CSC 215

Math and Computer Science



Permutations

- Heart of many brute force algorithms
- Generate all possible combinations
- How many are there?

$$P(n,k) = \frac{n!}{(n-k)!}$$

$$n = 3, k = 3$$

$$P(3,3) = \frac{n!}{(n-k)!} = 3! = 6$$



$$P(3,3)=3!=6$$

• Set $\{0,1,2\}$, n = 3

$$\{0,1,2\}$$
 $\{0,2,1\}$ $\{1,0,2\}$ $\{1,2,0\}$ $\{2,0,1\}$ $\{2,1,0\}$

6 ways to order the 3 numbers



$$P(4,4) = 4! = 24$$

• Set {0,1,2,3}

```
\{0,1,2,3\} \{0,1,3,2\} \{0,2,1,3\} \{0,2,3,1\} \{0,3,1,2\} \{0,3,2,1\} \{1,0,2,3\} \{1,0,3,2\} \{1,2,0,3\} \{1,2,3,0\} \{1,3,0,2\} \{1,3,2,0\} \{2,0,1,3\} \{2,0,3,1\} \{2,1,0,3\} \{2,1,3,0\} \{2,3,0,1\} \{2,3,1,0\} \{3,0,1,2\} \{3,0,2,1\} \{3,1,0,2\} \{3,1,2,0\} \{3,2,0,1\} \{3,2,1,0\}
```



Solving P(3,3) Iteratively



Solving P(4,4) Iteratively



Solving P(10,10) Iteratively

```
void iter_permute10()
{
    int i,j,k,1,m,n,o,p,q,r;
    for(i=0; i<10; i++)
        for(j=0; j<10; j++)
        for(l=0; l<10; l++)
        for(m=0; m<10; m++)
        for(n=0; n<10; n++)
        for(p=0; o<10; o++)
        for(p=0; q<10; q++)
        for(r=0; r<10; r++)</pre>
```

```
if( i!=j && i!=k && i!=1 && i!=m && i!=n && i!=n && i!=p && i!=p && i!=p && i!=r &&
    j!=k && j!=1 && j!=m && j!=n && j!=n && j!=p && j!=q && j!=r &&
    k!=1 && k!=m && k!=n && k!=n && k!=p && k!=q && k!=r &&
    l!=m && l!=n && l!=n && l!=p && l!=q && l!=r &&
    m!=n && m!=n && m!=p && m!=q && m!=r &&
    n!=n && m!=n && n!=p && m!=q && m!=r &&
    n!=n && n!=p && n!=q && n!=r &&
    n!=n && n!=p && n!=q && n!=r &&
    n!=n && n!=p && n!=q && n!=r &&
    n!=n && n!=n && n!=n && n!=n &&
    n!=n && n!=n && n!=n && n!=r &&
    n!=n && n!=n && n!=n && n!=n && n!=n &&
    n!=n && n!=n && n!=n && n!=r &&
    n!=n && n!=n && n!=n && n!=n &&
    n!=n && n!=n && n!=n
```



Iterative Approach

- n must be known before programming
- n number of for loops will be needed
- If statement to make sure no two numbers are the same.



Recursive Approach

- N not necessarily known at time of programming
- Each recursive call will nest a for loop inside of another for loop
- We will not put duplicate numbers into our solution



What is Needed – Set of n Integers

- An array of size (n) to hold the solution
- An array to keep track of the numbers used in the solution
- A variable to keep track of how many numbers are in the solutions so far
- A variable to store the size of our set (n)



The Solution Array

- Assume n = 3
- int p[3];

P Array	2	0	1
Index	0	1	2

• The P array holds a solution set of {2,0,1}



The Used Array

- Lets look a placing a number someplace in the P array
- Put the number 2 in the zeroth spot of the P array

Р	2		
Index	0	1	2
Used	0	0	1
Index	0	1	2

 Mark the 2 spot of the used array with a 1 to show it is in the solution already



Other Variables

- Need a variable to keep track of how many numbers we have put into the P array.
 - Tells us we have inserted k number of n
 - Tells us the location (index) of where to put the next number in the P array.
 - Call it pos for position
- Need a variable to tell us how many items are in the set.
 - Call it n,



Filling a Spot in P

- n = 10
- pos = 4 <= some spot in the P array

Р	X	Х	X	X	?					
Index	0	1	2	3	4	5	6	7	8	9

Must try and put every number in this location

• There are numbers in the previous spots and we can not duplicate.



Filling a Spot in P - continued

```
        P
        x
        x
        x
        x
        ?

        Index
        0
        1
        2
        3
        4
        5
        6
        7
        8
        9
```



Moving to the next position

```
        P
        x
        x
        x
        x
        ?

        Index
        0
        1
        2
        3
        4
        5
        6
        7
        8
        9
```



The Base Case

- When I have put n number of items in the P array, we have a solution.
- When this happens, do something with the solution, like print it out.



