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A Mini Project Report on

"Diffie-Hellman key exchange"

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ABSTRACT

Problem Statement:

Write a Java or C program to implement Diffie-Hellman key exchange Algorithm.

Description:

Diffie-Hellman is a way of generating a shared secret between two people in such a way that the secret can't be seen by observing the communication. That's an important distinction: You're not sharing information during the key exchange, you're creating a key together.

ACKNOWLEDGEMENT

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INTRODUCTION

1.1 INTRODUCTION TO THE PROJECT

Diffie—Hellman Key Exchange establishes a shared secret between two parties that can be used for secret communication for exchanging data over a public network. The following conceptual diagram illustrates the general idea of the key exchange by using colors instead of very large numbers.

The process begins by having the two parties, Alice and Bob, agree on an arbitrary starting color that does not need to be kept secret (but should be different every time[8]); in this example the color is yellow. Each of them selects a secret color that they keep to themselves. In this case, orange and blue-green. The crucial part of the process is that Alice and Bob now mix their secret color together with their mutually shared color, resulting in orange-tan and light-blue mixtures respectively, then publicly exchange the two mixed colors. Finally, each of the two mix together the color they received from the partner with their own private color. The result is a final color mixture yellow-brown that is identical to the partner's color mixture.

If a third party listened to the exchange, it would be computationally difficult for them to determine the secret colors. In fact, when using large numbers rather than colors, this action is computationally expensive for modern supercomputers to do in a reasonable amount of time.



1.2 PURPOSE OF THE PROJECT

The Diffie-Hellman protocol is a method for two computer users to generate a shared private key with which they can then exchange information across an insecure channel. Let the users be named Alice and Bob. First, they agree on two prime numbers g and p, where p is large (typically at least 512 bits) and g is a primitive root modulo p. (In practice, it is a good idea to choose p such that (p-1)/2 is also prime.) The numbers g and p need not be kept secret from other users. Now Alice chooses a large random number a as her private key and Bob similarly chooses a large number b. Alice then computes A=g^a (mod p), which she sends to Bob, and Bob computes B=g^b (mod p), which he sends to Alice.

Now both Alice and Bob compute their shared key K=g^(ab) (mod p), which Alice computes as

 $K=B^a \pmod{p}=(g^b)^a \pmod{p}$ and Bob computes as

 $K=A^b \pmod{p}=(g^a)^b \pmod{p}$.

Alice and Bob can now use their shared key K to exchange information without worrying about other users obtaining this information. In order for a potential eavesdropper (Eve) to do so, she would first need to obtain $K=g^{(ab)}$ (mod p) knowing only g, p, $A=g^{(ab)}$ (mod p) and $B=g^{(b)}$ (mod p).

This can be done by computing a from A=g^a (mod p) or b from B=g^b (mod p). This is the discrete logarithm



SYSTEM REQUIREMENTS

2.1 SOFTWARE REQUIREMENTS

- 1. Text editor
- 2. Python and JSON

2.2 HARDWARE REQUIREMENTS

1. RAM: 512 Mb

2. HARDDISK SPACE: 5GB

3. SYSTEM TYPE: 32-BIT OR 64-BIT PROCESSOR



DESIGN

DIFFIE HELLMEN ALGORITHM:

Diffie-Hellman key exchange, also called exponential key exchange, is a method of digital encryption that uses numbers raised to specific powers to produce decryption keys on the basis of components that are never directly transmitted, making the task of a would-be code breaker mathematically overwhelming.

To implement Diffie-Hellman, the two end users Alice and Bob, while communicating over a channel they know to be private, mutually agree on positive whole numbers p and q, such that p is a prime number and q is a generator of p. The generator q is a number that, when raised to positive whole-number powers less than p, never produces the same result for any two such whole numbers. The value of p may be large but the value of q is usually small.

