Hive and Shark
SQL and Rich Analytics at Scale

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Motivation

- MapReduce is hard to program.

- No schema, lack of query languages, e.g., SQL.

Solution

- Adding tables, columns, partitions, and a subset of SQL to unstructured data.

Hive

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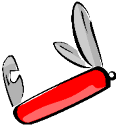
- Initially developed by Facebook.
- Focuses on scalability and extensibility.

Scalability

- Massive scale out and fault tolerance capabilities on commodity hardware.
- Can handle petabytes of data.

Extensibility

- **Data types**: primitive types and complex types.
- User Defined Functions (UDF).
- **Serializer/Deserializer**: text, binary, JSON ...
- **Storage**: HDFS, Hbase, S3 ...



RDBMS vs. Hive

	RDBMS	Hive
Language	SQL	HiveQL
Update Capabilities	INSERT, UPDATE, and DELETE	INSERT OVERWRITE; no UPDATE or DELETE
OLAP	Yes	Yes
OLTP	Yes	No
Latency	Sub-second	Minutes or more
Indexes	Any number of indexes	No indexes, data is always scanned (in parallel)
Data size	TBs	PBs

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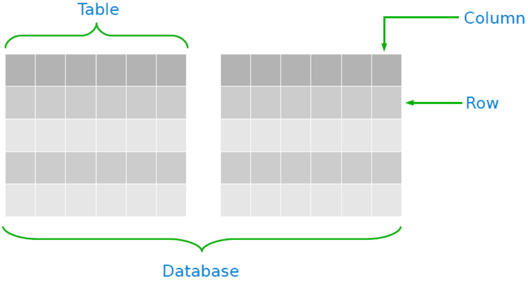
- Online Analytical Processing (OLAP): allows users to analyze database information from multiple database systems at one time.
- Online Transaction Processing (OLTP): facilitates and manages transaction-oriented applications.

Hive Data Model

- Re-used from RDBMS:
  - Database: Set of Tables.
  - Table: Set of Rows that have the same schema (same columns).
  - Row: A single record; a set of columns.
  - Column: provides value and type for a single value.



Hive Data Model - Table

- Analogous to tables in relational databases.
- Each table has a corresponding HDFS directory.
- For example data for table `customer` is in the directory
  `/db/customer`.

Hive Data Model - Partition

- A coarse-grained partitioning of a table based on the value of a column, such as a date.
- Faster queries on slices of the data.
- If `customer` is partitioned on column `country`, then data with a particular country value `SE`, will be stored in files within the directory `/db/customer/country=SE`.

Hive Data Model - Bucket

- Data in each partition may in turn be divided into buckets based on the hash of a column in the table.
- For more efficient queries.
- If `customer` country partition is subdivided further into buckets, based on `username` (hashed on username), the data for each bucket will be stored within the directories:

```
/db/customer/country=SE/000000_0
...
/db/customer/country=SE/000000_5
```

Column Data Types

- Primitive types
  - integers, float, strings, dates and booleans
- Nestable collections
  - array and map
- User-defined types
  - Users can also define their own types programmatically

Hive Operations

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  - Load and Insert (overwrite)
  - Does not support updating and deleting (later versions support update and delete)

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- SQL operations
  - Select, Filter, Join, Groupby

DDL Operations (1/3)

- Create tables

```
-- Creates a table with three columns
CREATE TABLE customer (id INT, name STRING, address STRING)
ROW FORMAT DELIMITED FIELDS TERMINATED BY '\t';
```

DDL Operations (1/3)

- Create tables

```
-- Creates a table with three columns
CREATE TABLE customer (id INT, name STRING, address STRING)
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```

- Create tables with partitions

```
-- Creates a table with three columns and a partition column
-- /db/customer2/country=SE;
-- /db/customer2/country=IR;
CREATE TABLE customer2 (id INT, name STRING, address STRING)
PARTITION BY (country STRING)
```

DDL Operations (2/3)

- Create tables with buckets

