Programming Abstractions

Lecture 19: MiniScheme B and C

What can MiniScheme do at this point?

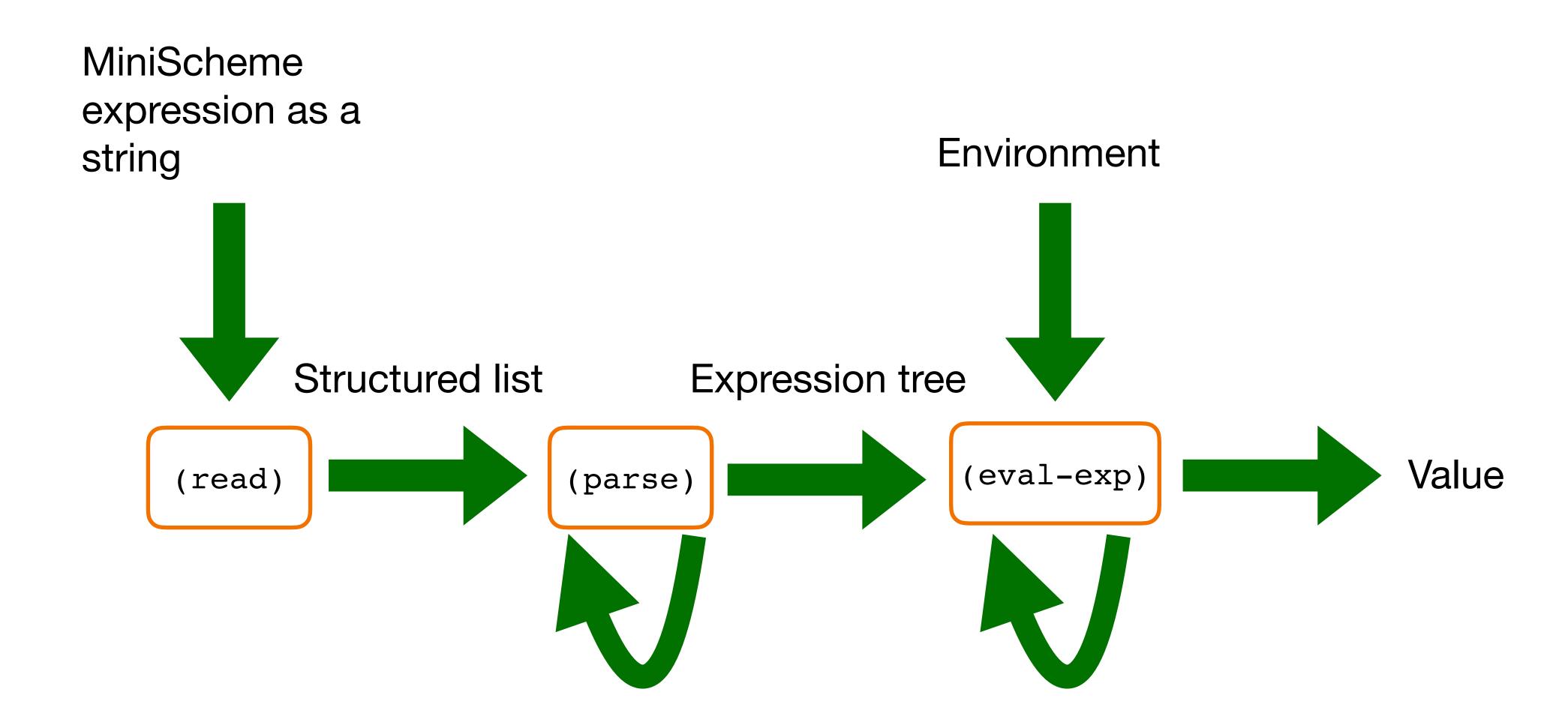
MiniScheme A has constant numbers

Recall

```
(parse input) — Parses the input, at this point only numbers, and returns a (lit-exp num)

(eval-exp tree e) — Evaluates the parse tree in the environment e, returning a value
```

Interpreter flow



Let's add some variables!

MiniScheme B

```
Grammar

EXP → number parse into lit-exp

symbol parse into var-exp
```

Data type for a variable reference expression

```
(struct var-exp (symbol) #:transparent)
  (var-exp symbol)
  (var-exp? exp)
  (var-exp-symbol exp)
```

Parsing symbols

MiniScheme B

```
(define (parse input)
  (cond [(number? input) (lit-exp input)]
        [(symbol? input) (var-exp input)]
        [else (error 'parse "Invalid syntax ~s" input)]))
When I run (parse 'foo), I get
  (var-exp 'foo)
```

Interpreting symbols

MiniScheme B

```
(define (eval-exp tree e)
  (cond [(lit-exp? tree) (lit-exp-num tree)]
        [(var-exp? tree)
         (env-lookup e (var-exp-symbol tree))]
        [else (error 'eval-exp "Invalid tree: ~s" tree)]))
You'll need a working env-lookup
> (env-lookup init-env 'x)
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> (eval-exp (var-exp 'x) init-env)
```

Assuming that x is bound to 10 and y to 25 in init-env, what does (parse 'x) return (assuming the implementation discussed so far)?

- A. 10
- B. 25
- C. (lit-exp 10)
- D. (var-exp 'x)
- E. It's an error of some sort

Assuming that x is bound to 10 and y to 25 in init-env, what does (eval-exp (parse 'x) init-env) return (assuming the implementation discussed so far)?

- A. 10
- B. 25
- C. (lit-exp 10)
- D. (var-exp 'x)
- E. It's an error of some sort

What can MiniScheme do at this point?

MiniScheme B has constant numbers

MiniScheme B has pre-bound symbols that are in the init-env

Let's add arithmetic and some list procedures MiniScheme C

Let's add +, -, *, /, car, cdr, cons, etc.

Students find this to be the hardest part of the project

- It's the first complex part
- It contains some things that make more sense later, once we add lambda expressions

Many ways to call procedures

```
(+ 2 3)
((lambda (x y) (+ x y)) 2 3)
(let ([f +]) (f 2 3))
```

The parser can't identify primitive procedures like + because symbols like f may be bound to primitive procedures

It can't tell because the parser does not have access to the environment

All that the parser can do is recognize a procedure application and parse

- the procedure; and
- the arguments

Enter lists

So far, the input to MiniScheme A and B has just been a number or a symbol

If the input is a list, then the kind of expression it represents depends on the first element

- If the first element is 'lambda, it's a lambda expression
- If the first element is 'let, it's a let expression
- ► If the first element is 'if, it's an if-then-else expression
- etc.

Applications don't have keywords, so any nonempty list for which the first element is not one of our supported keywords is an application

Procedure applications

MiniScheme C

An app-exp is a new data type that stores

- The parse tree for a procedure
- A list of parse trees for the arguments

Data type procedures

- (app-exp proc args)
- (app-exp? exp)
- (app-exp-proc exp)
- (app-exp-args exp)

Recursive implementation Parsing

Expressions are recursive: $EXP \rightarrow (EXP EXP^*)$

When parsing an application expression, you want to parse the sub expressions using parse

How should you parse the arguments?

```
Consider input that looks like ((lambda (x y) x) 2 3) or (f 4 5 6)
```

The procedure part can be parsed with (parse (first input))

How should you parse the arguments?

What is the result of (parse '(foo x y z))?

E. It's an error because the variables foo, x, y, and z aren't defined

What is the result of (parse '(foo (add1 x))?

D. It's an error