CSC215

Math and Computer Science



Recursion

- Related to mathematical induction
 - Definition of the nth value in terms of the first n-1 values
 - A stopping condition

- In C++, it is a function that calls itself
 - Has what is known as a base case for stopping the function calls



Ask Yourself the Following

- 1. How to define the problem in terms of a smaller problem of the same type?
- 2. How does each recursive call diminish the size of the problem?
- 3. What instance of the problem can serve as the base case?
- 4. As the problem diminishes, will it reach the base case?

The example here are best done iteratively.



Factorial

- N! = N * (n-1) * (n-2) * (n-3) *......
- How to define in terms of the first n-1 terms
- N! = N * (n-1)!

- What will serve as our base case?
 - 0! = 1
 - Could use 1! But the smallest possible factorial is defined as 0! = 1



The Factorial Criteria

• N = 0

• N > 0

answer is 1

answer is N * (N-1)!



Drawing out N!

- Assume n = 6
 - 6! = 6 * 5!
 - 5! = 5 * 4!
 - 4! = 4 * 3!
 - 3! = 3 * 2!
 - 2! = 2 * 1!
 - 1! = 1 * 0!
 - 0! = 1

← Good, I know the answer



Substitute Going Back Up

• Answer to n = 6, 6! = 720

```
6! = 6 * 5! = 6 * 120 = 720
5! = 5 * 4! = 5 * 24 = 120
4! = 4 * 3! = 4 * 6 = 24
3! = 3 * 2! = 3 * 2 = 6
2! = 2 * 1! = 2 * 1 = 2
1! = 1 * 0! = 1 * 1 = 1
0! = 1
```



Writing the function

```
unsigned int factorial( unsigned int n )
{
    // write the base case first
    if( n == 0 )
        return 1;
```

}



Writing the function

```
unsigned int factorial( unsigned int n )
    int ans;
    // write the base case first
    if( n == 0 )
        return 1;
    // compute the answer
    ans = n * factorial( n - 1 );
    return ans;
```



Stepping Through the Function

```
int main()
    int result;
    int n = 4;
    result = factorial ( n );
    cout << "4! = " << result << endl;</pre>
    return 0;
```



Memory	main
? = F(4)	result
4	n



Memory	main	F(4)
? = F(4)	Result	
4	n	
4		n
? = N *F(3)		ans



Memory	main	F(4)	F(3)
? = F(4)	Result		
4	n		
4		n	
? = n * F(3)		ans	
3			n
? = n * F(2)			ans



Memory	main	F(4)	F(3)	F(2)
? = F(4)	Result			
4	n			
4		n		
? = n * F(3)		ans		
3			n	
? = n * F(2)			ans	
2				n
? = n * F(1)				ans



Memory	main	F(4)	F(3)	F(2)	F(1)
? = F(4)	Result				
4	n				
4		n			
? = n * F(3)		ans			
3			n		
? = n * F(2)			ans		
2				n	
? = n * F(1)				ans	
1					n
? = n * F(0)					ans



Memory	main	F(4)	F(3)	F(2)	F(1)	F(0)
? = F(4)	Result					
4	n					
4		n				
? = n * F(3)		ans				
3			n			
? = n * F(2)			ans			
2				n		
? = n * F(1)				ans		
1					n	
? = n * F(0)					ans	
0						n
						ans

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Base Case Reached

Memory	main	F(4)	F(3)	F(2)	F(1)	F(0) = 1
? = F(4)	Result					
4	n					
4		n				
? = n * F(3)		ans				
3			n			
? = n * F(2)			ans			
2				n		
? = n * F(1)				ans		
1					n	
n * F(0) = 1					ans	

$$F(0) = 1$$



Memory	main	F(4)	F(3)	F(2)	F(1) = 1
? = F(4)	Result				
4	n				
4		n			
? = n * F(3)		ans			
3			n		
? = n * F(2)			ans		
2				n	
n * F(1) = 2				ans	



Memory	main	F(4)	F(3)	F(2) = 2
? = F(4)	Result			
4	n			
4		n		
? = n * F(3)		ans		
3			n	
n * F(2) = 6			ans	

