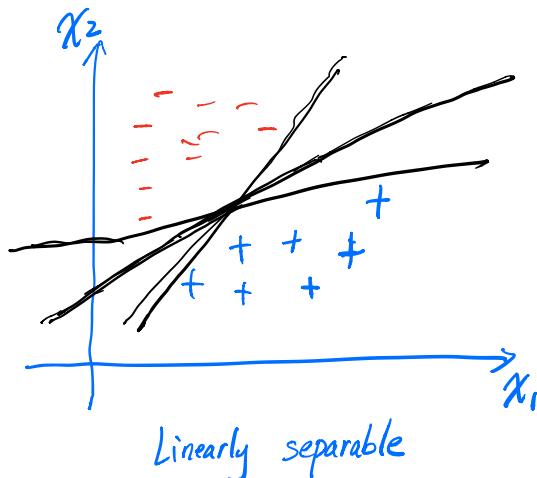


Logistic Loss ✓

Hinge loss.

Used by Support Vector Machine (SVM)

Main idea: Margin Maximization.



Infinite linear models
w/ perfect accuracy.

Linearly separable

Want to maximize the distance between
decision boundary & examples.

$w \in \mathbb{R}^d$: weight vector

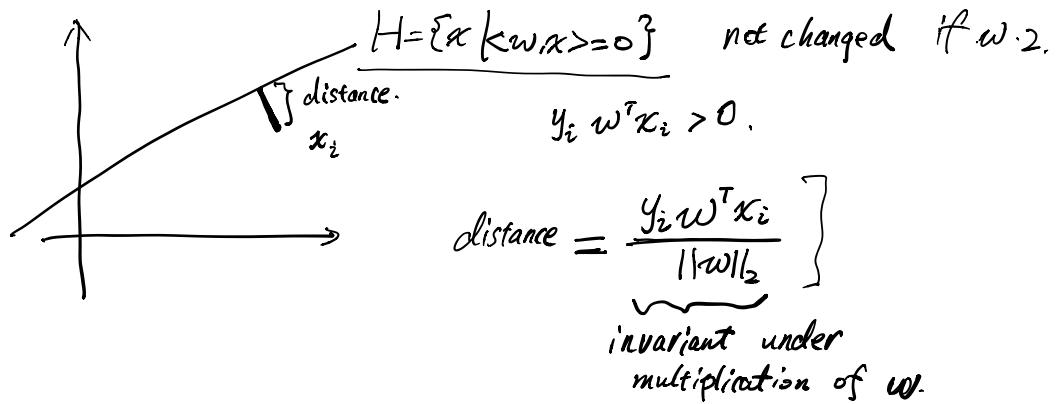
Decision boundary : $H = \{x \in \mathbb{R}^d \mid \underbrace{\langle w, x \rangle}_\text{Hyperplane} = 0\}$

Suppose w is perfectly accurate

for all examples $(x_i, y_i) \in \mathbb{R}^d \times \{\pm 1\}$

$$y_i w^T x_i > 0$$

Distance between x & H .



Idea of margin / distance maximization.

$$\boxed{\max_{\mathbf{w}} \min_i \frac{y_i \mathbf{w}^T \mathbf{x}_i}{\|\mathbf{w}\|_2}}$$

closest distance

optimization does not care about $\|\mathbf{w}\|_2$

might as well consider \mathbf{w} 's

such that $\boxed{\min_i y_i \mathbf{w}^T \mathbf{x}_i = 1}$

$\Leftrightarrow \boxed{\max_{\mathbf{w}} \frac{1}{\|\mathbf{w}\|_2}}$ such that $\min_i y_i \mathbf{w}^T \mathbf{x}_i = 1. \quad (2)$

$\Leftrightarrow \boxed{\min_{\mathbf{w}} \frac{1}{2} \|\mathbf{w}\|_2^2}$ such that $\boxed{y_i (\mathbf{w}^T \mathbf{x}_i) \geq 1 \quad \forall i}$

$\min_i y_i (\mathbf{w}^T \mathbf{x}_i) = 1 \quad \text{for all } i$

w such that

$\min_i y_i (\mathbf{w}^T \mathbf{x}_i) = 2 \rightarrow \text{take } \frac{w}{2}. \quad (3)$

margin maximization formulation

"In a perfect world"

What if not linearly separable?

$$\min_w \left[\frac{1}{2} \|w\|_2^2 \right] \text{ such that } \boxed{y_i(w^\top x_i) \geq 1} \quad \forall i$$

all

"Hard" Margin SVM (may not be feasible)

"Soft" Margin SVM.

$\min_w \frac{1}{2} \|w\|_2^2 + C \sum_{i=1}^n \xi_i$

such that

$\forall i, y_i(w^\top x_i) \geq 1 - \xi_i$, $\forall x_i$

$\xi_i \geq 0$.

get rid of?

