CS472 Assignment 07: Computational Complexity

1. Write a program that a) Converts a Sudoku puzzle stored in a 9x9 array into a graph, per the construction above, b) uses your favorite graph library to 9-color the resulting graph, c) and then uses the coloring to complete the puzzle.

```
import networkx as nx
from networkx.algorithms.coloring import greedy color
def sudoku to graph(sudoku):
```

```
def solve sudoku(sudoku):
  coloring = greedy color(G, strategy="largest first")
def print sudoku(sudoku):
sudoku = [
```

```
# Solve the Sudoku

solution = solve_sudoku(sudoku)

# Print the solved Sudoku

print("Solved Sudoku:")

print_sudoku(solution)
```

```
(cs472) $ python prog3_sudoku.py

Solved Sudoku:

4 5 10 3 7 1 9 2 6

9 3 1 2 8 4 5 10 7

6 8 2 9 10 5 4 3 1

1 10 6 5 3 2 7 4 9

3 4 5 10 6 7 8 1 2

7 2 9 4 1 8 3 6 5

5 11 3 6 4 10 1 9 8

8 7 4 1 2 9 6 5 3

2 1 12 8 5 3 10 7 4

(cs472) $
```

- **2.**Sorting n numbers can be reduced to convex hull computation by mapping so that the convex hull then yields points sorted by x. Since sorting requires $\Omega(n \log n)$, convex hull (in this case) has a lower bound of $\Omega(n \log n)$. Source: https://www.geeksforgeeks.org/convex-hull-algorithm/
- **3.** This is a *reduction* algorithm that transforms the Max-Cut problem into another problem. The goal is to visit each edge at least once. To reduce Max-Cut in a planar graph G (where G represents graph) to a shortest tour in its dual G^* , we follow these steps: 1) Embed G by drawing G planarly 2) Form G where G have G have G and G have G weighted edges in G^* . 3) Then we find the shortest tour visiting all edges of G^* . 4) Tour in G^* and G (partitioning vertices).5) Select the maximum-weight cut (Max-Cut step). So the complete algorithm is:
 - 1. Embed G.
 - 2. Form dual *G*.

- 3. Find shortest tour in *G*.
- 4. Tour \rightarrow cut in G.
- 5. Select max-weight cut.
- **4.** The complexity class P contains decision problems solvable in polynomial time. However, the given brute-force algorithm has a time complexity of O(n). Since the input size (b) is the number of bits to represent n, the algorithm's is actually exponential time complexity, meaning the algorithm does not run in polynomial time and therefore does not put the composite number problem into the class P. source: https://www.geeksforgeeks.org/composite-number/