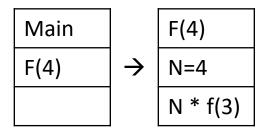


Memory	main	F(4)	F(3) = 6
? = F(4)	Result		
4	n		
4		n	
n * F(3) = 24		ans	

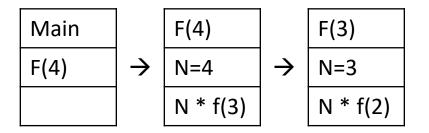


Memory	main	F(4) = 24
= F(4) = 24	Result	
4	n	

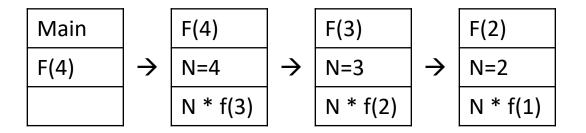




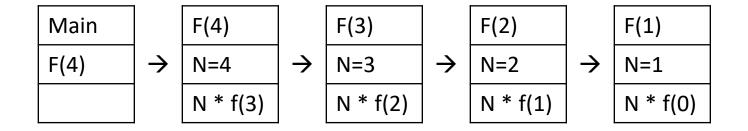




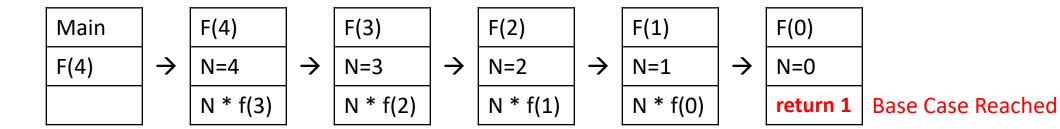




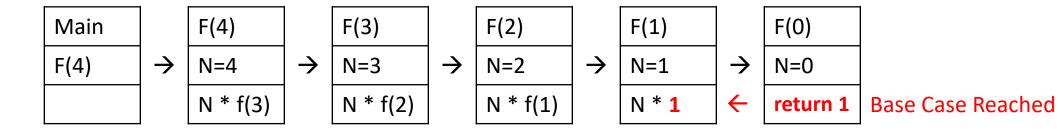




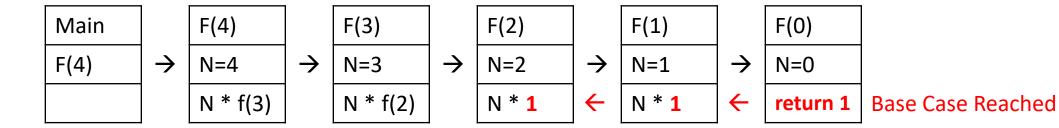




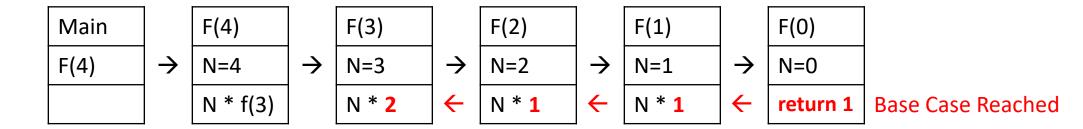




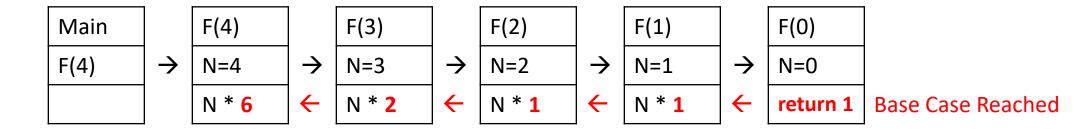




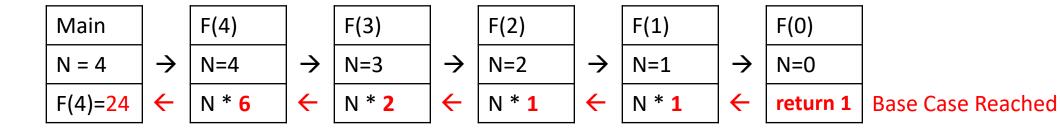














Factorial Iterative

```
unsigned int ifactorial( unsigned int n )
    int i;
    unsigned int ans = 1;
    for( i=1; i<=n; i++)</pre>
        ans *= i;
    return ans;
```



