

1 How do we check entailments?

1. Create set of “axioms” based on the program and the entailment to check
 - (a) Datatype Declarations
 - Enumeration types: Variables, Classes, Fields
 - ADT: Path
 - (b) Function Declarations
 - Declared: path-equivalence, instance-of, instantiated-by
 - Defined recursively: substitute
 - (c) Calculus rules as all quantified formulas
 - Static Rules: C-Refl, C-Class, C-Subst
 - Template Rules: C-Prog
 - (d) Assert entailment to be checked: $c_1, \dots, c_n \vdash c \Rightarrow \neg(c_1 \wedge \dots \wedge c_n \rightarrow c)$
2. Obtain Solution: Does the entailment contradict the rules?
 - If unsat: valid/correct entailment.
 - If sat: invalid entailment.

1.1 General First-Order Encoding Template

$$\text{Variable } (\dots) \quad (1)$$

$$\text{Field } (\dots) \quad (2)$$

$$\text{Class } (\dots) \quad (3)$$

$$\text{Path } ((\text{Variable}) (\text{Path Field})) \quad (4)$$

$$\equiv : \text{Path} \times \text{Path} \rightarrow \mathcal{B} \quad (5)$$

$$:: : \text{Path} \times \text{Class} \rightarrow \mathcal{B} \quad (6)$$

$$\text{.cls} \equiv : \text{Path} \times \text{Class} \rightarrow \mathcal{B} \quad (7)$$

$$_ \{\!\!\{ \mapsto \}\!\!\} : \text{Path} \times \text{Variable} \times \text{Path} \rightarrow \text{Path} \quad (8)$$

$$\forall p : \text{Path}. p \equiv p \quad (9)$$

$$\forall a : \mathcal{B}, p : \text{Path}, c : \text{Class}. \quad (10)$$

$$(a \rightarrow p.\text{cls} \equiv c) \rightarrow (a \rightarrow p :: c) \quad (11)$$

$$\forall a : \mathcal{B}, p : \text{Path}, q : \text{Path}, r : \text{Path}, s : \text{Path}, x : \text{Variable}. \quad (12)$$

$$(a \rightarrow p \{\!\!\{ x \mapsto r \}\!\!\} \equiv q \{\!\!\{ x \mapsto r \}\!\!\} \wedge (a \rightarrow s \equiv r)) \rightarrow \quad (13)$$

$$(a \rightarrow p \{\!\!\{ x \mapsto s \}\!\!\} \equiv q \{\!\!\{ x \mapsto s \}\!\!\}) \quad (14)$$

$$\forall a : \mathcal{B}, p : \text{Path}, c : \text{Class}, r : \text{Path}, s : \text{Path}, x : \text{Variable}. \quad (15)$$

$$(a \rightarrow p \{\!\!\{ x \mapsto r \}\!\!\} :: c \wedge (a \rightarrow s \equiv r)) \rightarrow \quad (16)$$

$$(a \rightarrow p \{\!\!\{ x \mapsto s \}\!\!\} :: c) \quad (17)$$

$$\forall a : \mathcal{B}, p : \text{Path}, c : \text{Class}, r : \text{Path}, s : \text{Path}, x : \text{Variable}. \quad (18)$$

$$(a \rightarrow p \{\!\!\{ x \mapsto r \}\!\!\}.\text{cls} \equiv c \wedge (a \rightarrow s \equiv r)) \rightarrow \quad (19)$$

$$(a \rightarrow p \{\!\!\{ x \mapsto s \}\!\!\}.\text{cls} \equiv c) \quad (20)$$

$$\forall a : \mathcal{B}, p : \text{Path}. (a \rightarrow \bigwedge _) \rightarrow (a \rightarrow p :: _) \quad (21)$$

1.2 Natural Numbers Program

$$\text{Zero}(x. \epsilon) \quad (22)$$

$$\text{Succ}(x. x.p :: \text{Nat}) \quad (23)$$

$$\forall x. x :: \text{Zero} \Rightarrow x :: \text{Nat} \quad (24)$$

$$\forall x. x :: \text{Succ}, x.p :: \text{Nat} \Rightarrow x :: \text{Nat} \quad (25)$$

$$\text{prev}(x. x :: \text{Nat}) : [y. y :: \text{Nat}] \quad (26)$$

$$\text{prev}(x. x :: \text{Zero}) : [y. y :: \text{Nat}] := \text{new Zero}() \quad (27)$$

$$\text{prev}(x. x :: \text{Succ}, x.p :: \text{Nat}) : [y. y :: \text{Nat}] := x.p \quad (28)$$

