# Team24\_Assignment4\_report

106034061 曾靖渝 106062116 黃晨 106062216 馮謙

#### **Environment**

- 1. OS: Mac OS Catalina 10.15.2
- 2. CPU: 2.7 GHz 四核心Intel Core i7
- 3. RAM: 8 GB 2133 MHz LPDDR3
- 4. GPU: Intel Iris Plus Graphics 655 1536 MB

# **Implementation**

- 1. FileMgr
  - a. All the critical sections in the FileMgr are removed.
- 2. Blockld
  - a. Add private variable named "cacheHashCode" to store calculated hash code.
  - b. Add private variable named "isNew" to indicate whether it has been accessed.
  - c. In hashCode function, use isNew to check if we have calculated hash number. If we haven't entered yet, assign isNew to false, calculate the hash number, and store it in cacheHashCode.
- 3. Page
  - a. We unplug all the synchronized lock except getVal and setVal.
- 4. BufferMgr
  - a. In BufferMgr, we can make the critical section smaller by moving the "synchronized (bufferPool)" before the statement involving "bufferPool".
- 5. BufferPoolMgr
  - a. Delete "synchronized" from the function flushAll and available.

### **Experiments**

#### Comparison of BufferPoolsize(Microbenchmark)

```
21 # The running time for warming up before benchmarking
22 org.vanilladb.bench.BenchmarkerParameters.WARM_UP_INTERVAL=60000
23 # The running time for benchmarking
24 org.vanilladb.bench.BenchmarkerParameters.BENCHMARK_INTERVAL=100000
25 # The number of remote terminal executors for benchmarking
26 org.vanilladb.bench.BenchmarkerParameters.NUM_RTES=10
27 # The IP of the target database server
28 org.vanilladb.bench.BenchmarkerParameters.SERVER_IP=127.0.0.1
29 # 1 = JDBC, 2 = Stored Procedures
30 org.vanilladb.bench.BenchmarkerParameters.CONNECTION_MODE=2
31 # 1 = Micro, 2 = TPC-C,
32 org.vanilladb.bench.BenchmarkerParameters.BENCH_TYPE=1
33 # Whether it enables the built-in profiler on the server
34 org.vanilladb.bench.BenchmarkerParameters.PROFILING_ON_SERVER=false
35 # The path to the generated reports
36 org.vanilladb.bench.StatisticMgr.OUTPUT_DIR=/Users/<u>cengjingyu</u>/Desktop/experiments/<u>microbench</u>
37 # The granularity for summarizing the performance of benchmarking
38 org.vanilladb.bench.StatisticMgr.GRANULARITY=1000
39 # Whether the RTEs display the results of each transaction
40 org.vanilladb.bench.rte.TransactionExecutor.DISPLAY_RESULT=false
42
43 #
44 # Micro-benchmarks Parameters
47 # The number of items in the testing data set
48 org.vanilladb.bench.benchmarks.micro.MicrobenchConstants.NUM_ITEMS=100000
49 # The ratio of read-write transactions during benchmarking
50 org.vanilladb.bench.benchmarks.micro.rte.MicrobenchmarkParamGen.RW_TX_RATE=0.5
51 # The ratio of long-read transactions during benchmarking
52 org.vanilladb.bench.benchmarks.micro.rte.MicrobenchmarkParamGen.LONG_READ_TX_RATE=0.0
53 # The number of read records in a transaction
54 org.vanilladb.bench.benchmarks.micro.rte.MicrobenchmarkParamGen.TOTAL_READ_COUNT=10
55 # The number of hot record in the read set of a transaction
56 org.vanilladb.bench.benchmarks.micro.rte.MicrobenchmarkParamGen.LOCAL_HOT_COUNT=1
57 # The ratio of writes to the total reads of a transaction
58 org.vanilladb.bench.benchmarks.micro.rte.MicrobenchmarkParamGen.WRITE_RATIO_IN_RW_TX=0.5
59 # The conflict rate of a hot record
60 org.vanilladb.bench.benchmarks.micro.rte.MicrobenchmarkParamGen.HOT_CONFLICT_RATE=0.08
```

