

Happy New Year 2006!

## Introduction and Basic Concepts

CS240: Data Structures and Data Management – Lecture 02

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Happy New Year 2006!

## Outline

## When and Where

CS240  
Tuesday, Thursday  
10am-11:20am  
MC 2054

## Who

- ▶ Instructor: J       Barbay
  - [email](#) – jbarbay (at) uwaterloo.ca
  - [contact](#) – DC 2332, x7824
  - [hours](#) – Thursday 1-3pm
- ▶ Tutor: Olga Miltchman
  - [email](#) – omiltchm (at) student.cs.uwaterloo.ca
  - [contact](#) – DC ???
  - [hours](#) – ???
- ▶ Administrative: Fenglian Qiu
  - [email](#) – f2qiu (at) cs.uwaterloo.ca
  - [contact](#) – DC 3115, Ext.2753

## Electronic Resources

[Web Page](#) — <http://www.student.cs.uwaterloo.ca/~cs240>

- ▶ Before the class: slide handout.
- ▶ After the class: lecture summary.
- ▶ Every two weeks: a new assignment.
- ▶ Useful links and policies.

[News Group](#) — <news:uw.cs.cs240>

## References

*Introduction to Algorithms* [CLRS]

- ▶ Covers 40% of our course material.
- ▶ Required readings for the course are given in the web page under the “Schedule” link.

*Algorithm Design* [Goodrich/Tamassia]

- ▶ Additional coverage for some specific topics.

## Mark Breakdown

Prospective Mark Breakdown:

- ▶ Assignments 30%
- ▶ Midterm 30%
- ▶ Final Exam 40%

You pass the course iff your total average is  $\geq 50\%$ .

## Assignments

- ▶ One assignment every two weeks, five in total.
- ▶ Hand-in and hand-out **electronically**.
- ▶ Release and retrieval on **Tuesdays**.  
Each assignment is worth 7 marks,  
**1 mark** removed per day late.
- ▶ Programming can be done in either Java or C++

**All programming assignments will be tested in the Undergrad Math/CS Unix Environment.**

## Policies – Academic Discipline

University Policy 71 (“Student Academic Discipline Policy”) contains relevant information and is available from the Web site of the University Secretariat at

<http://www.adm.uwaterloo.ca/infosec/>

- ▶ First offense:
  - ▶ –100% on the assignment,
  - ▶ at least –5% on your final course grade.
- ▶ Second offense: suspension for a term.

## Summary

- ▶ All the information is on the **webpage** of the course at  
<http://www.student.cs.uwaterloo.ca/~cs240/>
- ▶ Regularly check the **newsgroup** to be informed of last minute changes.  
<news://uw.cs.cs240>
- ▶ See me or the tutor if necessary:
  - ▶ J          , DC2332

## Outline

## Course Objectives.

- ▶ Sequel to CS134.
- ▶ Now focus on Data Structures, and Abstract Data Types.
- ▶ On the way:
  - ▶ Some more (light) mathematic analysis.
  - ▶ Some notions of the limits of computation.
  - ▶ some SQL.

## Course Topics

- ▶ Our Computational Model
- ▶ Time and Space Analysis
- ▶ Lists
- ▶ Graphs
- ▶ Search Trees
- ▶ Priority Queues
- ▶ Hashing
- ▶ Text Compression
- ▶ Pattern Matching
- ▶ Sorting
- ▶ Database Systems
- ▶ Memory Management

## Story Line

1. Analysis (Measuring Tape)
  - ▶ asymptotic Worst Case.
  - ▶ some Average Case.
2. Abstract Data Types (Kind of Tool: a hammer or a saw?)
  - ▶ Stack.
  - ▶ Queue.
  - ▶ Graph.
  - ▶ Tree.
  - ▶ Dictionary.
3. Data Structures (Material: Metal or Rubber hammer?)
  - ▶ array, matrix.
  - ▶ pointers, list.
  - ▶ hash
4. Applications (What to build: a chair or a table?)
  - ▶ Sorting
  - ▶ Text Compression
  - ▶ Text Search
  - ▶ Database
  - ▶ Memory Management

## Summary

- ▶ Abstract Data Types **ADT** and **Data Structures**.
- ▶ From the Theory to the Applications.
- ▶ Some **lower bounds**.
- ▶ Some **SQL**.

