## 编译原理 - 作业(4):语义分析

**Q1:** (P309, Exercise 5.1.1) For the SDD below, give annotated parse trees for the following expressions:

| PRODUCTIONS                     | SEMANTIC RULES                 |
|---------------------------------|--------------------------------|
| 1) $L \rightarrow E \mathbf{n}$ | L.val = E.val                  |
| 2) $E \rightarrow E_1 + T$      | $E.val = E_1.val + T.val$      |
| 3) $E \rightarrow T$            | E.val = T.val                  |
| 4) $T \rightarrow T_1 * F$      | $T.val = T_1.val \times F.val$ |
| 5) $T \rightarrow F$            | T.val = F.val                  |
| 6) $F \rightarrow (E)$          | F.val = E.val                  |
| 7) $F \rightarrow \text{digit}$ | F.val = digit.lexval           |

(1) 
$$(3+4)*(5+6)$$
 n

(2) 
$$(9+8*(7+6)+5)*4$$
 n

(1) 
$$(3+4)*(5+6)$$
 n (2)  $(9+8*(7+6)+5)*4$  n

(1)  $(3+4)*(5+6)$  n (2)  $(9+8*(7+6)+5)*4$  n

(1)  $(3+4)*(5+6)$  n (2)  $(9+8*(7+6)+5)*4$  n

(1)  $(2+8)*(7+6)+5)*4$  n

(1)  $(2+8)*(7+6)+5)*4$  n

(1)  $(2+8)*(7+6)+5)*4$  n

(2)  $(2+8)*(7+6)+5)*4$  n

(3)  $(2+8)*(7+6)+5)*4$  n

(4)  $(2+8)*(7+6)+5)*4$  n

(5)  $(2+8)*(7+6)+5)*4$  n

(6)  $(2+8)*(7+6)+5)*4$  n

(7)  $(2+8)*(7+6)+5)*4$  n

(8)  $(2+8)*(7+6)+5)*4$  n

(9)  $(2+8)*(7+6)+5)*4$  n

(1)  $(2+8)*(7+6)+5)*4$  n

(2)  $(2+8)*(7+6)+5)*4$  n

(3)  $(2+8)*(7+6)+5)*4$  n

(4)  $(2+8)*(7+6)+5)*4$  n

(5)  $(2+8)*(7+6)+5)*4$  n

(6)  $(2+8)*(7+6)+5)*4$  n

(7)  $(2+8)*(7+6)+5)*4$  n

(8)  $(2+8)*(7+6)+5)*4$  n

(9)  $(2+8)*(7+6)+5)*4$  n

(10)  $(2+8)*(7+6)+5)*4$  n

(11)  $(2+8)*(7+6)+5)*4$  n

(12)  $(2+8)*(7+6)+5)*4$  n

(13)  $(2+8)*(7+6)+5)*4$  n

(14)  $(2+8)*(7+6)+5)*4$  n

(15)  $(2+8)*(7+6)+5)*4$  n

(17)  $(2+8)*(7+6)+5)*4$  n

(18)  $(2+8)*(7+6)+5)*4$  n

(19)  $(2+8)*(7+6)+5$  n

Q2: (p323, Exercises 5.3.1) Below is a grammar for expressions involving operator + and integer

of floating-point operands. Floating-point numbers are distinguished by having a decimal point: 
$$E \rightarrow E + T \mid T$$

 $L \rightarrow \text{num} \cdot \text{num} \mid \text{num}$ 

(4) T > num T. type = int

Give an SDD to determine the type of each term 
$$T$$
 and expression  $E$ .

Production Rules Semantic Rules

(1)  $E \rightarrow E + T$   $[E, type = E, type = floot] | T.type = floot? floot: int

(2)  $E \rightarrow T$   $[E, type = T, type]$ 

(3)  $T \rightarrow num \cdot num$   $[E, type] = floot$$ 

**Q3:** (p317, Exercises 5.2.4) This grammar generates binary numbers with a "decimal" point:

$$S \to L \cdot L \mid L$$
$$L \to L \mid B \mid B$$

 $B \rightarrow 0 \mid 1$ 

(4) L -> B

(5) B >0

(6) B->1

- (1) Design an L-attributed SDD to compute S.val, the decimal number value of an input string. For example, the translation of string 101.101 should be the decimal number 5.625. Hint: use an inherited attribute L.side that tells which side of the decimal point a bit is on.
- (2) Draw the annotated parse tree of 101.101.

$$L.val = L.side = zleft? Li.val + 2+B.val$$

L. 1en=3

B . val= 1

L.val=0.10 B. val=/ L.side=vight

L. 
$$val = |o|$$

L.  $val = 0.10|$ 

L.  $side = left$ 

L.  $len = 3$ 

L.  $val = |o|$ 

L.  $val = |o|$ 

L.  $val = |o|$ 

L.  $val = |o|$ 

B.  $val = |o|$ 

L.  $side = vight$ 

L.  $len = 2$ 

L.  $len = 2$ 

L.  $val = |o|$ 

L.  $val = |o$ 

$$L.val = 1$$

$$L.side = left$$

5: 
$$y = x$$
;  
6:  $z = 0$ :

1: int x = 0;

8: 
$$z = z + y$$
;  
9:  $y = y + 1$ ;

Regarding the semantic analysis of variable type, we consider the following simplified grammar and syntax-directed translation (SDT):

(1) In the above SDT, both T and L have attribute 'type'. The type attribute is synthesized or inherited? Please explain.

(2) For Line 4 of the code snippet: int y, z; Construct the annotated parse tree based on the above SDT.

(2)

T. type = int L. type = int

int L. type = int

id. entry = 
$$Z$$

id. type = int

id. type = int

(3) For Lines 3, 7 and 11 of the code snippet, list the valid variables (name and type) in symbol table.