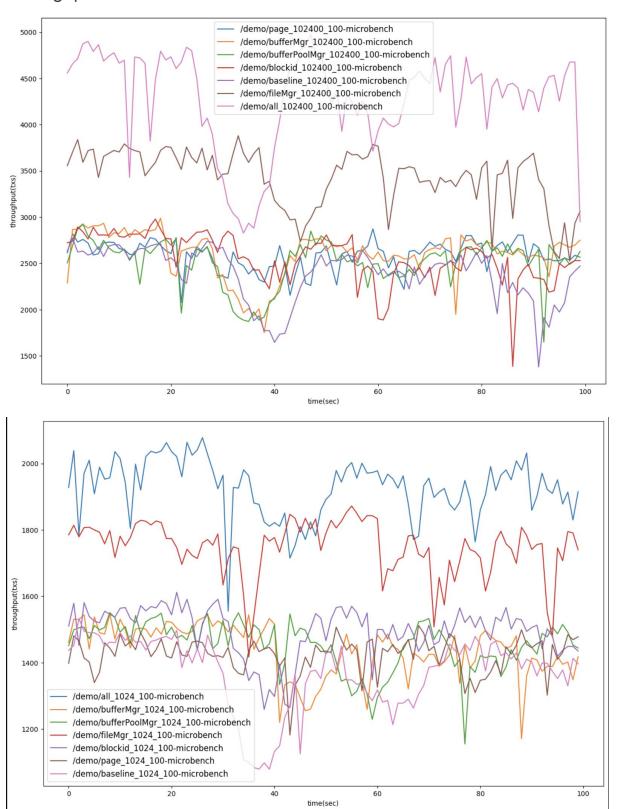
- Result (Time Interval= 100s)
  - a. BufferPoolSize=102400:

Speedup=427874(commits)/238933(commits)=1.790(79%)

b. BufferPoolSize=1024:

Speedup=192590(commits)/136930(commits)=1.406(40.6%)

