

# Wireless Sensor Network Denial of Sleep Attack

Presentation by Michael I Brownfield

## Paper Authors

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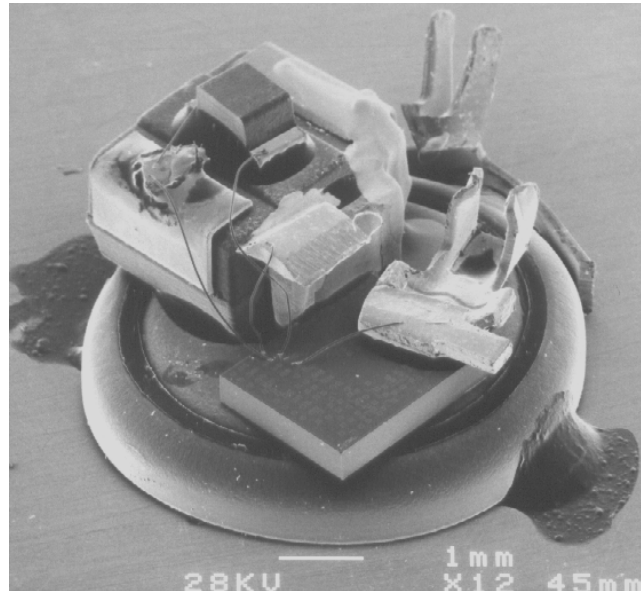
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Computer Engineering, Air Force Institute of Technology)

# Motivation

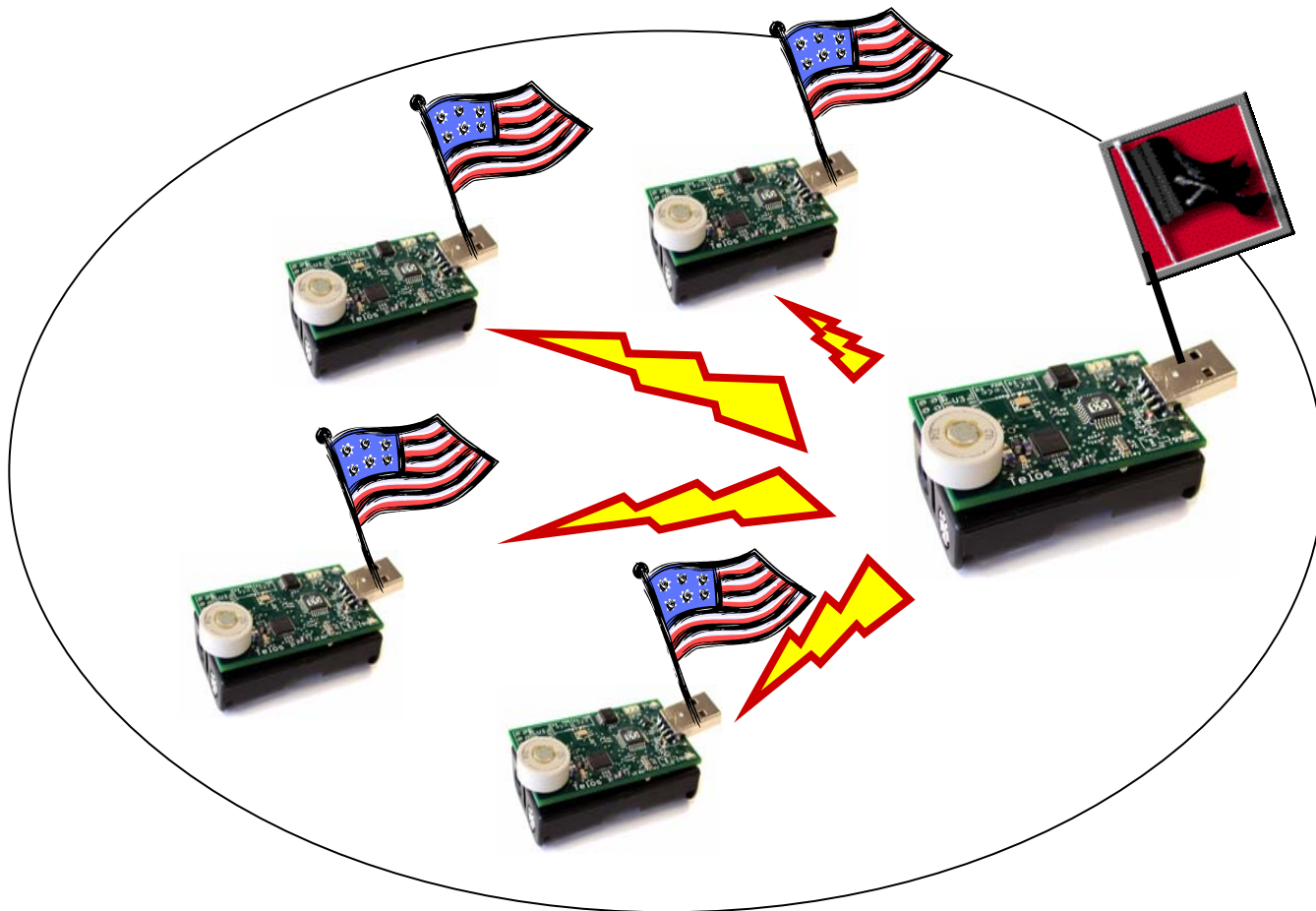


## FORTUNE

1. **SMART DUST KICKS UP A STORM**  
Tiny wireless sensors start monitoring the nation's food, workplaces, and welfare (02/2004).