Summation

$$\sum_{i=0}^{n} i = 0 + 1 + \dots + (n-2) + (n-1) + n$$



Summation

How to define in terms of the first n-1 terms

$$\sum_{i=0}^{n} i = n + \sum_{i=0}^{n-1} i$$

- What will serve as our base case?
 - N = 0, answer is 0



Summation Criteria

N = 0

answer is 0

N > 0

answer is
$$n + \sum_{i=0}^{n-1} i$$



Drawing it out

- Assume n = 6
 - Sum(6) = 6 + Sum(5)
 - Sum(5) = 5 + Sum(4)
 - Sum(4) = 4 + Sum(3)
 - Sum(3) = 3 + Sum(2)
 - Sum(2) = 2 + Sum(1)
 - Sum(1) = 1 + Sum(0)
 - Sum(0) = 0

← Good, I know the answer



Substitute Going Back Up

- Assume n = 6
 - Sum(6) = 6 + Sum(5) = 6 + 15 = 21
 - Sum(5) = 5 + Sum(4) = 5 + 10 = 15
 - Sum(4) = 4 + Sum(3) = 4 + 6 = 10
 - Sum(3) = 3 + Sum(2) = 3 + 3 = 6
 - Sum(2) = 2 + Sum(1) = 2 + 1 = 3
 - Sum(1) = 1 + Sum(0) = 1 + 0 = 1
 - Sum(0) = 0 \leftarrow Good, I know the answer



Writing the function

```
unsigned int summation( unsigned int n )
{
    // write the base case first
    if( n == 0 )
        return 0;
```

}



Writing the function

```
unsigned int summation( unsigned int n )
    int ans;
    // write the base case first
    if( n == 0 )
        return 0;
    // compute the answer
    ans = n + summation(n - 1);
    return ans;
```



Stepping Through the Function

```
int main()
    int result;
    int n = 4;
    result = summation ( n );
    cout << "sum of 0 to " << n << " = "
         << result << endl;
    return 0;
```



Memory	main
? = S(4)	result
4	n



Memory	main	S(4)
? = S(4)	Result	
4	n	
4		n
? = N + S(3)		ans



Memory	main	S(4)	S(3)
? = S(4)	Result		
4	n		
4		n	
? = n + S(3)		ans	
3			n
? = n + S(2)			ans



Memory	main	S(4)	S(3)	S(2)
? = S(4)	Result			
4	n			
4		n		
? = n + S(3)		ans		
3			n	
? = n + S(2)			ans	
2				n
? = n + S(1)				ans



Memory	main	S(4)	S(3)	S(2)	S(1)
? = S(4)	Result				
4	n				
4		n			
? = n + S(3)		ans			
3			n		
? = n + S(2)			ans		
2				n	
? = n + S(1)				ans	
1					n
? = n + S(0)					ans



Memory	main	S(4)	S(3)	S(2)	S(1)	S(0)
? = S(4)	Result					
4	n					
4		n				
? = n + S(3)		ans				
3			n			
? = n + S(2)			ans			
2				n		
? = n + S(1)				ans		
1					n	
? = n + S(0)					ans	
0						n
						ans

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Base Case Reached

Memory	main	S(4)	S(3)	S(2)	S(1)	S(0) = 0
? = S(4)	Result					
4	n					
4		n				
? = n + S(3)		ans				
3			n			
? = n + S(2)			ans			
2				n		
? = n + S(1)				ans		
1					n	
n + S(0) = 1					ans	



Memory	main	S(4)	S(3)	S(2)	S(1) = 1
? = S(4)	Result				
4	n				
4		n			
? = n + S(3)		ans			
3			n		
? = n + S(2)			ans		
2				n	
n + S(1) = 3				ans	



Memory	main	S(4)	S(3)	S(2) = 3
? = S(4)	Result			
4	n			
4		n		
? = n + S(3)		ans		
3			n	
n + S(2) = 6			ans	