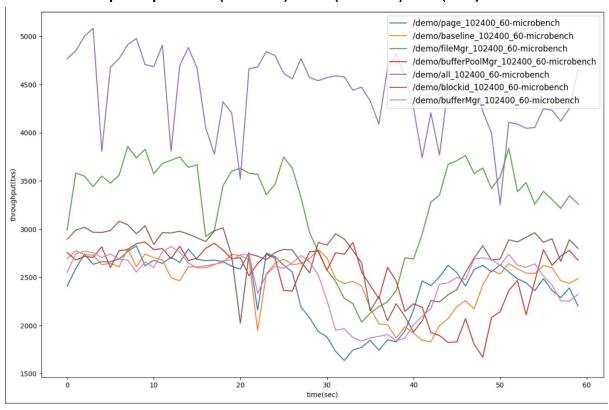
### Throughput

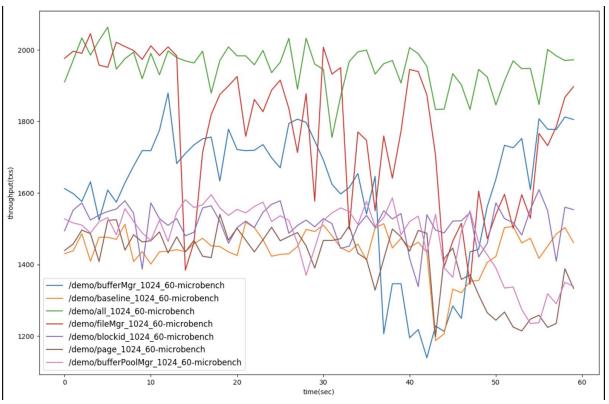


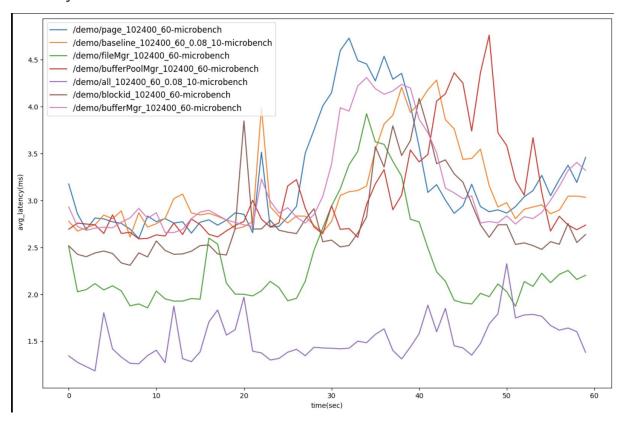
- 2. Result(Time Interval = 60s)
  - a. BufferPoolSize=102400:

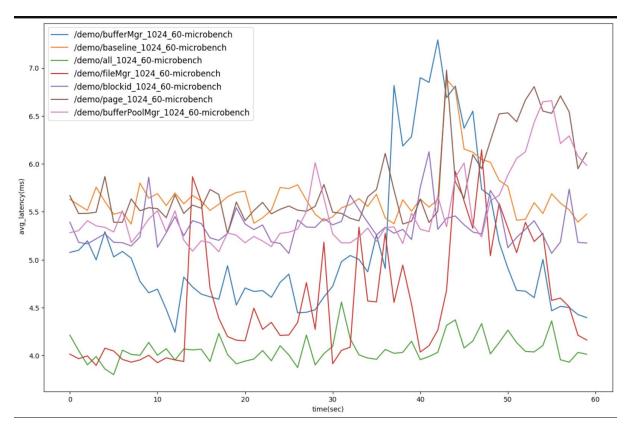
Speedup=265853(commits)/148399(commits)=1.791(79.1%)

