# CSCI 210: Computer Architecture Lecture 25: Datapath

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# Today's Class

• The datapath!

# CS History: The Manchester Baby



- First stored-program computer
- Ran its first program on June 21, 1948
- Designed as a testbed for the first random-access memory
- Only arithmetic operations were addition and subtraction
- Its first program calculated the highest proper divisor of 2<sup>18</sup> (262,144), by testing every integer from 2<sup>18</sup> downwards
- This program was 17 instructions and took 52 minutes to run

# The Processor: Datapath & Control

- We're ready to look at an implementation of MIPS simplified to contain only:
  - memory-reference instructions: lw, sw
  - arithmetic-logical instructions: add, sub, and, or, slt
  - control flow instructions: beq

# Generic implementation

#### Fetch

- Use the program counter (PC) to supply instruction address
- Get the instruction from memory
- Update the program counter to the next instruction
- Decode instruction
  - Read registers
  - Read the instruction to decide how to execute

#### Execute

- Perform necessary data manipulation
- Write to registers

### To fetch an instruction, what hardware do we need?

- Fetch
  - Use the program counter (PC)
     to supply instruction address
  - Get the instruction from memory
  - Update the program counter to the next instruction

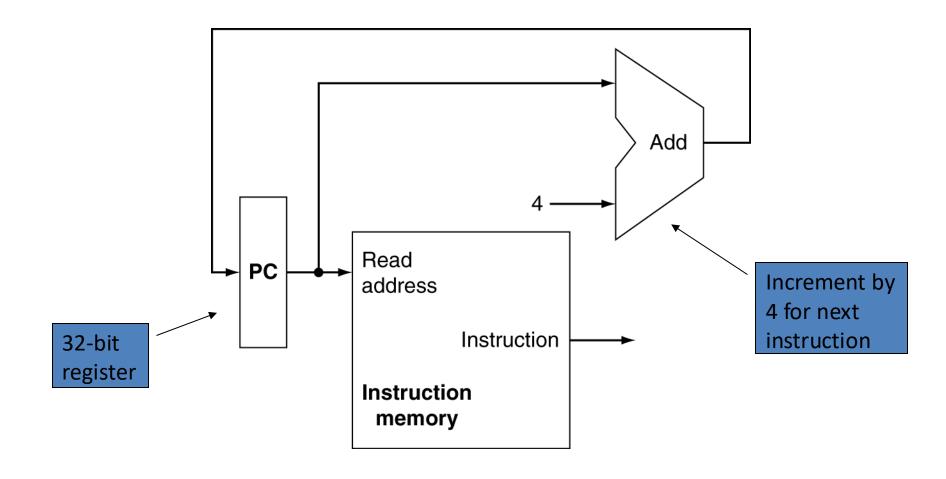
A. Register(s), Memory

B. Register(s), Adder, Memory

C. Register(s), ALU, Memory

D. More than this

### Instruction Fetch



# Generic implementation

- Fetch
  - Use the program counter (PC) to supply instruction address
  - Get the instruction from memory
  - Update the program counter to the next instruction
- Decode instruction
  - Read registers
  - Read the instruction to decide how to execute
- Execute
  - Perform necessary data manipulation
  - Write to registers

# Registers for instructions

- add \$t0, \$t1, \$t2 needs to read the values of registers \$t1 and \$t2 and write to register \$t0
- Iw \$t0, 4(\$t8) needs to read one register and write one register
- sw \$t0, -8(\$s0) needs to read two registers and write zero registers

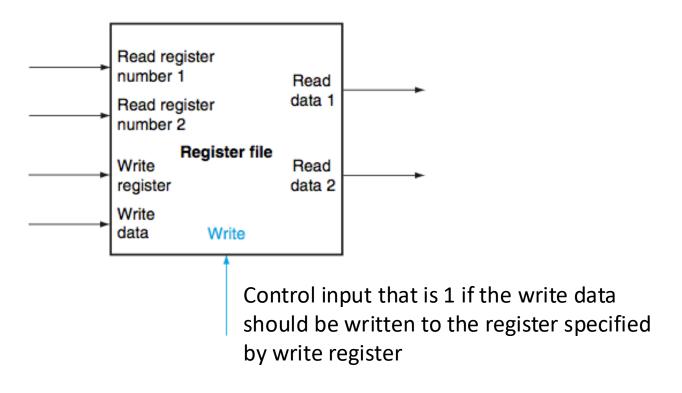
# Interface for the register file

 We need the ability to read from up to 2 registers and write up to 1 register

#### Interface:

- Three 5-bit register select inputs (rs, rt, rd)
- Two 32-bit data outputs (data in rs and rt)
- One 32-bit data input (data to write to rd)
- One 1-bit control input (should input data be written to rd or not)