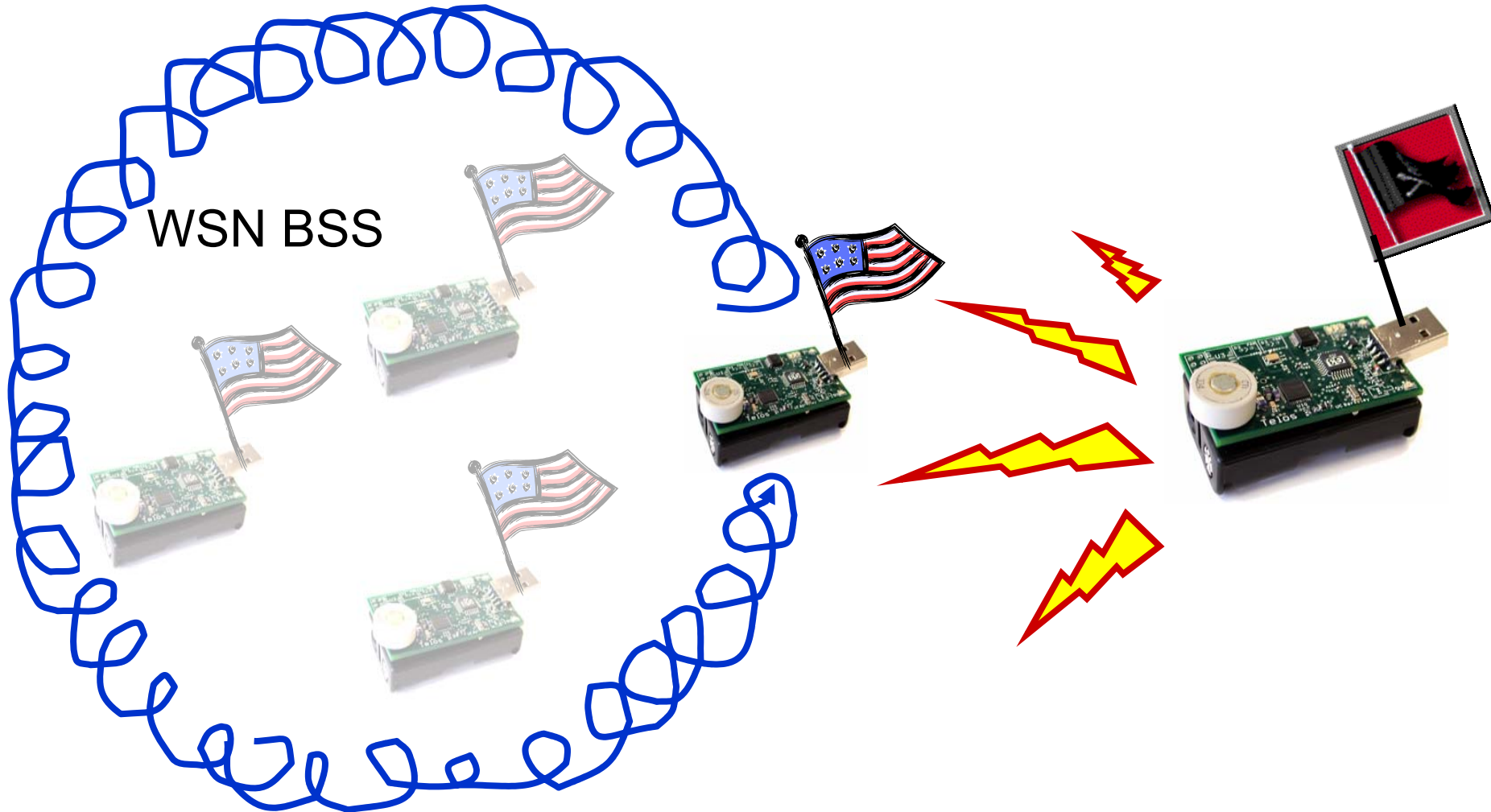


# Even More Motivation



## Wireless Sensor BSS

# Yes, Even More Motivation



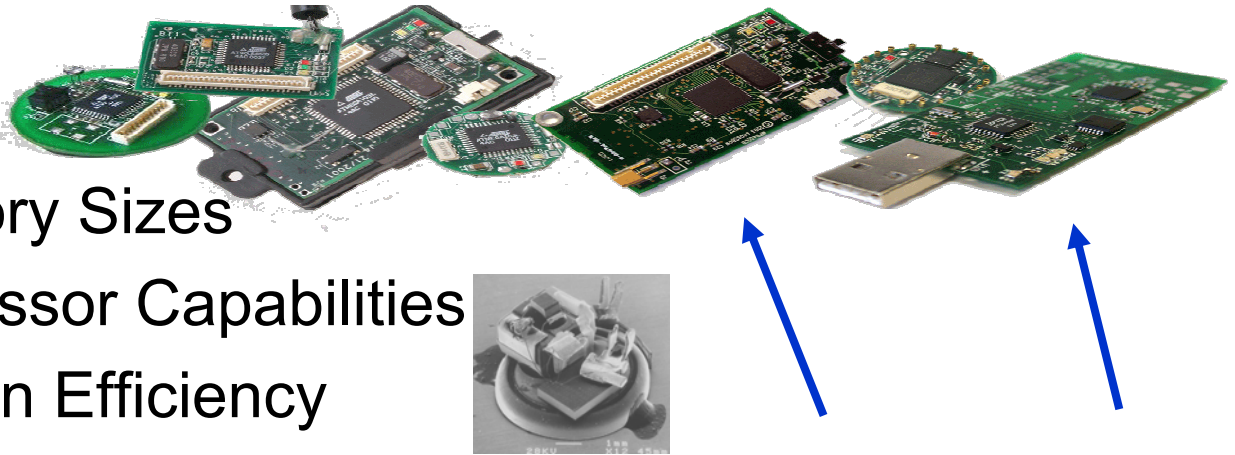


# Agenda

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- Wireless Network Challenges
- Related Security Work
  - Software Solutions
  - Hardware Solutions
- MAC Layer Sources of Energy Loss
- Comparison Wireless Sensor Network Protocols
- Gateway MAC
- Protocol Analysis
- Conclusion

# Sensor MAC Protocol Challenges

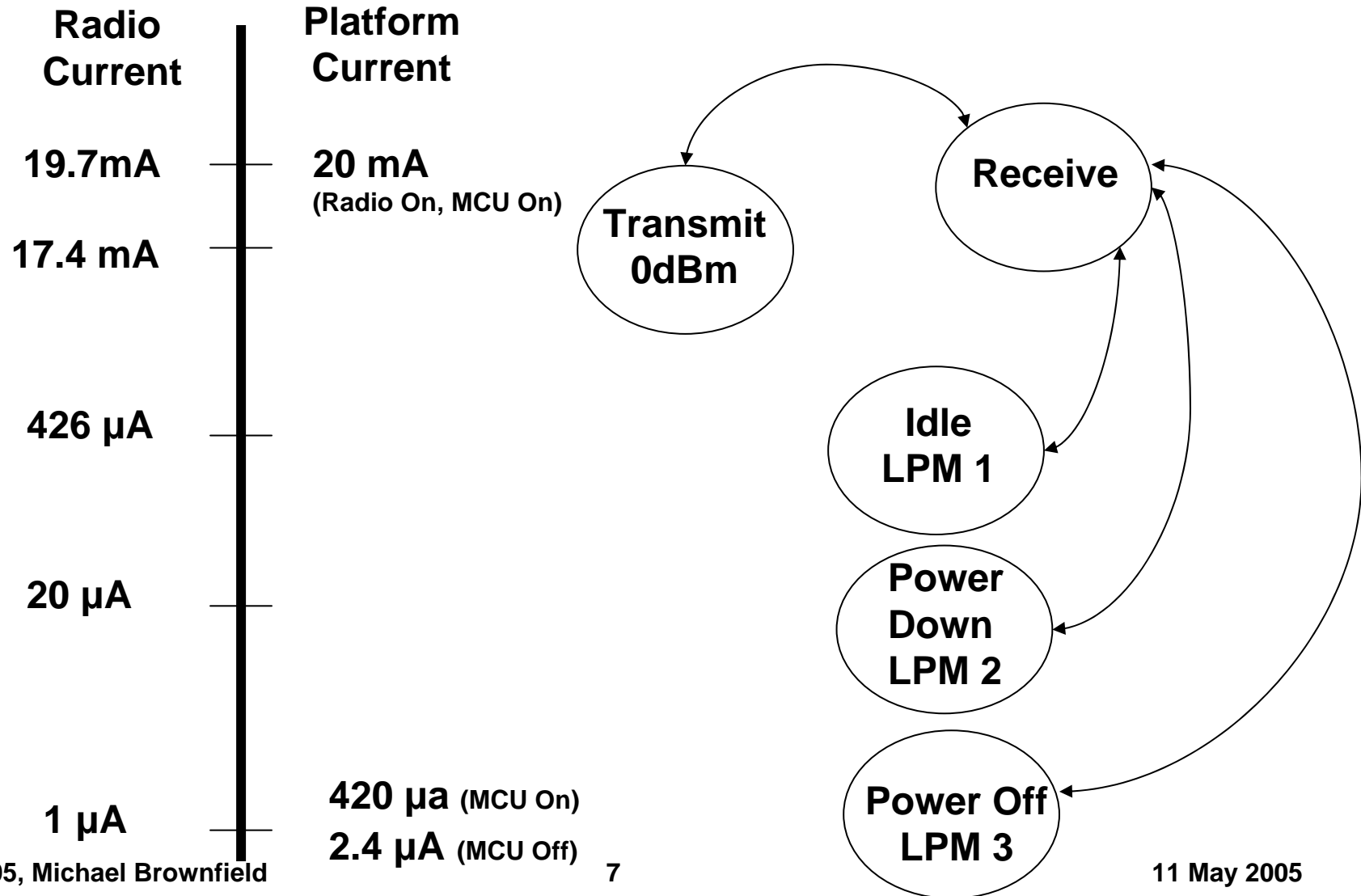


- Energy Efficiency
- RAM/ROM Memory Sizes
- Embedded Processor Capabilities
- Channel Utilization Efficiency
  - Throughput
  - Delay
  - Fairness
- Self-Configuration
- Network Scalability

Platform	Smartdust	Mica2	TelosA
$\mu$ -controller	8-bit	16-bit	16-bit
MCU RAM	512 B	4 kB	2 kB
EEPROM	512 B	128 kB	60 kB
Radio	916 MHz	868 MHz	2.4 GHz
Data Rate	10 kbps	76.8 kbps	250 kbps



# Sensor Motes Believe it Or Not!



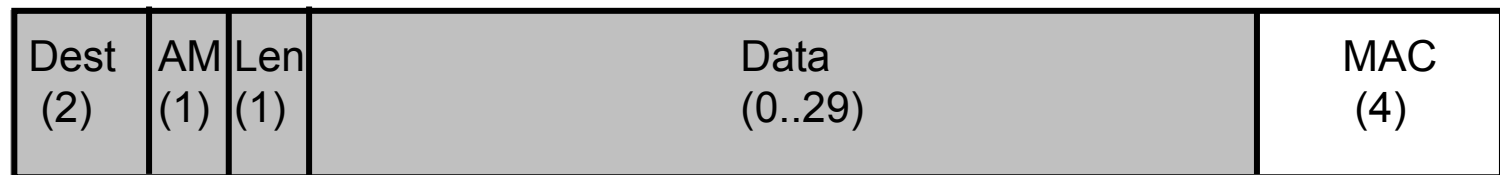
## ● Tiny Sec

- Fully Implemented with only 10% energy, latency, and BW overhead
- Proves hardware assistance is unnecessary
- Provides message authentication, integrity, and confidentiality w/out hardware acceleration