# b. BufferPoolSize=1024:Speedup=117089(commits)/86460(commits)=1.35(35%)





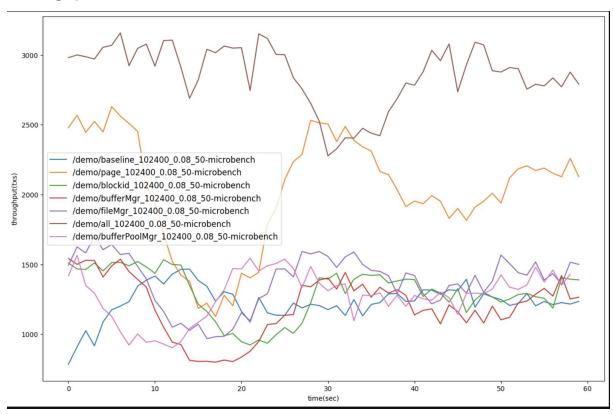


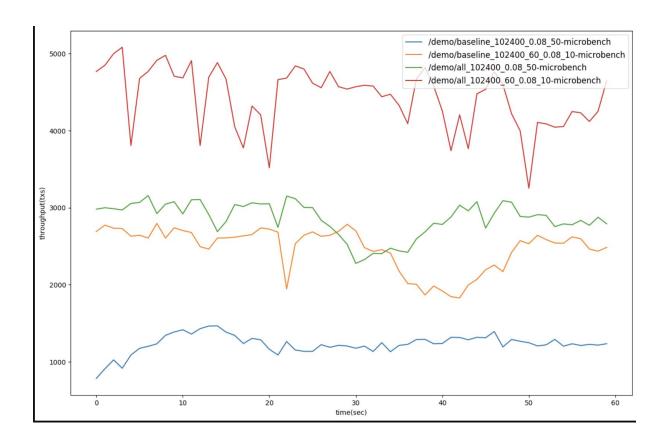


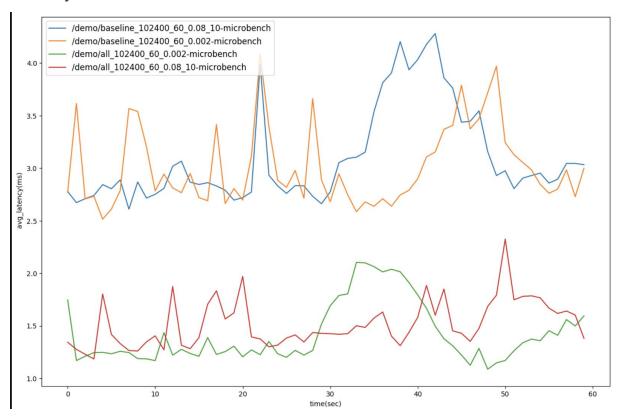
## Comparison of NUM\_RTES(Microbenchmark)

- 1. Result (BufferPoolSize=102400)
  - a. NUM\_RTES=10:
  - c. Speedup=265853(commits)/148399(commits)=1.791(79.1%)
  - b. NUM\_RTES=50 : Speedup=171426(commits)/73962(commits)=2.317(131.7%)

### Throughput





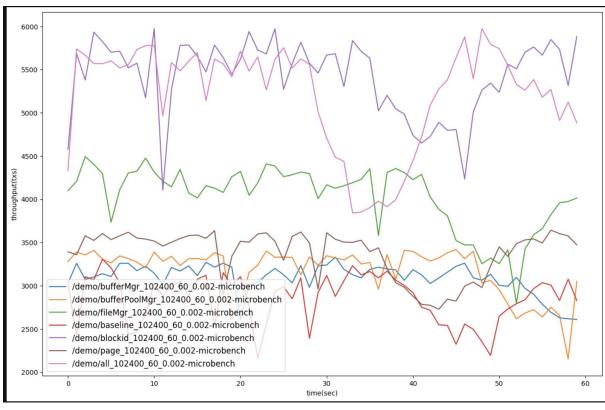


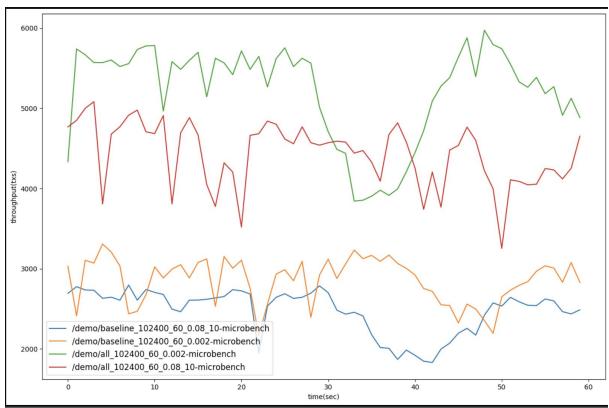
#### Comparison of HOT\_CONFLICT\_RATE(Microbenchmark)