```
-- Specify the columns to bucket on and the number of buckets
-- /db/customer3/000000_0
-- /db/customer3/000000_1
-- /db/customer3/000000_2
set hive.enforce.bucketing = true;
CREATE TABLE customer3 (id INT, name STRING, address STRING)
CLUSTERED BY (id) INTO 3 BUCKETS;
```

DDL Operations (2/3)

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```

- Browsing through tables

```
-- lists all the tables
SHOW TABLES;

-- shows the list of columns
DESCRIBE customer;
```

DDL Operations (3/3)

- Altering tables

```sql
-- rename the customer table to alaki
ALTER TABLE customer RENAME TO alaki;

-- add two new columns to the customer table
ALTER TABLE customer ADD COLUMNS (job STRING);
ALTER TABLE customer ADD COLUMNS (grade INT COMMENT 'some
comment');
```

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- Dropping tables

```
DROP TABLE customer;
```

DML Operations

- Loading data from flat files.

```
-- if 'LOCAL' is omitted then it looks for the file in HDFS.
-- the 'OVERWRITE' signifies that existing data in the table is deleted.
-- if the 'OVERWRITE' is omitted, data are appended to existing data sets.
LOAD DATA LOCAL INPATH 'data.txt' OVERWRITE INTO TABLE customer;

-- loads data into different partitions
LOAD DATA LOCAL INPATH 'data1.txt' OVERWRITE INTO TABLE
customer2 PARTITION (country='SE');
LOAD DATA LOCAL INPATH 'data2.txt' OVERWRITE INTO TABLE
customer2 PARTITION (country='IR');
```

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```

- Store the query results in tables

```
INSERT OVERWRITE TABLE customer SELECT * From old_customers;
```

SQL Operations (1/3)

- Selects and filters

```
SELECT id FROM customer2 WHERE country='SE';

-- selects all rows from customer table into a local directory
INSERT OVERWRITE LOCAL DIRECTORY '/tmp/hive-sample-out' SELECT *
FROM customer;

-- selects all rows from customer2 table into a directory in hdfs
INSERT OVERWRITE DIRECTORY '/tmp/hdfs_ir' SELECT * FROM customer2
WHERE country='IR';
```

SQL Operations (2/3)

- Aggregations and groups

```
SELECT MAX(id) FROM customer;

SELECT country, COUNT(*), SUM(id) FROM customer2 GROUP BY
country;

INSERT TABLE high_id_customer SELECT c.name, COUNT(*) FROM
customer c WHERE c.id > 10 GROUP BY c.name;
```

SQL Operations (3/3)

- Join

```
CREATE TABLE customer (id INT, name STRING, address
STRING) ROW FORMAT DELIMITED FIELDS TERMINATED BY
'\t';

CREATE TABLE order (id INT, cus_id INT, prod_id INT, price
INT) ROW FORMAT DELIMITED FIELDS TERMINATED BY '\t';

SELECT * FROM customer c JOIN order o ON (c.id =
o.cus_id);
```

User-Defined Function (UDF)

```
package com.example.hive.udf;

import org.apache.hadoop.hive.ql.exec.UDF; import
org.apache.hadoop.io.Text;

public final class Lower extends UDF {
  public Text evaluate(final Text s) {
    if (s == null) {
      return null;
    }
     return new Text(s.toString().toLowerCase());
   }
}
```

```
-- Register the class
CREATE FUNCTION my_lower AS 'com.example.hive.udf.Lower';