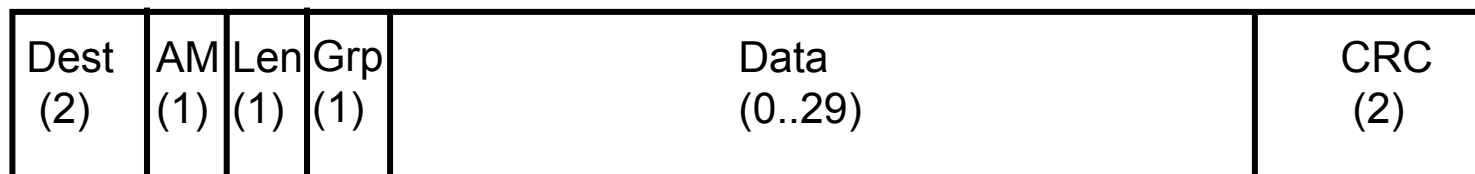
+5

TinySec-AE packet format



+1

TinySec-Authentication packet format



Standard TinyOS packet format





# CC2420 WSN Mote Transceiver Hardware Security Solutions

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## AES-128 Hardware Encryption with 4 Security Modes

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1. Disabled
2. Cipher block chaining – Message Authentication Code (CBC-MAC) authentication
3. Counter (CTR) encryption / decryption
4. Counter with CBC-MAC (CCM) authentication and encryption / decryption

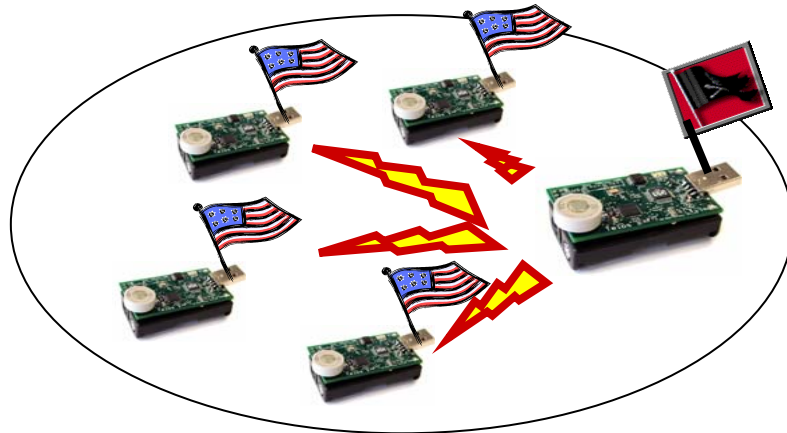
# Link Layer Denial of Service ATTACKS!

- Link Layer Collision Attack

- Attacker sends a signal at same time and frequency as legitimate traffic

- Link Layer Exhaustion Attack

- Manipulate the MAC protocol to force nodes to stay awake (repeated RTS messages)





# MAC Layer Sources of Energy Loss

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- Idle Listening
- Frame Collisions
  - Propagation Delay
  - Hidden Terminal
  - Capture Effect
- Overhearing
- Protocol / Control Packet Overhead
  - Control Packets (RTS, CTS, ACK)
  - Management Packets (Association requests, Beacon)
  - MAC Header transmission , Interframe Spacing
  - Contention Backoff



# Current WSN MAC Approaches

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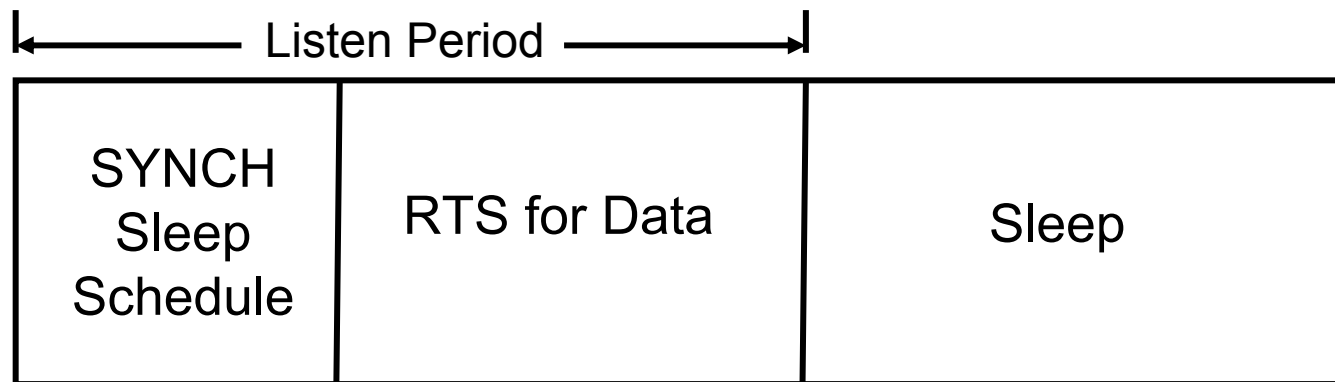
- Slotted CSMA:

- Sensor MAC (S-MAC)      implemented
- Timeout MAC (T-MAC)      implemented
- Berkeley-MAC (B-MAC)      implemented

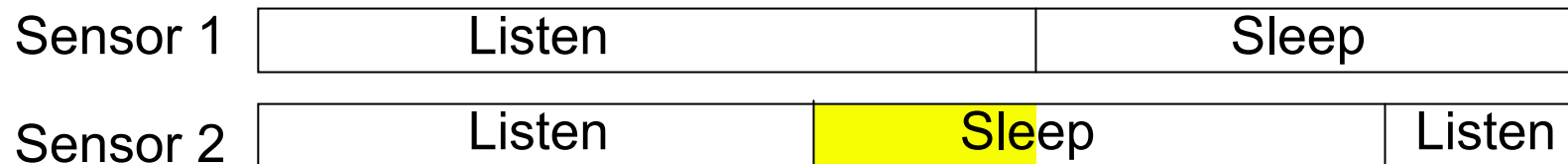
- TDMA

- Traffic-Adaptive MAC (TRAMA)
- Low-Energy Adaptive Clustering Heirarchy (LEACH)
- Power Aware Clustered TDMA (PACT)
- Bit-Map Assisted (BMA)
- Proposed Gateway MAC (G-MAC)