Р Index 0 1 2 Used 0 0 0 Index 0 1 2

for(item=0; item<n; item++) // start new for loop</pre>

make recursive call to pos+1

Ρ Index

0		
0	1	2

Used Index

1	0	0
0	1	2



Р 0 Index 0 1 2 Used 0 0 1 Index 0 1 2

for(item=0; item<n; item++)</pre> Item = 0, used[item] = 1
← already in the solution Item = 1, used[item] = 0

// start new for loop

← move into P at pos, mark used make recursive call to pos+1

Ρ Index

0	1	
0	1	2

Used Index

1	1	0
0	1	2



Р Index

0	1	
0	1	2

Used Index

1	1	0
0	1	2

for(item=0; item<n; item++)</pre>

Item = 1, used[item] = 1

Item = 2, used[item] = 0

Ρ Index

0	1	2
0	1	2

// start new for loop

Item = 0, used[item] = 1
← 0 is in solution already

← 1 is in solution already

← move into P at pos, mark used make recursive call to pos+1

Used

1	1	1
0	1	2



P Index

0	1	2
0	1	2

Used Index

1	1	1
0	1	2

Base case is reached (pos == n)
print out solution to screen.

Return to try other numbers

0 1 2



Р 0 1 2 Index 0 1 2

Used 1 1 1 Index 0 2 1

for(item=0; item<n; item++)</pre>

Item = 1, used[item] = 1

Item = 2, used[item] = 0

Item = 3,

Р Index

0	1	
0	1	2

Item = 0, used[item] = 1
← 0 is in solution already

← 1 is in solution already

← Mark 2 as unused

← return, tried all numbers

Used	1	1	0
ndex	0	1	2



P 0 1 1 Index 0 1 2

 Used
 1
 1
 0

 Index
 0
 1
 2

for(item=0; item<n; item++)</pre>

Item = 0, used[item] = 1

Item = 1, used[item] = 0

Item = 2, used[item] = 0

P Index

0	2	
0	1	2

← 0 is in solution already

← Mark 1 as unused

← move into P at pos, mark used make recursive call to pos+1

Jsed	1	0	
ndex	0	1	



Р Index

0	2	
0	1	2

Used Index

1	0	1
0	1	2

for(item=0; item<n; item++) // start new for loop</pre>

Item = 1, used[item] = 0

Item = 0, used[item] = 1
← 0 is in solution already

← move into P at pos, mark used make recursive call to pos+1

Р Index

0	2	1
0	1	2

Used

1	1	1
0	1	2



P Index

0	2	1
0	1	2

Used Index

1	1	1
0	1	2

Base case is reached (pos == n)
print out solution to screen.

Return to try other numbers

0 1 2

0 2 1



Р Index

0	2	1
0	1	2

Used Index

d	1	1	1
X	0	1	2

for(item=0; item<n; item++)</pre>

Item = 1, used[item] = 0

Item = 2, used[item] = 1

Item = 3

Р Index

0	2	
0	1	2

Item = 0, used[item] = 1
← 0 is in solution already

← Mark 1 as unused

← return

Used

1	0	1
0	1	2



Р 2 0 Index 0 1 2 Used 1 0 1 Index 0 2 1

for(item=0; item<n; item++)</pre>

Item = 1, used[item] = 0

Item = 2, used[item] = 0

Item = 3

Р Index

0		
0	1	2

Item = 0, used[item] = 1
← 0 is in solution already

← Mark 1 as unused

← Mark 2 as unused

← return

Used Index

1	0	0
0	1	2



Р Index

0		
0	1	2

Used Index

	1	0	0
(0	1	2

← Mark 0 as unused

← move into P at pos, mark used make recursive call to pos+1

Р		
In	de	X

1		
0	1	2

Use

Used	0	1	0
Index	0	1	2



Р 1 Index 0 1 2 Used 0 0 1 Index 0 1 2

for(item=0; item<n; item++) // start new for loop</pre>

Item = 0, used[item] = 1
← move into P at pos, mark used make recursive call to pos+1

Ρ Index

1	0	
0	1	2

Used Index

1	1	0
0	1	2



Р Index

1	0	
0	1	2

Used Index

1	1	0
0	1	2

for(item=0; item<n; item++) // start new for loop</pre>

Item = 1, used[item] = 1

Item = 2, used[item] = 0

Ρ Index

1	0	2
0	1	2

Item = 0, used[item] = 1
← 0 is in solution already

← 1 is in solution already

1

← move into P at pos, mark used make recursive call to pos+1

1

1

Used



P Index

1	0	2
0	1	2

Used Index

1	1	1
0	1	2

Base case is reached (pos == n) print out solution to screen.

