

Assignment 4

CIS 621: Algorithms and Complexity

Problem 1 (10 points) For the graph $G = (\mathcal{U}, \mathcal{E})$, where \mathcal{U} is the set of vertices and \mathcal{E} is the set of edges, we define the following nonlinear integer program, where $w_{i,j} \geq 0, \forall (i,j) \in \mathcal{E}$ and k is a nonnegative integer:

$$\begin{aligned} \max \quad & \sum_{(i,j) \in \mathcal{E}} w_{i,j} (x_i + x_j - 2x_i x_j) \\ \text{s. t.} \quad & \sum_{i \in \mathcal{U}} x_i = k, \\ & x_i \in \{0, 1\}, \forall i \in \mathcal{U}. \end{aligned}$$

- Show that the following linear program is a relaxation of the above problem:

$$\begin{aligned} \max \quad & \sum_{(i,j) \in \mathcal{E}} w_{i,j} z_{i,j} \\ \text{s. t.} \quad & z_{i,j} \leq x_i + x_j, \forall (i,j) \in \mathcal{E}, \\ & z_{i,j} \leq 2 - x_i - x_j, \forall (i,j) \in \mathcal{E}, \\ & \sum_{i \in \mathcal{U}} x_i = k, \\ & 0 \leq x_i \leq 1, \forall i \in \mathcal{U}, \\ & 0 \leq z_{i,j} \leq 1, \forall (i,j) \in \mathcal{E}. \end{aligned}$$

- Let $F(x) = \sum_{(i,j) \in \mathcal{E}} w_{i,j} (x_i + x_j - 2x_i x_j)$ be the objective function of the nonlinear integer program. Show that for any (x, z) that is feasible to the linear program, $F(x) \geq \frac{1}{2} \sum_{(i,j) \in \mathcal{E}} w_{i,j} z_{i,j}$.

Problem 2 (10 points) For the directed graph $G = (\mathcal{U}, \mathcal{E})$, where \mathcal{U} is the set of vertices and \mathcal{E} is the set of directed edges, we want to partition \mathcal{U} into two sets \mathcal{V} and $\mathcal{W} = \mathcal{U} \setminus \mathcal{V}$ in order to maximize the total weight of the edges going from \mathcal{V} to \mathcal{W} (i.e., the edges (i,j) with $i \in \mathcal{V}$ and $j \in \mathcal{W}$).

- Give a randomized $\frac{1}{4}$ -approximation algorithm for this problem.
- Show that the following linear program is a relaxation of this problem:

$$\begin{aligned} \max \quad & \sum_{(i,j) \in \mathcal{E}} w_{i,j} z_{i,j} \\ \text{s. t.} \quad & z_{i,j} \leq x_i, \forall (i,j) \in \mathcal{E}, \\ & z_{i,j} \leq 1 - x_j, \forall (i,j) \in \mathcal{E}, \\ & 0 \leq x_i \leq 1, \forall i \in \mathcal{U}, \\ & 0 \leq z_{i,j} \leq 1, \forall (i,j) \in \mathcal{E}. \end{aligned}$$

- For the above linear program, give a randomized $\frac{1}{2}$ -approximation algorithm based on rounding $x_i, \forall i \in \mathcal{U}$ to 1 with the probability of $\frac{1}{2}x_i + \frac{1}{4}$.