```
l, r = 0, len(nums) ----> This is a "left-closed, right-open" interval: [l, r) l, r = 0, len(nums) - 1 ----> This is a "fully closed" interval: [l, r]
```

This means you'll need to **update the loop condition and possibly other parts** of the code.

Changes Required for [l, r] Interval:

```
class Solution:
    def searchInsert(self, nums: List[int], target: int) -> int:
        l, r = 0, len(nums) - 1

    while l <= r:
        mid = l + (r - l) // 2

    if nums[mid] == target:
        return mid
    elif target < nums[mid]:
        r = mid - 1
    else:
        l = mid + 1

return l</pre>
```

***** Key Differences:

```
[l, r] version (r = len(nums) -
                  [l, r) version (r = len(nums))
     Aspect
                                                                       1)
                while l < r
                                                     while l <= r
Loop Condition
Shrink Right
                 r = mid
                                                     r = mid - 1
Bound
                 Works for lower-bound type insert
Purpose
                                                     Works directly with index insert logic
                 position
When Loop Ends l == r is exit point
                                                     1 > r is exit point
```

Final Output Logic:

In both cases, returning 1 gives you the correct **insertion index** — the position where target should go to maintain sorted order.

Would you like a visual step-through example comparing both versions?

You said:

what is the condition for "left closed, right open" and "fully closed"

ChatGPT said:

Great! Let's clearly break down what "left-closed, right-open" ([l, r)) and "fully closed" ([l, r]) intervals mean — especially in binary search logic.

Interval Terminology:

```
1. Left-Closed, Right-Open — [l, r)
```

- Includes the **left index** (1)
- Excludes the right index (r)
- Meaning: you consider all indices i such that l <= i < r

✓ Binary Search Logic for [1, r):

```
l = 0
r = len(nums)  # `r` is one past the last index
while l < r:  # because r is exclusive
    mid = l + (r - l) // 2
    if nums[mid] == target:
        return mid
    elif target < nums[mid]:
        r = mid  # exclude mid
    else:
        l = mid + 1  # exclude mid
return l</pre>
```

2. Fully Closed — [l, r]

- Includes both 1 and r
- Meaning: you consider all indices i such that l <= i <= r

▶ Binary Search Logic for [1, r]

```
l = 0
r = len(nums) - 1  # because r is inclusive now

while l <= r:  # must include r
    mid = l + (r - l) // 2
    if nums[mid] == target:
        return mid
    elif target < nums[mid]:
        r = mid - 1  # exclude mid
    else:
        l = mid + 1  # exclude mid

return l</pre>
```

Summary Table:

Feature [l, r) (Left-closed, right-open) [l, r] (Fully closed) len(nums) len(nums) - 1Initial r l < rl <= r Loop condition r = mid - 1When shrinking r r = midWhen shrinking l l = mid + 1l = mid + 1Loop ends when l == rl > rLower bound search (insertion) Typical use Exact match or insert index

Rule of Thumb:

• Use [1, r) if you prefer math-style "range-end exclusive" logic (like Python's range).

• Use [1, r] if you prefer more **intuitive range** (includes both ends), but must be careful with loop boundaries.

```
☑ Q.35: Search Insert Position

Return the index of target if found, otherwise the index where it should be inserted to maintain order.
Find the first position i where nums[i] >= target
Always return a valid index, whether or not the target is found.
Left-closed, Right-open [left, right)
Works perfectly for lower bound problems
def searchInsert(nums, target):
  left, right = 0, len(nums)
  while left < right:
     mid = (left + right) // 2
     if nums[mid] < target:
       left = mid + 1
       right = mid
  return left

☑ Q.704: Binary Search

Return the index of the target if found, else return -1.
Find the exact position of target
Return -1 if not found
Fully closed [left, right]
Includes the last index, safely checks exact matches
def search(nums, target):
  left, right = 0, len(nums) - 1
  while left <= right:
     mid = (left + right) // 2 if nums[mid] == target:
       return mid
     elif nums[mid] < target:
       left = mid + 1
     else:
       right = mid - 1
  return -1
Summary Table
Feature / Question
                            Q.35: Insert Position
                                                                     Q.704: Binary Search
Target found?
                            Return index
                                                                     Return index
Target not found?
                            Return insert position
                                                                     Return -1
                            Find first i such that nums[i] ≥ target
Goăl
                                                                     Find exact target
Return Range
                            Always [0, len(nums)]
                                                                     -1 if not found
Preferred Interval
                            [left, right)
                                                                     [left, right]
```

Loop Condition

Suitable for Lower Bound

left < right

✓ Yes

left <= right

× No

Rule of Thumb: Task Type Insert Position / Lower Bound Exact Match Search

Use Interval [left, right) [left, right]