

CSCI 210: Computer Architecture

Lecture 30: Pipelining the Datapath

Stephen Checkoway

Oberlin College

December 17, 2021

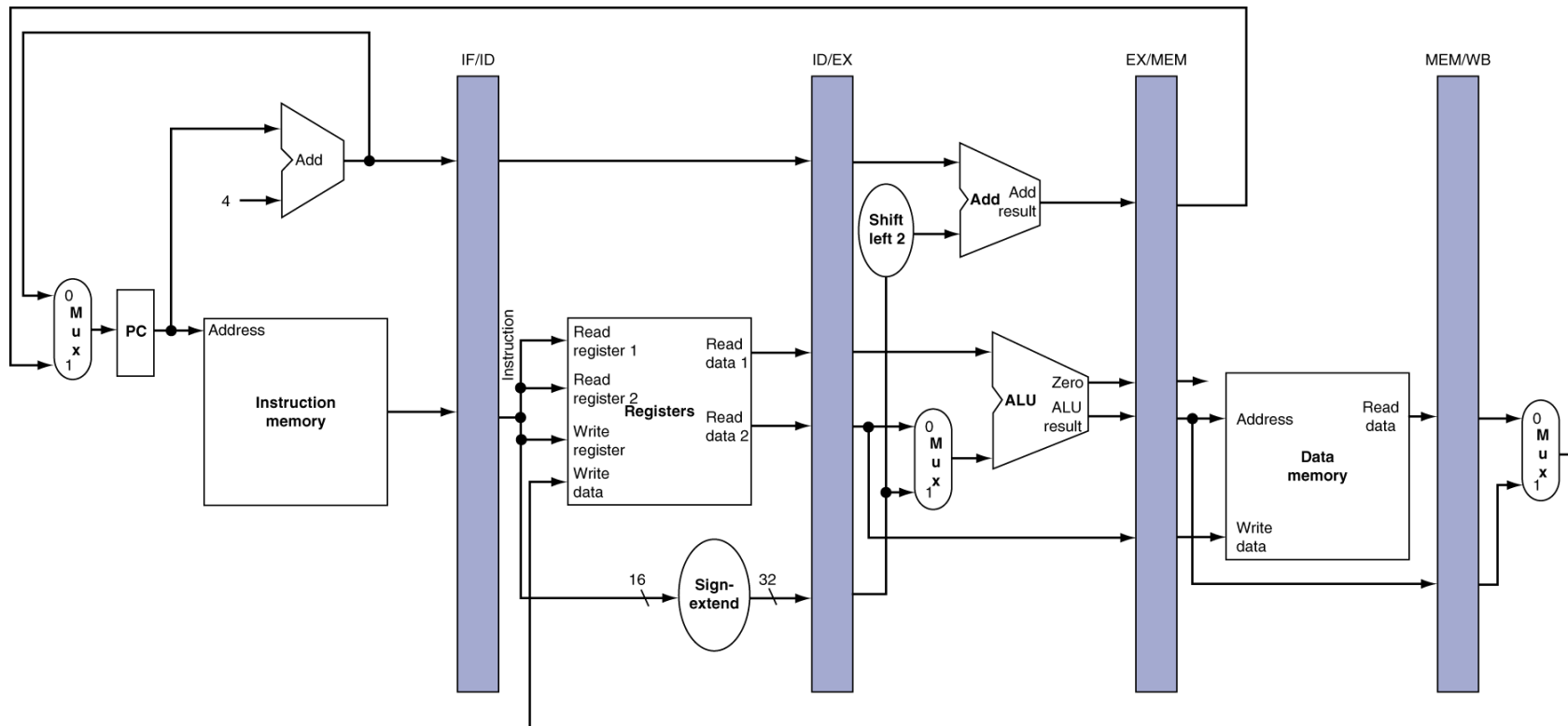
Slides from Cynthia Taylor

Announcements

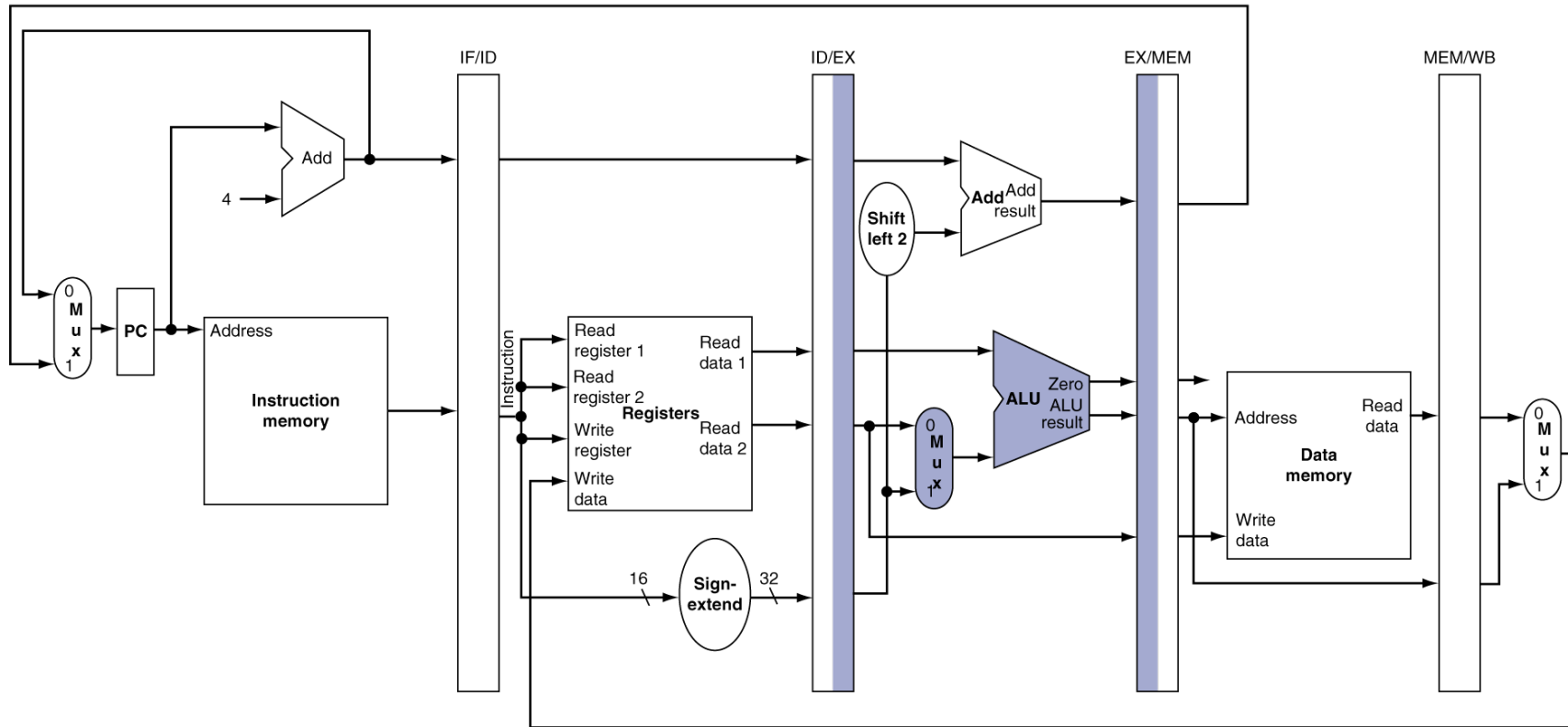
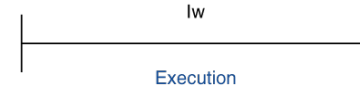
- Problem Set 10 due Jan. 2
- Lab 8 due Jan. 2
- Office Hours Friday 13:30 – 14:30