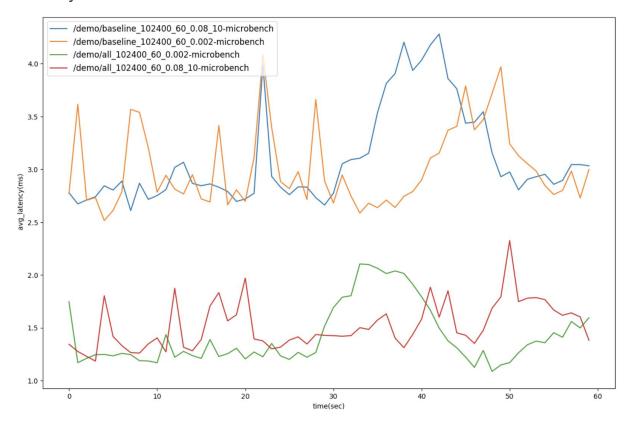
```
21 # The running time for warming up before benchmarking
22 org.vanilladb.bench.BenchmarkerParameters.WARM_UP_INTERVAL=60000
23 # The running time for benchmarking
24 org.vanilladb.bench.BenchmarkerParameters.BENCHMARK_INTERVAL=60000
25 # The number of remote terminal executors for benchmarking
26 org.vanilladb.bench.BenchmarkerParameters.NUM_RTES=10
27 # The IP of the target database server
28 org.vanilladb.bench.BenchmarkerParameters.SERVER_IP=127.0.0.1
29 # 1 = JDBC, 2 = Stored Procedures
30 org.vanilladb.bench.BenchmarkerParameters.CONNECTION_MODE=2
31 # 1 = Micro, 2 = TPC-C,
32 org.vanilladb.bench.BenchmarkerParameters.BENCH_TYPE=1
33 # Whether it enables the built-in profiler on the server
34 org.vanilladb.bench.BenchmarkerParameters.PROFILING_ON_SERVER=false
35 # The path to the generated reports
36 org.vanilladb.bench.StatisticMgr.OUTPUT_DIR=/Users/cengjingyu/Desktop/experiments/microbench
37 # The granularity for summarizing the performance of benchmarking
38 org.vanilladb.bench.StatisticMgr.GRANULARITY=1000
39 # Whether the RTEs display the results of each transaction
40 org.vanilladb.bench.rte.TransactionExecutor.DISPLAY_RESULT=false
42
43 #
44 # Micro-benchmarks Parameters
47 # The number of items in the testing data set
48 org.vanilladb.bench.benchmarks.micro.MicrobenchConstants.NUM_ITEMS=100000
49 # The ratio of read-write transactions during benchmarking
50 org.vanilladb.bench.benchmarks.micro.rte.MicrobenchmarkParamGen.RW_TX_RATE=0.5
51 # The ratio of long-read transactions during benchmarking
52 org.vanilladb.bench.benchmarks.micro.rte.MicrobenchmarkParamGen.LONG_READ_TX_RATE=0.0
53 # The number of read records in a transaction
54 org.vanilladb.bench.benchmarks.micro.rte.MicrobenchmarkParamGen.TOTAL_READ_COUNT=10
55 # The number of hot record in the read set of a transaction
56 org.vanilladb.bench.benchmarks.micro.rte.MicrobenchmarkParamGen.LOCAL_HOT_COUNT=1
57 # The ratio of writes to the total reads of a transaction
58 org.vanilladb.bench.benchmarks.micro.rte.MicrobenchmarkParamGen.WRITE_RATIO_IN_RW_TX=0.5
59 # The conflict rate of a hot record
60 org.vanilladb.bench.benchmarks.micro.rte.MicrobenchmarkParamGen.HOT_CONFLICT_RATE=0.002
```

- 1. Result (BufferPoolSize=102400)
  - a. HOT CONFLICT RATE=0.08:
  - d. Speedup=265853(commits)/148399(commits)=1.791(79.1%)
  - b. HOT\_CONFLICT\_RATE=0.002: Speedup=312686(commits)/171295(commits)=1.825(82.5%)

