

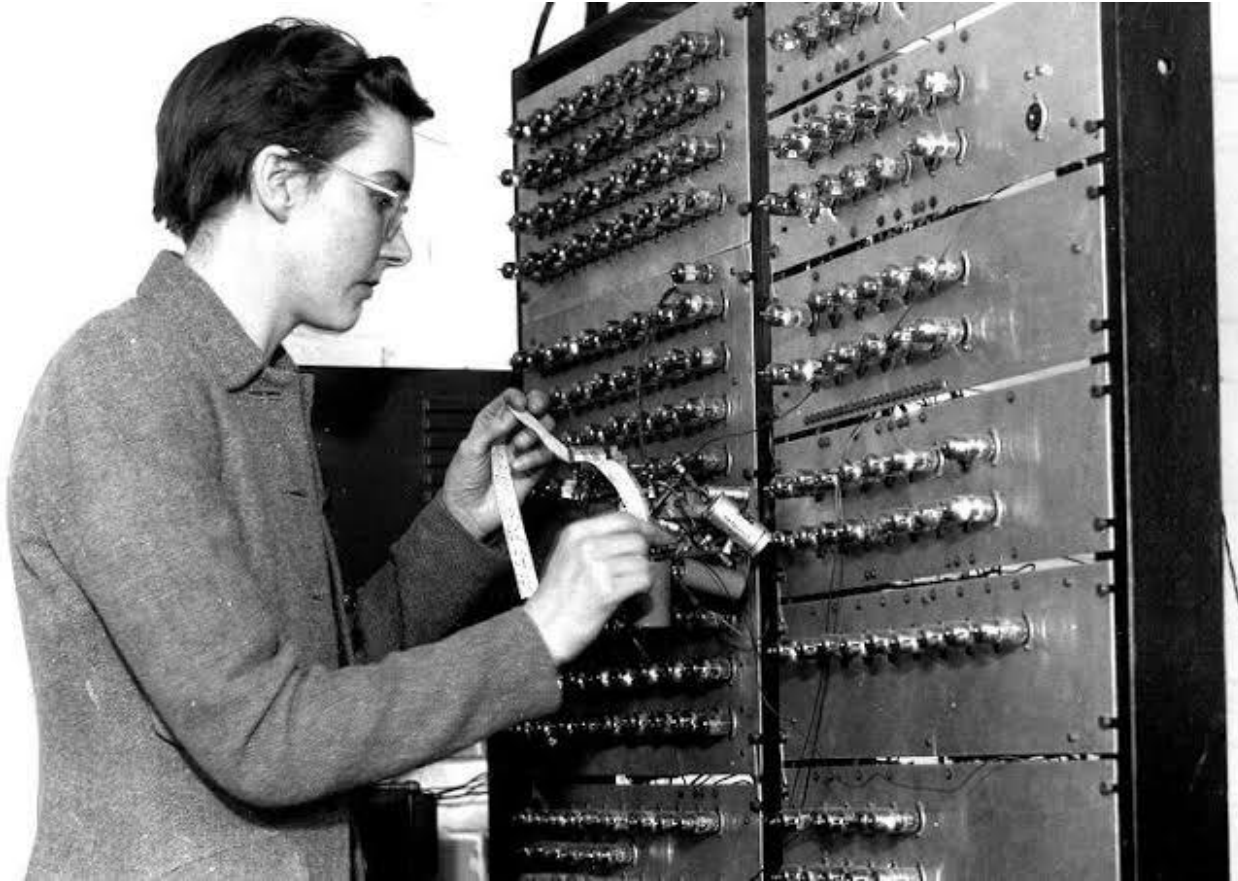
# CSCI 210: Computer Architecture

## Lecture 9: Logical Operations

Stephen Checkoway

Slides from Cynthia Taylor

# CS History: Kathleen Britton



- Applied mathematician and computer scientist
- Wrote the first assembly language and assembler in 1947
- Collaborated with Andrew Booth to develop three early computers: the ARC (Automatic Relay Calculator), SEC (Simple Electronic Computer), and APE(X)C
- Later worked with neural nets