-- Using the function
SELECT my_lower(title), sum(freq) FROM titles GROUP BY
my_lower(title);
```

Executing SQL Questions

- Processes HiveQL statements and generates the execution plan through three-phase processes.

  1. Query parsing: transforms a query string to a parse tree representation.
  2. Logical plan generation: converts the internal query representation to a logical plan, and optimizes it.
  3. Physical plan generation: split the optimized logical plan into multiple map/reduce and HDFS tasks.

Optimization (1/2)

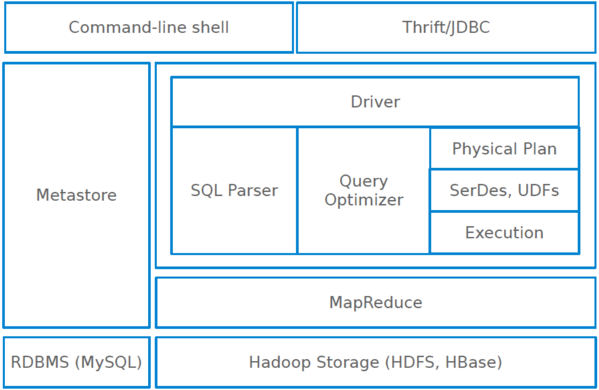
- Column pruning
  - Projecting out the needed columns.
- Predicate pushdown
  - Filtering rows early in the processing, by pushing down predicates to the scan (if possible).
- Partition pruning
  - Pruning out files of partitions that do not satisfy the predicate.

Optimization (2/2)

- Map-side joins
  - The small tables are replicated in all the mappers and joined with other tables.
  - No reducer needed.

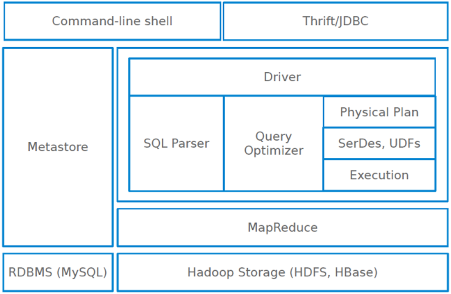
- Join reordering
  - Only materialized and kept small tables in memory.
  - This ensures that the join operation does not exceed memory limits on the reducer side.

Hive Components (1/8)



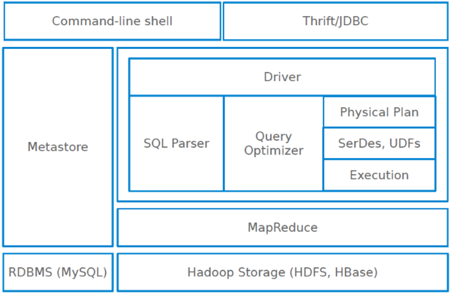
Hive Components (2/8)

- External interfaces
  - User interfaces, e.g., CLI and web UI
  - Application programming interfaces, e.g., JDBC and ODBC
  - Thrift, a framework for cross-language services.



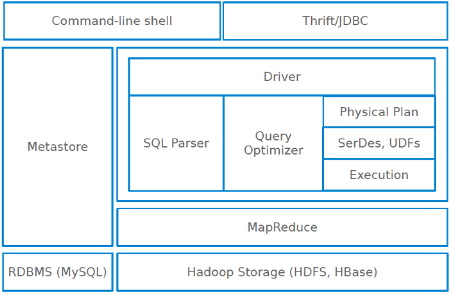
Hive Components (3/8)

- Driver
  - Manages the life cycle of a HiveQL statement during compilation, optimization and execution.



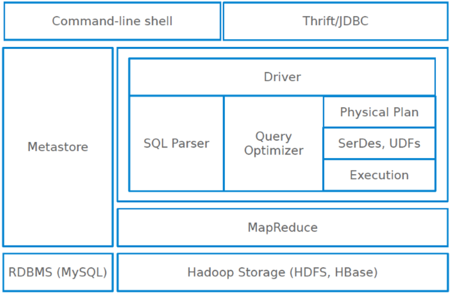
Hive Components (4/8)

- Compiler (Parser/Query Optimizer)
  - Translates the HiveQL statement into a a logical plan, and optimizes it.



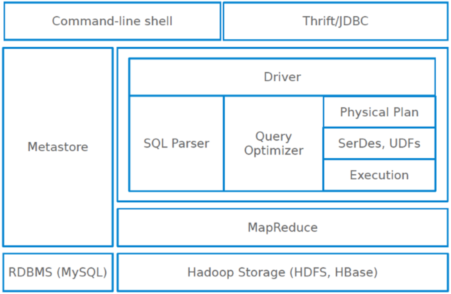
