

Escape Analysis

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What is Escape Analysis?

- Escape analysis tries to **statically** reason about the **dynamic scope** of an **object**.

Traditional Use Cases

- Stack-allocations
- Scalar-replacement
- Unnecessary synchronization removal
- Identify pure methods (see purity analysis slides)
- Identify immutable data-structures
- Generate more precise call graphs

Stack Allocation

```
public void foo() {  
    Object o = new Object();  
    processObject(o);  
}
```

```
public void processObject(Object o) {  
    // ...  
}
```

- The object o can be allocated on the stack instead of the heap.
- No garbage collection is needed for o.

Scalar Replacement

```
class Circle {  
    double r;  
    public Circle(double r) { this.r = r; }  
    public double area() { return Math.PI * r * r; }  
}
```

```
public double foo() {  
    Circle c = new Circle(1);  
    return c.area();  
}
```

- No object needs to be created for c.
- The call to c.area() can be inlined.
- May also require inlining analysis.

Removing Unnecessary Synchronization

```
public class AppletViewer {  
    // ...  
    public URL getCodeBase() {  
        Object o = new Object();  
        synchronized (o) {  
            // ...  
        }  
        return this.baseURL;  
    }  
}
```

Synchronization is useless on o, as only this call of getCodeBase has access to it.

This code was found in JDK 8 update 151

What is Escape Analysis (cont.)?

- Escape analysis makes a statement about the object's **lifetime** and **accessibility**
- In contrast to lifetime analysis, escape analysis is more coarse grain and binds the object's lifetime/accessibility to other entities (in particular **methods** or **threads**)

Dimensions of Escaping: Lifetime

```
public void m() { // in thread t  
    Object o = new Object();  
}
```

1. $\text{Lifetime}(t) \geq \text{Lifetime}(o) \leq \text{Lifetime}(m)$

Stack Allocation: ✓ Scalar Replacement: ? Synchronization Removal: ?

2. $\text{Lifetime}(t) \geq \text{Lifetime}(o) > \text{Lifetime}(m)$

Stack Allocation: ✗ Scalar Replacement: ✗ Synchronization Removal: ?

3. $\text{Lifetime}(o) > \text{Lifetime}(t) \geq \text{Lifetime}(m)$

Stack Allocation: ✗ Scalar Replacement: ✗ Synchronization Removal: ✗

Dimensions of Escaping: Accessibility

```
public void m() { // in thread t  
    Object o = new Object();  
}
```

1. $\text{Access}(o) = \{m, t\}$

Stack Allocation: ✓ Scalar Replacement: ✓ Synchronization Removal: ✓

2. $(\exists m' \neq m . m' \in \text{Access}(o)) \wedge (\nexists t \neq t' . t' \in \text{Access}(o))$

Stack Allocation: ? Scalar Replacement: ✗/? Synchronization Removal: ✓

3. $\exists t' \neq t . t' \in \text{Access}(o)$

Stack Allocation: ✗ Scalar Replacement: ✗ Synchronization Removal: ✗

Method Calls

```
public void foo() {  
    Object o = new Object();  
    bar(o);  
}
```

```
public void bar(Object o) {  
    // ...  
}
```

On Method calls (**c**):

- Access(o) = {foo, bar, t}
- Lifetime(o) \leq Lifetime(foo)
- Stack Allocation: ✓
- Scalar Replacement: ✗
- Synchronization Removal: ✓

Method Returns

```
public Object foo() {  
    Object o = new Object();  
    return o;  
}
```

```
public void bar() {  
    foo();  
}
```

On returns (**r**):

- $\text{Access}(o) = \{\text{foo}, \text{bar}, t\}$
- $\text{Lifetime}(o) > \text{Lifetime}(\text{foo})$
- $\text{Lifetime}(o) < \text{Lifetime}(t)$
- Stack Allocation: **X**
- Scalar Replacement: **X**
- Synchronization Removal: **✓**

Unhandled Exceptions

```
public void foo() {  
    throw new Exception();  
}
```

```
public void bar() {  
    foo();  
}
```

On abnormal returns (**a**):

- $\text{Access}(o) = \{\text{foo}, \text{bar}, t\}$
- $\text{Lifetime}(o) > \text{Lifetime}(\text{foo})$
- $\text{Lifetime}(o) < \text{Lifetime}(t)$
- Stack Allocation: **X**
- Scalar Replacement: **X**
- Synchronization Removal: **✓**

Mutating the State of Given Parameters

```
public class MyClass {  
    Object f;
```

```
    public void foo(MyClass p) {  
        p.f = new Object();  
    }
```

```
    public void bar() {  
        foo(new MyClass);  
    }  
}
```

Parameter field writes(**p**):

