HW 10
Due: 27 September 2011

COSC 3015

## 1 Datatypes and Structural Induction

As discussed in class today, Haskell datatypes have corresponding structural induction princples.

#### 1.1 Nat

The datatype definition for natural numbers was given as:

```
datatype Nat = Zero \mid Succ Nat
```

The induction principle for the type Nat is

$$(P(Zero) \land \forall k :: Nat. \ P(k) \Rightarrow P(Succ \ k)) \Rightarrow \forall n :: Nat. \ P(n)$$

Thus to show a property P holds for every finite natural number, show two things:

Case Zero: Show P(Zero).

Case Succ: For arbitrary k :: Nat assume P(k) and show P(Succ k).

Note, if we want to reason about possibly infinte elements of the type we add the case to show  $P(\perp)$ . This is true for each of the induction principles below.

#### 1.2 Lists

Recall the defintion of lists containing elements of type a.

```
datatype \ List \ a = Nil \ | \ Cons \ a \ (List \ a)
```

The induction principle for finite lists is

$$(P(Nil) \land \forall x :: a. \forall xs :: List \ a. \ P(xs) \Rightarrow P(Cons \ x \ xs)) \Rightarrow \forall ys :: List \ a. \ P(ys)$$

Thus to show a property P holds for every list, show two things:

Case Nil: Show P(Nil).

Case Cons: For arbitrary x :: a and xs :: List a assume P(xs) and show  $P(Cons \ x \ xs)$ .

### 1.3 Binary Trees

Recall the definition of the datatype for binary trees containing elements of type a.

```
datatype \ BTree \ a = Leaf \ | \ Node \ a \ (BTree \ a) \ (BTree \ a)
```

The induction principle for finite binay trees is

$$(P(Leaf) \land \forall x :: a. \ \forall t_l, t_r :: BTree \ a. \ (P(t_l) \land P(t_r)) \Rightarrow P(Node(v, t_l, t_r))) \Rightarrow \forall t :: BTree \ a. \ P(t)$$

Thus to show a property P holds for every BTree, show two things:

Case Leaf: Show P(Leaf).

Case Node: For arbitrary x :: a and  $t_l, t_r :: BTree \ a$  assume  $P(t_l)$  and assume  $P(t_r)$  and show  $P(Node(v, t_l, t_r))$ .

# 2 Assignment

Write the structural induction principles for the finite instances of the following types.

1. Trees having nodes with both one and with two children.

 $One Two Tree\ a = Empty \mid Single\ a\ (One Two Tree\ a) \mid Branch\ a\ (One Two Tree\ a)$ 

2. Trees whose nodes have exactly three children.

Three  $Tree\ a = Leaf \mid Node\ a$  (Three  $Tree\ a$ ) (Three  $Tree\ a$ ) (Three  $Tree\ a$ )

3. Formulas of the following form.

 $Formula = Bottom \mid Not Formula \mid And Formula Formula$