

QEC 101

Introduction to quantum error correction

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Invitation: QR code



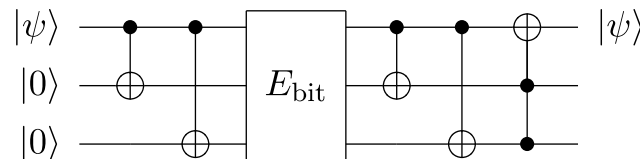
- Basic idea: using **redundant** information
- Error correction level: 7–30%
- Reed–Solomon codes
 - Polynomials over finite fields (Galois fields)
- 回形针 PaperClip : Vol.120 二维码的秘密

Quantum errors

- Bloch sphere state: $|\psi\rangle = \cos \frac{\theta}{2} |0\rangle + e^{i\phi} \sin \frac{\theta}{2} |1\rangle$
- General coherent error: $U(\delta\theta, \delta\phi) |\psi\rangle = \cos \frac{\theta+\delta\theta}{2} |0\rangle + e^{i(\phi+\delta\phi)} \sin \frac{\theta+\delta\theta}{2} |1\rangle$
 - Rotation on Bloch sphere
- Write in Pauli matrices: $U(\delta\theta, \delta\phi) = \alpha_I I + \alpha_X X + \alpha_Z Z + \alpha_Y Y (= \alpha_{XZ} XZ)$
- Two types: X -error ($|0\rangle \leftrightarrow |1\rangle$) and Z -error ($|0\rangle \rightarrow |0\rangle, |1\rangle \rightarrow -|1\rangle$)
- Challenges:
 - No-cloning theorem: $U_{\text{clone}}(|\psi\rangle \otimes |0\rangle) \nrightarrow |\psi\rangle \otimes |\psi\rangle$
 - Detect both X and Z errors simultaneously
 - Wavefunction collapse
 - Continuous error in θ and ϕ
- General procedure:
 - Syndrome extraction: measure $U |\psi\rangle$ and project to subspace (e.g. $X |\psi\rangle$)
 - Apply corresponding operator (e.g. X^{-1}) to recover $|\psi\rangle$

Bit flip code (1)

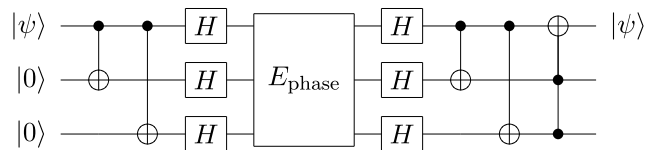
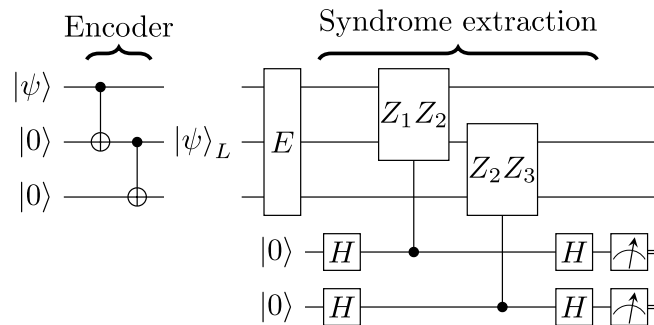
- Initial state: $|\psi\rangle = \alpha|0\rangle + \beta|1\rangle$
- Bit flip error: $P(|0\rangle \leftrightarrow |1\rangle) = p$
- No-cloning theorem:
 - $|\psi\rangle \not\rightarrow |\psi\psi\rangle$
 - $|0\rangle \rightarrow |0_L\rangle = |000\rangle, |1\rangle \rightarrow |1_L\rangle = |111\rangle$
 - Logical qubits: $|x_L\rangle$, ancilla qubits: $|x00\rangle$
- Syndrome diagnosis:
 - $P_0 = |000\rangle\langle 000| + |111\rangle\langle 111|$: no error
 - $P_1 = |100\rangle\langle 100| + |011\rangle\langle 011|$: bit flip on qubit 1
 - $P_2 = |010\rangle\langle 010| + |101\rangle\langle 101|$: bit flip on qubit 2
 - $P_3 = |001\rangle\langle 001| + |110\rangle\langle 110|$: bit flip on qubit 3
- Correct error: apply $M = \sum_i X_i P_i$



- Only single bit flip error can be corrected
- Fidelity:
 - $F_\psi = 1 - p$
 - $F_{\psi'} = (1 - p)^3 + 3p(1 - p)^2 > F_\psi$ for $p < \frac{1}{2}$

Bit flip code (2)

- Another way to do syndrome measurement:
 - $Z_1 Z_2 = Z \otimes Z \otimes I$: compare first two qubits
 - $Z_2 Z_3 = I \otimes Z \otimes Z$: compare last two qubits
- Syndrome table:
 - 00: no error
 - 01: bit flip on qubit 1
 - 10: bit flip on qubit 2
 - 11: bit flip on qubit 3
- Measurement **does not** change the state
- Sign flip code:
 - Dual between X and Z



Quantum code distance

- Classical:
 - Minimum Hamming distance between any two codewords
 - $d_{\text{classical}}(0000, 1111) = 4$
- Quantum:
 - Shortest path to get from one state to another by using Pauli operators
 - Minimum size error that will go undetected
- $[[n, k, d]]$ notation:
 - n : number of **physical** qubits
 - k : number of **logical** qubits
 - d : code distance

Stabilizer formalism (1): we need some group theory!