Pipeline registers

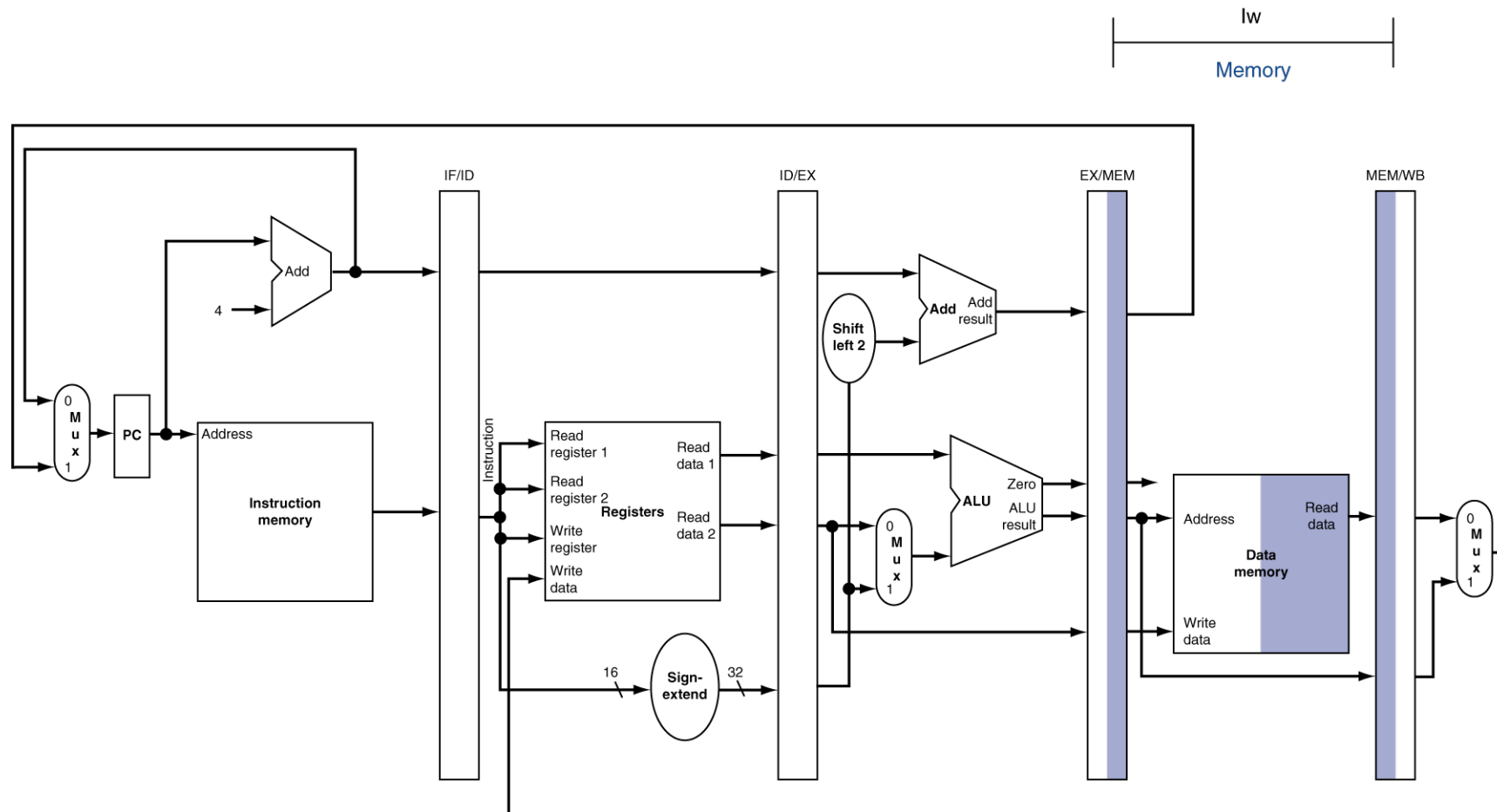
- Need registers between stages
 - To hold information produced in previous cycle