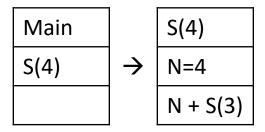


Memory	main	S(4)	S(3) = 6
? = S(4)	Result		
4	n		
4		n	
n + S(3) = 10		ans	

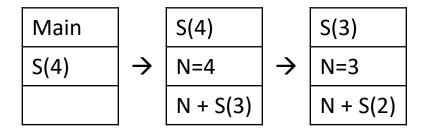


Memory	main	S(4) = 10
= S(4) = 10	Result	
4	n	

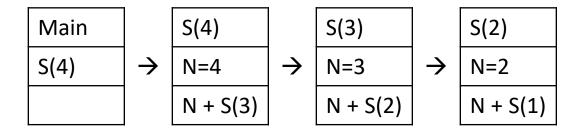




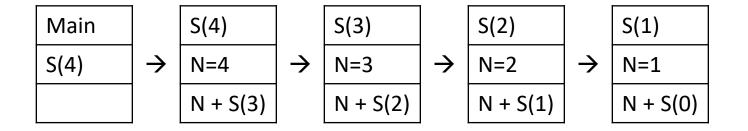




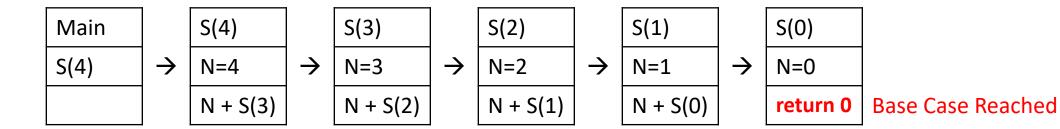




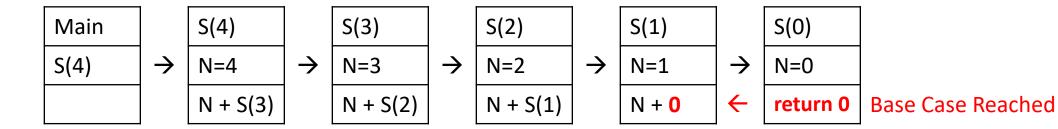




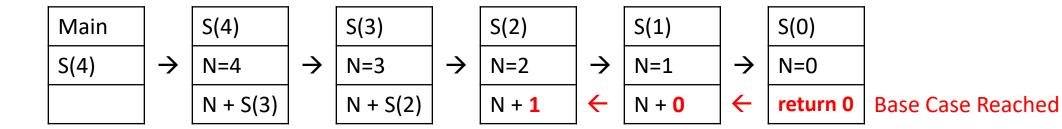




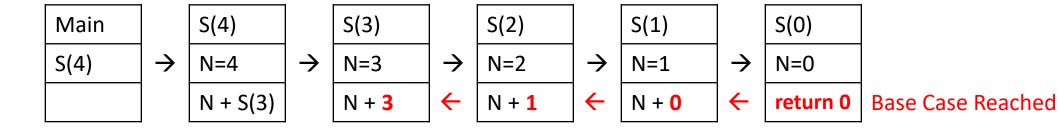




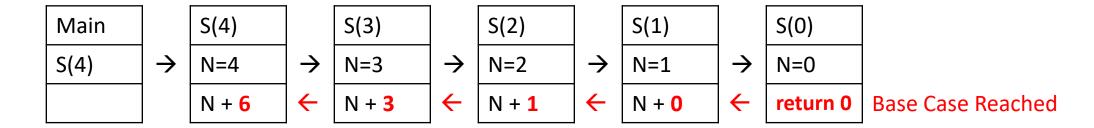




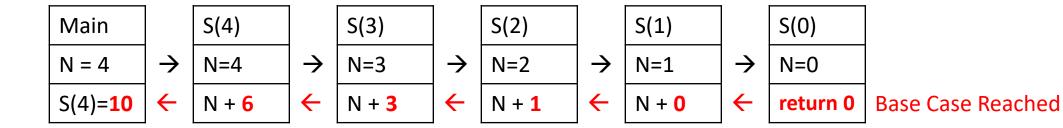














Summation Iterative

```
unsigned int isummation( unsigned int n )
    int i;
    unsigned int ans = 0;
    for( i=0; i<=n; i++)</pre>
        ans += i;
    return ans;
```



You Try It at Home

• Given two integer number a and b where a is less than b and a is greater than or equal to 0. Sum the range from a to b.

le a = 3, b = 7
Sum(3,7) =
$$3 + 4 + 5 + 6 + 7 = 25$$

Careful, Sum(3,3) = 3

To get started:

unsigned int sumRange(unsigned int a, unsigned int b);