- Reduce energy through sleeping

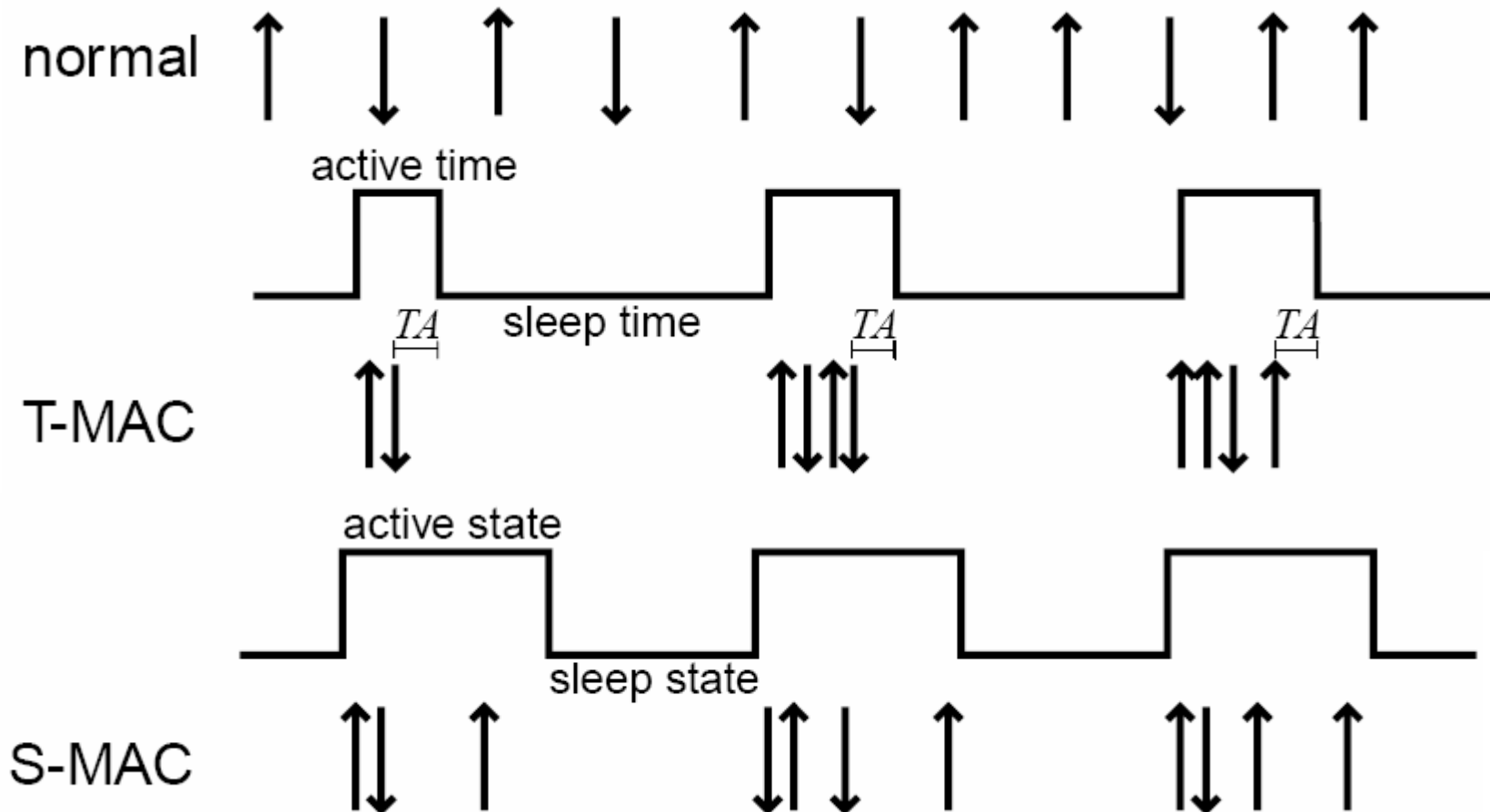


- Reduce energy by synchronizing sleep schedules in clusters

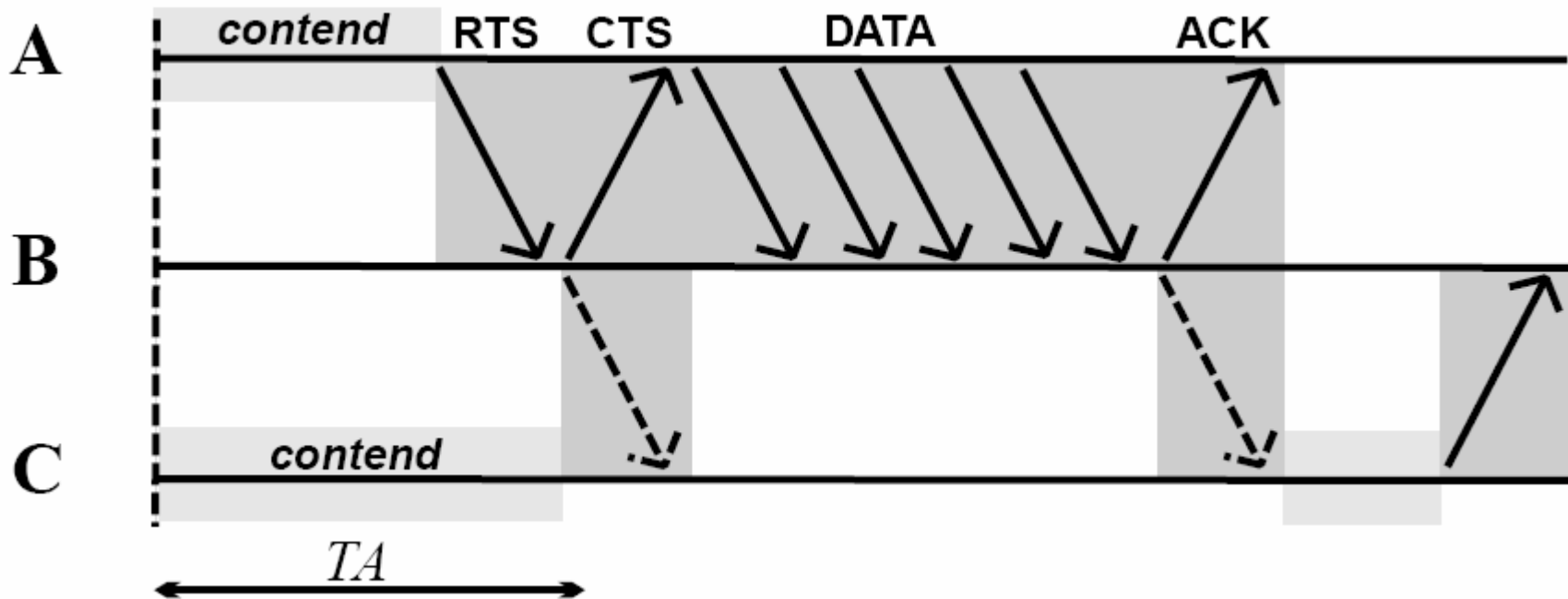


# Idle Listening Slotted Approaches

- Reduce energy through sleeping after timeout



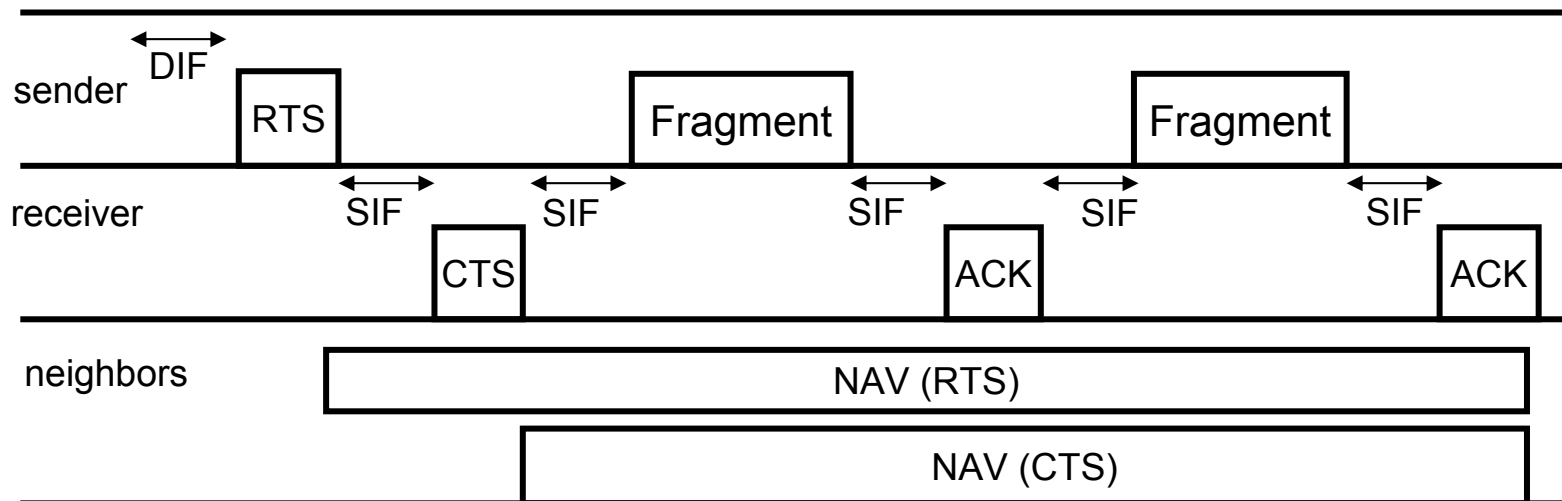
# T-MAC Adaptive Timeout (TA)



$$TA = 1.5 * (T_{\text{contention\_interval}} + T_{\text{RTS\_interval}} + T_{\text{SIFS}})$$

# Frame Collisions / Overhearing Avoidance

- Frame Collisions: RTS-CTS-Data-ACK
- Overhearing Avoidance: In-channel signaling to sleep during other transmissions