# Logical Operations

- Instructions for bitwise manipulation

Operation	C	Java	MIPS
Shift left	<<	<<	sll
Shift right	>>	>>>	srl
Bitwise AND	&	&	and, andi
Bitwise OR			or, ori
Bitwise NOT	~	~	nor

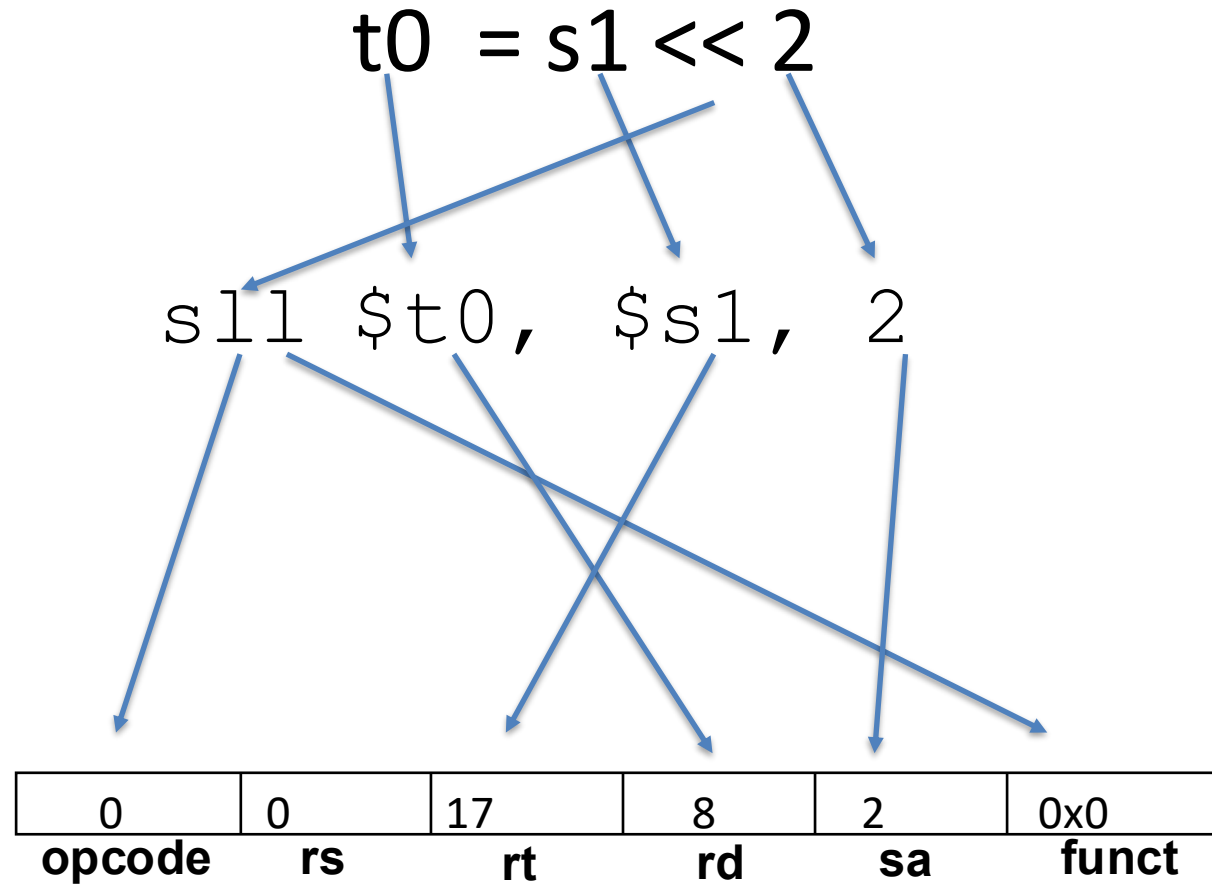
- Useful for extracting and inserting groups of bits in a word

