

# COMP30026 Models of Computation

## Lecture 2: Propositional Logic

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# Our Goals for the Next Few Lectures

Introduce formal propositional logic.

Use as vehicle for more general logic concepts.

Use for simple mechanised proof.

**Pay attention!**

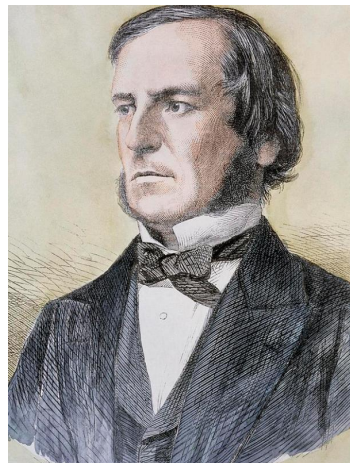
Even if you have seen this before, we need the concepts for later.

# Propositional = Boolean Logic

Until the mid-19th century, “logic” meant Aristotelian logic.

George Boole took an **algebraic** view of logic.

Deep connection between logic and arithmetic.



**Figure:** George Boole, circa 1864

# Intro Puzzle

Heidi, Dina and Louise are being questioned by their aunt.

Here is what they say:

*Heidi: "Dina and Louise had equal share in it; if one is guilty, so is the other."*

*Dina: "If Heidi is guilty, then so am I."*

*Louise: "Dina and I are not both guilty."*

Their aunt, knowing that they are honest kids, realises that they cannot tell a lie.

Has she got sufficient information to decide who (if any) are guilty?

# Syntax

We shall build propositional formulas from this set of symbols:

$$\underbrace{A, B, C, \dots, Z}_{\text{prop. letters}}, \underbrace{\neg, \wedge, \vee, \rightarrow, \leftrightarrow}_{\text{connectives}}, (, ).$$

Well-formed formulas (wffs) are generated by this grammar:

$$\begin{aligned} wff &:= A \mid B \mid C \mid \dots \mid Z \\ &\mid \neg wff \\ &\mid (wff \wedge wff) \\ &\mid (wff \vee wff) \\ &\mid (wff \rightarrow wff) \\ &\mid (wff \leftrightarrow wff) \end{aligned}$$

# Some Well-Formed Formulas

$$P \quad (1)$$

$$(P \rightarrow Q) \quad (2)$$

$$(P \vee \neg P) \quad (3)$$

$$\neg(P \wedge \neg P) \quad (4)$$

$$(P \leftrightarrow \neg P) \quad (5)$$

$$(((P \rightarrow Q) \rightarrow P) \rightarrow P) \quad (6)$$

# Pronunciation Guide

Symbol	Pronunciation
$\neg$	not
$\vee$	or; vee
$\wedge$	and; wedge
$\rightarrow$	if ... then; implies; only if; arrow
$\leftrightarrow$	if and only if; biimplies; double arrow

**Warning:** This is **only** pronunciation!

These symbols are **not** shorthands for English words!

# Notational Conveniences

For sake of readability, we will follow these **informal** rules:

- Drop outermost parentheses.
- Drop inner parentheses in nested uses of  $\wedge$  and  $\vee$ .
  - $P \wedge Q \wedge R$  is short for either of:
    - $((P \wedge Q) \wedge R)$
    - $(P \wedge (Q \wedge R))$
  - Same is true if you replace every “ $\wedge$ ” with “ $\vee$ ”.
  - **Warning:**  $P \wedge Q \vee R$  is nonsense.



# What is Truth?

What does it mean for  $P \wedge Q$  to be true?

What about just  $P$ ?

Is  $P \vee \neg P$  true?

# Boolean Semantics: Connectives

## Definition (Truth function)

A function from truth values to truth values.

Boolean truth values: **t** and **f** (also written **1** and **0**).

Connectives are **truth-functional**.