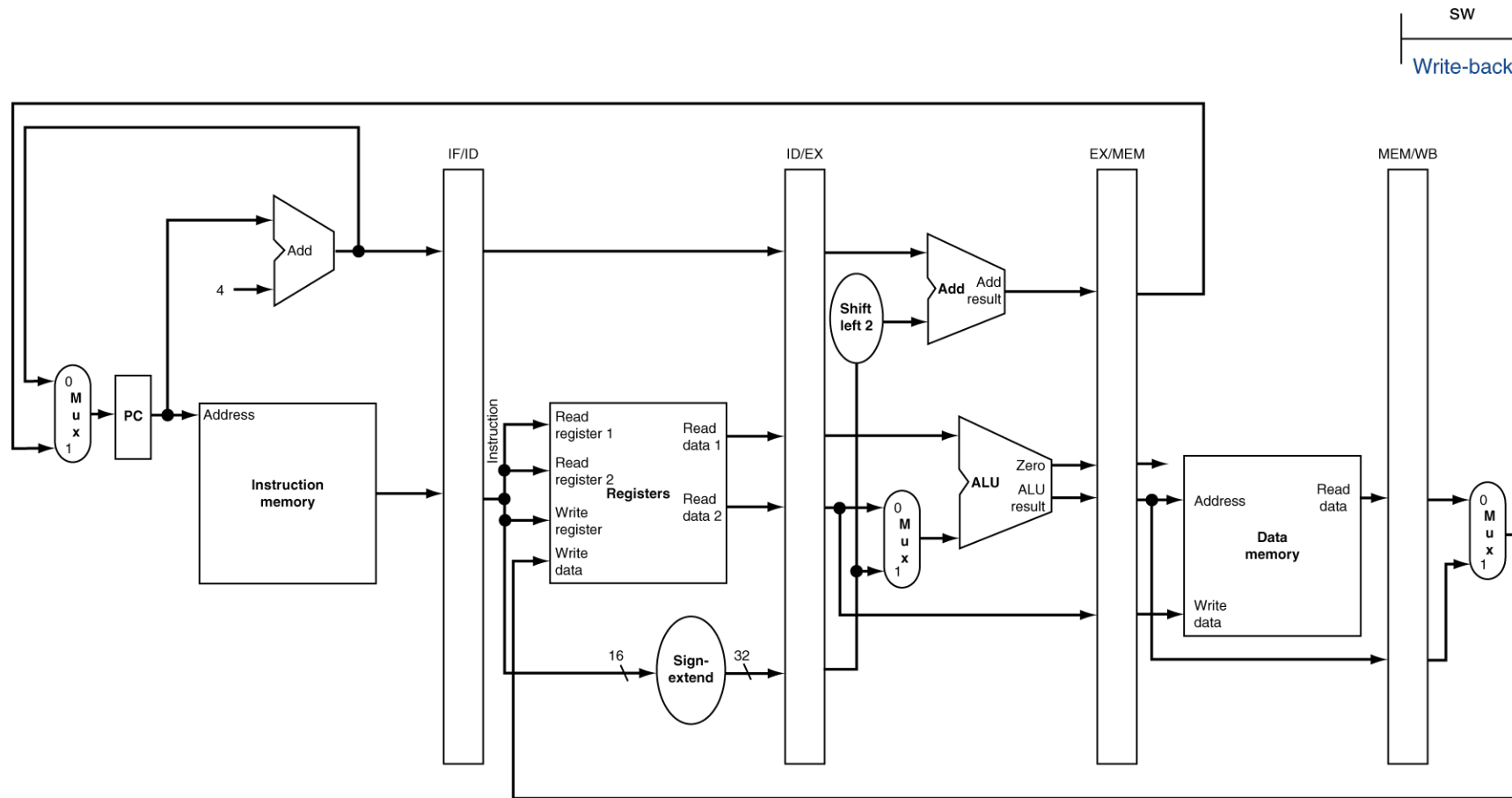
EX for Load



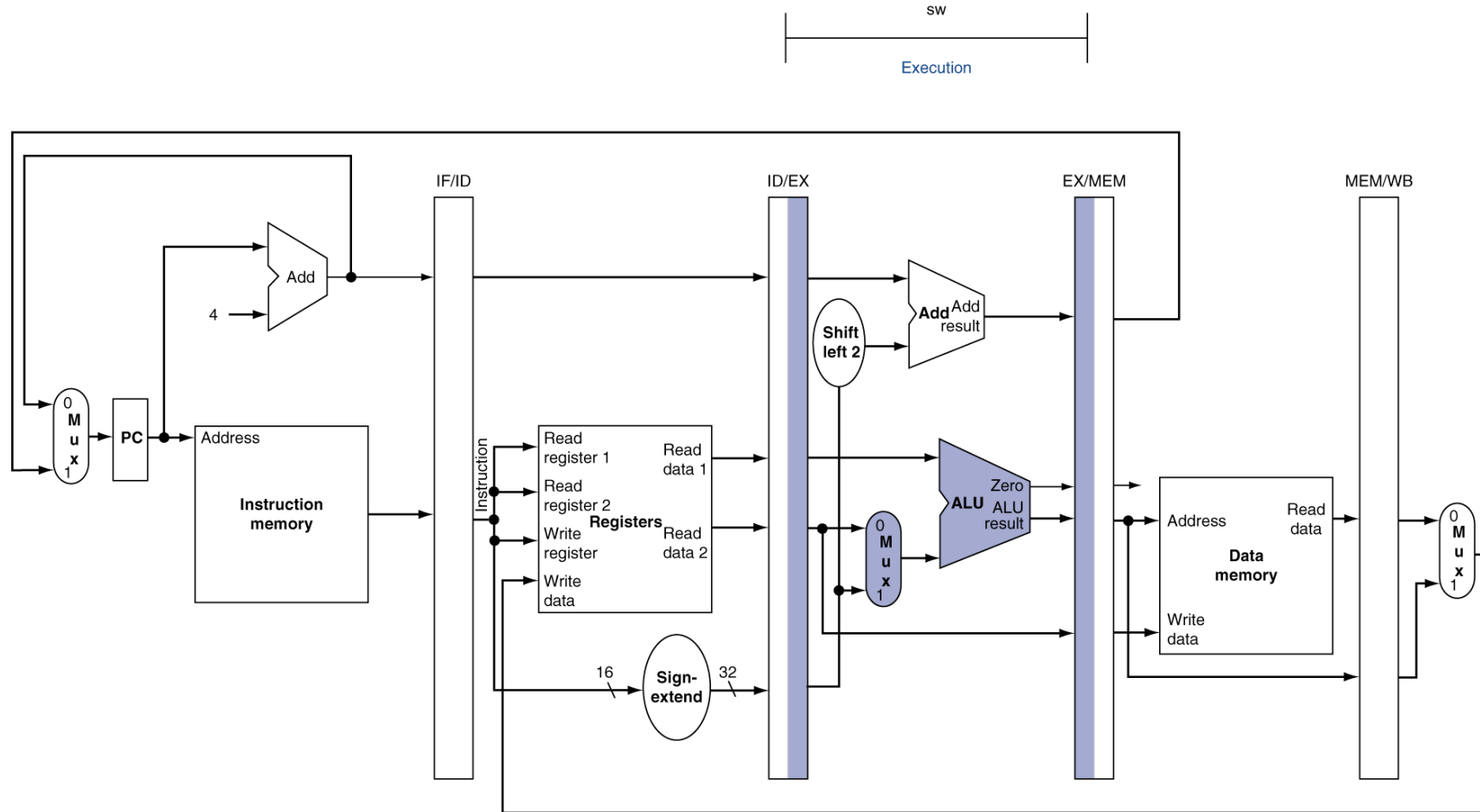
MEM for Load



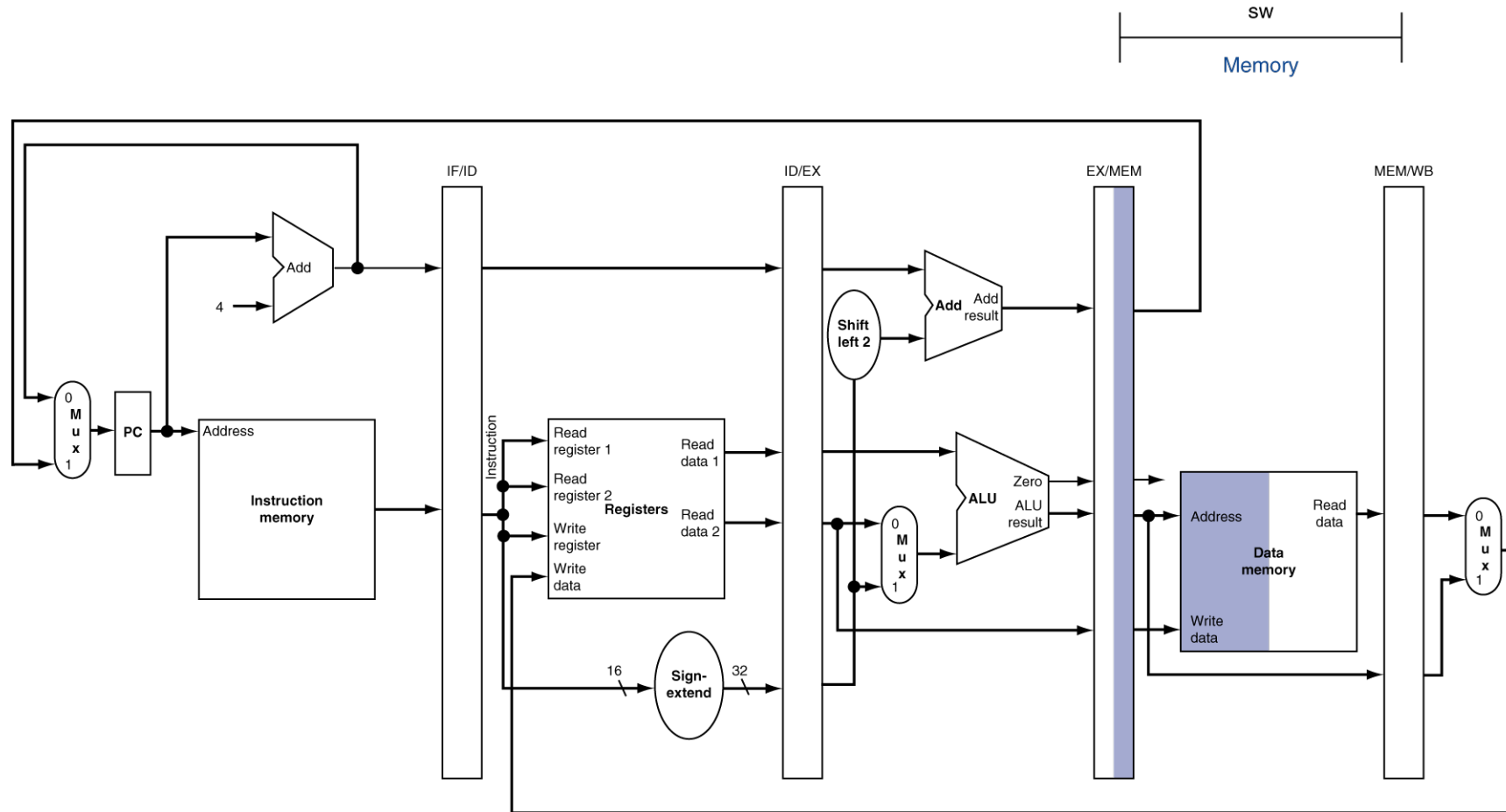
WB for Load



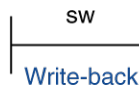
EX for Store



MEM for Store



SW
Write-back

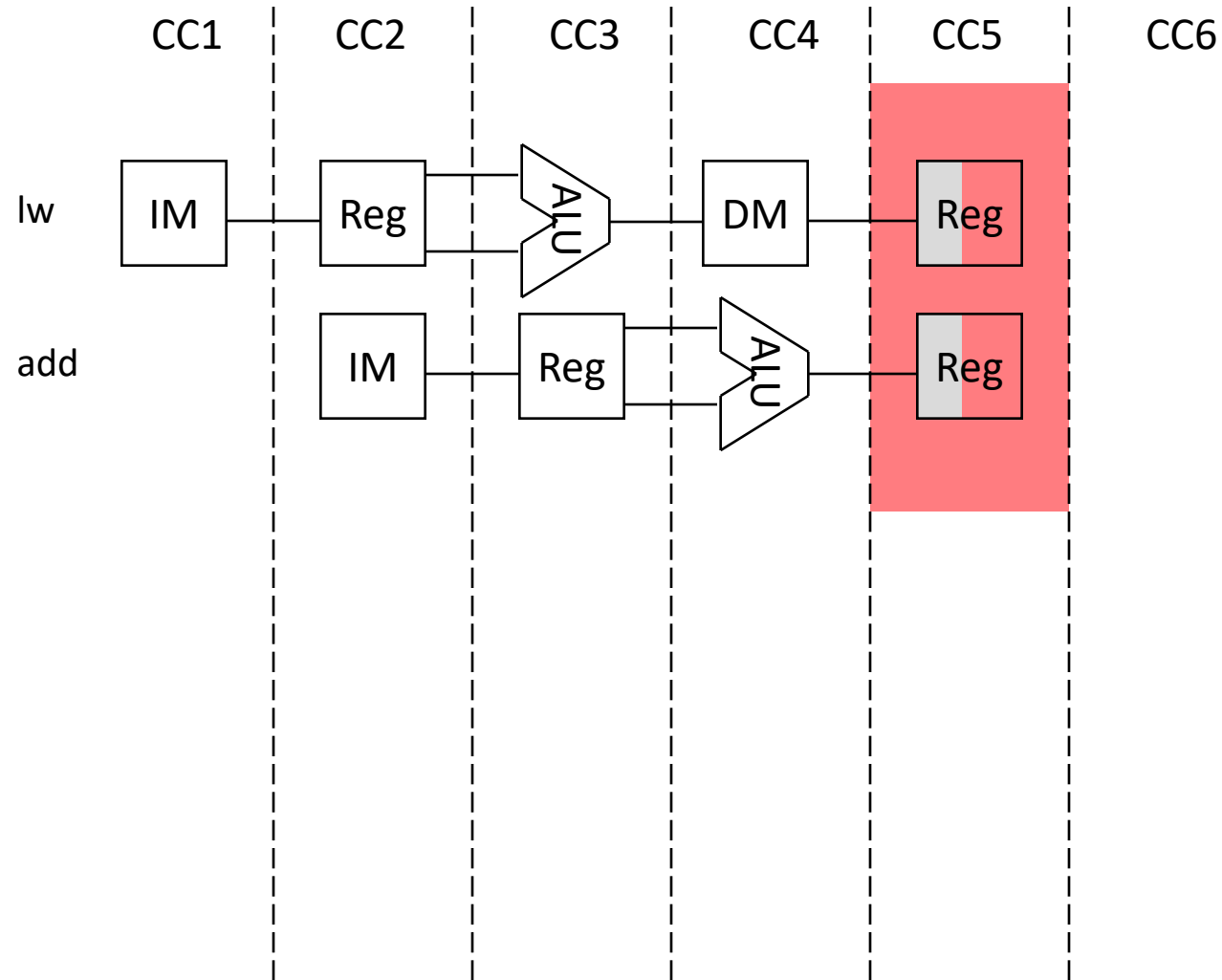


Pipeline Stages

Should we force every instruction to go through all 5 stages? Can we break it up, with R-type taking 4 cycles instead of 5?

Selection	Yes/No	Reason (Choose BEST answer)
A	Yes	Decreasing R-type to 4 cycles improves instruction throughput
B	Yes	Decreasing R-type to 4 cycles improves instruction latency
C	No	Decreasing R-type to 4 cycles causes hazards
D	No	Decreasing R-type to 4 cycles causes hazards and doesn't impact throughput
E	No	Decreasing R-type to 4 cycles causes hazards and doesn't impact latency