# Berkeley MAC (B-MAC)

## (Contention-based)



Versatile Low Power Media Access for Wireless Sensor Networks

Joseph Polastre, Jason Hill and David Culler

Sensys November 2004



# B-MAC GOALS



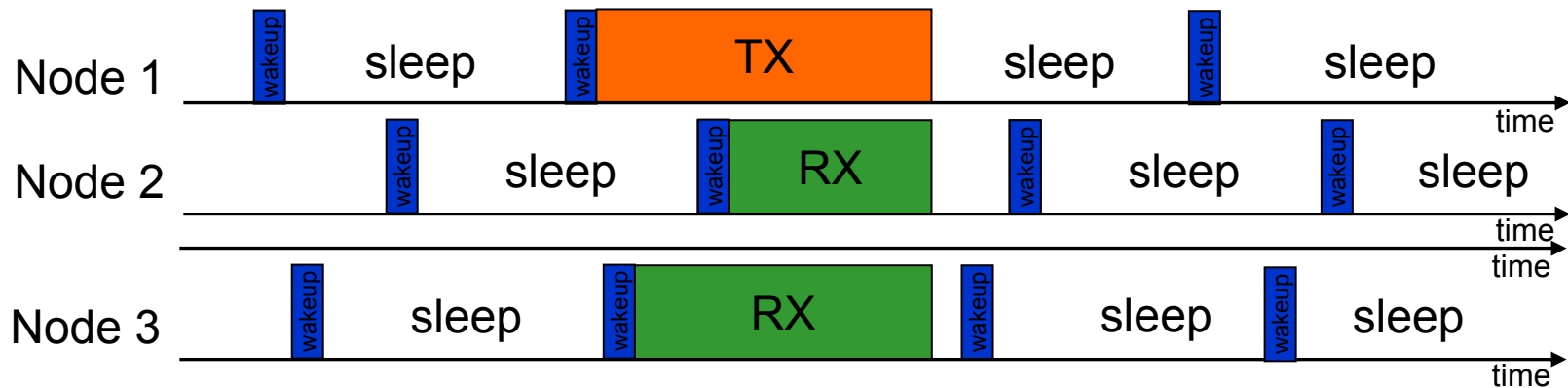
- Low Power operation
- Effective collision avoidance
- Simple and predictable
- Small code size and RAM usage
- Tolerable to changing RF/networking conditions
- Scalable to large numbers of nodes

# BMAC Design Approach



- Minimalistic (BMAC)

- Small core functionality: media access control
- RTS/CTS, ACKs, etc are considered higher layer functionality (services) that Applications can turn on and off





# Denial of Sleep Threat Analysis

- S-MAC

$$T_{networklifetime} = T_{sensorlifetime} = \frac{C_{battery(mAhr)}}{(D)(I_{active(mA)}) + (1-D)(I_{sleep(mA)})}$$

- T-MAC and B-MAC

$$T_{networklifetime} = T_{sensorlifetime} = \frac{C_{battery(mAhr)}}{I_{active(mA)}}$$

- Proposed Gateway MAC

$$T_{networklifetime} = T_{sensorlifetime} = \frac{n_{nodes} * C_{battery(mAhr)}}{(I_{active(mA)})}$$

# Gateway MAC (G-MAC)

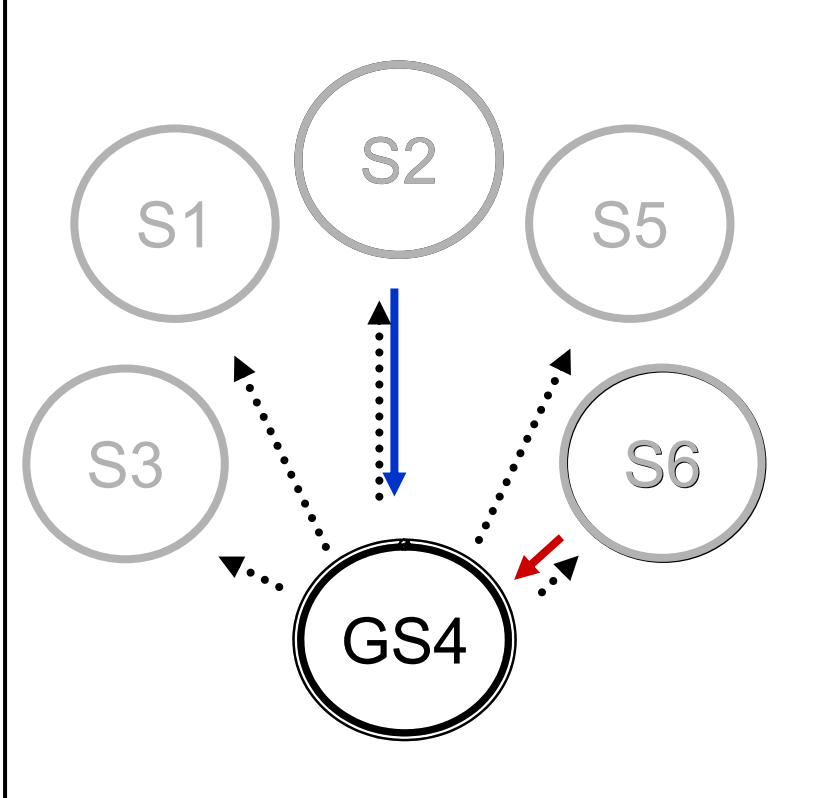
## Design Goals

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- Flexible and Scalable
  - Allocate Transmission Slots for Active Communication
  - *Distribute Clusterhead/Gateway Duties*
- Energy-efficient
  - *Eliminate Network-wide Idle Listening*
  - Create Traffic Rhythm for Increased Sleep Duration
- Secure
  - *Gateway Sentry*
  - *Embedded AES-128*
- Dynamic
  - *Self-configure, Self-recover*

# GMAC Design: Phase 1 Data Collection

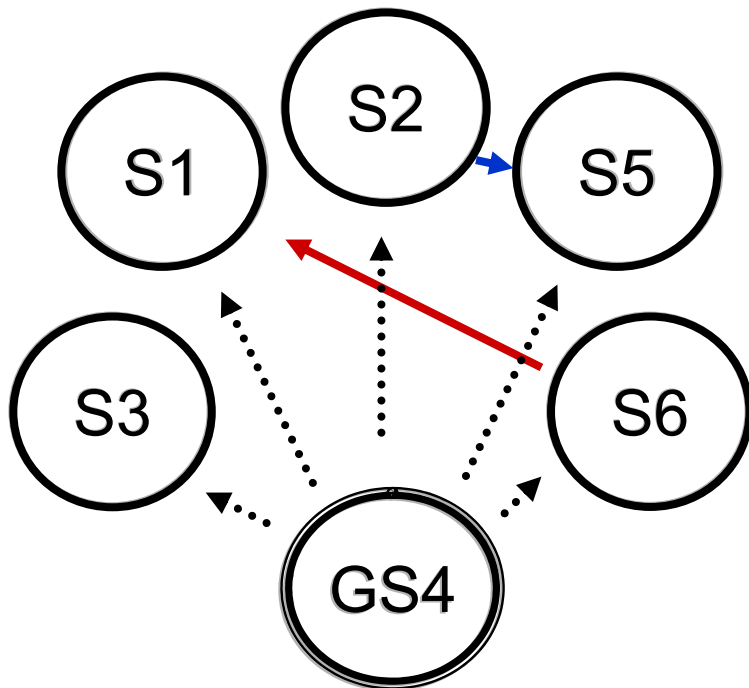
## WSN: Wireless Sensor Network



- Collects Intra-Network Traffic:
  - Local Gossip reservations
  - Cluster Association Requests
- Collects Inter-Network Traffic:
  - Converge-cast (out to WDS)
  - Network Routing
    - ◆ Tandem routing
    - ◆ Terminal