# Shift Operations



- shamt: how many positions to shift
- Shift left logical
  - Shift left and fill with 0 bits
  - **sll by n bits multiplies by  $2^n$**
- Shift right logical
  - Shift right and fill with 0 bits
  - **srl by n bits divides by  $2^n$  (unsigned only)**

# MIPS shift instructions



# Shift left logical

- 0110 1001 << 2 in 8 bits
  - Most significant 2 bits are **dropped**
  - 2 0s are **added** to become the least significant bits
  - Result: ~~01~~ 1010 01**00** => 1010 0100

# Shift right logical

- 1010 1001 >>> 3 in 8 bits
  - Least significant 3 bits are **dropped**
  - 3 0s are **added** to become the most significant bits
  - Result: **000**1 0101 ~~001~~ => 0001 0101

# Shift right arithmetic

- sra rd, rt, shamt
  - Shift right and copy the sign bit
- 1010 1001 >> 3 in 8 bits
  - Least significant 3 bits are **dropped**
  - 3 1s are **added** because the MSB is 1 to become the most significant bits
  - Result: **1111** 0101 ~~001~~ => 1111 0101



A new op HEXSHIFTRIGHT shifts hex numbers right by a digit. HEXSHIFTRIGHT  $i$  times is equivalent to

- A. Dividing by  $i$
- B. Dividing by  $2^i$
- C. Dividing by  $16^i$
- D. Multiplying by  $16^i$

# Remember Boolean Operations?

- and, or, not . . .
- Now we'll apply them to bits!
- Just think of 1 as True, and 0 as False

# And Truth Table

	<i><b>0</b></i>	<i><b>1</b></i>
<i><b>0</b></i>	0	0
<i><b>1</b></i>	0	1

# AND Operations

- Useful to mask bits in a word
  - Select some bits, clear others to 0

and \$t0, \$t1, \$t2

\$t2    0000 0000 0000 0000 0000 1101 1100 0000

\$t1    0000 0000 0000 0000 0011 1100 0000 0000

\$t0    0000 0000 0000 0000 0000 1100 0000 0000

# AND identities (for a single bit)

- $x \& 0 =$
- $x \& 1 =$

$$\begin{array}{r} 01101001 \\ \& 11000111 \end{array}$$

A. 00010000

B. 01000001

C. 10101110

D. 11101111

If we want to zero out bits\* 3 – 0 in a byte we should AND with

A. 00000000

\*MSB (bit 7) is on the left,  
LSB (bit 0) is on the right

B. 00001111

C. 11110000

D. 11111111

One way to represent colors is to specify the amount of **red**, **green**, and **blue** (RGB) that makes up the color. The three color “channels” are often packed into a 32-bit integer as follows:

00000000 **rrrrrrrr** **gggggggg** **bbbbbbbb**

where each color is represented in 8 bits. If `color` is a 32-bit RGB color, which high-level expression extracts the green channel as an 8-bit value in the range 0–255?

A. `green = (color >> 8) & 0xFF`

B. `green = (color >> 16) & 0xFF`

C. `green = (color & 0xFF) >> 8`

D. `green = (color & 0xFF00) >> 16`

E. More than one of the above (which?)



Assume `color` is stored in `$t0` and `green` should be extracted to `$t1`. Which sequence of MIPS instructions corresponds to  $\text{green} = (\text{color} \gg 8) \ \& \ 0\text{xFF}$

A. `sll $t1, $t0, 8`  
    `andi $t1, $t1, 0xFF`

B. `srl $t1, $t0, 8`  
    `andi $t1, $t1, 0xFF`

C. `andi $t1, $t0, 0xFF`  
    `sll $t1, $t1, 8`

D. `andi $t1, $t0, 0xFF`  
    `srl $t1, $t1, 8`

# Or Truth Table

	<i><b>0</b></i>	<i><b>1</b></i>
<i><b>0</b></i>	0	1
<i><b>1</b></i>	1	1

# OR Operations

- Useful to set bits in a word
  - Set some bits to 1, leave others unchanged

or \$t0, \$t1, \$t2

\$t2    0000 0000 0000 0000 0000 1101 1100 0000

\$t1    0000 0000 0000 0000 0011 1100 0000 0000

\$t0    0000 0000 0000 0000 0011 1101 1100 0000

# OR Identities (for a single bit)

- $x \mid 0 =$

- $x \mid 1 =$

01101001  
| 11000111

A. 00010000

B. 01000001

C. 10101110

D. 11101111

Recall RGB:

