



Structural Induction on Trees

Principles of Functional Programming

Structural Induction on Trees

Structural induction is not limited to lists; it applies to any tree structure.

The general induction principle is the following:

To prove a property $P(t)$ for all trees t of a certain type,

- ▶ show that $P(l)$ holds for all leaves l of a tree,
- ▶ for each type of internal node t with subtrees s_1, \dots, s_n , show that $P(s_1) \wedge \dots \wedge P(s_n)$ *implies* $P(t)$.

Example: IntSets

Recall our definition of IntSet with the operations contains and incl:

```
abstract class IntSet:  
  def incl(x: Int): IntSet  
  def contains(x: Int): Boolean  
  
object Empty extends IntSet:  
  def contains(x: Int): Boolean = false  
  def incl(x: Int): IntSet = NonEmpty(x, Empty, Empty)
```

Example: IntSets (2)

```
case class NonEmpty(elem: Int, left: IntSet, right: IntSet) extends IntSet:
```

```
  def contains(x: Int): Boolean =  
    if x < elem then left.contains(x)  
    else if x > elem then right.contains(x)  
    else true
```

```
  def incl(x: Int): IntSet =  
    if x < elem then NonEmpty(elem, left.incl(x), right)  
    else if x > elem then NonEmpty(elem, left, right.incl(x))  
    else this
```

The Laws of IntSet

What does it mean to prove the correctness of this implementation?

One way to define and show the correctness of an implementation consists of proving the laws that it respects.

In the case of IntSet, we have the following three laws:

For any set s , and elements x and y :

```
Empty.contains(x)      = false
s.incl(x).contains(x)  = true
s.incl(x).contains(y)  = s.contains(y)    if  $x \neq y$ 
```

(In fact, we can show that these laws completely characterize the desired data type).

Proving the Laws of IntSet (1)

How can we prove these laws?

Proposition 1: `Empty.contains(x) = false`.

Proof: According to the definition of `contains` in `Empty`.

Proving the Laws of IntSet (2)

Proposition 2: `s.incl(x).contains(x) = true`

Proof by structural induction on `s`.

Base case: `Empty`

`Empty.incl(x).contains(x)`

Proving the Laws of IntSet (2)

Proposition 2: `s.incl(x).contains(x) = true`

Proof by structural induction on `s`.

Base case: `Empty`

`Empty.incl(x).contains(x)`

`= NonEmpty(x, Empty, Empty).contains(x) // by definition of Empty.incl`

Proving the Laws of IntSet (2)

Proposition 2: `s.incl(x).contains(x) = true`

Proof by structural induction on `s`.

Base case: `Empty`

`Empty.incl(x).contains(x)`

`= NonEmpty(x, Empty, Empty).contains(x) // by definition of Empty.incl`

`= true // by definition of NonEmpty.contains`

Proving the Laws of IntSet (3)

Induction step: `NonEmpty(x, l, r)`

`NonEmpty(x, l, r).incl(x).contains(x)`

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`NonEmpty(x, l, r).incl(x).contains(x)`

`= NonEmpty(x, l, r).contains(x) // by definition of NonEmpty.incl`

Proving the Laws of IntSet (3)

Induction step: `NonEmpty(x, l, r)`

`NonEmpty(x, l, r).incl(x).contains(x)`

`= NonEmpty(x, l, r).contains(x)` *// by definition of NonEmpty.incl*

`= true` *// by definition of NonEmpty.contains*

Proving the Laws of IntSet (4)

Induction step: `NonEmpty(y, l, r)` **where** $y < x$

`NonEmpty(y, l, r).incl(x).contains(x)`

Proving the Laws of IntSet (4)

Induction step: `NonEmpty(y, l, r)` where $y < x$

`NonEmpty(y, l, r).incl(x).contains(x)`

`= NonEmpty(y, l, r.incl(x)).contains(x) // by definition of NonEmpty.incl`

Proving the Laws of IntSet (4)