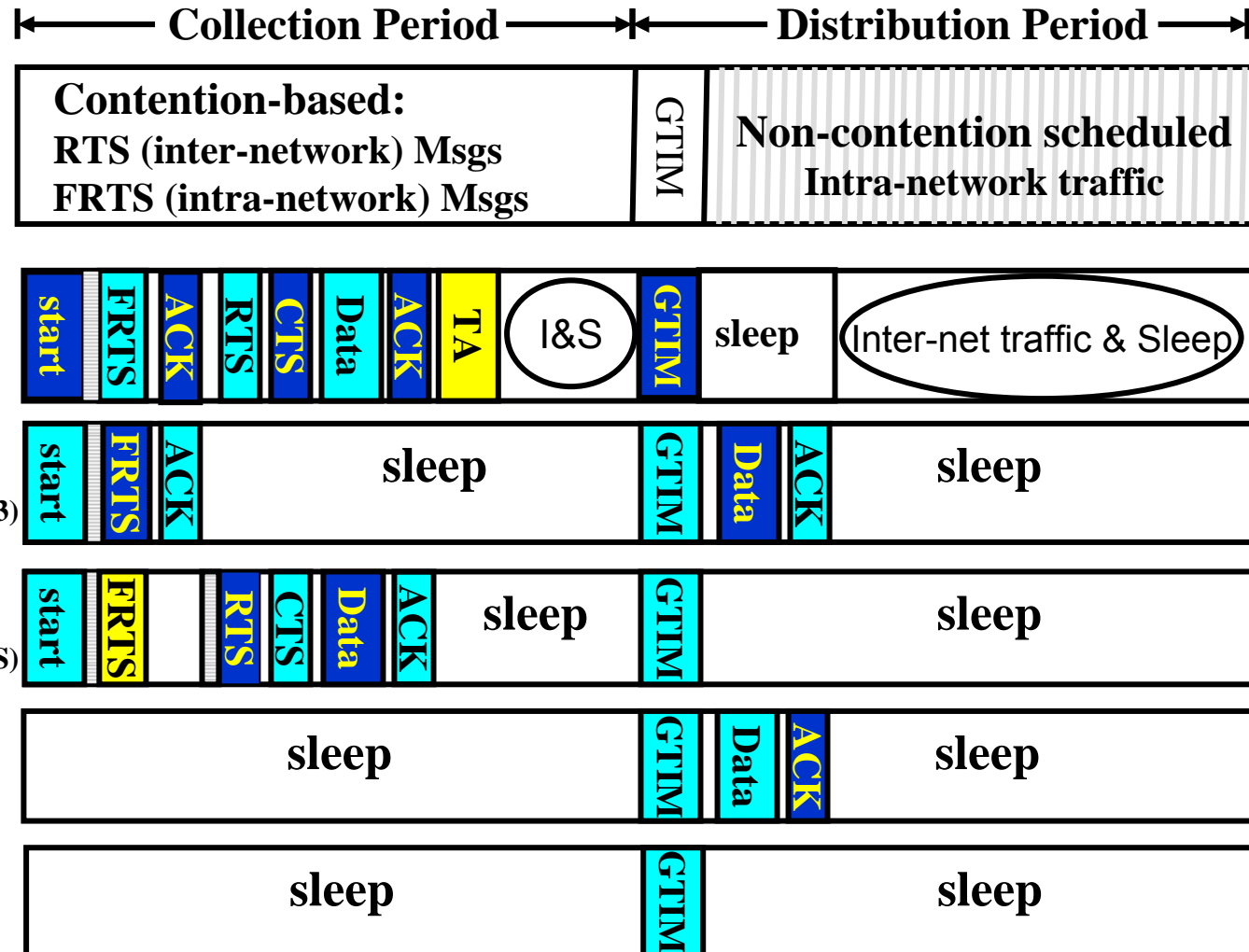
# GMAC Design: Phase 1 Data Distribution

WSN: Wireless Sensor Network



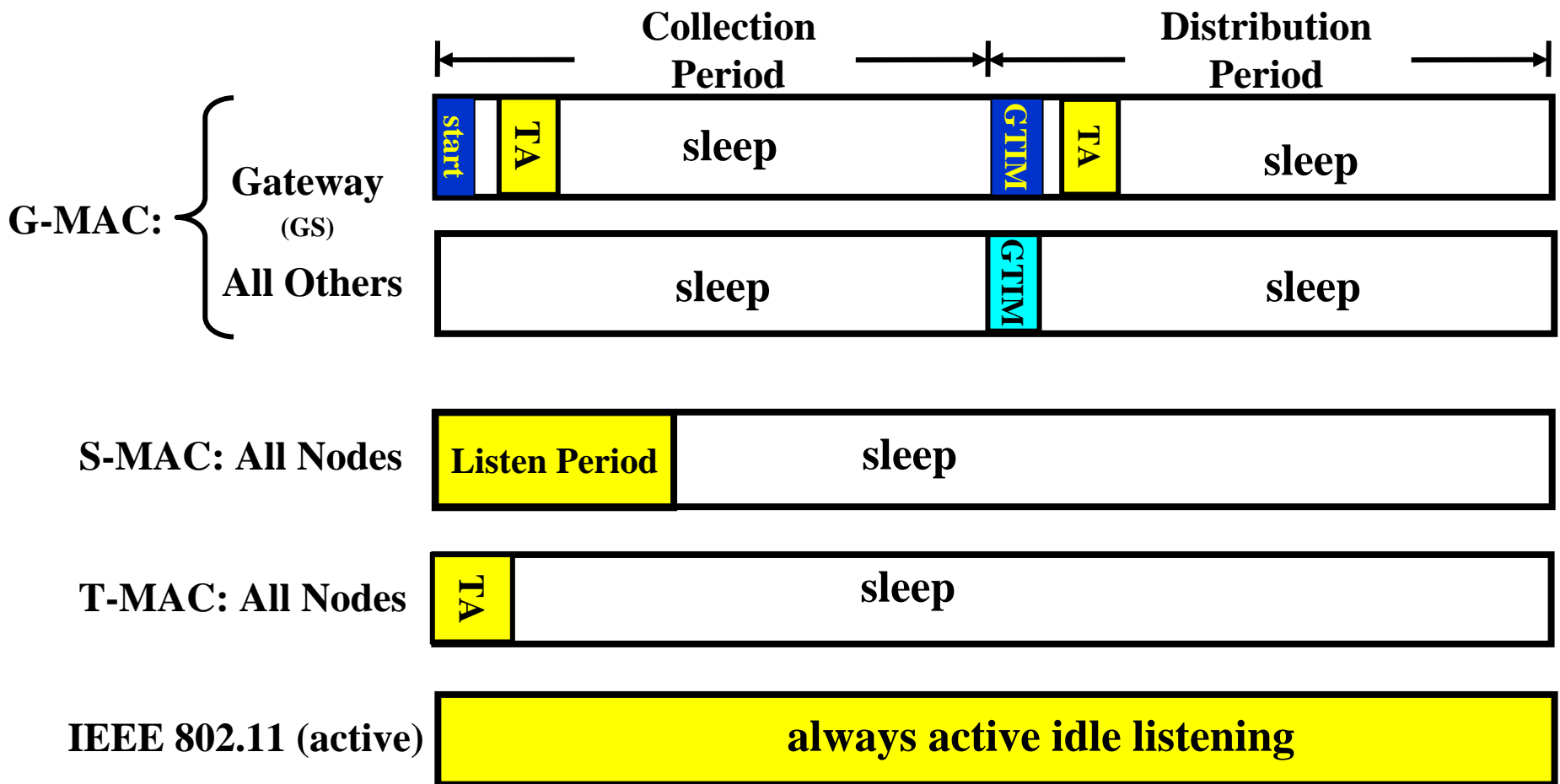
- Gateway Node broadcasts Traffic Indication Map (GTIM)
  - Local Broadcast (dotted lines)
  - Gossip exchanges
  - Cluster Synchronization Beacon