1.3 C-Prog Rules for Natural Numbers Program

$$\forall a:\mathcal{B}, p:\text{Path}. \quad (29)$$

$$(a \rightarrow p :: \text{Zero}) \rightarrow \quad (30)$$

$$(a \rightarrow p :: \text{Nat}) \quad (31)$$

$$\forall a:\mathcal{B}, p:\text{Path}. \quad (32)$$

$$(a \rightarrow p :: \text{Succ} \wedge x.p\llbracket x \mapsto p \rrbracket :: \text{Nat}) \rightarrow \quad (33)$$

$$(a \rightarrow p :: \text{Nat}) \quad (34)$$

2 Example Entailments

2.1 Working valid entailment

$$p.\text{cls} \equiv \text{Zero} \vdash p :: \text{Nat}$$

- Checks unsatisfiable
- Unsat Core: C-Class, C-Prog-Zero

2.2 Working invalid entailment

$$\vdash x \equiv y$$

- Checks satisfiable
- Model: $p \equiv q \doteq p = q$

$$a \equiv b \vdash a \equiv c$$

Checks satisfiable

Model:

$$\begin{aligned}
& \text{helper}(q:\text{Path}) := \\
& \quad q = a \text{ ? } a : \\
& \quad \quad q = c \text{ ? } c : \\
& \quad \quad \quad q = x \text{ ? } x : \\
& \quad \quad \quad \quad q = b \text{ ? } b : \\
& \quad \quad \quad \quad \quad q = b.p \text{ ? } b.p : \\
& \quad \quad \quad \quad \quad \quad q = a.b \text{ ? } a.b : c.p \\
& p:\text{Path} \equiv q:\text{Path} := \\
& \quad \text{let } a_1 := \text{helper}(p) = a.p \wedge \text{helper}(q) = b.p \\
& \quad \quad a_2 := \text{helper}(p) = a.p \wedge \text{helper}(q) = a.p \\
& \quad \quad a_3 := \text{helper}(p) = c.p \wedge \text{helper}(q) = c.p \\
& \quad \quad a_4 := \text{helper}(p) = b.p \wedge \text{helper}(q) = b.p \\
& \quad \quad a_5 := \text{helper}(p) = b.p \wedge \text{helper}(q) = a.p \\
& \quad \text{in } \bigvee_{i \in \{j \mid 1 \leq j \leq 5\}} a_i \\
& \quad \vee \text{helper}(p) = a \wedge \text{helper}(q) = a \\
& \quad \vee \text{helper}(p) = b \wedge \text{helper}(q) = b \\
& \quad \vee \text{helper}(p) = b \wedge \text{helper}(q) = a \\
& \quad \vee \text{helper}(p) = c \wedge \text{helper}(q) = c \\
& \quad \vee \text{helper}(p) = x \wedge \text{helper}(q) = x \\
& \quad \vee \text{helper}(p) = a \wedge \text{helper}(q) = b
\end{aligned}$$

2.3 Non-working invalid entailment

$x :: \text{Succ}, x.p :: \text{Zero} \vdash x :: \text{Zero}$

- Location: /paper/dep-classes/smt/semantic_entailment/fieldAccessTimeout.smt
- Solver has “infinite” runtime.
- Wanted behavior: Checks satisfiable.
- Current Solution: Impose a timeout on the solver.
 - Assume an entailment to be invalid if the solver times out (to retain soundness).
 - Procedure isn’t complete anymore.
 - Balance the timeout between the amount of checkable valid entailments and tolerable waiting time. (Set it such that the majority of valid entailments can still be checked.)

3 Undefinedness

3.1 adding $\neg x \equiv x.p$ results in ???

3.2 asserting $x :: \text{Zero} \wedge x :: \text{Nat}$ is SAT

But with further investigation, again doesn't seem to be a problem, as we are searching conflicts

4 Faithful encoding

5 Change entailment translation

Instead $\bigwedge c_i \rightarrow \text{conclusion}$, use multiple asserts. for c_i in *context*. assert c_i
assert $\neg \text{conclusion}$

6 Explicitly define constraint predicates

6.1 Non-working invalid entailment

$x \equiv y \vdash x \equiv z$

- Checks unsatisfiable, but shouldn't
- Unsat Core: C-Subst-PathEq
- TODO: possible reason? (check instantiation of subst rule)