### Throughput







#### Comparison of NUM WAREHOUSES(TPCC)

```
18 # Basic Parameters
19#
20
21 # The running time for warming up before benchmarking
22 org.vanilladb.bench.BenchmarkerParameters.WARM_UP_INTERVAL=60000
23 # The running time for benchmarking
24 org.vanilladb.bench.BenchmarkerParameters.BENCHMARK_INTERVAL=60000
25 # The number of remote terminal executors for benchmarking
26 org.vanilladb.bench.BenchmarkerParameters.NUM_RTES=2
27 # The IP of the target database server
28 org.vanilladb.bench.BenchmarkerParameters.SERVER_IP=127.0.0.1
29 # 1 = JDBC, 2 = Stored Procedures
30 org.vanilladb.bench.BenchmarkerParameters.CONNECTION_MODE=2
31 # 1 = <u>Micro</u>, 2 = TPC-C
32 org.vanilladb.bench.BenchmarkerParameters.BENCH_TYPE=2
33# Whether it enables the built-in profiler on the server
34 org.vanilladb.bench.BenchmarkerParameters.PR0FILING_ON_SERVER=<mark>false</mark>
35 # The path to the generated reports
36 org.vanilladb.bench.StatisticMgr.OUTPUT_DIR=/Users/<u>cengjingyu/</u>Desktop/experiments/TPCC
37 # The granularity for summarizing the performance of benchmarking
38 org.vanilladb.bench.StatisticMgr.GRANULARITY=1000
39 # Whether the RTEs display the results of each transaction
40 org.vanilladb.bench.rte.TransactionExecutor.DISPLAY_RESULT=false
4 # TPC-C Parameters
5 #
```

```
# # TPC-C Parameters
# # TPC-C Parameters
# # The number of warehouses in the testing data set
# org.vanilladb.bench.benchmarks.tpcc.TpccConstants.NUM_WAREHOUSES=5
# The total number of frequency
# org.vanilladb.bench.benchmarks.tpcc.TpccConstants.FREQUENCY_TOTAL=100
# The frequency of new-order transactions
# org.vanilladb.bench.benchmarks.tpcc.TpccConstants.FREQUENCY_NEW_ORDER=100
# The frequency of payment transactions
# org.vanilladb.bench.benchmarks.tpcc.TpccConstants.FREQUENCY_PAYMENT=0
# The frequency of order-status transactions
# XXX: Not implemented
# The frequency of delivery transactions
# XXX: Not implemented
# org.vanilladb.bench.benchmarks.tpcc.TpccConstants.FREQUENCY_ORDER_STATUS=0
# The frequency of delivery transactions
# XXX: Not implemented
# org.vanilladb.bench.benchmarks.tpcc.TpccConstants.FREQUENCY_DELIVERY=0
# The frequency of stock-level transactions
# XXX: Not implemented
# org.vanilladb.bench.benchmarks.tpcc.TpccConstants.FREQUENCY_STOCK_LEVEL=0
# Whether it enables the thinking and keying time defined in TPC-C specification
# org.vanilladb.bench.benchmarks.tpcc.rte.TpccTxExecutor.ENABLE_THINK_AND_KEYING_TIME=false
```

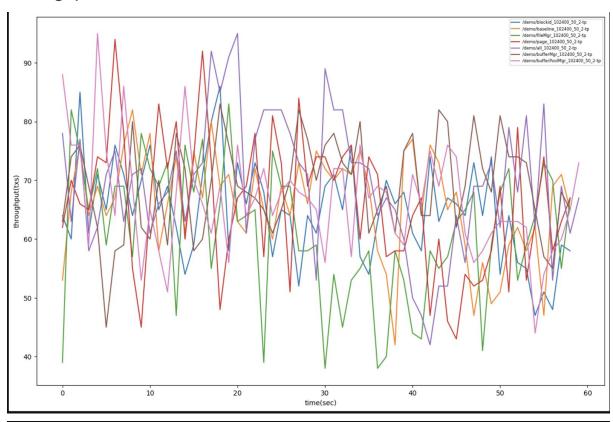
- Result (BufferPoolSize=512)
  - a. NUM WAREHOUSES=5:

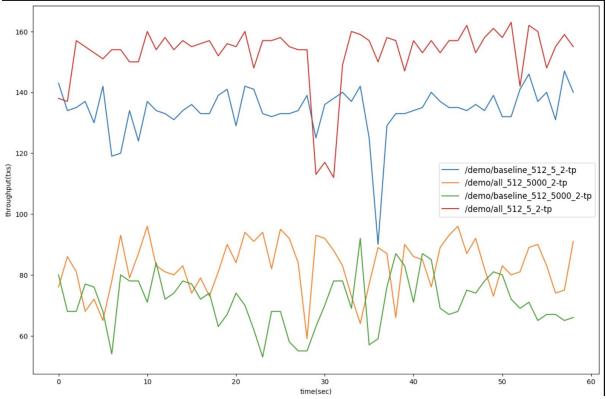
Speedup=9124(commits)/8095(commits)=1.127(12.7%)