# GMAC Design





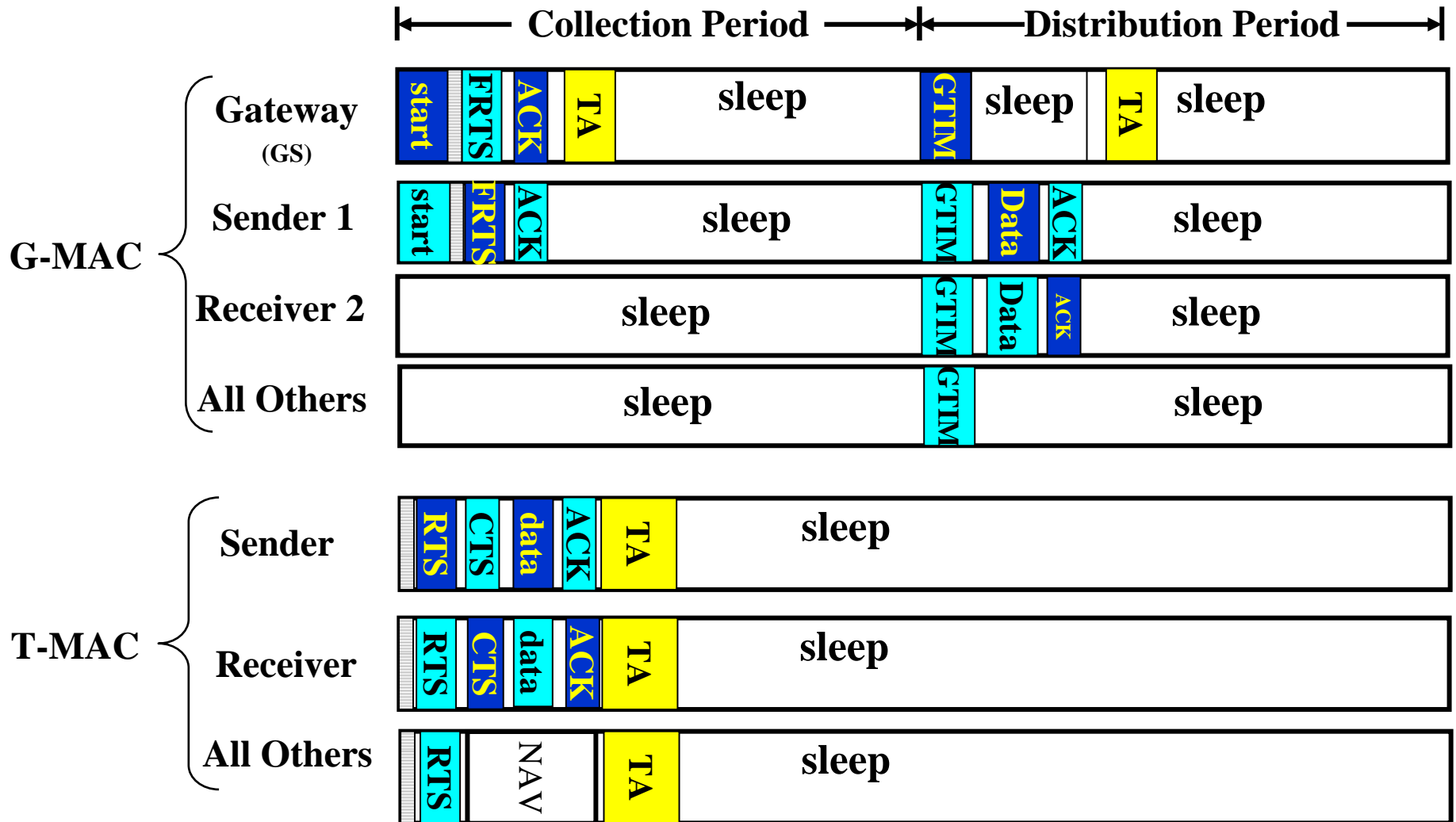
# Protocols in Empty Network





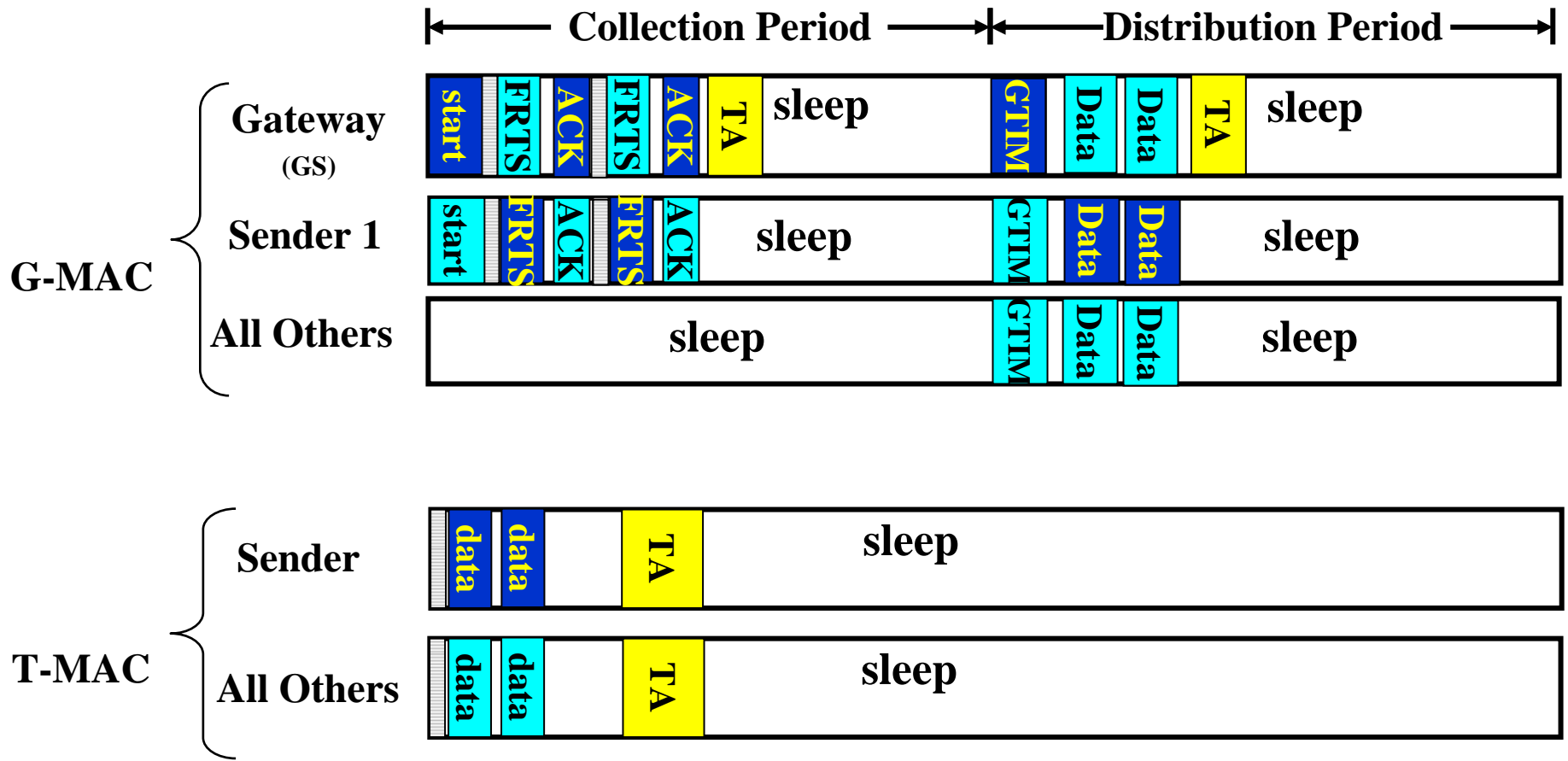
# G-MAC vs T-MAC

## Unicast Message



# G-MAC vs T-MAC

## Broadcast Message





# GMAC Point Coordinator/ Gateway Election

Battery Pwr Level		Voltage Range
11	High	$2.6 < \text{Pwr} \quad (3.0-3.6)$
10	Med	$2.4 < \text{Pwr} \leq 2.6$
01	Low	$2.1 < \text{Pwr} \leq 2.4$
00	Min	$\text{Pwr} \leq 2.1 \text{ volt}$

Memory Capacity Level		Percentage Avail. Capacity
11	High	$30\% < \text{Mem}$
10	Med	$20\% < \text{Mem} \leq 30\%$
01	Low	$10\% < \text{Mem} \leq 20\%$
00	Min	$\text{Mem} \leq 10\%$

## Critical Resource Level Algorithm

if Pwr = **Min** or Mem = **Min**

then Resource Level = 3

elseif Pwr = **Low** or Mem = **Low**

then Resource Level = 2

elseif Pwr = **Med** or Mem = **Med**

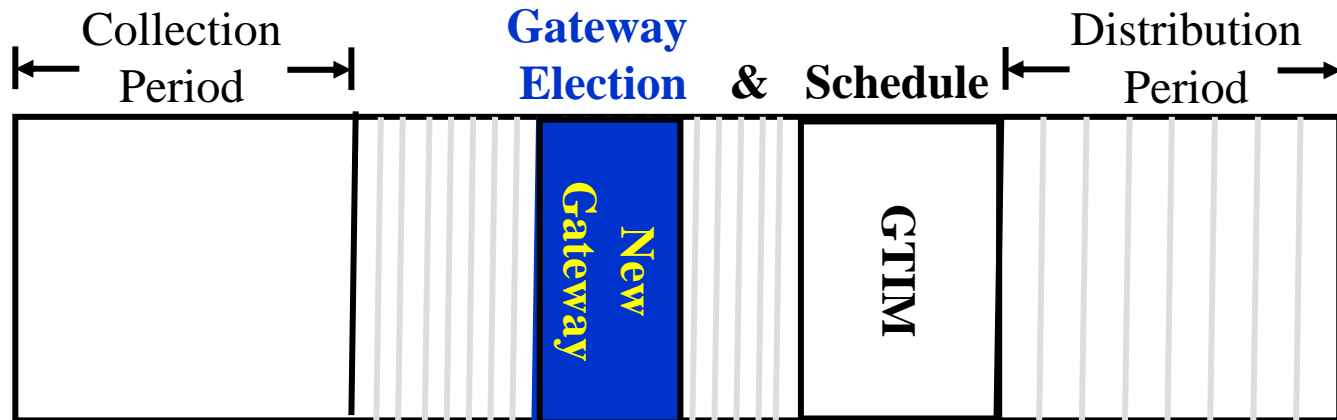
then Resource Level = 1

Elseif Pwr = **High** or Mem = **High**

then Resource Level = 0



# GMAC Point Coordinator/ Gateway Election



Election Contention Backoff =  $(16\text{-bit Random Number mod } 2^7) + (RL * 128)$

