

ECSE 420 Assignment 2 Report

Group 19

Stuart Mashaal
260639962

Oliver Tse Sakkwun
260604362

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Question 1

1.2

Some stuff

1.4

Some stuff

1.5

Some stuff

Question 2

2.1 - LockOne

No. If the shared `flag` atomic registers are replaced by regular registers, `LockOne` does not still provide mutual exclusion. **Proof:**

Thread 1	Thread 2
1 <code>public void lock() {</code>	1 <code>//</code>
2 <code>int i = ThreadID.get();</code>	2 <code>//</code>
3 <code>int j = i - 1;</code>	3 <code>//</code>
4 <code>flag[i] = true // only local</code>	4 <code>//</code>
5 <code>//</code>	5 <code>public void lock() {</code>
6 <code>//</code>	6 <code>int i = ThreadID.get();</code>
7 <code>//</code>	7 <code>int j = i - 1;</code>
8 <code>//</code>	8 <code>flag[i] = true // only local</code>
9 <code>//</code>	9 <code>while (flag[j]) {} // false</code>
10 <code>//</code>	10 <code>} // enter critical section</code>
11 <code>while (flag[j]) {} // false</code>	11 <code>//</code>
12 <code>} // enter critical section</code>	12 <code>//</code>

The problem is that if we don't use atomic registers, then the writes to `flag` are not necessarily shared right away. This means that both threads can write to `flag` without the other being able to detect it. Then, both can enter the critical section.

2.2 - LockTwo

No, if shared variable `victim` uses a regular register instead of an atomic one `LockTwo` does not still provide mutual exclusion, for similar reasons to those of 2.1. **Proof:**

Thread 1	Thread 2
1 <code>public void lock() {</code>	1 <code>//</code>
2 <code>int i = ThreadID.get();</code>	2 <code>//</code>
3 <code>victim = i //only local</code>	3 <code>//</code>
4 <code>//</code>	4 <code>public void lock() {</code>
5 <code>//</code>	5 <code>int i = ThreadID.get();</code>
6 <code>//</code>	6 <code>victim = i //only local</code>
7 <code>//</code>	7 <code>while (victim == i){} //see t1 write</code>
8 <code>//</code>	8 <code>} //enter critical section</code>
9 <code>while (victim == i){} //see t2 write</code>	9 <code>//</code>
10 <code>} // enter critical section</code>	10 <code>//</code>

Since the `victim` register is no longer atomic, writes to it by either thread are not only *not shared immediately* but also are *not sequentially consistent*. The Java Concurrency Model does not guarantee sequential consistency without use of synchronization primitives (such as the `volatile` keyword, which makes a register atomic).

Question 3

Some stuff