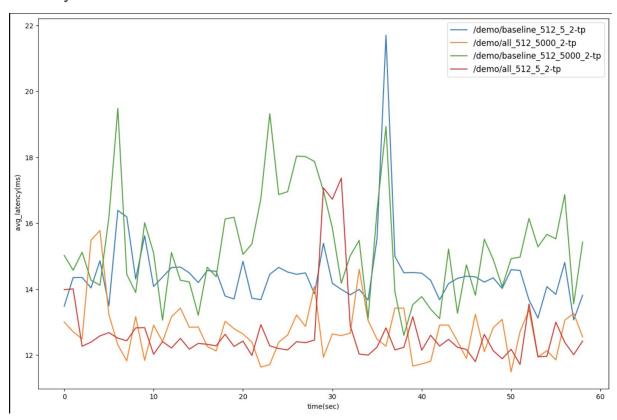
b. NUM WAREHOUSES=5000:

Speedup=4981(commits)/4272(commits)=1.165(16.5%)

# Throughput







# **Analysis**

#### **Smaller Critical Section**

Using smaller critical section can improve the performance of DBMS. We try to reduce the regions of critical sections. After tracing codes, we found that there are several critical sections in files.(BufferMgr, BufferPoolMgr, Page, FileMgr). To being with, in order to realize multi-threading, we can rescind different levels' "synchronize" parameters. It mainly can be divided into three parts. Buffer level, Page level and FileMgr level. Initially, all levels' functions are with "synchronize" parameters. We try to reduce the size of critical sections, the performance became better, due to maintaining data consistency, we can modify many functions. Then, we try to decrease the size of Page, FileMgr levels, since we maintains data consistency of threads in buffer level, we can repeal most parts of each level. Finally the performance becomes much better.

#### Never Do It Again

When using pinningbuffer hashmap, hashcode() method will be called to calculate the hash value of the blockID. In origin version, the value has to be calculated whenever it is called. We add a variable to store the calculated value when the method is first called. So afterward, it can instantly return the stored value.

#### Size of BufferPool

In our experiments, we choose two buffer pool size to implement, which is 1024 and 102400. According to the result, when using larger buffer pool size, our optimization has more significant improvement on performance. We guess it is because our optimization is done by decreasing critical section and reducing duplicated calculation, but not replacement strategy. So the improvement on performance when using smaller buffer pool will be constrained by page swapping.

#### Num of RTEs

We compare two experiments, which has 10 RTEs, 50 RTEs respectively. Result shows that the optimization is better when having 50 RTEs. We think it is because when more clients are accessing a DBMS concurrently, our optimization on concurrency (smaller critical section) will be more effective.

#### Rate of Conflicts

To find out the influence of conflict rate on optimization performance, we compare two conflict rate, 0.002 and 0.08. The result shows that with conflict rate 0.002, the performance (throughput and latency) and the improvement ratio is a better than with conflict rate 0.08. We believe the reason is that the smaller the conflict rate is, the less possible that clients will access the same block. Accessing the same block would be blocked since all buffer method are synchronized. So smaller conflict rate leads to better performance.

#### TPC-C

When using TPC-C benchmarking, we first want to use the same bufferpool size like micro benchmark. However, since TPC-C database is large, it would run out of memory. So we finally choose 512 as our bufferpool size. Unfortunately, due to the small size of bufferpool, many transaction is aborted by deadlock. In both two experiment with warehouse number 5 and 5000, the performance and improvement are bad compared to micro benchmarking.