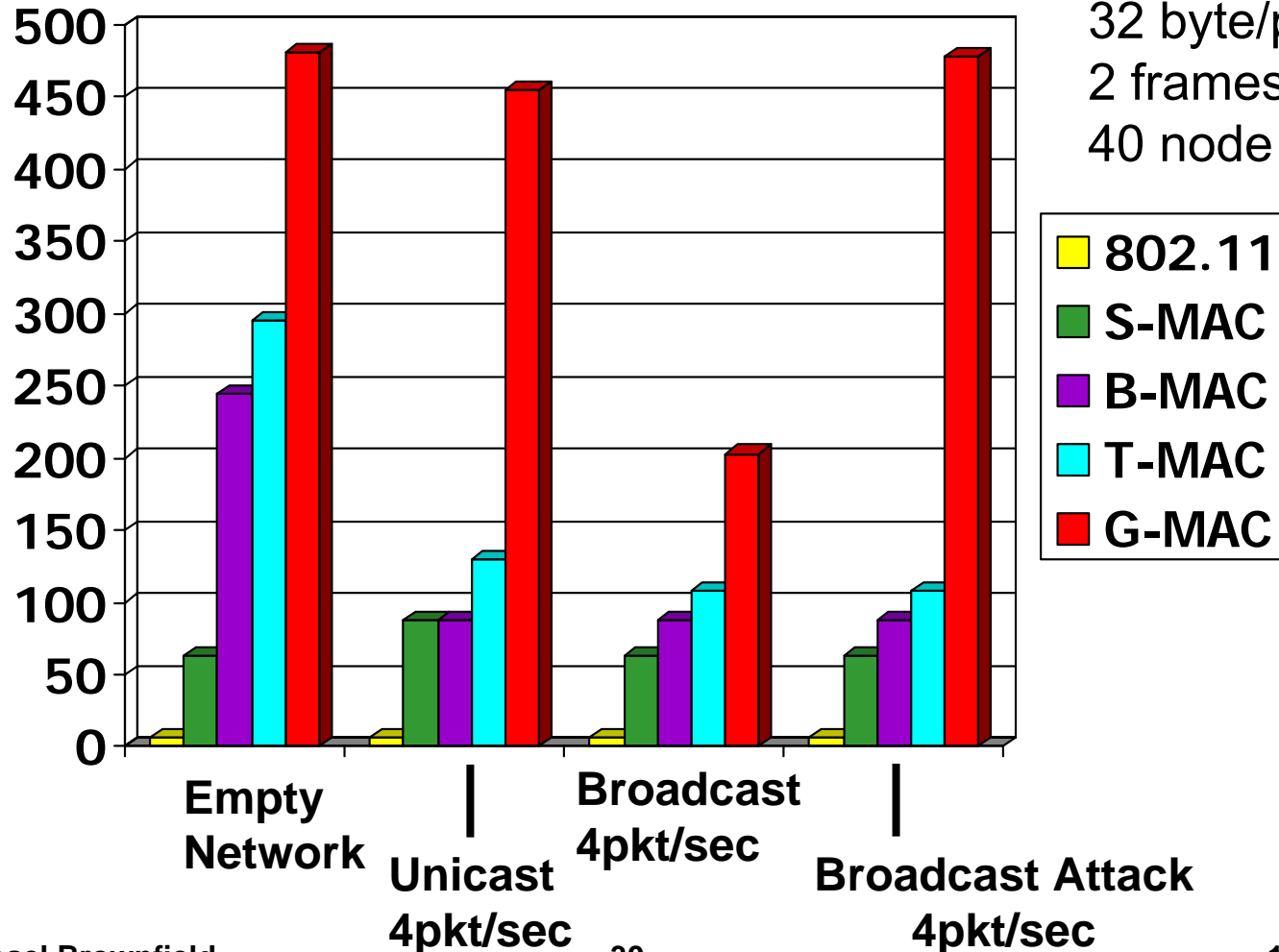
Resource Level (RL)	Contention Window
	$\text{Random}(2^7) + RL * 128$
0 (Max)	0 to 127 Slots (16 $\mu$ s to 2ms)
1	128 to 255 Slots (2ms to 4ms)
2	256 to 511 Slots (4ms to 8ms)
3 (Low)	512 to 1023 Slots (8ms to 16ms)

# Preliminary Results: G-MAC

VIRGINIA POLYTECHNIC INSTITUTE  
AND STATE UNIVERSITY

Data Rate 62.6kbps  
32 byte/pkt  
2 frames/second  
40 node network

Network  
Lifetime  
(days)





# Denial of Sleep Summary

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- WSN networks are vulnerable to Denial of Sleep Attacks
- Design Security Measures at the System Architectural Level (when possible)
- Centralized network coordination functions (Gateway) can reduce the effects of Denial of Sleep Attacks

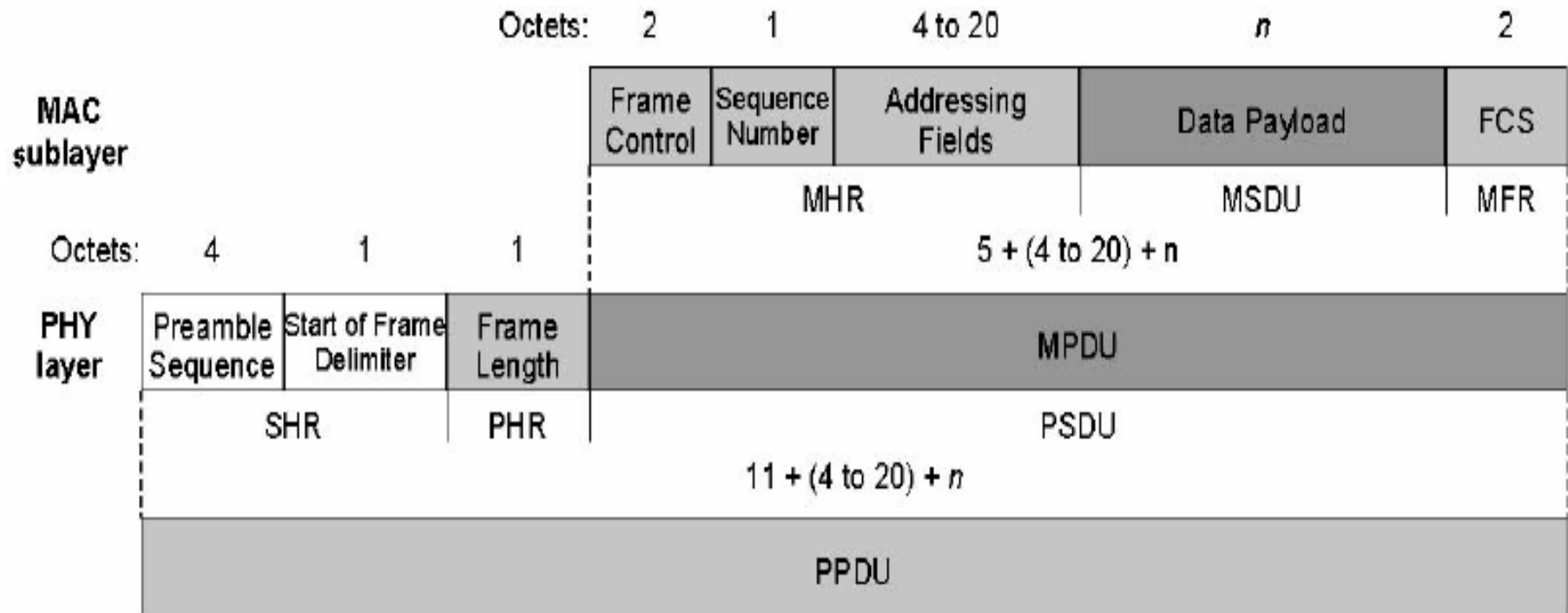
# Questions

**If we knew what it was we were doing,  
it would not be called research, would it ?  
--Albert Einstein**





# 802.15.4 Frame Format



# Symbiotic Network