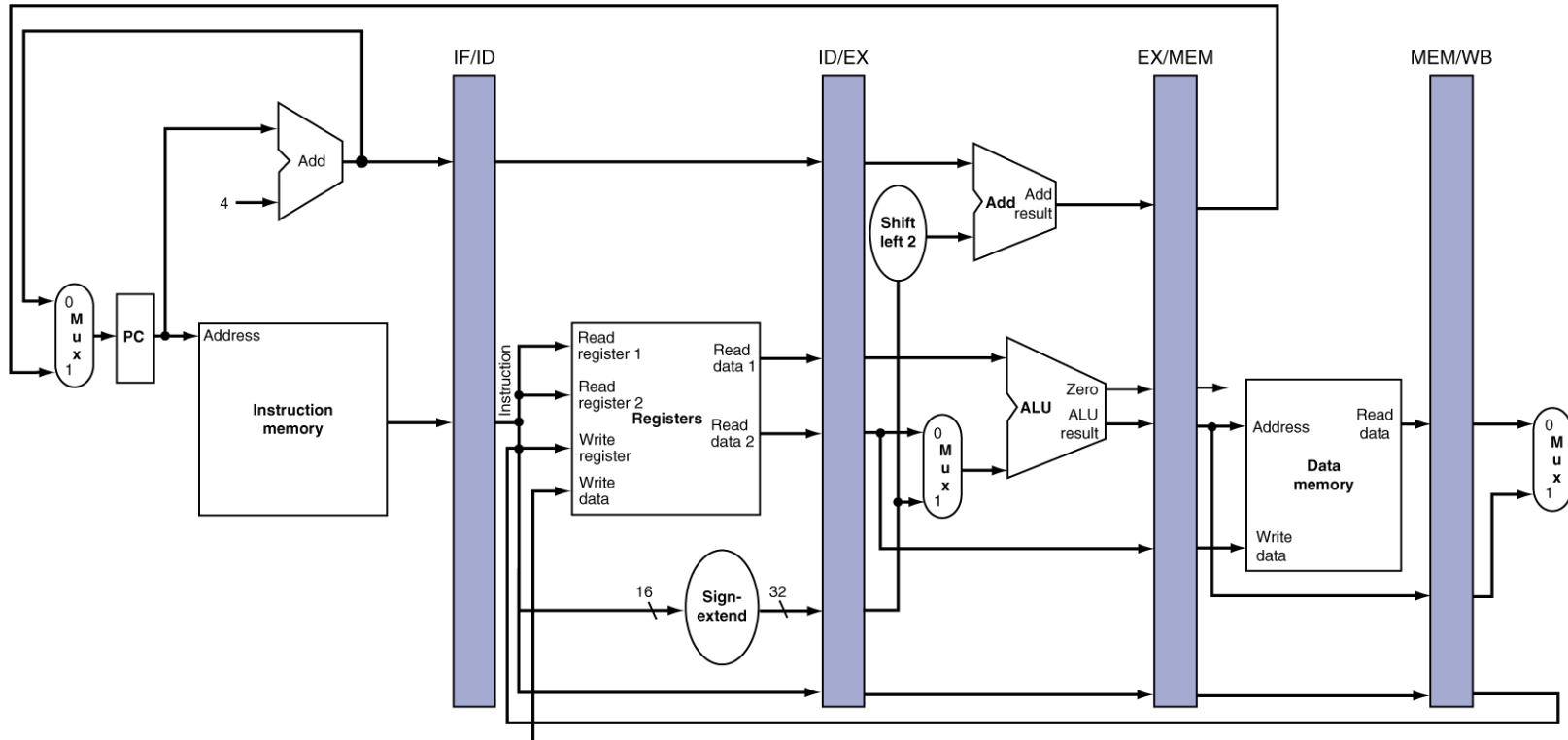
Mixed Instructions in the Pipeline



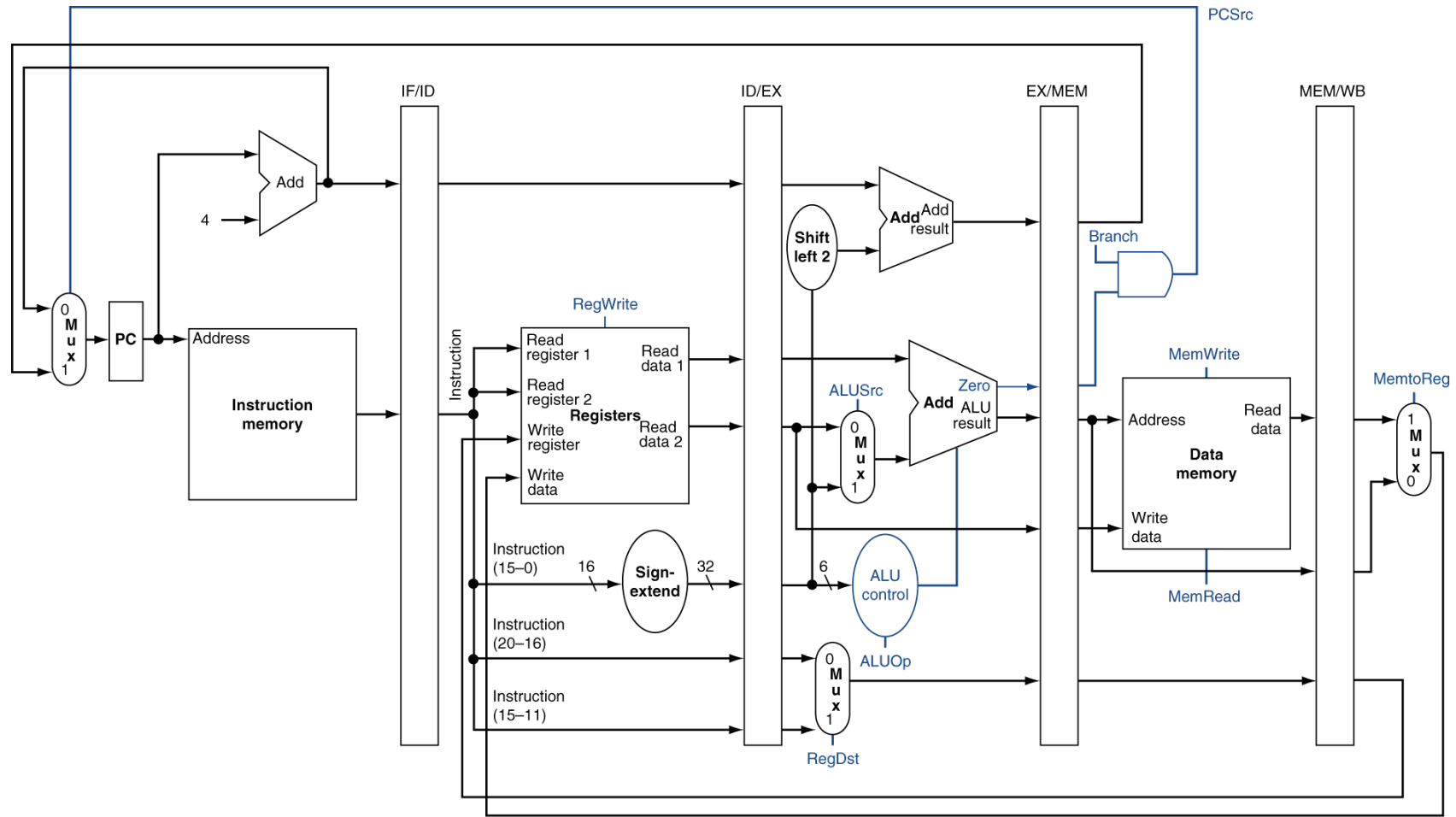
Single-Cycle Pipeline Diagram

- State of pipeline in a given cycle

add \$14, \$5, \$6	lw \$13, 24(\$1)	add \$12, \$3, \$4	sub \$11, \$2, \$3	lw \$10, 20(\$1)
Instruction fetch	Instruction decode	Execution	Memory	Write-back



Pipelined Control



