Notes on **Sha1-Prefix**'s Probabilities

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1 The Sha1-Prefix Program

Sha1-Prefix is a solution program to the following *SHA1 prefix problem*:

Given an input UTF-8 string and a difficulty number $n\in\mathbb{N}$ with $1\leq n\leq 9$, do find an UTF-8 string that appended to the input string would produce a SHA1 digest of which hexadecimal representation has its leftmost n digits all equal to 0.

Sha1-Prefix is structured on two tiers:

- shal-prefix obtained from shal-prefix.c and a few other source files is the main program tackling the core of the SHA1 prefix problem. It gets the input string (not necessarily UTF-8) from stdin or from a named file and the difficulty number as a command line argument. shal-prefix will run until it finds the first suffix string that satisfies the specified prefix condition. It can be told to terminate cleanly the execution when signaled with SIGHUP if `-e|--sighup-exits' was passed on the invoking command line.
- shal-prefix.sh is a quite involved bash shell script that drives the main program's execution. The shell script's main use case is that of running shal-prefix in series controlled by time outs.

Bellow we will present simple theoretical arguments that founds the approach taken. Under reasonably acceptable assumptions, we show that running sha-prefix in series increases **Sha1-Prefix**'s probability of success producing required outcome.

2 Mathematical Evaluations

The basic theoretical assumption we're considering hereafter is that successive runs of shal-prefix by shal-prefix.sh are to be modeled as Bernoulli trials [1] with constant or, by case, non-constant probability of success.

1 Fact. If the random variable X follows the binomial distribution with parameters $n \in \mathbb{N}^*$ and $0 \le p \le 1$, then the probability of getting at least one success in n trials is given by $1 - (1 - p)^n$.

Proof. As per [2], the probability of exactly k successes in n trials is given by:

$$\Pr(X = k) = \binom{n}{k} p^k (1 - p)^{n - k}$$

Then follows easily that:

$$Pr(X \ge 1) = 1 - Pr(X < 1)$$

$$= 1 - Pr(X = 0)$$

$$= 1 - \binom{n}{0} p^0 (1 - p)^{n-0}$$

$$= 1 - (1 - p)^n$$

2 Fact. If $0 , for <math>B_n \stackrel{\text{def}}{=} 1 - (1 - p)^n$, it holds true that $p < B_n < B_{n+1}$, for all $n \in \mathbb{N}^*$.

Proof. We have that:

$$0
$$\implies (1 - p)^{n-1} < 1$$

$$\implies (1 - p)^n < 1 - p$$

$$\implies p < B_n$$$$

Moreover:

$$0
$$\implies (1 - p)^{n+1} < (1 - p)^n$$

$$\implies B_n < B_{n+1}$$$$

3 Fact. If the random variable X follows the Poisson's binomial distribution with parameters $n \in \mathbb{N}^*$ and $0 \le p_i \le 1$, for $1 \le i \le n$, then the probability of getting at least one success in n trials is given by $1 - \prod_{i=1}^n (1-p_i)$.

Proof. As per [3], the probability of exactly 0 successes in n trials is given by:

$$\Pr(X = 0) = \prod_{i=1}^{n} (1 - p_i)$$

Then:

$$Pr(X \ge 1) = 1 - Pr(X < 1)$$

$$= 1 - Pr(X = 0)$$

$$= 1 - \prod_{i=1}^{n} (1 - p_i)$$

4 Fact. If $0 < p_i < 1$ for all $i \in \mathbb{N}$, $1 \le i \le n$, for $P_n \stackrel{\mathrm{def}}{=} 1 - \prod_{i=1}^n (1-p_i)$, it holds true that $p_i < P_n < P_{n+1}$, for all $i \in \mathbb{N}$, $1 \le i \le n$, and $n \in \mathbb{N}^*$.

Proof. For arbitrary but fixed $i \in \mathbb{N}$, $1 \le i \le n$, we have that:

$$0 < p_i < 1 \implies 0 < 1 - p_i < 1$$

$$\implies \prod_{k=1}^{n} (1 - p_k) < 1 - p_i$$

$$\implies p_i < P_n$$

Moreover:

$$\begin{array}{ll} 0 < p_i < 1 & \Longrightarrow & 0 < 1 - p_i < 1 \\ & \Longrightarrow & \prod_{k=1}^{n+1} (1 - p_k) = \left(\prod_{k=1}^n (1 - p_k) \right) \cdot (1 - p_{n+1}) < \prod_{k=1}^n (1 - p_k) \\ & \Longrightarrow & P_n < P_{n+1} \end{array}$$

References

[1] Bernoulli Trial

https://en.wikipedia.org/wiki/Bernoulli_trial

[2] Binomial Distribution

https://en.wikipedia.org/wiki/Binomial_distribution

[3] Poisson Binomial Distribution

https://en.wikipedia.org/wiki/Poisson_binomial_distribution

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