Induction step: `NonEmpty(y, l, r)` where $y < x$

`NonEmpty(y, l, r).incl(x).contains(x)`

= `NonEmpty(y, l, r.incl(x)).contains(x)` // by definition of `NonEmpty.incl`

= `r.incl(x).contains(x)` // by definition of `NonEmpty.contains`

Proving the Laws of IntSet (4)

Induction step: `NonEmpty(y, l, r)` where $y < x$

`NonEmpty(y, l, r).incl(x).contains(x)`

`= NonEmpty(y, l, r.incl(x)).contains(x) // by definition of NonEmpty.incl`

`= r.incl(x).contains(x) // by definition of NonEmpty.contains`

`= true // by the induction hypothesis`

Proving the Laws of IntSet (4)

Induction step: `NonEmpty(y, l, r)` where $y < x$

`NonEmpty(y, l, r).incl(x).contains(x)`

= `NonEmpty(y, l, r.incl(x)).contains(x)` // by definition of `NonEmpty.incl`

= `r.incl(x).contains(x)` // by definition of `NonEmpty.contains`

= `true` // by the induction hypothesis

Induction step: `NonEmpty(y, l, r)` where $y > x$ is analogous

Proving the Laws of IntSet (5)

Proposition 3: If $x \neq y$ then

$$xs.incl(y).contains(x) = xs.contains(x).$$

Proof by structural induction on s . Assume that $y < x$ (the dual case $x < y$ is analogous).

Base case: Empty

`Empty.incl(y).contains(x)`

`// to show: = Empty.contains(x)`

Proving the Laws of IntSet (5)

Proposition 3: If $x \neq y$ then

$$xs.incl(y).contains(x) = xs.contains(x).$$

Proof by structural induction on s . Assume that $y < x$ (the dual case $x < y$ is analogous).

Base case: Empty

```
Empty.incl(y).contains(x)           // to show: = Empty.contains(x)  
  
= NonEmpty(y, Empty, Empty).contains(x) // by definition of Empty.incl
```

Proving the Laws of IntSet (5)

Proposition 3: If $x \neq y$ then

$$xs.incl(y).contains(x) = xs.contains(x).$$

Proof by structural induction on s . Assume that $y < x$ (the dual case $x < y$ is analogous).

Base case: Empty

`Empty.incl(y).contains(x)` // to show: `= Empty.contains(x)`

`= NonEmpty(y, Empty, Empty).contains(x)` // by definition of `Empty.incl`

`= Empty.contains(x)` // by definition of `NonEmpty.contains`

Proving the Laws of IntSet (6)

For the inductive step, we need to consider a tree $\text{NonEmpty}(z, l, r)$. We distinguish five cases:

1. $z = x$
2. $z = y$
3. $z < y < x$
4. $y < z < x$
5. $y < x < z$

First Two Cases: $z = x$, $z = y$

Induction step: `NonEmpty(x, l, r)`

`NonEmpty(x, l, r).incl(y).contains(x)` // to show: `= NonEmpty(x, l, r).contains(x)`

First Two Cases: $z = x$, $z = y$

Induction step: `NonEmpty(x, l, r)`

`NonEmpty(x, l, r).incl(y).contains(x)` // to show: `= NonEmpty(x, l, r).contains(x)`

`= NonEmpty(x, l.incl(y), r).contains(x)` // by definition of `NonEmpty.incl`

First Two Cases: $z = x$, $z = y$

Induction step: `NonEmpty(x, l, r)`

`NonEmpty(x, l, r).incl(y).contains(x)` // to show: `= NonEmpty(x, l, r).contains(x)`

`= NonEmpty(x, l.incl(y), r).contains(x)` // by definition of `NonEmpty.incl`

`= true` // by definition of `NonEmpty.contains`

First Two Cases: $z = x$, $z = y$

Induction step: `NonEmpty(x, l, r)`

`NonEmpty(x, l, r).incl(y).contains(x)` // to show: `= NonEmpty(x, l, r).contains(x)`