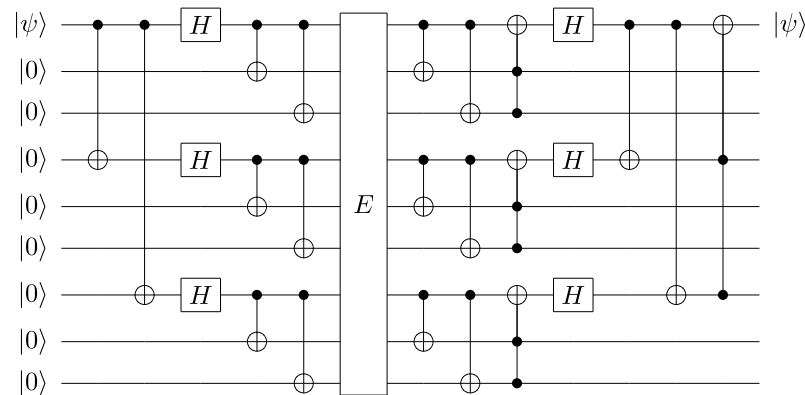
- Pauli group:
 - 1 qubit: $G_1 = \{\pm I, \pm iI, \pm X, \pm iX, \pm Y, \pm iY, \pm Z, \pm iZ\} = \langle X, Y, Z \rangle$
 - n qubits: $G_n = \{g_i \otimes \cdots \otimes g_n | g_i \in G_1\}$
- Stabilizer group: all elements that leave a state invariant
 - $\mathcal{S} = \{g \in G_n | g|\psi\rangle = |\psi\rangle\}$
 - \mathcal{S} is Abelian group
- Example:
 - 3 qubits, $\mathcal{S} = \{I, Z_1 Z_2, Z_2 Z_3, Z_3 Z_1\} = \langle Z_1 Z_2, Z_2 Z_3 \rangle$
 - $Z_1 Z_2$ can stabilize $\{|000\rangle, |001\rangle, |110\rangle, |111\rangle\}$
 - $Z_2 Z_3$ can stabilize $\{|000\rangle, |100\rangle, |011\rangle, |111\rangle\}$
 - Hence \mathcal{S} can stabilize $\{|000\rangle, |111\rangle\}$

Stabilizer formalism (2)

- Error $E \in \mathcal{E} \subset G_n$, measurement $M \in \mathcal{S}$
 - M and E anti-commute: $M(E|\psi\rangle) = -EM|\psi\rangle = -E|\psi\rangle$
 - Error will corrupt the encoded state
 - Can detect error E
 - M and E commute: $M(E|\psi\rangle) = EM|\psi\rangle = E|\psi\rangle$
 - Can't correct error E
- Stabilizer error-correcting condition:
 - Any errors $E_1, E_2 \in \mathcal{E}$ can be corrected, if
 - $E_1^\dagger E_2 \notin \mathcal{Z}(\mathcal{S})$ or
 - $E_1^\dagger E_2 \in \mathcal{S}$
 - Centralizer: $\mathcal{Z}(\mathcal{S}) = \{g \in G_n | gM = Mg, \forall M \in \mathcal{S}\}$
 - Commute with all elements in \mathcal{S}

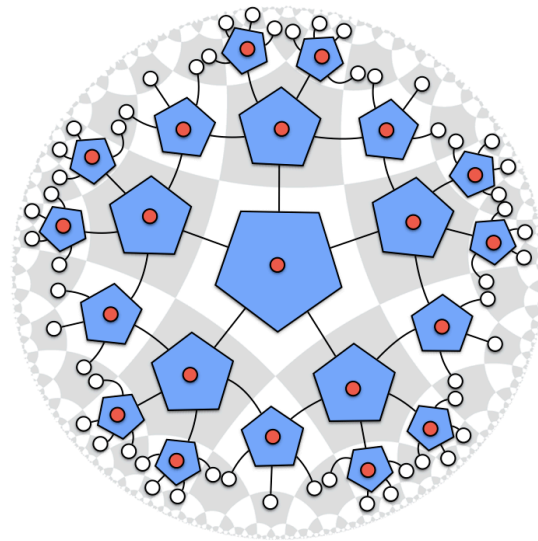
Shor code

- 9 physical qubits for 1 logical qubit
- Combine bit flip and sign flip codes
 - $Z_1 Z_2, Z_2 Z_3, \dots, Z_8 Z_9$ will detect X errors
 - $X_1 X_2 X_3 X_4 X_5 X_6, X_4 X_5 X_6 X_7 X_8 X_9$ will detect Z errors
- Stabilizer language ($[[9, 1, 3]]$ code):
 - $\mathcal{S} = \{Z_1 Z_2, \dots, Z_8 Z_9, X_1 \cdots X_6, X_4 \cdots X_9\}$
 - Example:
 - $E = X_1 Y_4 \notin \mathcal{Z}(\mathcal{S})$: anti-commute with $Z_1 Z_2$, can be corrected
 - $E = Z_1 Z_3 \in \mathcal{S}$: can be corrected



And more...

- Surface code: topological QEC
 - 2D lattice of qubits
 - Stabilizer: Pauli chains on the surface
 - Anyons, topological order, string-net condensation, categorical symmetries, etc
 - See [arXiv:quant-ph/9707021](https://arxiv.org/abs/quant-ph/9707021) etc
- HaPPY code: holographic QEC
 - Encode k "bulk" logical qubits into n "boundary" physical qubits
 - Bulk-boundary correspondence, AdS/CFT duality, Ryu–Takayanagi formula, tensor networks, p -adic physics, etc
 - See [arXiv:1503.06237](https://arxiv.org/abs/1503.06237) and [1802.01040](https://arxiv.org/abs/1802.01040)



References

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