Hive Components (5/8)

- Physical plan
  - Transforms the logical plan into a DAG of Map/Reduce jobs.



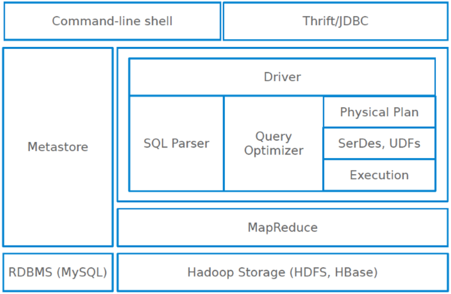
Hive Components (6/8)

- Execution engine
  - The driver submits the individual mapreduce jobs from the DAG to the execution engine in a topological order.



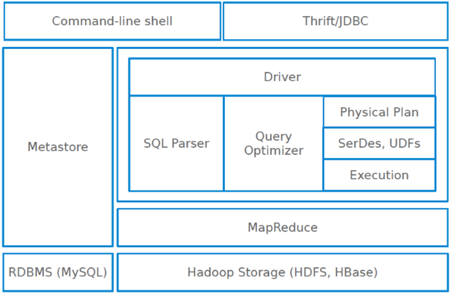
Hive Components (7/8)

- SerDe
  - Serializer/Deserializer allows Hive to read and write table rows in any custom format.



Hive Components (8/8)

- Metastore
  - The system catalog.
  - Contains metadata about the tables.
  - Metadata is specified during table creation and reused every time the table is referenced in HiveQL.
  - Metadatas are stored on either a traditional relational database, e.g., MySQL, or file system and not HDFS.



Hive on SHARK

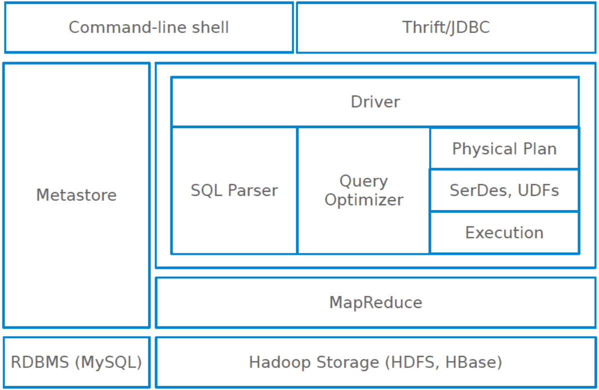
Spark RDD - Reminder

- RDDs are immutable, partitioned collections that can be created through various transformations, e.g., map, groupByKey, join.

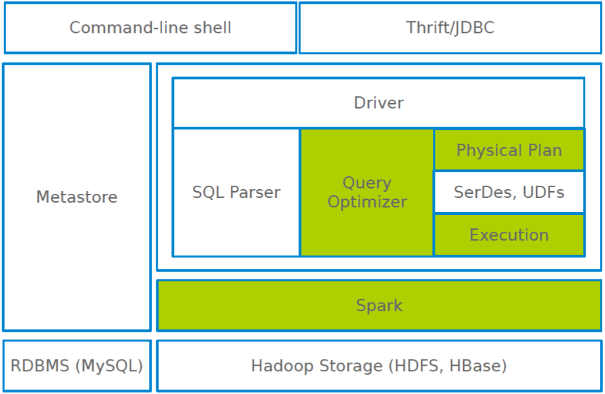
Executing SQL over Spark RDDs

- Shark runs SQL queries over Spark using three-step process:
  1. Query parsing: Shark uses Hive query compiler to parse the query and generate a parse tree.
  2. Logical plan generation: the tree is turned into a logical plan and basic logical optimization is applied.
  3. Physical plan generation: Shark applies additional optimization and creates a physical plan consisting of transformations on RDDs.

Hive Components



Shark Components



Shark and Spark

- Shark extended RDD execution model:

- Partial DAG Execution (PDE): to re-optimize a running query after running the first few stages of its task DAG.
- In-memory columnar storage and compression: to process relational data efficiently.
- Control over data partitioning.

Partial DAG Execution (1/2)

- How to optimize the following query?