- Access(o) = {foo, bar, t}
- Lifetime(o) > Lifetime(foo)
- Lifetime(o) < Lifetime(t)
- Stack Allocation: **X**
- Scalar Replacement: **X**
- Synchronization Removal: **?**

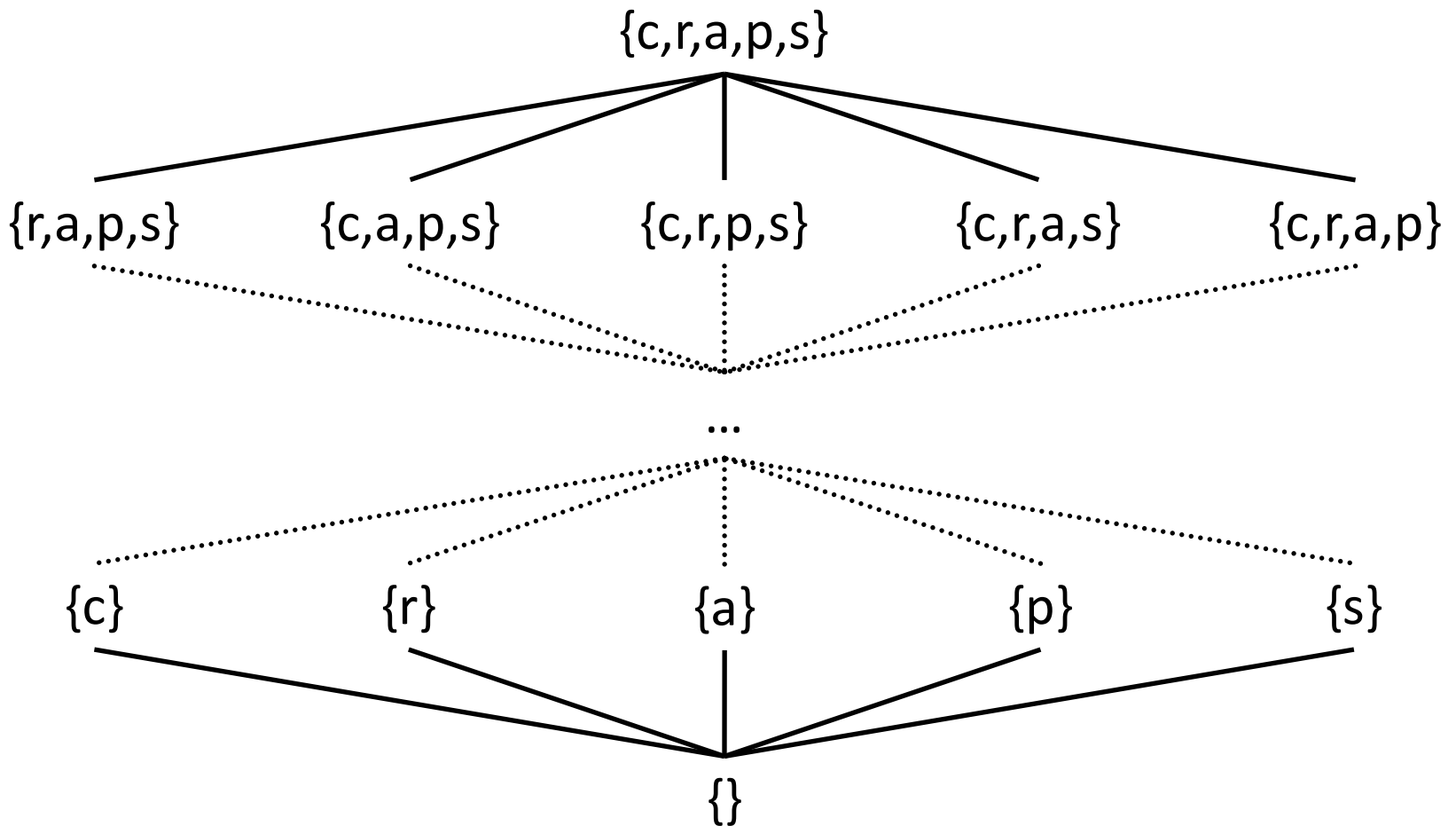
Writing Static Fields

```
public class MyClass {  
    public static Object f;  
  
    public void foo() {  
        Object o = new Object();  
        MyClass.f = o;  
    }  
  
    public void bar() {  
        foo();  
    }  
}
```

