



A *tabular* method!

## Basics of Dynamic Programming

### Dynamic Programming

- Not a specific algorithm, but a technique (like divide-and-conquer).
- Developed back in the day when “programming” meant “**tabular method**” (like linear programming). Doesn’t really refer to computer programming.
- Used for optimization problems (a set of choices must be made in order to arrive at an optimal solution):
  - Find a solution with *the* optimal value.
  - Minimization or maximization.

## Dynamic Programming

- Another strategy for designing algorithms is *dynamic programming*
  - A metatechnique, not an algorithm (like divide & conquer)
  - The word “programming” is historical and predates computer programming
- Use when problem breaks down into recurring small subproblems

## Optimization Problems

- In an optimization problem, there are typically many *feasible* solutions for any input instance  $I$
- For each solution  $S$ , we have a *cost* or *value* function  $f(S)$
- Typically, we wish to find a feasible solution  $S$  such that  $f(S)$  is either *minimized* or *maximized*
- Thus, when designing an algorithm to solve an optimization problem, we must prove the algorithm produces a best possible solution.

## Principle of Optimality

- In book, this is termed “Optimal substructure”
- An optimal solution contains within it optimal solutions to subproblems.
- More detailed explanation
  - Suppose solution  $S$  is optimal for problem  $P$ .
  - Suppose we decompose  $P$  into  $P_1$  through  $P_k$  and that  $S$  can be decomposed into pieces  $S_1$  through  $S_k$  corresponding to the subproblems.
  - Then solution  $S_i$  is an optimal solution for subproblem  $P_i$

Divide-and-Conquer	Dynamic Programming
Combines solutions of subproblems to solve the original problem	
<i>Disjoint</i> subproblems	<i>Overlapping</i> subproblems

## Dynamic Programming

- Dynamic programming is a divide-and-conquer technique at heart
- That is, we solve larger problems by patching together solutions to smaller problems
- Dynamic programming can achieve efficiency by storing solutions to subproblems to avoid redundant computations
  - We typically avoid redundant computations by computing solutions in a bottom-up fashion

## Dynamic Programming (DP)

- Like divide-and-conquer, solve problem by combining the solutions to sub-problems.
- Differences between divide-and-conquer and DP:
  - **Independent** sub-problems, solve sub-problems **independently** and **recursively**, (so same sub(sub)problems solved **repeatedly**)
  - Sub-problems are **dependent**, i.e., sub-problems **share** sub-sub-problems, every sub(sub)problem solved **just once**, solutions to sub(sub)problems are **stored in a table** and used for solving higher level sub-problems.

## Optimal Substructure

- Optimal substructure
- Overlapping subproblems

## Optimal Substructure – contd.

- Optimal (sub)structure
  - An optimal solution to the problem contains within it optimal solutions to subproblems.
- Overlapping subproblems
  - The space of subproblems is “small” in that a recursive algorithm for the problem solves the same subproblems over and over. Total number of distinct subproblems is typically polynomial in input size.
- (Reconstruction an optimal solution)

## Optimal Substructure – contd.

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- We say that a problem exhibits **optimal substructure** if an optimal solution to the problem contains within it optimal solution to subproblems.
- Example: Matrix-multiplication problem

## Optimal Substructure – contd.

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Optimal substructure varies across problem domains in two ways:

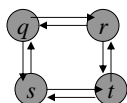
1. how many subproblems are used in an optimal solution to the original problem, and
2. how many choices we have in determining which subproblem(s) to use in an optimal solution.

## Subtleties when Determining Optimal Structure

- Take care that optimal structure does not apply even it looks like to be in first sight.
- Unweighted shortest path:
  - Find a path from  $u$  to  $v$  consisting of fewest edges.
  - Can be proved to have optimal substructures.

## Subtleties when Determining Optimal Structure – contd.

- Unweighted longest simple path:
  - Find a simple path from  $u$  to  $v$  consisting of most edges.
  - Figure 15.4 shows it does not satisfy optimal substructure.
- Independence (no share of resources) among subproblems if a problem has optimal structure.



$q \rightarrow r \rightarrow t$  is the longest simple path from  $q$  to  $t$ .  
But  $q \rightarrow r$  is not the longest simple path from  $q$  to  $r$ .

## Efficient Top-Down Implementation

- We can implement any dynamic programming solution top-down by storing computed values in the table
  - If all values need to be computed anyway, bottom up is more efficient
  - If some do not need to be computed, top-down may be faster

## Rules for Dynamic Programming

1. **Characterize** the structure of an optimal solution.
2. **Derive** a **recursive formula** for computing the values of optimal solutions.
3. **Compute the value** of an optimal solution **in a bottom-up fashion** (top-down is also applicable).
4. Construct an optimal solution from computed information.



## When is dynamic programming effective?

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- Dynamic programming works best on objects that are linearly ordered and cannot be rearranged
  - characters in a string
  - files in a filing cabinet
  - points around the boundary of a polygon
  - the left-to-right order of leaves in a search tree.
- Whenever your objects are ordered in a left-to-right way, dynamic programming must be considered.

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The End