and sets errno

In what scenario would we return less than the number of bytes specified?

- A. Reading until the end of the file
- B. Writing to the end of the file
- C. Reading from a network
- D. Writing to a network
- E. More than one of the above (which ones?)

Seek in file: Iseek(2)

```
#include <sys/types.h>
#include <unistd.h>

off_t lseek(int fd, off_t offset, int whence);
    whence is one of SEEK_SET, SEEK_CUR, SEEK_END
```

- On success, returns the resultant offset in terms of bytes from the beginning of the file
- On error, returns (off_t)-1 and sets errno

Close files: close(2)

Reading a file with system calls

- 1. Open the file with libc::open() and handle errors
- 2. Reserve space in a Vec<u8>
- 3. Read some data with libc::read() and handle errors
- 4. If all of the data was not read, go back to step 2
- 5. Close the file with libc::close()

Why do we have an open() call, as opposed to just read()ing or write()ing a file path and checking permissions each time?

- A. To improve security
- B. To improve performance
- C. It makes interacting with the hard drive easier
- D. Two of the above
- E. All of the above

Opening the file

```
use std::ffi::CString;
use std::io;
fn read_file(path: &str) -> io::Result<Vec<u8>>> {
    let path = CString::new(path)?;
                                                     Construct a 0-terminated
    let mut data: Vec<u8> = Vec::new();
                                                            C string
    unsafe {
        let fd = libc::open(path.as_ptr(), libc::0_RDONLY);
        if fd == -1  {
            return Err(io::Error::last_os_error());
        // Read the data here
        libc::close(fd);
    Ok(<u>data</u>)
```

Reserve space

Read some data

```
fn read_file(path: &str) -> io::Result<Vec<u8>>> {
    // ...
        loop {
             let ptr: *mut libc::c_void = data.as_mut_ptr()
                 .offset(data.len() as isize)
                 .cast();
             let amount = libc::read(fd, ptr, <u>data</u>.capacity() - <u>data</u>.len());
             if amount < 0 {</pre>
                 let err = io::Error::last_os_error();
                 libc::close(fd);
                                                  Easy to forget to close
                 return Err(err);
                                                         the file!
                amount == 0 \{
                 break;
             data.set_len(data.len() + amount as usize);
    Ok(data)
                                            51
```

```
fn read_file(path: &str) -> io::Result<Vec<u8>>> {
    let path = CString::new(path)?;
    let mut data: Vec<u8> = Vec::new();
    unsafe {
        let fd = libc::open(path.as_ptr(), libc::0_RDONLY);
        if fd == -1 {
            return Err(io::Error::last_os_error());
        loop {
            if data.capacity() - data.len() < 4096 {</pre>
                 data_reserve(4096);
             let ptr: *mut libc::c_void = data.as_mut_ptr().offset(data.len() as isize).cast();
             let amount = libc::read(fd, ptr, <u>data</u>.capacity() - <u>data</u>.len());
            if amount < 0 {</pre>
                 let err = io::Error::last_os_error();
                 libc::close(fd);
                 return Err(err);
               amount == 0 \{
                 break;
            data.set_len(data.len() + amount as usize);
        libc::close(fd);
    Ok(data)
```

Contrast with normal Rust

```
fn read_file(path: &str) -> io::Result<Vec<u8>>> {
    use std::io::Read;
    let mut file = File::open(path)?;
    let mut data = Vec::new();
    file.read_to_end(&mut data)?;
    Ok(data)
}

1 or more read system calls
```

close system call when file is dropped

Or even easier

```
fn main() {
    let data1 = read_file("example.txt").unwrap();
    let data2 = std::fs::read("example.txt").unwrap();
    assert_eq!(data1, data2);
}
```

One function to call.

It'll call open(), read(), and close()

```
#include <errno.h>
                                                               errno = old_errno;
#include <fcntl.h>
                                                               return NULL;
#include <stdio.h>
#include <stdlib.h>
                                                          data = new data;
#include <unistd.h>
                                                      ssize_t amount = read(fd,
void *read_file(char const *path,
                                                                             &data[len],
                                                                             cap - len);
                size t *len ptr)
                                                      if (amount < 0) {
    int fd = open(path, 0_RDONLY);
                                                           int old_errno = errno;
                                                          free(data);
    if (fd == -1) {
                                                          close(fd);
        return NULL;
                                                          errno = old_errno;
    *len_ptr = 0;
                                                          return NULL;
    char *data = NULL;
                                                      if (amount == 0) {
    size_t len = 0;
    size_t cap = 0;
                                                          break;
    while (1) {
        if (cap - len < 4096) {
                                                       len += amount;
           cap += 4096;
           char *new_data = realloc(data,
                                                  close(fd);
                                                  *len_ptr = len;
                                     cap);
            if (new_data == NULL) {
                                                  return data;
                int old_errno = errno;
                free(data);
                close(fd);
                                             55
```

File system manipulation system calls

Delete files: unlink(2)

Rename files: rename(2)

```
#include <stdio.h>
```

```
int rename(char const *oldpath, char const *newpath);
```

- Renames oldpath to newpath, returns 0 on success
- ► Returns –1 and sets errno on error
- This can change directories, but not file systems!

Get current directory: getcwd(3)

```
#include <unistd.h>
```

```
char *getcwd(char *buf, size_t size);
```

- Copies absolute path of current working directory to buf
 - length of array is "size"
 - if path is too long (including null byte), NULL/ERANGE
- Linux allows NULL for buf for dynamic allocation, see man page

Basically just a wrapper around the getcwd system call plus some memory allocations

Change directories: chdir(2)

```
#include <unistd.h>
int chdir(const char *path);
int fchdir(int fildes);
```

Change working directory of calling process

- How "cd" is implemented
- fchdir() is only in certain standards, but widely available
- fchdir() lets you return to a directory referenced by a file descriptor from open(2)ing a directory

0 on success, -1/errno on error

Create/delete a directory

```
#include <sys/stat.h>
#include <sys/types.h>
int mkdir(char const *path, mode_t mode);

    Create a directory called path

  Don't forget execute bits in mode!
#include <unistd.h>
int rmdir(char const *path);

    Delete the directory specified by path

0 for success, -1/errno on error
```

Reading directories

```
opendir(3), readdir(3), closedir(3)
```

Enables the application to read the contents of directories

These are actually just higher-level wrappers around open(2), getdirents(2), and close(2) which are themselves wrappers around the corresponding system calls

Libc crate and normal Rust

The libc crate declares all of these functions

The std::fs module has Rust-versions

- remove_file() for unlink(2)
- rename()
- create_dir() for mkdir(2)
- remove_dir() for rmdir(2)
- read_dir() for opening, reading, and closing directories

The std::env module has some other related functions

- current_dir() for getcwd(2)
- set_current_dir() for chdir(2)

What are some reasons you would programmatically create, delete, and change directories?

A. Discuss and select any letter on your clicker when you are done