Return to try other numbers



Р Index

1	0	2
0	1	2

Used Index

1	1	1
0	1	2

for(item=0; item<n; item++)</pre>

Item = 1, used[item] = 1

Item = 2, used[item] = 0

Item = 3

Р Index

1	0	
0	1	2

Item = 0, used[item] = 1
← 0 is in solution already

← 1 is in solution already

← Mark 2 as unused

← return

Used

1	1	0
0	1	2



P Index

1	0	
0	1	2

Used Index

1	1	0
0	1	2

for(item=0; item<n; item++)</pre>

Item = 0, used[item] = 0

Item = 1, used[item] = 1

Item = 2, used[item] = 0

P Index

1	2	
0	1	2

← Mark 0 as unused

← 1 is in solution already

← move into P at pos, mark used make recursive call to pos+1

Used

0	1	1
0	1	2



P 1 2 Index 0 1 2

 Used
 0
 1
 1

 Index
 0
 1
 2

for(item=0; item<n; item++)
Item = 0, used[item] = 0</pre>

// start new for loop

← move into P at pos, mark used make recursive call to pos+1

P Index

1	2	0
0	1	2

Used

1	1	1
0	1	2



P Index

1	2	0
0	1	2

Used Index

ed	1	1	1
lex	0	1	2

Base case is reached (pos == n) print out solution to screen.

Return to try other numbers



P 1 2 0 Index 0 1 2
 Used
 1
 1
 1

 Index
 0
 1
 2

for(item=0; item<n; item++)</pre>

Item = 0, used[item] = 0

Item = 1, used[item] = 1

Item = 2, used[item] = 1

Item = 3

P Index

1	2	
0	1	2

← Mark 0 as unused

← Return

Used

0	1	1
0	1	2



Р 2 1 Index 0 1 2 Used 0 1 1 Index 0 2 1

for(item=0; item<n; item++)</pre>

Item = 0, used[item] = 0
← Mark 0 as unused

Item = 1, used[item] = 1

Item = 2, used[item] = 0

Item = 3

Р Index

1		
0	1	2

← 1 is in solution already

← Mark 2 as unused

← return

d	0	1	0
ex	0	1	2



P 1 1 1 1 1 2

 Used
 0
 1
 0

 Index
 0
 1
 2

for(item=0; item<n; item++)</pre>

Item = 0, used[item] = 0

Item = 1, used[item] = 0

Item = 2, used[item] = 0

P Index

2		
0	1	2

← Mark 0 as unused

← Mark 1 as unused

move into P at pos, mark used
make recursive call to pos+1

0	0	1
0	1	2



P 2 1 1 2

 Used
 0
 0
 1

 Index
 0
 1
 2

for(item=0; item<n; item++)
Item = 0, used[item] = 0</pre>

// start new for loop

← move into P at pos, mark used make recursive call to pos+1

P Index

2	0	
0	1	2

1	0	1
0	1	2



Ρ Index

2	0	
0	1	2

Used Index

d	1	0	1
X	0	1	2

for(item=0; item<n; item++)</pre>

Item = 1, used[item] = 0

// start new for loop

Item = 0, used[item] = 1
← 0 already in solution

← move into P at pos, mark used make recursive call to pos+1

1

Р Index

2	0	1
0	1	2

Used

Oscu	1	1
Index	0	1



P Index

2	0	1
0	1	2

Used Index

1	1	1
0	1	2

Base case is reached (pos == n)
print out solution to screen.

Return to try other numbers



Р Index

2	0	1
0	1	2

Used Index

1	1	1
0	1	2

for(item=0; item<n; item++)</pre>

Item = 1, used[item] = 0

Item = 2, used[item] = 1

Item = 3

Р Index

2	0	
0	1	2

← Mark 1 as unused

← 2 already in solution

← return

Used

Index

1	0	1
0	1	2



P Index

2	0	
0	1	2

Used Index

1	0	1
0	1	2

for(item=0; item<n; item++)</pre>

Item = 0, used[item] = 1

Item = 1, used[item] = 0

← Mark 0 as unused

← move into P at pos, mark used make recursive call to pos+1

1

1

1

P Index

2	1	
0	1	2

Used

Index	0

0



P 2 1 1 Index 0 1 2

 Used
 0
 1
 1

 Index
 0
 1
 2

for(item=0; item<n; item++)
Item = 0, used[item] = 0</pre>

// start new for loop

← move into P at pos, mark used make recursive call to pos+1

P Index

2	1	0
0	1	2

1	1	1
0	1	2



P Index

2	1	0
0	1	2

Used Index

1	1	1
0	1	2

Base case is reached (pos == n)
print out solution to screen.

0 1 2

Return to try other numbers



Р Index

2	1	0
0	1	2

Us Ind

ed	1	1	1
dex	0	1	2

for(item=0; item<n; item++)</pre>

Item = 3

Р Index

2	1	
0	1	2

← 1 already in solution

← 2 already in solution

← return

Used

Index

0	1	1
0	1	2



P 2 1 1 Index 0 1 2

 Used
 0
 1
 1

 Index
 0
 1
 2

for(item=0; item<n; item++)</pre>

Item = 0, used[item] = 0

Item = 1, used[item] = 0

Item = 2, used[item] = 1

Item = 3

P Index

2		
0	1	2

← Mark 0 as unused

← Mark 1 as Unused

← 2 already in solution

← Return

0	0	1
0	1	2



Item
$$= 3$$

Р	
Index	

0	1	2

- ← Mark 0 as unused
 - ← Mark 1 as unused
 - ← Mark 2 as unused
 - ← return

Used	
Index	

0	0	0
0	1	2

