

# CSE 210: Computer Architecture

## Lecture 7: Negative Numbers, Overflow

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# Announcements

- Problem Set 2 due Friday 23:59
- Lab 1 due Sunday 23:59
- Office Hours Tuesday 13:30 – 14:30

# How do we indicate a negative number?

- Sign and magnitude
- Ones' Complement
- Two's Complement

# Ones' Complement

- To make a number negative, just flip all its bits!
- Need to know how many bits: -5 in
  - 4 bits: 0101 => 1010
  - 8 bits: 00000101 => 11111010

A byte representing  $-6_{10}$  in Ones' Complement is

- A. 00000110
- B. 10000110
- C. 11111001
- D. 11110110
- E. None of the above

# Ones' complement

- Two zeros: 00000000 and 11111111 (in 8 bits)
- Addition:
  - Perform normal n-bit addition
  - Add the carryout bit back to the result

# Two's Complement

- Flip all the bits and add 1
- For  $n$  bits, the unsigned version of  $-x = 2^n - x$
- Can represent  $-128$  to  $127$  in 8 bits
  - In  $n$  bits, can represent  $-2^{n-1}$  to  $2^{n-1} - 1$
- Only one zero (00000000 in 8 bits)
- Used in modern computers

## -6 in Two's Complement

- A. 11110110
- B. 11111001
- C. 11111010
- D. 11111110
- E. None of the above



Two's Complement:  $1111101_2 = ?_{10}$

- A. -2
- B. -3
- C. -4
- D. -5
- E. None of the above

The negation of  $11110001_2$  is \_\_\_\_\_<sub>2</sub>

- A. 00001110
- B. 00001111
- C. 00011110
- D. 01110001
- E. None of the above

# Addition and Subtraction

- Positive and negative numbers are handled in the same way.
- The carry out from the most significant bit is ignored.
- To perform the subtraction  $A - B$ , compute  $A +$  (two's complement of  $B$ )

For  $n$  bits, the sum of a number and its negation will be

A.  $0_{n-1} \dots 0_0$

B.  $1_{n-1} 0_{n-2} \dots 0_0$

C.  $1_{n-1} \dots 1_0$

D. It will vary

E. None of the above

$$11110110_2 + 00001100_2 = ?_2$$

- A. 00000010
- B. 00001100
- C. 11110010
- D. 11111110
- E. None of the above

$$1111 + 1000 = \underline{\hspace{1cm}}_2$$

A. 0111

B. 1000

C. 1111

D. 0000

E. None of the above

# Overflow

- Overflow occurs when an addition or subtraction results in a value which cannot be represented using the number of bits available.
- In that case, the algorithms we have been using produce incorrect results.

# Is overflow a problem in modern programs?

A. Nope, we have totally solved this business!

B. Yep, still a problem.



# Handling Overflow

- Hardware can detect when overflow occurs
- Software may or may not check for overflow
  - Java guarantees two's complement behavior!
  - In C, overflow is “undefined behavior” meaning, it can do anything
  - In Rust, overflow is checked in debug builds but not optimized builds!

# How To Detect Overflow

- On an addition, an overflow occurs if and only if the carry into the sign bit differs from the carry out from the sign bit.
- Overflow occurs if adding two negative numbers produces a positive result or if adding two positive numbers produces a negative result.

Will  $01111111_2 + 00000101_2$  result in overflow?

A. Yes

B. No

C. It depends

# Unsigned Numbers

- Some types of numbers, such as memory addresses, will never be negative
- Some programming languages reflect this with types such as “unsigned int”, which only hold positive numbers
  - `uint32_t` in C99
  - `u32` in Rust
  - Java only has signed types
- In an unsigned byte, values will range from 0 to 255

# In MIPS

- add, sub, addi instructions cause exceptions on (signed) overflow
- addu, subu, addiu instructions do not
- Rationale: In C, unsigned types never cause overflow, they're defined to wrap (produce a value modulo  $2^n$ )
- In practice: Since overflow is undefined behavior, it is assumed to never happen so compilers always use addu/subu/addiu

# Reading

- Next lecture: How Instructions Are Represented
  - Section 2.5
- Problem Set 2 due Friday
- Lab 1 due Sunday