How do we control our pipelined CPU?

- A. We need to add new control signals.
- B. We need to forward the control values to the correct stage.
- C. We don't need to do anything special; it will work the way it is.

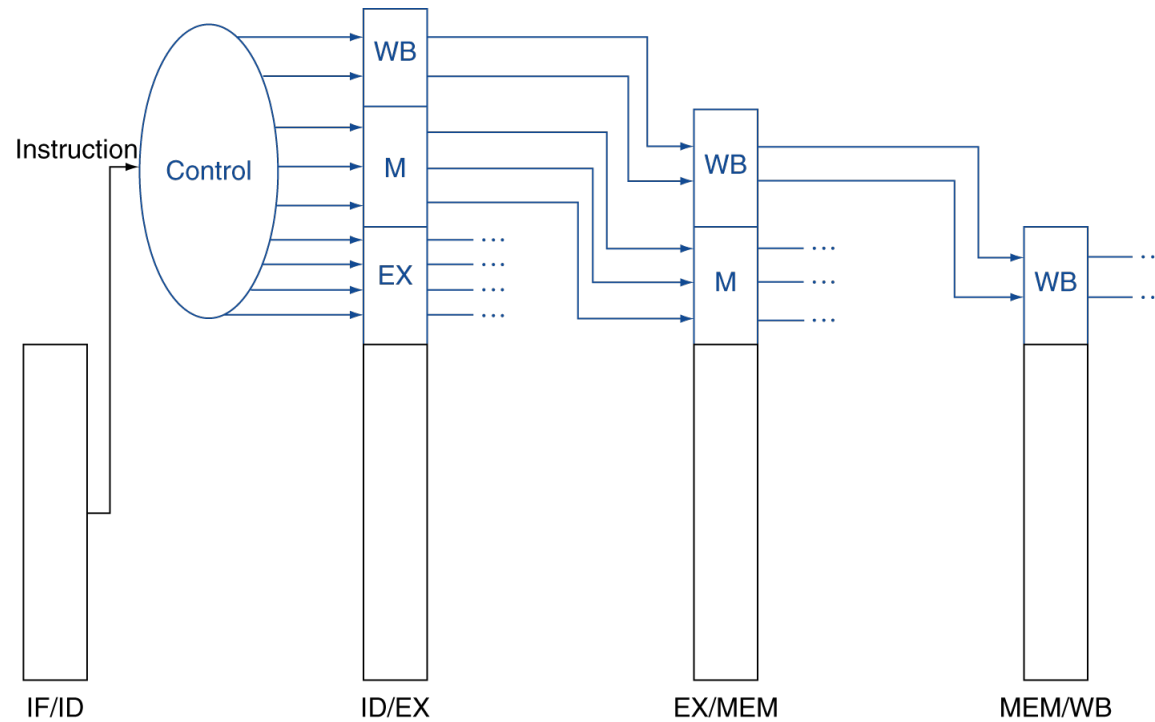
Pipeline Control

- IF Stage: read Instr Memory (always asserted) and write PC (on System Clock)
- ID Stage: no optional control signals to set