00000000 rrrrrrrr gggggggg bbbbbbbb

If  $r$ ,  $g$ , and  $b$  are values in the range 0–255, how can we construct the RGB value,  $c$ , whose channels are  $r$ ,  $g$ , and  $b$ ?

A.  $c = (r \ll 16) \mid (g \ll 8) \mid b$

B.  $c = (r \ll 24) \mid (g \ll 16) \mid (b \ll 8)$

C.  $c = (r \gg 8) \mid (g \gg 16) \mid (b \gg 24)$

D.  $c = (r \gg 16) \mid (g \gg 24) \mid (b \gg 32)$

E. More than one of the above (which?)

# Nor Truth Table

	<i><b>0</b></i>	<i><b>1</b></i>
<i><b>0</b></i>	1	0
<i><b>1</b></i>	0	0

# NOR Operations

- MIPS has NOR 3-operand instruction
  - $a \text{ NOR } b = \text{NOT } (a \text{ OR } b)$

```
nor $t0, $t1, $t2
```

\$t2    0000 0000 0000 0000 0000 1101 1100 0000

\$t1    0000 0000 0000 0000 0011 1100 0000 0000

\$t0    1111 1111 1111 1111 1100 0010 0011 1111



01101001  
NOR 11000111

A. 00010000

B. 01000001

C. 10101110

D. 11101111

# NOT operations

- Inverts all the bits in a word
  - Change 0 to 1, and 1 to 0

MIPS does not need a NOT operation because we can use \_\_\_\_\_ for NOT \$t1, \$t2

A. `nor $t1, $t2, $zero`

B. `nor $t1, $t2, $t3`, where all bits in \$t3 are set to 1

C. `nori $t1, $t2, 0b11111111_11111111`, where nori is Nor Immediate

D. It does require a NOT operation

E. None of the above are correct

# XOR Truth Table

	<i><b>0</b></i>	<i><b>1</b></i>
<i><b>0</b></i>	0	1
<i><b>1</b></i>	1	0

# XOR Operations

- Exclusive OR (written  $x \oplus y$  or  $x \wedge y$ )
  - Set bits to one only if they are not the same

`xor $t0, $t1, $t2`

\$t2    0000 0000 0000 0000 0000 1101 1100 0000

\$t1    0000 0000 0000 0000 0011 1100 0000 0000

\$t0    0000 0000 0000 0000 0011 0001 1100 0000

01101001  
XOR 11000111

A. 00010000

B. 01000001

C. 10101110

D. 11101111

# XOR Identities (for a single bit)

- $x \text{ XOR } 0 =$
- $x \text{ XOR } 1 =$

10 & 7

- A. 0
- B. 2
- C. 7
- D. 10
- E. None of the above



Set bit 4 in byte x to 1, leaving the rest of the bits unchanged

A.  $x = x \text{ AND } 00010000$

B.  $x = x \text{ AND } 11101111$

C.  $x = x \text{ OR } 00010000$

D.  $x = x \text{ NOR } 11101111$

# Invert bits 2–0 of x

A.  $x = x \text{ AND } 00000111$

B.  $x = x \text{ OR } 00000111$

C.  $x = x \text{ NOR } 00000111$

D.  $x = x \text{ XOR } 00000111$

Find the ones' complement of x (in 8 bits)

A.  $x \text{ XOR } 00000000$

B.  $x \text{ XOR } 11111111$

C.  $x \text{ XOR } 11111110$

D.  $x \text{ OR } 11111111$

# Reading

- Next lecture: Branching instructions