

# Abstract Algebra and Type Classes

Principles of Functional Programming

# Doing Abstract Algebra with Type Classes

Type classes let one define concepts that are quite abstract, and that can be instantiated with many types. For instance:

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trait SemiGroup[T]:
  extension (x: T) def combine (y: T): T
```

This models the algebraic concept of a semigroup with an associative operator combine.

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This models the algebraic concept of a semigroup with an associative operator combine.

We can then define methods that work for all semigroups. For instance:

```
def reduce[T: SemiGroup](xs: List[T]): T =
    xs.reduceLeft(_.combine(_))
```

### Type Class Hierarchies

Algebraic type classes often form natural hierarchies. For instance, a *monoid* is defined as a semigroup with a left and right unit element.

Here's its natural definition:

```
trait Monoid[T] extends SemiGroup[T]:
  def unit: T
```

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```
def reduce[T](xs: List[T])(using m: Monoid[T]): T =
    xs.foldLeft(m.unit)(_.combine(_))
```

#### Using Context Bounds

In the previous example we had to pass an explicitly named type class instance m: Monoid[T] to reduce, so that we could refer to m.unit.

One could alternatively use a context bound and a summon.

```
def reduce[T: Monoid](xs: List[T]): T =
    xs.reduceLeft(summon[Monoid[T]].unit)(_.combine(_))
```

#### Streamlining Access

A simpler calling syntax can be obtained if we do some preparation in the Monoid trait itself.

```
trait Monoid[T] extends SemiGroup[T]:
  def unit: T
object Monoid:
  def apply[T](using m: Monoid[T]): Monoid[T] = m
```

This defines a global function Monoid.apply[T] that returns the Monoid[T] instance that is currently visible.

With this helper, reduce can be written like this:

```
def reduce[T: Monoid](xs: List[T]): T =
    xs.reduceLeft(Monoid[T].unit)(_.combine(_))
```

## Multiple Typeclass Instances

It's possible to have several given instances for a typeclass/type pair. For instance, Int could be a Monoid in (at least) two ways:

```
with + as combine and 0 as unit, or
with * as combine and 1 as unit.
given sumMonoid: Monoid[Int] with
  extension (x: Int) def combine(y: Int) : Int = x + y
  def unit: Int = 0

given prodMonoid: Monoid[Int] with
  extension (x: Int) def combine(y: Int) : Int = x * y
  def unit: Int = 1
```

Define the sum and product functions on List[Int] in terms of reduce.

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```
def sum(xs: List[Int]): Int = reduce(xs)(using sumMonoid)
def product(xs: List[Int]): Int = reduce(xs)(using prodMonoid)
```

What happens if you leave out the using arguments?

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```

What happens if you leave out the using arguments?

An ambiguity error.

## Typeclass Laws

Algebraic type classes are not just defined by their type signatures but also by the laws that hold for them.

For example, any given instance of Monoid[T] should satisfy the laws:

```
x.combine(y).combine(z) == x.combine(y.combine(z))
unit.combine(x) == x
x.combine(unit) == x
```

where x, y, z are arbitrary values of type T and unit = Monoid.unit[T].

The laws can be verified either by a formal or informal proof, or by testing them.

A good way to test that an instance is *lawful* is using randomized testing with a tool like ScalaCheck.