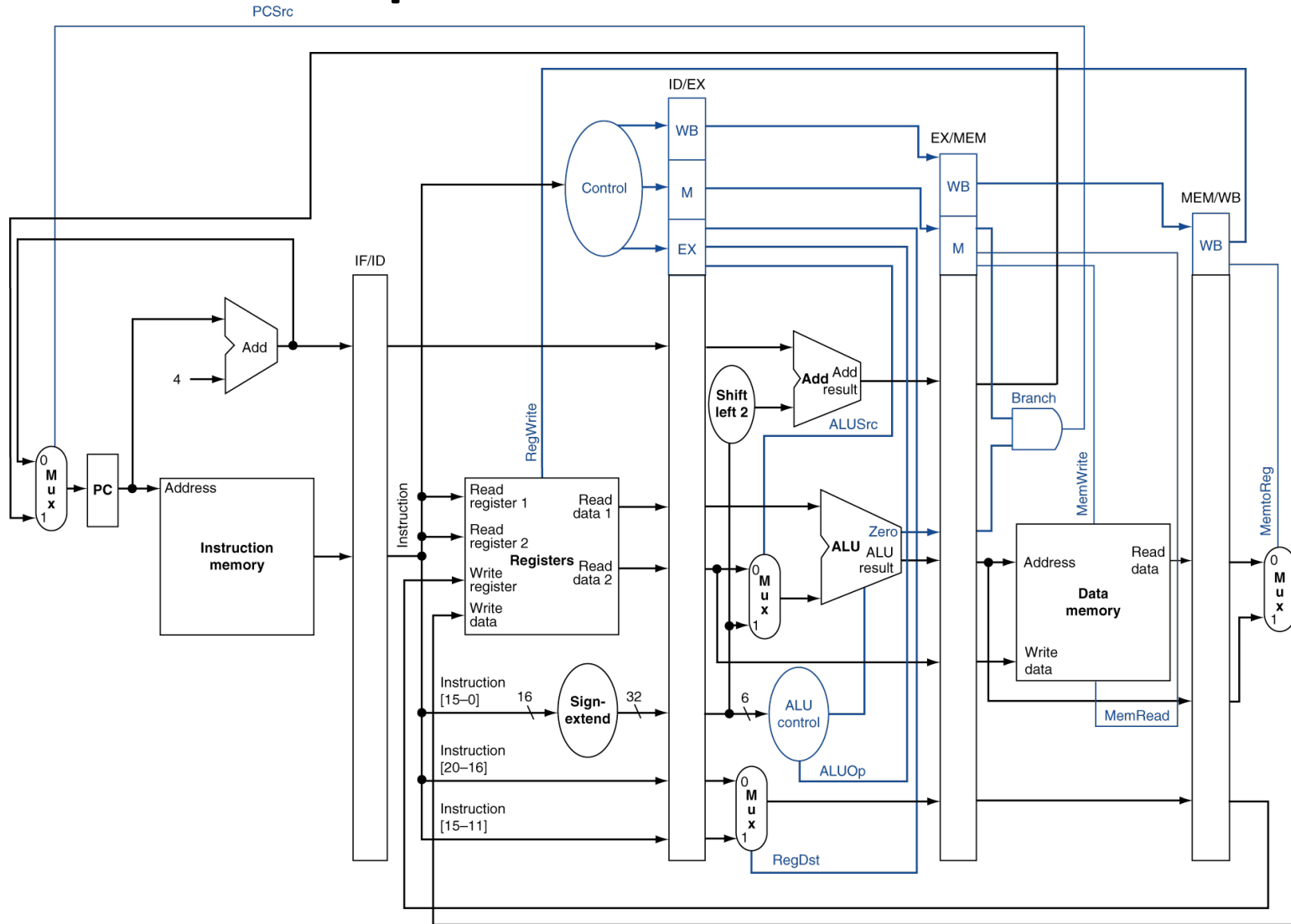
	EX Stage				MEM Stage			WB Stage	
	Reg Dst	ALU Op1	ALU Op0	ALU Src	Brch	Mem Read	Mem Write	Reg Write	Mem toReg
R	1	1	0	0	0	0	0	1	0
lw	0	0	0	1	0	1	0	1	1
sw	X	0	0	1	0	0	1	0	X
beq	X	0	1	0	1	0	0	0	X

