

Assignment 3

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Question 1

1.1

1.2

1.3

1.4

1.5

Question 2

2.1

See in code

2.2

The contains method scans the list to find the pair of nodes (pred,curr) reachable from head such that $\text{pred.next} == \text{curr}$, $\text{pred.key} \leq \text{key}$ and $\text{curr.key} < \text{key}$. The traversal uses hand-over-hand locking.

- Item is not in the list
When $\text{curr.key} == \text{key}$ is false that is $\text{curr.key} < \text{key}$. From the sortedness invariant of the list, $\text{pred.next} == \text{curr}$, and $\text{pred.key} \leq \text{key}$ we conclude that item cannot be in the set.
- Item is in the list
When $\text{curr.key} == \text{key}$ is true, then from the uniqueness of keys that $\text{curr.item} = \text{item}$. Hence, item is in the set.

Question 3

3.1

Working on that

3.2

The difficulty is to define the criteria where the indexes are valid to be read or written. We had to add two more atomic variables: the first representing how many elements can still be inserted; the second represents the valid indices for reading.

Question 4

4.3

4.4

First, the implementation breaks up the matrix-by-vector multiplication into its various vector-vector dot products. It does this by breaking the matrix into two halves, then halving again and again until it reaches a half that consists of only one row. At this point, it initiates a dot product on that row.

Note that the number of *work nodes* of this step is therefore $\Theta(2N)$, the number of nodes in a fully-balanced binary tree containing N leaves.

Each “leaf” of this binary tree is a dot product to be computed. The dot product is broken up into the sum of two *sub-dot-products*, halving all the way down as before. At the bottom, there is a single two-integer multiplication to do. Once again, this is $\Theta(2N)$ work nodes.

Because $2N \cdot 2N \in \Theta(N^2)$, the implementation achieves work $\Theta(N^2)$.

Also note that the longest possible path through this graph is all the way to bottom of one tree, then another, then back up both of them. That's $\Theta(4 \cdot \log_2(N)) \in \Theta(\log_2(N))$, so the implementation meets both the *work* and *critical path* requirements.

As for the parallelism,