

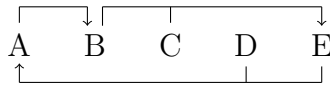
# CIS 551: Assignment 5

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## Problem 1: 19.2

Consider a relation  $R$  with five attributes  $ABCDE$ . You are given the following dependencies.



### Part 1)

List all keys for  $R$

We can see that  $A \rightarrow B$  which means we can rewrite the relationships as  $ACDE$ . Giving us



With  $A$  and  $C$  we are given  $E$ . Thus we have

$A \quad C \quad D$

And we no longer have any relationships.

Similarly we can get the other keys. Giving us the answer:

$ACD, BCD, CDE$

### Part 2)

Is  $R$  in 3NF?

Yes.  $B, E, A$  (the right sides from the FDs) are all part of keys.

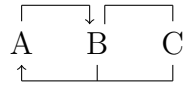
### Part 3)

Is  $R$  in BCNF?

No. None of the left sides of the FDs contain a key and there is no super key.

## Problem 2: 19.6

Suppose that we have the following three tuples in a legal instance of a relation scheme  $S$  with three attributes  $ABC$  (listed in order):  $(1, 2, 3)$ ,  $(4, 2, 3)$ , and  $(5, 3, 3)$ .

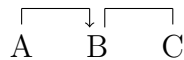
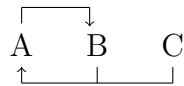
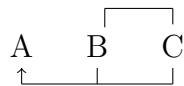


### Part 1)

Which of the following dependencies can you infer does *not* hold over scheme  $S$ ?

(a)  $A \rightarrow B$ , (b)  $BC \rightarrow A$ , (c)  $B \rightarrow C$

We'll look at the following three pictures



We see that the third diagram is the only legal dependency, which is where the dependency  $BC \rightarrow A$  is removed.  $\therefore$  (b) does **not** hold.

### Part 2)

Can you identify any dependencies that hold over  $S$ ?

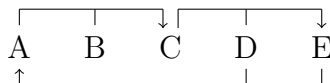
We are only given an instance of  $S$  and can't determine what relationships hold for **all** instances of  $S$ .

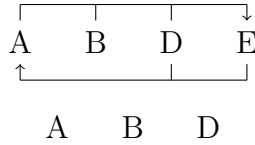
## Problem 3

Removed from assignment

## Problem 4

Consider relation  $R = ABCDE$  with the below relationships. Convert to BCNF.





## Problem 5: 19.13

Consider a relation  $R$  with five attributes  $ABCDE$

### Part 1)

For each of the following instances of  $R$ , state whether it violates (a) the FD  $BC \rightarrow D$  and (b) the MVD  $BC \twoheadrightarrow D$

- (a)  $\{\}$  No.
- (b)  $\{(a, 2, 3, 4, 5), (2, a, 3, 5, 5)\}$   $BC \rightarrow D$  is violated if  $a = 2$
- (c)  $\{(a, 2, 3, 4, 5), (2, a, 3, 5, 5), (a, 2, 3, 4, 6)\}$   $BC \rightarrow D$  and  $BC \twoheadrightarrow D$  are both violated if  $a = 2$ . We need another tuple that ends in 5, 6.
- (d)  $\{(a, 2, 3, 4, 5), (2, a, 3, 4, 5), (a, 2, 3, 6, 5)\}$   $BC \rightarrow D$  is violated because if  $a = 2$  then there is a duplicate row.
- (e)  $\{(a, 2, 3, 4, 5), (2, a, 3, 7, 5), (a, 2, 3, 4, 6)\}$  Both  $BC \rightarrow D$  and  $BC \twoheadrightarrow D$  are violated.  $BC \rightarrow D$  is violated because of the second tuple.  $BC \twoheadrightarrow D$  is violated because we must have another tuple that pairs with the second tuple.
- (f)  $\{(a, 2, 3, 4, 5), (2, a, 3, 4, 5), (a, 2, 3, 6, 5), (a, 2, 3, 6, 6)\}$   $BC \rightarrow D$  is violated because there is not a unique value defined by  $BC$ .  $BC \twoheadrightarrow D$  is violated because the entry for the  $D$ th column in the 4<sup>th</sup> tuple isn't the same as in the 1<sup>st</sup> tuple.
- (g)  $\{(a, 2, 3, 4, 5), (a, 2, 3, 6, 5), (a, 2, 3, 6, 6), (a, 2, 3, 4, 6)\}$  Just  $BC \rightarrow D$  is violated because the third tuple's last entry isn't also a 5.

### Part 2)

If each instance for  $R$  listed above is legal, what can you say about the DF  $A \rightarrow B$ ? We cannot say anything, because there are not unique relationships.

## Problem 6

Let  $R$  be a relation,  $X$  a set of attributes of  $R$ , and  $A$  an attribute of  $R$ . (Also denote that  $XA$  the result of adding  $A$  to  $X$ ) Define the support of  $X$ , written as  $\#X$ , as the number of distinct tuples in  $R|X$  ( $R$  restricted to, or projected onto, the attributes of  $X$ ). Prove that if  $X \rightarrow A$ , then  $\#X = \#XA$ .

Proof by contradiction:

Let  $X = BC$  such that  $BC \rightarrow A$ .

A	B	C
1	2	3
4	5	6
7	2	3

We see here that there is a violation in the third row because  $\#X = 2$  and  $\#XA = 3$ . This is not a legal relationship. We note that we only have legal schemas if there is an injective relationship. If we have an injective relationship (onto) then we would have the relationship  $\#X \leq \#XA$  (allows for surjection). For strict equality to hold, there needs to be a bijective relationship.

## Problem 7

Prove that if  $R$  is in 3NF and has only one candidate key, then  $R$  is in BCNF.

## Problem 8

Insert the following values into an initially empty B+ tree with parameter  $d = 2$  and values 17, 11, 50, 22, 5, 35, 42, 60, 15, 30, 25, 27, 37, 40, 20.

## Problem 9

Repeat the previous problem with  $d = 3$

## Problem 10: 10.8 part 1

Assume that you have just built a dense B+ tree index using Alternative (2) on a heap file containing 20,000 records. The key field for this B+ tree index is a 40-byte string, and it is a candidate key. Pointers (le., record ids and page ids) and (at most) 10-byte values. The size of one disk page is 1000 bytes. The index was built in a bottom-up fashion using the bulk loading algorithm, and the nodes at each level were filled up as much as possible.

### Part 1)

How many levels does the resulting tree have?