(Equivalent form)

$\min_w \frac{1}{2} \|w\|_2^2 + C \sum_{i=1}^n \max\{0, 1 - y_i w^\top x_i\}$ "Hinge loss"

$\Leftrightarrow \min_w \underbrace{\frac{1}{2} \|w\|_2^2}_{\text{Regularization.}} + \underbrace{\sum_{i=1}^n \max\{0, 1 - y_i w^\top x_i\}}_{\text{Hinge loss ERM}}$

Constrained OPT

~~Unconstrained OPT~~

$$\min_w \frac{1}{2} \|w\|_2^2 + C \sum_{i=1}^n \xi_i$$

such that

$$\forall i, y_i(w^T x_i) \geq 1 - \xi_i,$$

slack variable

$$\forall i, \xi_i \geq 0.$$

$$\min_w \frac{1}{2} \|w\|_2^2 + C \sum_{i=1}^n \max\{0, 1 - y_i w^T x_i\}$$

"Hinge loss"

Detour : Lagrange Duality.

$$\min_w F(w) \quad \text{such that } h_j(w) \leq 0, \quad \forall j \in [m]$$

Constrained OPT

$$\xi_1, \dots, \xi_m$$

"Put constraints into obj" $w \leftarrow$ "primal —"

For each j , introduce $\lambda_j \geq 0 \leftarrow$ "dual variable"

$$\text{Lagrangian : } L(w, \lambda) = F(w) + \sum_{j=1}^m \lambda_j h_j(w)$$

$$\begin{array}{c} \text{primal} \\ \text{control } w \\ \min L(w, \lambda) \end{array} \quad \begin{array}{c} \text{a game.} \\ \text{Dual.} \\ \text{control } \lambda \\ \max L(w, \lambda) \end{array}$$

forced to satisfy
 $h_j(w) \leq 0 \quad \forall j$

Suppose w is
 not a
 feasible solution
 $\underline{h_j(w) > 0}$
 $\lambda_j \leftarrow$ "infinity"
 $L(w, \lambda) \rightarrow \infty$

Zero-sum game
view

$$\begin{array}{ccccc} \min_w & \max_{\lambda} & L(w, \lambda) & \text{"minmax"} \\ \text{left} \longrightarrow \text{right} & & & \end{array}$$

$\min \max \Leftarrow$ minimization goes first
maximization "best response"

maxmin : opposite

Play second is "usually" better.

Weak Duality : $\underbrace{\max_{\lambda}}_{\text{favors min player.}} L(w, \lambda) \leq \underbrace{\min_w}_{\text{favors max player.}} \max_{\lambda} L(w, \lambda)$

Strong Duality : $\max_{\lambda} \min_w L(w, \lambda) = \min_w \max_{\lambda} L(w, \lambda)$

under "mild" condition (e.g., soft-SVM, Slaters Condition)

$$\lambda^* = \arg \max_{\lambda} \left(\min_w L(w, \lambda) \right)$$

$$w^* = \arg \min_w \left(\max_{\lambda} L(w, \lambda) \right)$$

Strong Duality \Rightarrow

$$\begin{aligned} F(w^*) &= \min_w \max_{\lambda} L(w, \lambda) && \text{(Strong Duality)} \\ &= \max_{\lambda} \min_w L(w, \lambda) \\ &= \min_w L(w, \lambda^*) && \text{(Defn } \lambda^*) \\ &\leq L(w^*, \lambda^*) && \text{(Defn min)} \\ &= F(w^*) + \sum_j \lambda_j^* k_j(w^*) && \text{(Defn of } L) \end{aligned}$$

$$\sum_j \lambda_j^* h_j(w^*) \geq 0 \Leftrightarrow$$

$\min_w F(w) \text{ such that } \begin{cases} h_j(w) \leq 0 \\ \forall j \in [m] \end{cases}$

w^* is feasible $\lambda_j^* \geq 0$

$$\sum_j \lambda_j^* h_j(w^*) \leq 0 \Leftrightarrow$$

$$\sum_j \lambda_j^* h_j(w^*) = 0$$

Each $\lambda_j^* h_j(w^*) \leq 0$

$$\Rightarrow \forall j, \lambda_j^* h_j(w^*) = 0$$

KKT conditions

(Characterization of w^*, λ^*)

- $\forall j, \lambda_j^* h_j(w^*) = 0$ (complementary slackness)
- w^* is the minimizer of $L(w, \lambda^*)$
- $\nabla_w L(w^*, \lambda^*) \rightarrow 0$ (Stationarity)
- $\lambda_j \geq 0, h_j(w^*) \leq 0$ for all j (Feasibility)

Dated back Kuhn-Tucker '1951

Karush '1939 in unpublished master thesis