

One Loop Does Not Fit All

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Motivation

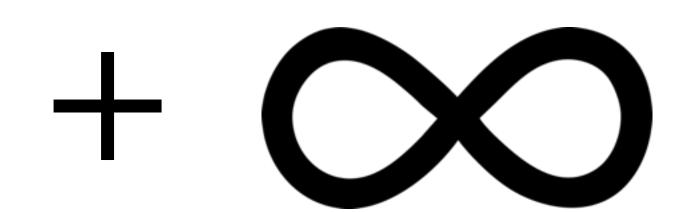


Just In Time

compilation is

a key technology in

column stores



Infinite Possibilities



One solution does not fit all



No clear conclusion given the various setups

Our Approach

Query 1 Select Query SELECT R.d FROM R WHERE R.a < v1 AND R.b < v2 AND R.c < v3 AND R.d < v4

Algorithm 1 4 Loops Strategy for Query 1

- 1. inter = select(R.a, v1)
- 2. inter = select_fetch(R.b, v2, inter)
- 3. inter = select_fetch(R.c, v3, inter) 4. dest = select_fetch(R.d, v4, inter)

Algorithm 2 1 Loop Strategy for Query 1

1. dest = select(R.a, v1, R.b, v2, R.c, v3, R.d, v4)

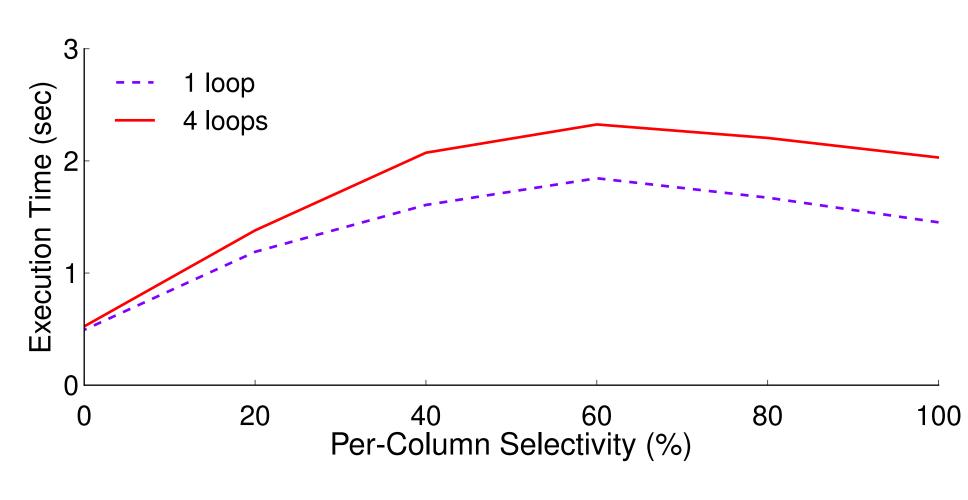
Experimental Setup

- Multi Threaded Execution: 8 available cores
- Evaluation with branching and branchless execution
- Loop unrolling to optimize the branchless version
- Vectorized Processing
- 4-way Intel Xeon E7-4820 configuration with 64 hardware threads
- 1 TB of main memory
- •GCC 4.7.2 (Debian 4.7.2-5).

Intuition

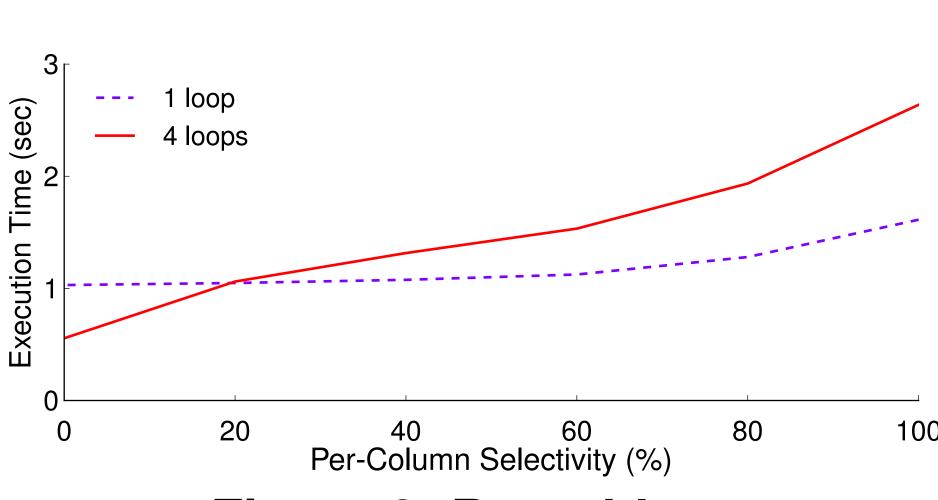
- Using separate loops for each predicate results in scanning fewer cache lines when the first filter selects very few tuples.
- For low selectivities it is the case that in the 1 Loop Strategy we are reading data that is not going to be used.

Results



Motivation: Optimizing for performance? Method: Branchless Execution High Selectivity: 1 Loop Strategy Low Selectivity: 1 Loop Strategy

Figure 1: With Branches



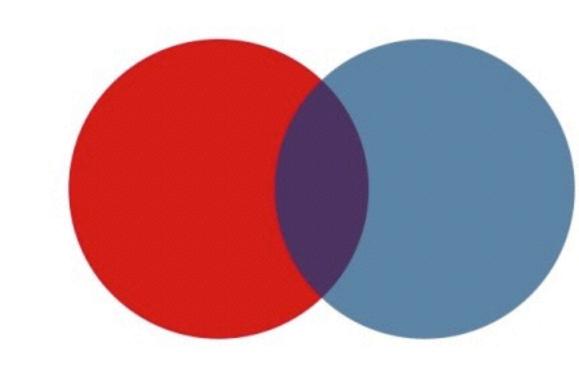
Motivation: Optimizing for robustness? Method: Branchless Execution High Selectivity: 1 Loop Strategy Low Selectivity: 4 Loop Strategy

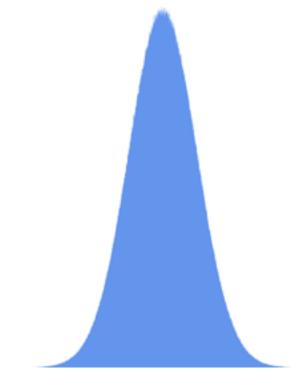
Figure 2: Branchless

Future Work

Disjunctive Selects

Alternate Distributions





Correlations