# Register File



	6 bits	5 bits	5 bits	5 bits	5 bits	6 bits
R-type	opcode	rs	rt	rd	sa	funct

# Generic implementation

#### Fetch

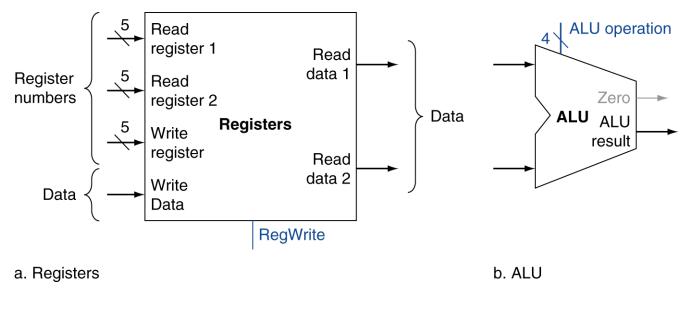
- Use the program counter (PC) to supply instruction address
- Get the instruction from memory
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- Decode instruction
  - Read registers
  - Read the instruction to decide how to execute

#### Execute

- Perform necessary data manipulation
- Write to registers

#### **R-Format Instructions**

- Read two register operands
- Perform arithmetic/logical operation
- Write register result



	6 bits	5 bits	5 bits	5 bits	5 bits	6 bits
R-type	opcode	rs	rt	rd	sa	funct

## Data memory

- sub \$t0, \$t1, \$t2 does not read or write memory
- Iw \$t0, 0(\$s0) reads 32-bits from memory
- sw \$t0, 0(\$s0) writes 32-bits to memory

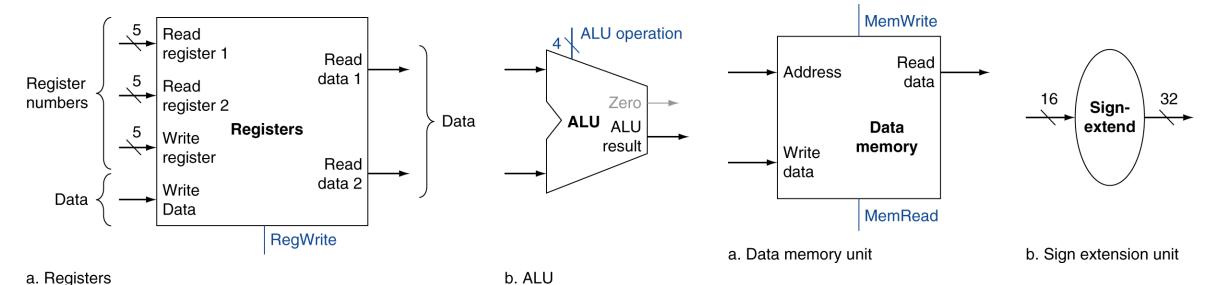
# Which of these describes our interface for data memory? What do we need for lw 0(\$t3), \$t2 or sw 4(\$t5), \$t4

- A. One 32-bit data output, 1 5-bit select input, 1 32-bit data input, 2 bits of control input
- B. One 32-bit data output, 2 5-bit select inputs, 2 bits of control input
- C. One 32-bit data output, 1 32-bit select input, 1 32-bit data input, 2 bits of control input
- D. One 32-bit data output, 1 32-bit select input, 2 bits of control input
- E. None of the above

Data is what we read/write to memory, select is our address, control is what the unit does