`= NonEmpty(x, l.incl(y), r).contains(x)` // by definition of `NonEmpty.incl`

`= true` // by definition of `NonEmpty.contains`

`= NonEmpty(x, l, r).contains(x)` // by definition of `NonEmpty.contains`

First Two Cases: $z = x$, $z = y$

Induction step: `NonEmpty(x, l, r)`

```
NonEmpty(x, l, r).incl(y).contains(x) // to show: = NonEmpty(x, l, r).contains(x)

= NonEmpty(x, l.incl(y), r).contains(x) // by definition of NonEmpty.incl

= true // by definition of NonEmpty.contains

= NonEmpty(x, l, r).contains(x) // by definition of NonEmpty.contains
```

Induction step: `NonEmpty(y, l, r)`

```
NonEmpty(y, l, r).incl(y).contains(x) // to show: = NonEmpty(y, l, r).contains(x)
```

First Two Cases: $z = x$, $z = y$

Induction step: `NonEmpty(x, l, r)`

```
NonEmpty(x, l, r).incl(y).contains(x) // to show: = NonEmpty(x, l, r).contains(x)

= NonEmpty(x, l.incl(y), r).contains(x) // by definition of NonEmpty.incl

= true // by definition of NonEmpty.contains

= NonEmpty(x, l, r).contains(x) // by definition of NonEmpty.contains
```

Induction step: `NonEmpty(y, l, r)`

```
NonEmpty(y, l, r).incl(y).contains(x) // to show: = NonEmpty(y, l, r).contains(x)

= NonEmpty(y, l, r).contains(x) // by definition of NonEmpty.incl
```

Case $z < y$

Induction step: `NonEmpty(z, l, r)` **where** $z < y < x$

`NonEmpty(z, l, r).incl(y).contains(x)` // to show: `= NonEmpty(z, l, r).contains(x)`

Case $z < y$

Induction step: `NonEmpty(z, l, r)` where $z < y < x$

`NonEmpty(z, l, r).incl(y).contains(x)` // to show: `= NonEmpty(z, l, r).contains(x)`

`= NonEmpty(z, l, r.incl(y)).contains(x)` // by definition of `NonEmpty.incl`

Case $z < y$

Induction step: `NonEmpty(z, l, r)` where $z < y < x$

`NonEmpty(z, l, r).incl(y).contains(x)` // to show: `= NonEmpty(z, l, r).contains(x)`

`= NonEmpty(z, l, r.incl(y)).contains(x)` // by definition of `NonEmpty.incl`

`= r.incl(y).contains(x)` // by definition of `NonEmpty.contains`

Case $z < y$

Induction step: `NonEmpty(z, l, r)` where $z < y < x$

`NonEmpty(z, l, r).incl(y).contains(x)` // to show: `= NonEmpty(z, l, r).contains(x)`

`= NonEmpty(z, l, r.incl(y)).contains(x)` // by definition of `NonEmpty.incl`

`= r.incl(y).contains(x)` // by definition of `NonEmpty.contains`

`= r.contains(x)` // by the induction hypothesis

Case $z < y$

Induction step: `NonEmpty(z, l, r)` where $z < y < x$

```
NonEmpty(z, l, r).incl(y).contains(x) // to show: = NonEmpty(z, l, r).contains(x)

= NonEmpty(z, l, r.incl(y)).contains(x) // by definition of NonEmpty.incl

= r.incl(y).contains(x) // by definition of NonEmpty.contains

= r.contains(x) // by the induction hypothesis

= NonEmpty(z, l, r).contains(x) // by definition of NonEmpty.contains
```


Case $y < z < x$

Induction step: `NonEmpty(z, l, r)` **where** $y < z < x$

`NonEmpty(z, l, r).incl(y).contains(x)` // to show: `= NonEmpty(z, l, r).contains(x)`