```
SELECT * FROM table1 a JOIN table2 b ON (a.key = b.key) WHERE
my_crazy_udf(b.field1, b.field2) = true;
```

Partial DAG Execution (1/2)

- How to optimize the following query?

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```

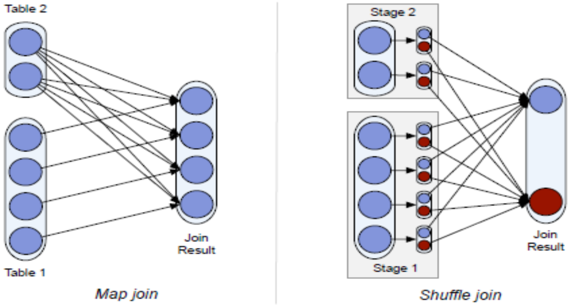
- It can not use cost-based optimization techniques that rely on accurate a priori data statistics.
- They require dynamic approaches to query optimization.
- Partial DAG Execution (PDE): dynamic alteration of query plans based on data statistics collected at run-time.

Partial DAG Execution (2/2)

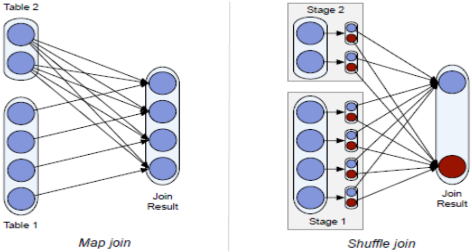
- The workers gather customizable statistics at global and per-partition granularities at run-time.
- Each worker sends the collected statistics to the master.
- The master aggregates the statistics and alters the query plan based on such statistics.

Join Optimization

Partial DAG execution can be used to perform several run-time optimizations for join queries

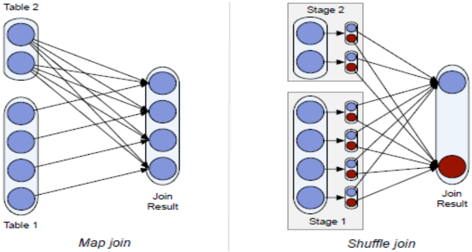


Join Optimization



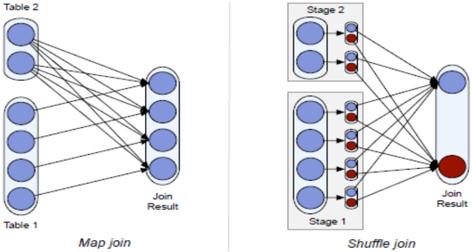
In shuffle join, both join tables are hash-partitioned by the join key. Each reducer joins corresponding partitions using a local join algorithm, which is chosen by each reducer based on runtime statistics

Join Optimization



In map join (also known as broadcast join) a small input table is broadcast to all nodes, where it is joined with each partition of a large table. This approach can result in significant cost savings by avoiding an expensive repartitioning and shuffling phase

Join Optimization



Map join is only worthwhile if some join inputs are small, so Shark uses partial DAG execution to select the join strategy at runtime based on its inputs' exact sizes

Columnar Memory Store

- Simply caching Hive records as JVM objects is inefficient.
- 12 to 16 bytes of overhead per object in JVM implementation:
  - e.g., storing a 270MB table as JVM objects uses approximately 971 MB of memory.
- Shark employs column-oriented storage using arrays of primitive objects.

Data Partitioning

- Shark allows co-partitioning two tables, which are frequently joined together, on a common key for faster joins in subsequent queries.

Shark/Spark Integration

- Shark provides a simple API for programmers to convert results from SQL queries into a special type of RDDs: `sql2rdd`.

```
val youngUsers = sql2rdd("SELECT * FROM users WHERE age < 20")
println(youngUsers.count)
val featureMatrix = youngUsers.map(extractFeatures(_))
kmeans(featureMatrix)
```

Summary

- Operators: DDL, DML, SQL
- Hive architecture vs. Shark architecture
- Add advance features to Spark, e.g., PDE, columnar memory store

Questions
?