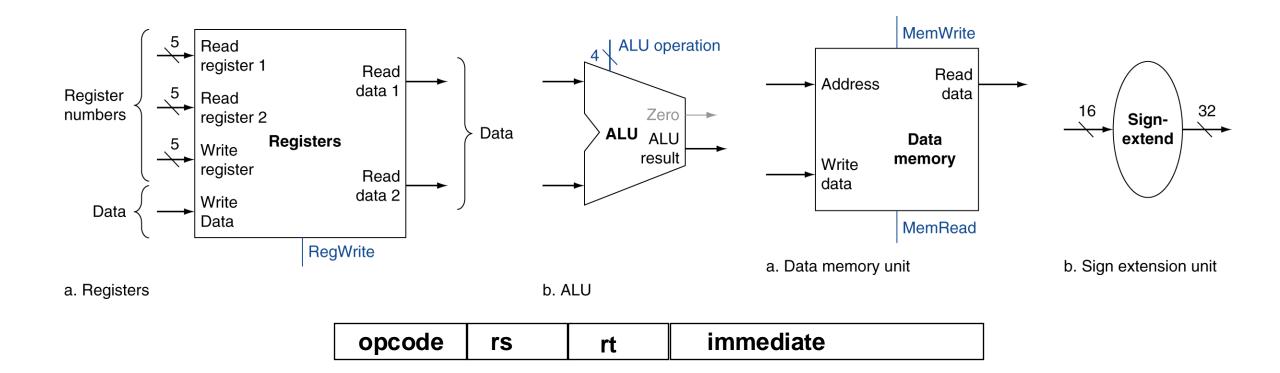
# Load/Store Instructions

- Read register operands
- Calculate address using 16-bit offset
  - Use ALU, but sign-extend offset
- Load: Read memory and update register
- Store: Write register value to memory



opcode rs rt immediate

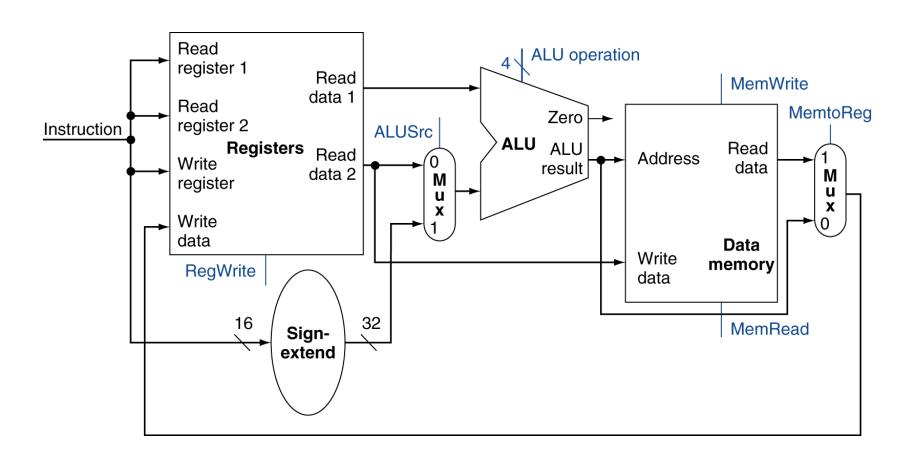
# Memory



# Which is true about the ALU and the register file in MIPS?

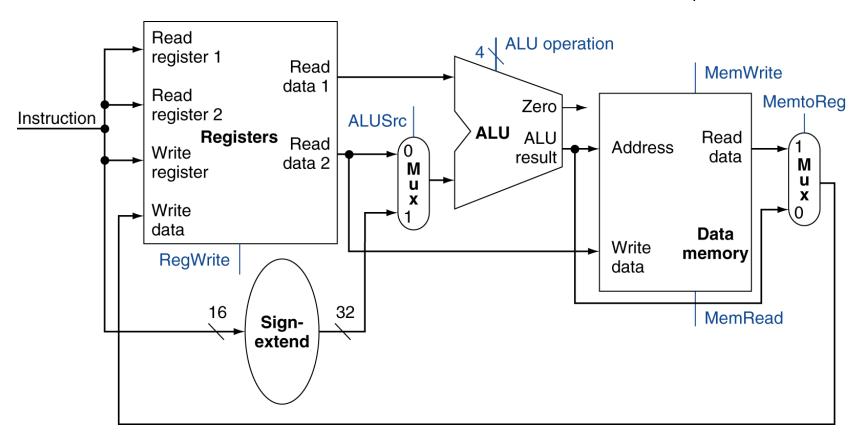
- A. The ALU always performs an operation before accessing the register file
- B. The ALU *sometimes* performs an operation before accessing the register file
- C. The register file is *always* accessed before performing an ALU operation
- D. The register file is *sometimes* accessed before performing an ALU operation
- E. None of the above.