## Outline

## Log and Exponent Identities

The more common identities you will likely use:

- ▶  $\log_b b^a =$
- ▶  $b^{\log_b a} =$
- ▶  $(b^a)^c =$
- ▶  $b^a b^c =$
- ▶  $\log_b(ac) =$
- ▶  $\log_b(a^c) =$
- ▶  $\log_b a =$
- ▶  $b^{\log_c a} =$

For short, note  $\log_2 a$  as  $\lg a$

## Log/Exponent Identities (Cont')

- ▶ **Example.** Simplify:

$$\begin{aligned}\lg(2^n) + n^2 2^{3 \lg n} &= \\ &= \\ &= \\ &= \\ &= \end{aligned}$$

- ▶ May need floors and ceilings:

- ▶  $\lfloor 3.14159265 \rfloor =$
- ▶  $\lceil 3.14159265 \rceil =$

## Common Summations

- ▶ Arithmetic series:

$$\sum_{i=1}^n i =$$
$$\sum_{i=1}^n i^2 =$$

## Common Summations (cont')

- ▶ Useful approximation:

$$\sum_{i=1}^n i^k \approx \frac{n^{k+1}}{k+1}$$

- ▶ Geometric series (where  $a \neq 1$ ):

$$\sum_{i=0}^n a^i = \frac{a^{n+1} - 1}{a - 1}$$

- ▶ Infinite series (where  $0 < a < 1$ ):

$$\sum_{i=0}^{\infty} a^i = \frac{1}{1 - a}$$

## Derivations

**Example:**

$$\sum_{i=0}^{\infty} \frac{1}{2^i} =$$
$$=$$
$$=$$

## Factorial

- ▶ The number of arrangements of  $n$  distinct objects is  $n!$ .
- ▶  $n! = n \times (n-1) \times \cdots \times 2$
- ▶ Stirling's approximations:
  - ▶  $n! \approx \sqrt{2\pi n} \left(\frac{n}{e}\right)^n$
  - ▶  $\log(n!) \approx n \lg(n) - n + \frac{\ln n}{2} + \frac{\ln 2\pi}{2}$

## Comparison of Algorithms

How do we find out which of these algorithms is the best?

- ▶ Selection Sort
- ▶ Merge Sort
- ▶ Counting Sort

## Selection sort

```
for i:=1 to n
  | min:=i;
  | for j:=i+1 to n
  | | if a[min]>a[j] min:=j;
  | tmp:=a[i]; a[i]:=a[min]; a[min]:=tmp;
```

## Merge sort

```
function sort(from,to)
  | if (from>to)
  | | mid:=floor((from+to)/2);
  | | sort(from,mid); sort(mid+1,to);
  | | merge(from,mid,to);

function merge(from,mid,to)
  | copy a[from..mid] to a new array b
  | copy a[mid+1..to] to a new array c
  | add infinity to both b and c as the last element
  | k:=1; m:=1;
  | for j:=from to to
  | | if b[k]<c[m] then
  | | | a[j]:=b[k]; k++;
  | | else
  | | | a[j]:=c[m]; m++;
```

## Counting sort

Assume that all numbers in the array are between 1 and 1000:

```
clear array count[1..1000];
for i:=1 to n
  | count[a[i]]++;
k:=0;
for i:=1 to 1000
  | for j:=1 to count[i];
  | | a[k]:=i; k++;
```

## Comparison of Algorithms

How do we find out which of these algorithms is the best?

- ▶ Study worst/average case on instances of same size
- ▶ Measure time by key operations
- ▶ Suppose Uniform Time Access (RAM model)

## Summary

- ▶ There are some Math Formula worth remembering.
- ▶ There are some algorithms worth studying.