Language Used: Python



IMPLEMENTATION

Alice Side code:

```
prime = 1
while True:
       prime = random.randint(100,1000)
       if isPrime(prime) == True:
              break
       else:
              prime = prime
while True:
       generator = random.randint(2,prime-2)
       if isPrime(generator) == True:
              break
       else:
              generator = generator
print "prime =",prime
                      #prime
print "generator =",generator #generator
a = random.randint(1,35535)
tcp = socket(AF_INET,SOCK_STREAM)
tcp.setsockopt(SOL_SOCKET,SO_REUSEADDR,1)
tcp.bind(('localhost',2222))
tcp.listen(5)
TEMP = \{\}
FOO = "
BAR = "
TEMP = {'p':prime,'base':generator}
conn, addr = tcp.accept()
while 1:
```



```
msg = conn.recv(1024)
  if "NEGOCIATION" in msg:
    conn.send(json.dumps(TEMP))
    print "Sending base and prime number..."
    break
A = (generator ** a)% prime
print "the shared secret of alice is: %d " % (A)
msg = conn.recv(1024)
FOO = json.loads(msg)
B = FOO['B']
s = (B ** a)\% prime
TEMP = \{'A':A\}
conn.send(json.dumps(TEMP))
print("wait for calculations to be completed !!!!<<<<<>>>>")
time.sleep(3)
print "The secret is %d " % (s)
msg = conn.recv(1024)
BAR = json.loads(msg)
secret2= BAR['S']
if s == secret2:
       print "keys are verified successfully!"
```



Bob Side code:

```
a = random.randint(1,35535)
p = TEMP['p']
g = TEMP['base']
B = (int(g)**a)\%int(p)
print "The shared secret of Bob is %d: " % (B)
DICT = \{'B':B\}
tcp.send(json.dumps(DICT))
msg = tcp.recv(1024)
TEMP = json.loads(msg)
A = TEMP['A']
s = (A^{**}a)\%p
print("wait for calculations to be completed !!!!<<<<<>>>>")
time.sleep(3)
print "The secret key is %d " % (s)
SECT = \{'S':s\}
tcp.send(json.dumps(SECT))
```



DISCUSSIONS AND SCREENSHOTS

5.1 Demo of differ-hellman

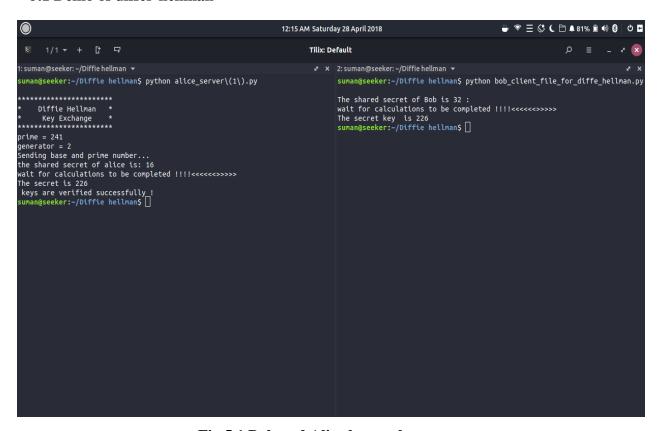


Fig 5.1 Bob and Alice key exchange



CONCLUSION

Ideally, Diffie-Hellman should be used in conjunction with a recognized authentication method such as digital signatures to verify the identities of the users over the public communications medium. Diffie-Hellman is well suited for use in data communication but is less often used for data stored or archived over long periods of time.



FUTURE SCOPE

Encryption

Public key encryption schemes based on the Diffie–Hellman key exchange have been proposed. The first such scheme is the ElGamal encryption. A more modern variant is the Integrated Encryption Scheme.

Forward secrecy

Protocols that achieve forward secrecy generate new key pairs for each session and discard them at the end of the session. The Diffie-Hellman key exchange is a frequent choice for such protocols, because of its fast key generation.

Password-authenticated key agreement

When Alice and Bob share a password, they may use a password-authenticated key agreement (PK) form of Diffie–Hellman to prevent man-in-the-middle attacks. One simple scheme is to compare the hash of s concatenated with the password calculated independently on both ends of channel. A feature of these schemes is that an attacker can only test one specific password on each iteration with the other party, and so the system provides good security with relatively weak passwords. This approach is described in ITU-T Recommendation X.1035, which is used by the G.hn home networking standard.



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