# R-Type/Load/Store Datapath



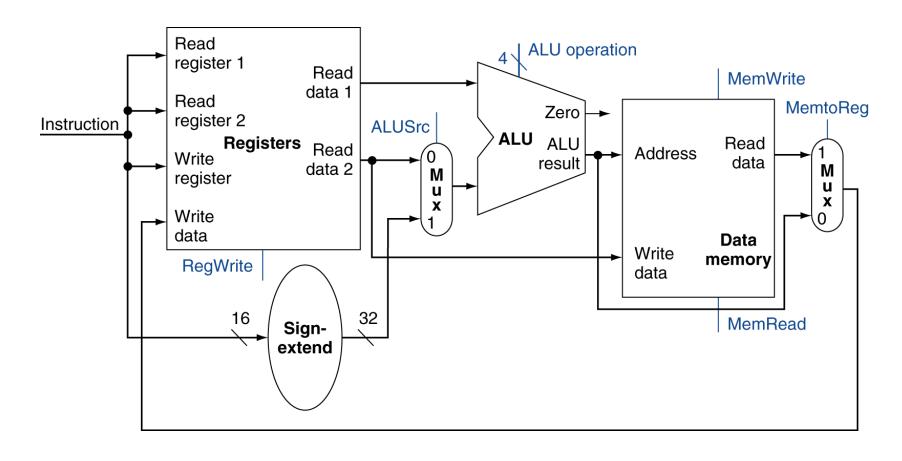
# Add \$t0, \$t0, \$t1

\$t0 is register 8, \$t1 is register 9 \$t0 holds 5 \$t1 holds 6



# lw \$t1, 4(\$t0)

\$t0 is register 8, \$t1 is register 9 \$t0 holds 0x07AB8110 0x07AB8114 holds 12



#### **Branch Instructions**

- Read register operands
- Compare operands
  - Use ALU, subtract and check Zero output
- Calculate target address
  - Sign-extend offset
  - Shift left 2 bits (word offset)
  - Add to PC + 4
    - Already calculated during instruction fetch

# What do we need to add Conditional Branch Instructions to our datapath?

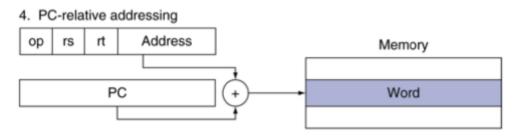
A. ALU

B. Registers and an ALU

C. Registers, ALU and Memory

D. Registers, an ALU and an Adder

beq \$t2, \$t3, 0x4F35



Read register operands

Compare operands

Use ALU, subtract and check Zero output

Calculate target address

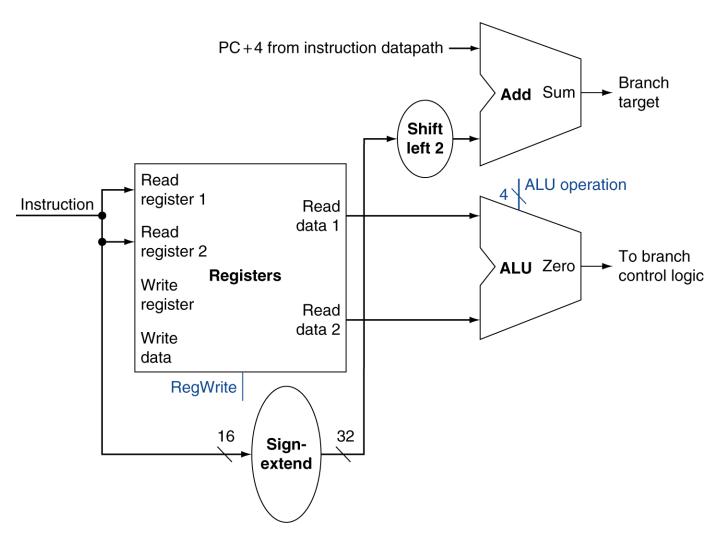
Sign-extend offset

Shift left 2 bits (word offset)