Usually presented as a **truth table**:

$A$	$B$	$\neg A$	$A \wedge B$	$A \vee B$	$A \rightarrow B$	$A \leftrightarrow B$
<b>0</b>	<b>0</b>	<b>1</b>	<b>0</b>	<b>0</b>	<b>1</b>	<b>1</b>
<b>0</b>	<b>1</b>	<b>1</b>	<b>0</b>	<b>1</b>	<b>1</b>	<b>0</b>
<b>1</b>	<b>0</b>	<b>0</b>	<b>0</b>	<b>1</b>	<b>0</b>	<b>0</b>
<b>1</b>	<b>1</b>	<b>0</b>	<b>1</b>	<b>1</b>	<b>1</b>	<b>1</b>

# Boolean Semantics: Letters

Propositional letters are **Boolean variables**.

## Definition (Truth assignment)

A function from propositional letters to truth values.

Usual notation:

$$v = \{P \mapsto \mathbf{1}, Q \mapsto \mathbf{0}\}.$$

We then have:

$$v(P) = \mathbf{1}$$

$$v(Q) = v(R) = \dots = v(Z) = \mathbf{0}.$$

# Truth of a Formula

Let  $v = \{P \mapsto \mathbf{1}, Q \mapsto \mathbf{0}\}$ .

**Poll:** Which of these formulas are true under  $v$ ?

1.  $P \wedge Q$
2.  $(P \vee Q) \wedge (P \vee R)$
3.  $P \rightarrow Q$
4.  $\neg P \rightarrow \neg Q$

**Shorthand:** “ $v \models \phi$ ” means “ $\phi$  is true under  $v$ ”.

# Truth Tables for Formulas

$P$	$Q$	$R$	$((P \wedge Q) \vee R)$				
0	0	0	0	0	0	<b>0</b>	0
0	0	1	0	0	0	<b>1</b>	1
0	1	0	0	0	1	<b>0</b>	0
0	1	1	0	0	1	<b>1</b>	1
1	0	0	1	0	0	<b>0</b>	0
1	0	1	1	0	0	<b>1</b>	1
1	1	0	1	1	1	<b>1</b>	0
1	1	1	1	1	1	<b>1</b>	1

Which of these have the same truth tables?

1.  $P \rightarrow Q$  and  $\neg Q \rightarrow \neg P$
2.  $(P \rightarrow Q) \wedge (P \rightarrow R)$  and  $P \rightarrow (Q \wedge R)$
3.  $(P \rightarrow R) \wedge (Q \rightarrow R)$  and  $(P \wedge Q) \rightarrow R$

**Hint:**  $P \rightarrow Q$  has the same truth table as  $\neg P \vee Q$ .

# Logical Equivalence

## Definition

Formulas are *logically equivalent* iff they have equal truth values under **every** truth assignment.

**Shorthand:** " $F \equiv G$ " means " $F$  is logically equivalent to  $G$ ".

# Material Conditional

**Warning:** “ $\rightarrow$ ” is weird!

Often read as “implies”, but causality is not required!

$A$	$B$	$A \rightarrow B$
0	0	1
0	1	1
1	0	0
1	1	1

1. If volume increases, then pressure falls.
2. If Melbourne is in Queensland, then Brisbane is in Victoria.
3. Melbourne and Brisbane are in different states **and** if Melbourne is in Queensland then so is Brisbane.



# Modus Ponens

$$\frac{P \rightarrow Q \quad P}{Q}$$

A rule is **sound** if every model of the premises is a model of the conclusion.

**Challenge:** prove that modus ponens is sound.

# Exit Puzzle

On the island of Knights and Knaves, everyone is a knight or knave. Knights always tell the truth. Knaves always lie.

Today there is a census on the island!

You are a census taker, going from house to house. Fill in what you know about each of these three houses.

- **In house 1:** We are both knaves.
- **In house 2:** At least one of us is a knave.
- **In house 3:** If I am a knight then so is my wife.

If you like these puzzles, Raymond Smullyan has written lots of books that you will like.