Pipelined Control

Control signals derived from instruction



Pipelined Control



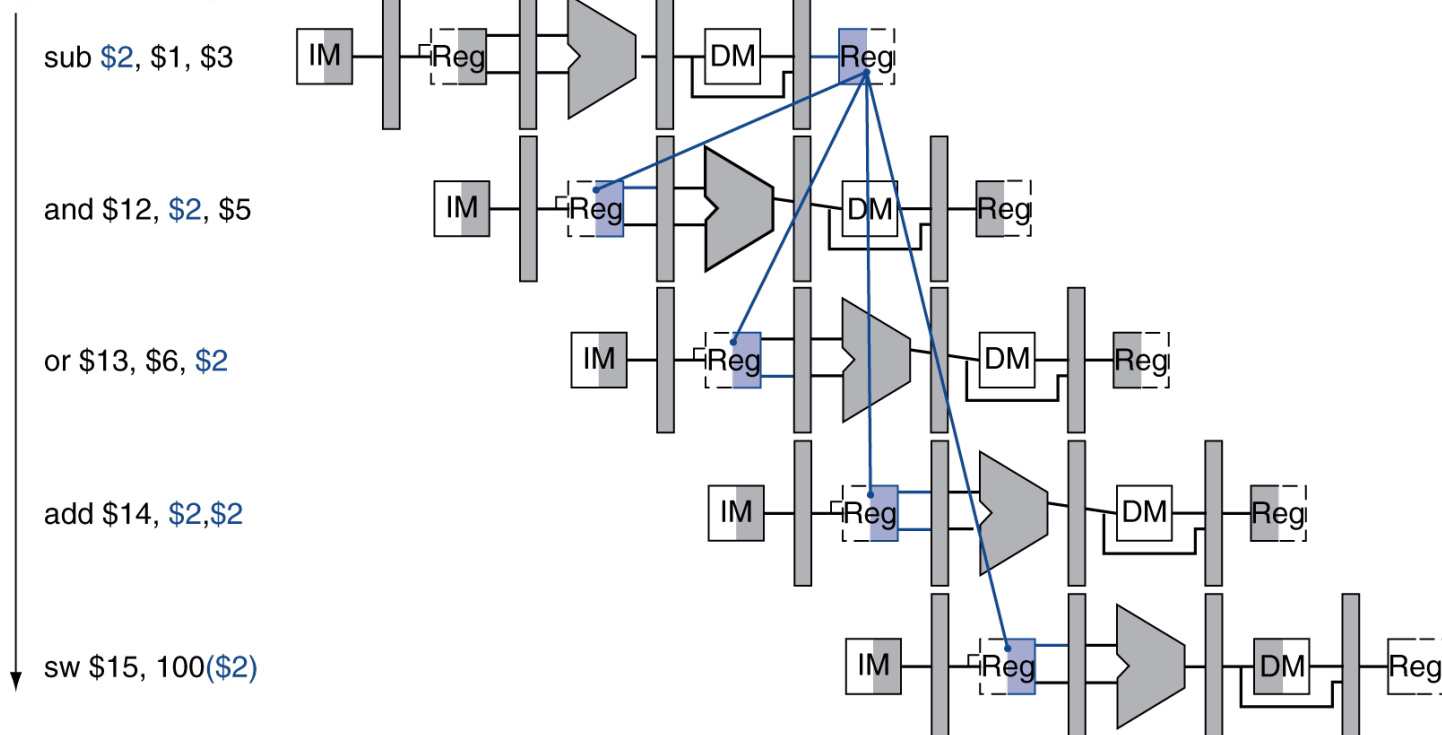
Will we need to add write control signals to our pipeline registers?

- A. No, we can automatically write them on the clock cycle
- B. Yes, we need to know if there's new information to write
- C. Yes, we might not want to overwrite previous values

Dependencies & Forwarding

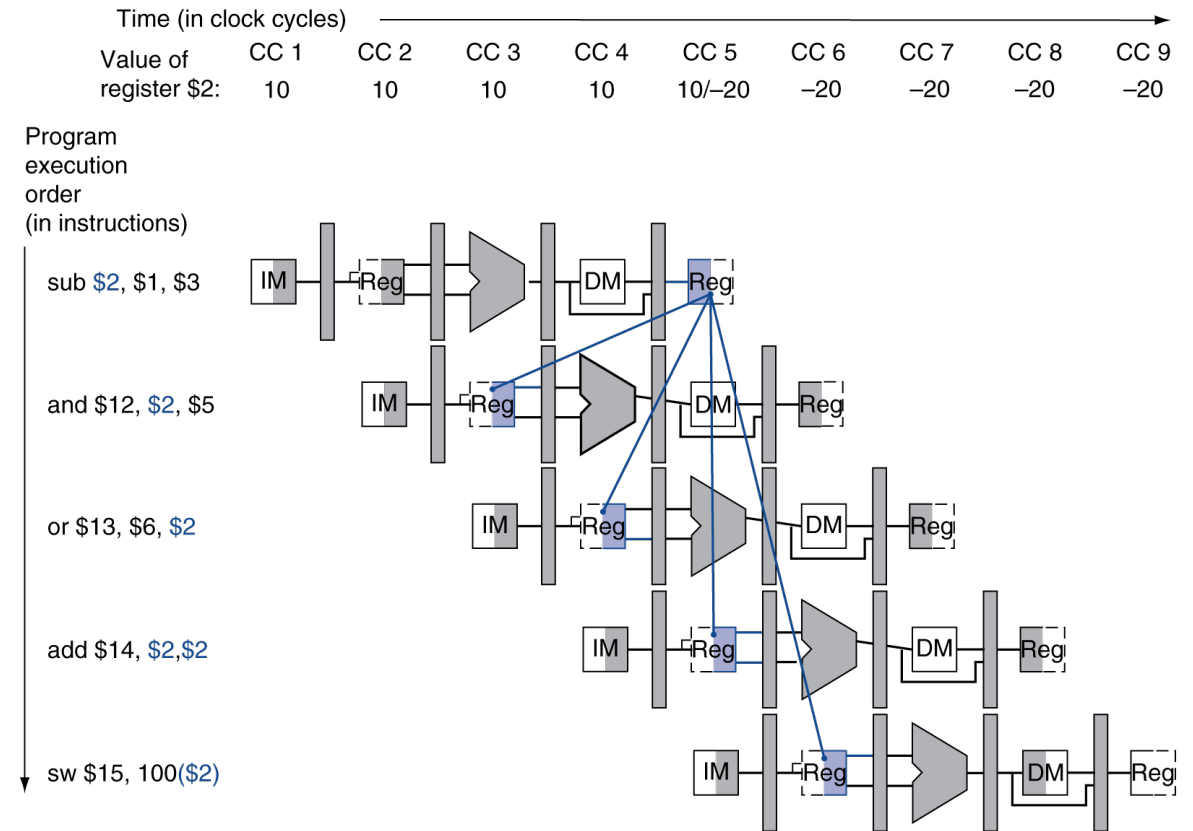
	Time (in clock cycles) →								
Value of register \$2:	CC 1	CC 2	CC 3	CC 4	CC 5	CC 6	CC 7	CC 8	CC 9
	10	10	10	10	10/-20	-20	-20	-20	-20

Program
execution
order
(in instructions)



We can BEST solve these data hazards

- A. By stalling.
- B. By forwarding.
- C. By combining forwards and stalls.
- D. By doing something else.



Data Hazards in ALU Instructions

- Consider this sequence:

```
sub $2, $1, $3
and $12, $2, $5
or  $13, $6, $2
add $14, $2, $2
sw  $15, 100($2)
```

- We can resolve hazards with forwarding
 - How do we detect when to forward?

Reading

- Next lecture: Data Hazards
 - Section 5.8