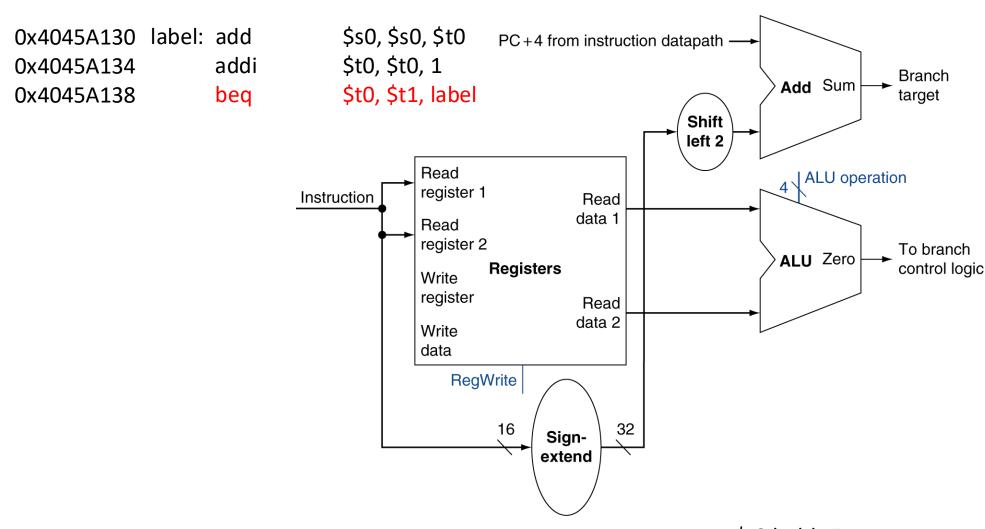
Add to PC + 4

Already calculated during instruction fetch

### **Branch Instructions**



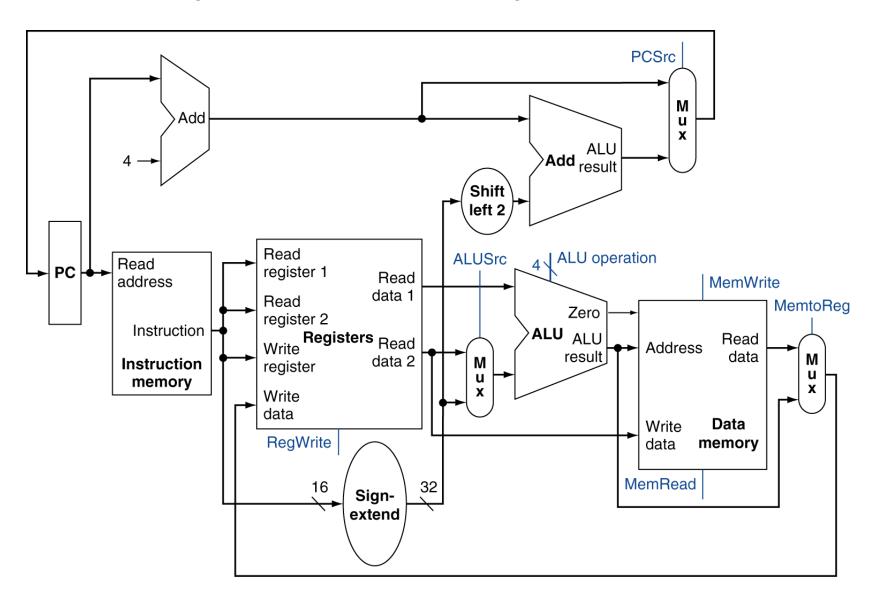
#### **Branch Instructions**



op = 0x04 rs = 8 rt = 9 imm = 0xFFFD

\$t0 holds 5 \$t1 holds 5

# Datapath (still simplified a bit)



# Reading

- Next lecture: Data path
  - Section 5.2