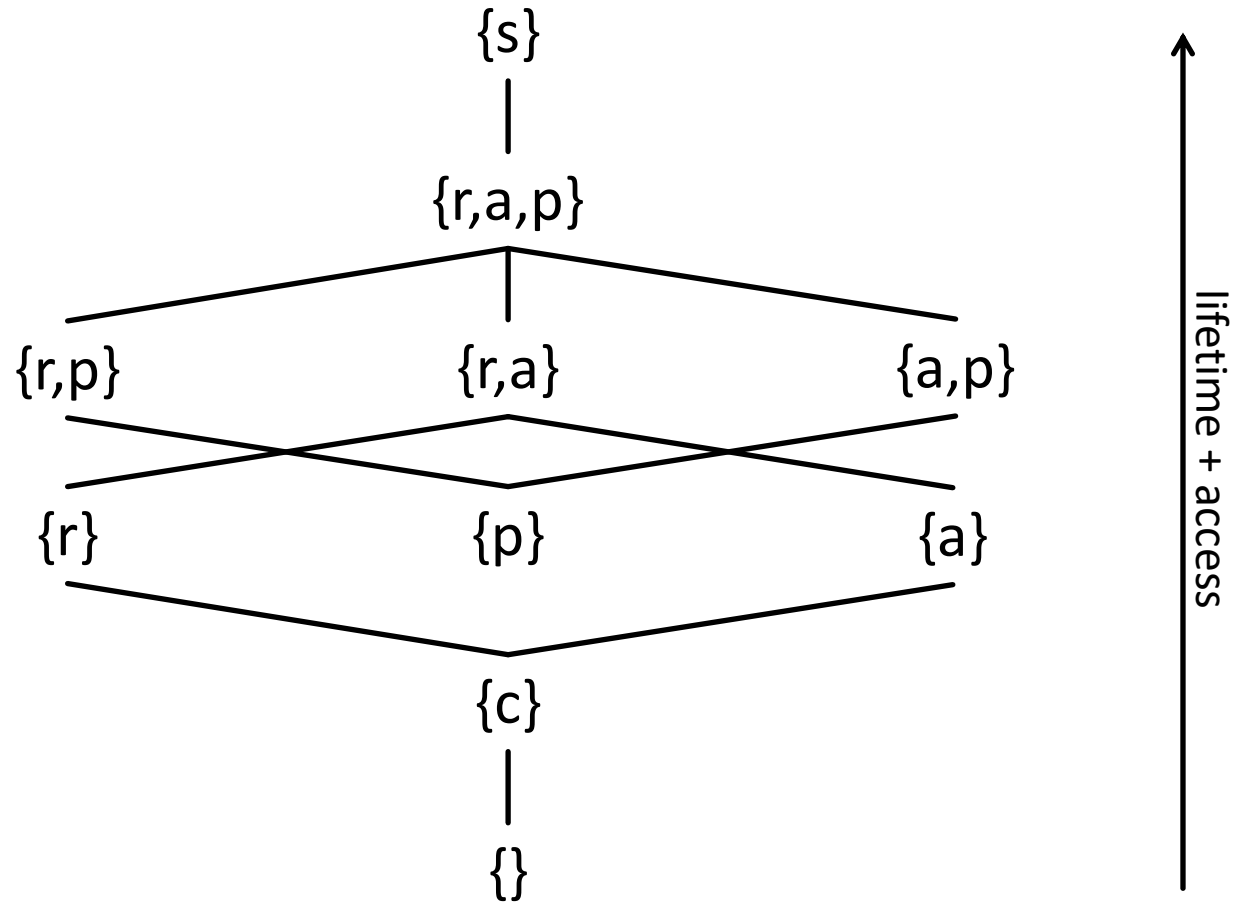
Static field writes(s):

- $\forall m . m \in \text{Access}(o)$
- $\forall t' . t' \in \text{Access}(o)$
- $\text{Lifetime}(o) > \text{Lifetime}(\text{foo})$
- $\text{Lifetime}(o) > \text{Lifetime}(t)$
- Stack Allocation: **X**
- Scalar Replacement: **X**
- Synchronization Removal: **X**

Naïve Subset Lattice



Simplified Lattice



Handling Constructor Calls

Object o = **new** Object();

NEW(java.lang.Object)
DUP
INVOKESPECIAL(java.lang.Object{ **void** <init>() })

- As known from the lecture, the JVM distinguishes between object allocation and invocations of <init>.
- Distinguish between constructor and other calls as <init> does not let escape the object per-se.
- Requires (at least some) inter-procedural analysis if it is not the constructor of java.lang.Object

May-Alias Analysis and Fields

```
public class MyClass {  
    Object f;  
    public static Object global;  
  
    public void foo() {  
        MyClass c1 = new MyClass();  
        MyClass.global = c1; // c1 escapes  
        MyClass c2 = c1; // requires may-alias detection  
        Object o = new Object();  
        c2.f = o; // c2 does escape, therefore o does as well  
    }  
}
```

Research Directions

- Use escape information for more precise call graphs for libraries.
- Modularization and soundness of call-graph and points-to analysis
- For lab, seminar or thesis topics as well as for HiWi positions contact me (kuebler@cs.tu-darmstadt.de).

References

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