

COMP30026 Models of Computation

Lecture 5: Predicate Logic

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Why Predicate Logic?

Express statements about objects more faithfully.

In particular, *finitely* express:

- statements about **infinite** collections (e.g. the integers);
- transitive verbs (e.g. “loves”) and relative pronouns (e.g. “whose”).

Some Translations

No emus fly: $\forall x(Emu(x) \rightarrow \neg Flies(x))$

There are black swans: $\exists x(Black(x) \wedge Swan(x))$

If all push the cart, the donkey will be happy: $(\forall x P(x, c)) \rightarrow H(d)$

If the cart is pushed, the donkey will be happy: $(\exists x P(x, c)) \rightarrow H(d)$

Building Blocks

Five new kinds of symbols:

- ① constants (e.g. c for “**the** cart”, d for “**the** donkey”)
- ② predicates (e.g. P for “__ is pushing __”, H for “__ is happy”)
 - Identity, $=$, is a special case.
- ③ quantifiers
 - \forall (pronounced “for all”)
 - \exists (pronounced “there exists”)
- ④ variables (e.g. x , y , z)
- ⑤ functions (e.g. $+$, \cdot , f)

More Translations

Tina found Rover and returned him to Anne:

$$Found(tina, rover) \wedge Gave(tina, rover, anne)$$

Tina found a dog and gave **it** to Anne:

$$\exists x(Dog(x) \wedge Found(tina, x) \wedge Gave(tina, x, anne))$$

Lea inhabits the house **that** Jackie built:

$$\exists x(House(x) \wedge Inhabits(lea, x) \wedge BuilderOf(jackie, x))$$

Mothers' mothers are grandmothers:

$$\forall x \forall y \forall z ((Mother(x, y) \wedge Mother(y, z)) \rightarrow Grandmother(x, z))$$

Existential Quantification

Existential quantification, \exists , is generalised \vee .

“Tina found some money and gave **it** to the Red Cross”:

$$\exists x (Money(x) \wedge Found(tina, x) \wedge Gave(tina, x, redcross))$$

means:

$$(Found(tina, \$1) \wedge Gave(tina, \$1, redcross)) \vee$$

$$(Found(tina, \$2) \wedge Gave(tina, \$2, redcross)) \vee$$

$$\vdots$$

Universal Quantification

Universal quantification, \forall , is generalised \wedge .

“The square of every integer is nonnegative”:

$$\forall x(x \in \mathbb{Z} \rightarrow (x \cdot x \geq 0))$$

means:

$$0 \times 0 \geq 0 \wedge$$

$$1 \times 1 \geq 0 \wedge$$

$$\vdots$$

Quiz: Translate This

Translate “Every Melburnian barracks for a frisbee team”.

Use these predicates:

Predicate	Interpretation
$M(x)$	x is a Melburnian
$T(x)$	x is a frisbee team
$B(x, y)$	x barracks for y

Terms

A **term** represents an **individual** object.

Terms are not formulas.

Examples:

- *redcross*
- 1
- x
- $f(x + y)$

Atomic Formulas

Atomic formulas represent statements about **individual** objects.

Examples:

- $Flies(x)$
- $P(x, c)$
- $Gave(tina, x, redcross)$
- $2 + 2 = 5$
- $x \in y$

A Notational Convention

A predicate starts with an upper case letter; nothing else does.

- “*parent(rhonda)*” is a **term**.
 - Likely refers to “the parent of Rhonda”.
- “*Parent(rhonda)*” is a **formula**.
 - Likely represents the proposition “Rhonda is a parent”.

Well-formed formulas (wffs) are generated by the grammar

$$\begin{aligned} wff &\rightarrow atom \\ &| \neg wff \\ &| (wff \wedge wff) \\ &| (wff \vee wff) \\ &| (wff \rightarrow wff) \\ &| (wff \leftrightarrow wff) \\ &| \forall var\ wff \\ &| \exists var\ wff \end{aligned}$$

Quantifier Scope

The subformula attached to a quantifier is its **scope**.

Example: $\forall x \underbrace{(P(x) \vee Q(c))}_{\text{scope of } \forall x}$

A quantifier with variable x **binds** x within its scope.

Example: x is bound in $\forall x P(x)$.

If a variable is not bound, it is **free**.

Example: x is free in $P(x)$.

Which variables are bound, and which are free?

- ① $\forall z(P(x, y, z) \wedge \forall y(P(f(x), z, y)))$
- ② $\forall x \exists y(x < y \wedge \exists z(y < z))$

Renaming Variables

Bound variables can be renamed unless there is a clash:

$\exists x \forall y (x < y)$ means the same as $\exists x \forall z (x < z)$.

But $\exists x \forall y (x \leq y)$ is very different to $\exists x \forall x (x \leq x)$.

Renaming **free** variables changes meaning:

$P(x)$ is different to $P(y)$.

From English to Predicate Logic

A rough guide:

- ① Identify nouns, verbs, pronouns, adjectives, relative clauses.
- ② Assign:
 - Constant symbols to singular objects (e.g. “Rhonda”).
 - Predicate symbols to verbs and adjectives (e.g. “loves”).
 - Function symbols to relative clauses (e.g. “the parent of”).
 - Variables to indefinite pronouns (e.g. “someone”).
- ③ Replicate logical structure of sentence.

Example Translations

Let $L(x, y)$ stand for “ x loves y ”.

$L(alex, eva)$

Alex loves Eva

$\forall x L(x, eva)$

Everyone loves Eva (incl. Eva)

$\forall x (\neg(x = eva) \rightarrow L(x, eva))$

Eva is loved by everyone else

$\exists x (\neg(x = alex) \wedge L(x, alex))$

Someone other than Alex loves Alex

$\forall x \exists y L(x, y)$

Everybody loves somebody

$\exists y \forall x L(x, y)$

Someone is loved by everybody

$\exists x \forall y L(x, y)$

Someone loves everybody

Word Order

Consider word order with care:

- “There is something which is not P ”: $\exists y \neg P(y)$
- “There is not something which is P ” (“nothing is P ”):
 $\neg \exists y P(y)$
- “All S are not P ” **vs** “not all S are P ”:
 $\forall x (S(x) \rightarrow \neg P(x))$ **vs** $\neg \forall x (S(x) \rightarrow P(x))$

Quantifier Order

Order of different quantifiers is important!!!

$\forall x \exists y L(x, y)$ says “everyone has someone they love”.

$\exists y \forall x L(x, y)$ says “there is someone who is loved by everyone”.

But $\forall x \forall y$ is the same as $\forall y \forall x$ and $\exists x \exists y$ is the same as $\exists y \exists x$.

Implicit Quantifiers

Often quantifiers are implicit in English.

Look for nouns (especially plural) without determiners.

“Humans are mortal” means “**all** humans are mortal”:

$$\forall x (Human(x) \rightarrow Mortal(x))$$

“A horse is stronger than a dog” would usually mean:

$$\forall x \forall y ((Horse(x) \wedge Dog(y)) \rightarrow Stronger(x, y))$$

“If a child owns a dog, the child spoils it”:

$$\forall x \forall y ((Child(x) \wedge Dog(y) \wedge Owns(x, y)) \rightarrow Spoils(x, y))$$