

1 How do we check entailments?

1. Create set of “axioms” based on the program and the entailment to check
 - (a) Datatype Declarations
 - Enumeration types: Variables, Classes, Fields
 - ADT: Path
 - (b) Function Declarations
 - Declared: path-equivalence, instance-of, instantiated-by
 - Defined recursively: substitute
 - (c) Calculus rules as all quantified formulas
 - Static Rules: C-Refl, C-Class, C-Subst
 - Dynamic Rules: C-Prog
 - (d) Assert entailment to be checked: $c_1, \dots, c_n \vdash c \Rightarrow \neg(c_1 \wedge \dots \wedge c_n \rightarrow c)$
2. Get Solution
 - If unsat: valid/correct entailment.
 - If sat: invalid entailment.

1.1 General First-Order Encoding Template

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|--|------|
| Variable (...) | (1) |
| Field (...) | (2) |
| Class (...) | (3) |
| Path ((Variable) (Path Field)) | (4) |
| $\equiv : \text{Path} \times \text{Path} \rightarrow \mathcal{B}$ | (5) |
| $:: : \text{Path} \times \text{Class} \rightarrow \mathcal{B}$ | (6) |
| $\text{.cls} \equiv : \text{Path} \times \text{Class} \rightarrow \mathcal{B}$ | (7) |
| $_ \llbracket \mapsto \rrbracket : \text{Path} \times \text{Variable} \times \text{Path} \rightarrow \text{Path}$ | (8) |
| $\forall p : \text{Path}. p \equiv p$ | (9) |
| $\forall a : \mathcal{B}, p : \text{Path}, c : \text{Class}.$ | (10) |
| $(a \rightarrow p.\text{cls} \equiv c) \rightarrow (a \rightarrow p :: c)$ | (11) |
| $\forall a : \mathcal{B}, p : \text{Path}, q : \text{Path}, r : \text{Path}, s : \text{Path}, x : \text{Variable}.$ | (12) |
| $(a \rightarrow p \llbracket x \mapsto r \rrbracket \equiv q \llbracket x \mapsto r \rrbracket \wedge (a \rightarrow s \equiv r)) \rightarrow$ | (13) |
| $(a \rightarrow p \llbracket x \mapsto s \rrbracket \equiv q \llbracket x \mapsto s \rrbracket)$ | (14) |
| $\forall a : \mathcal{B}, p : \text{Path}, c : \text{Class}, r : \text{Path}, s : \text{Path}, x : \text{Variable}.$ | (15) |
| $(a \rightarrow p \llbracket x \mapsto r \rrbracket :: c \wedge (a \rightarrow s \equiv r)) \rightarrow$ | (16) |
| $(a \rightarrow p \llbracket x \mapsto s \rrbracket :: c)$ | (17) |
| $\forall a : \mathcal{B}, p : \text{Path}, c : \text{Class}, r : \text{Path}, s : \text{Path}, x : \text{Variable}.$ | (18) |
| $(a \rightarrow p \llbracket x \mapsto r \rrbracket.\text{cls} \equiv c \wedge (a \rightarrow s \equiv r)) \rightarrow$ | (19) |
| $(a \rightarrow p \llbracket x \mapsto s \rrbracket.\text{cls} \equiv c)$ | (20) |
| $\forall a : \mathcal{B}, p : \text{Path}. (a \rightarrow \bigwedge _) \rightarrow (a \rightarrow p :: _)$ | (21) |

1.2 Natural Numbers Program

$$\text{Zero}(x. \epsilon) \quad (22)$$

$$\text{Succ}(x. x.p :: \text{Nat}) \quad (23)$$

$$\forall x. x :: \text{Zero} \Rightarrow x :: \text{Nat} \quad (24)$$

$$\forall x. x :: \text{Succ}, x.p :: \text{Nat} \Rightarrow x :: \text{Nat} \quad (25)$$

$$\text{prev}(x. x :: \text{Nat}) : [y. y :: \text{Nat}] \quad (26)$$

$$\text{prev}(x. x :: \text{Zero}) : [y. y :: \text{Nat}] := \text{new Zero}() \quad (27)$$

$$\text{prev}(x. x :: \text{Succ}, x.p :: \text{Nat}) : [y. y :: \text{Nat}] := x.p \quad (28)$$

1.3 C-Prog Rules for Natural Numbers Program

$$\forall a:\mathcal{B}, p:\text{Path}. \quad (29)$$

$$(a \rightarrow p :: \text{Zero}) \rightarrow \quad (30)$$

$$(a \rightarrow p :: \text{Nat}) \quad (31)$$

$$\forall a:\mathcal{B}, p:\text{Path}. \quad (32)$$

$$(a \rightarrow p :: \text{Succ} \wedge x.p\llbracket x \mapsto p \rrbracket :: \text{Nat}) \rightarrow \quad (33)$$

$$(a \rightarrow p :: \text{Nat}) \quad (34)$$

2 Example Entailments

2.1 Working valid entailment

$p.\text{cls} \equiv \text{Zero} \vdash p :: \text{Nat}$

- Checks unsatisfiable
- Unsat Core: C-Class, C-Prog-Zero

2.2 Working invalid entailment

TODO: search typechecking tests for (non-trivial) SAT entailment
 $\vdash x \equiv y$

- Checks satisfiable
- Model: $p \equiv q \doteq p = q$

2.3 Non-working invalid entailment

$x \equiv y \vdash x \equiv z$

- Checks unsatisfiable, but shouldn't
- Unsat Core: C-Subst-PathEq
- TODO: possible reason? (check instantiation of subst rule)

2.4 Non-working invalid entailment

fieldAccessTimeout.smt

- Neither checks satisfiable nor unsatisfiable.
- Solver has “infinite” runtime.
- TODO: check manually for a model?

3 Undefinedness

3.1 adding $\neg x \equiv x.p$ results in ???

3.2 asserting $x :: \text{Zero} \wedge x :: \text{Nat}$ is SAT

But with further investigation, again doesn't seem to be a problem, as we are searching conflicts