Case $y < z < x$

Induction step: `NonEmpty(z, l, r)` where $y < z < x$

`NonEmpty(z, l, r).incl(y).contains(x)` // to show: `NonEmpty(z, l, r).contains(x)`

= `NonEmpty(z, l.incl(y), r).contains(x)` // by definition of `NonEmpty.incl`

Case $y < z < x$

Induction step: NonEmpty(z, l, r) **where** $y < z < x$

```
NonEmpty(z, l, r).incl(y).contains(x)    // to show: = NonEmpty(z, l, r).contai
```

```
= NonEmpty(z, l.incl(y), r).contains(x) // by definition of NonEmpty.incl
```

```
= r.contains(x) // by definition of NonEmpty.contains
```

Case $y < z < x$

Induction step: `NonEmpty(z, l, r)` where $y < z < x$

`NonEmpty(z, l, r).incl(y).contains(x)` // to show: `= NonEmpty(z, l, r).contains(x)`

`= NonEmpty(z, l.incl(y), r).contains(x)` // by definition of `NonEmpty.incl`

`= r.contains(x)` // by definition of `NonEmpty.contains`

`= NonEmpty(z, l, r).contains(x)` // by definition of `NonEmpty.contains`

Case $x < z$

Induction step: `NonEmpty(z, l, r)` **where** $y < x < z$

`NonEmpty(z, l, r).incl(y).contains(x)` // to show: `= NonEmpty(z, l, r).contains(x)`

Case $x < z$

Induction step: `NonEmpty(z, l, r)` where $y < x < z$

`NonEmpty(z, l, r).incl(y).contains(x)` // to show: `= NonEmpty(z, l, r).contains(x)`

`= NonEmpty(z, l.incl(y), r).contains(x)` // by definition of `NonEmpty.incl`

Case $x < z$

Induction step: `NonEmpty(z, l, r)` where $y < x < z$

`NonEmpty(z, l, r).incl(y).contains(x)` // to show: `= NonEmpty(z, l, r).contains(x)`

`= NonEmpty(z, l.incl(y), r).contains(x)` // by definition of `NonEmpty.incl`

`= l.incl(y).contains(x)` // by definition of `NonEmpty.contains`

Case $x < z$

Induction step: `NonEmpty(z, l, r)` where $y < x < z$

`NonEmpty(z, l, r).incl(y).contains(x)` // to show: `= NonEmpty(z, l, r).contains(x)`

`= NonEmpty(z, l.incl(y), r).contains(x)` // by definition of `NonEmpty.incl`

`= l.incl(y).contains(x)` // by definition of `NonEmpty.contains`

`= l.contains(x)` // by the induction hypothesis

Case $x < z$

Induction step: `NonEmpty(z, l, r)` where $y < x < z$

```
NonEmpty(z, l, r).incl(y).contains(x)    // to show: = NonEmpty(z, l, r).contains(x)

= NonEmpty(z, l.incl(y), r).contains(x)  // by definition of NonEmpty.incl

= l.incl(y).contains(x)                  // by definition of NonEmpty.contains

= l.contains(x)                          // by the induction hypothesis

= NonEmpty(z, l, r).contains(x)         // by definition of NonEmpty.contains
```

These are all the cases, so the proposition is established.

Exercise (Hard)

Suppose we add a function `union` to `IntSet`:

```
abstract class IntSet:  
  ...  
  def union(other: IntSet): IntSet  
  
object Empty extends IntSet:  
  ...  
  def union(other: IntSet) = other  
  
class NonEmpty(x: Int, l: IntSet, r: IntSet) extends IntSet:  
  ...  
  def union(other: IntSet): IntSet = l.union(r.union(other)).incl(x)
```

Exercise (Hard)

The correctness of union can be translated into the following law:

Proposition 4:

$$xs.union(ys).contains(x) = xs.contains(x) \vee ys.contains(x)$$

Show proposition 4 by using structural induction on xs .