# Fall 2024 B561 Assignment 1 Relational Databases, Expressing Constraints and Queries in SQL, and in Tuple Relational Calculus (TRC)

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## 1 Introduction

The goals for this assignment are to

- 1. become familiar with the PostgreSQL system
- 2. create a relational database and populate it with data;
- 3. examine the side-effects on the state of a database by inserts and deletes in the presence or absence of primary and foreign key constraints;
- 4. formulate some queries in SQL and evaluate them in PostgreSQL; and
- 5. translate TRC queries to SQL and formulate queries and constraints in TRC

To turn in your assignment, you will need to upload to Canvas a file with name assignment1.sql which contains the necessary SQL statements that solve the graded problems in this assignment.

The assignment1.sql file must be such that the AI's can run it in their Post-greSQL environment. In addition, you will need to upload a separate assignment1.txt file that contains the results of running your queries. We have posted the exact requirements and an example for uploading your solution files. (See the module Instructions for turning in assignments.) Finally, you will need to upload an assignment1.pdf file that contains the solutions for problems related to TRC.

#### Note:

1. To solve this assignment, you will need to download and install PostgreSQL (version 12 or higher) on your computer.

- 2. To solve problems related to TRC, follow the syntax and semantics described in the lecture slides Tuple Relational Calculus and SQL.
- 3. For a good way to learn about Latex, look at <a href="https://www.overleaf.com/learn/latex/Free\_online\_introduction\_to\_LaTeX\_(part\_1)">https://www.overleaf.com/learn/latex/Free\_online\_introduction\_to\_LaTeX\_(part\_1)</a>.

For the problems in this assignment we will use the following database schema:  $^{1}$ 

```
Student(sid, sname, homeCity)
Department(deptName, mainOffice)
Major(major)
employedBy(sid, deptName, salary)
departmentLocation(deptName, building)
studentMajor(sid, major)
hasFriend(sid1, sid2)
```

In this database we maintain a set of students (Student), a set of departments (Department), and a set of (major) majors (Major). The sname attribute in Student is the name of the student. The homeCity attribute in Student specifies the home city of the student. The deptName attribute in Department is the name of the department. The mainOffice attribute in Department is the name of the building where the department has its main office. The major attribute in Major is the name of a (major) major.

A student can be employed by at most one department. This information is maintained in the employedBy relation. (We permit that a student is not employed by any department and that a department may not employ any students.) The salary attribute in employedBy specifies the salary made by the student.

The building attribute in departmentLocation indicates a building in which the department is located. (A department may be located in multiple buildings.)

A student can have multiple majors. This information is maintained in the studentMajor relation. A major can be the major of multiple students. (A student may not have any majors, and a major may have no students with that major.)

A pair  $(s_1, s_2)$  in hasFriend indicates that student  $s_1$  considers student  $s_2$  as a friend. It is possible that  $s_1$  considers  $s_2$  as a friend, but not necessarily the other way around. In other words, the hasFriend relation can not be assumed to be symmetric.

The domain for the attributes sid, salary, sid1, and sid2 is integer. The domain for all other attributes is text.

We assume the following foreign key constraints:

<sup>&</sup>lt;sup>1</sup>The primary key, which may consist of one or more attributes, of each of these relations is underlined.

- sid is a foreign key in employedBy referencing the primary key sid in Student;
- deptName is a foreign key in employedBy referencing the primary key deptName in Department;
- deptName is a foreign key in departmentLocation referencing the primary key deptName in Department;
- sid is a foreign key in studentMajor referencing the primary key sid in Student;
- major is a foreign key in studentMajor referencing the primary key major in Major;
- sid1 is a foreign key in hasFriend referencing the primary key sid in Student; and
- sid2 is a foreign key in hasFriend referencing the primary key sid in Student;

The data for the assignment is included in the schema.sql.

**Remark 1** We will typically use the primary key of an object to specify that object. This should not cause any confusion since an object can be referenced uniquely with its primary key value. E.g., we will write '... student s...' rather than '... the student with sid s...'. When posing a query, we may write 'Find each student who ...' instead of 'Find the sid of each student who ...'. The expected answer for such a query should be the set of sids of students who meet the criteria of the query.

# 2 Database creation and impact of constraints on insert and delete statements.

Create a database in PostgreSQL that stores the data provided in the schema.sql file. Make sure to specify primary and foreign keys.

1. Provide 4 conceptually different examples that illustrate how the presence or absence of primary and foreign keys affect inserts or deletes in these relations. To solve this problem, you will need to experiment with the relation schemas and instances for this assignment. For example, you should consider altering primary keys and foreign key constraints and then consider various sequences of insert and delete operations. You may need to change some of the relation instances to observe the desired effects. Certain inserts and deletes should succeed but others should generate error conditions. (Consider the lecture notes about keys, foreign keys, and inserts and deletes as a guide to solve this problem.) [5 Points]

## 3 Expressing queries in SQL

For this assignment, you are required to use tuple variables in your SQL statements. For example, in formulating the query "Find the sid and sname of each student who lives in Bloomington" you should write the query

```
SELECT s.sid, s.sname
FROM Student s
WHERE s.homeCity = 'Bloomington'
```

rather than

```
SELECT sid, sname
FROM Student
WHERE homeCity = 'Bloomington'
```

Write SQL statements for the following queries. Make sure that each of your queries returns a set but not a bag. In other words, make appropriate use of the DISTINCT clause where necessary.

You can **not** use the SQL JOIN operations or SQL aggregate functions such as COUNT, SUM, MAX, MIN, etc in your solutions.

- 2. For each department, list its name along with the highest salary made by students who work for it. [5 Points]
- 3. Find the sid and sname of each student s who is employed by a department d and who has a salary that is strictly higher than the salary of each of his or her friends who is employed by that department d. (Student should only be considered if indeed he or she has a friend who is employed by department d.) [5 Points]
- 4. Find the deptName of each department that not only employs students whose home city is Indianapolis. (In other words, there exists at least one student who is employed by such a department whose home city is not Indianapolis.) [5 Points]
- 5. Find the sid, sname of each student who (a) has home city Bloomington, (b) works for a department where he or she earns a salary that is higher than 20000, and (c) has at least one friend. [5 Points]
- 6. Find the pairs  $(d_1, d_2)$  of names of different departments whose main offices are located in the same building. [5 Points]
- 7. Find the sid and sname of each student whose home city is different than those of his or her friends. [5 Points]
- 8. Find each major that is the major of at most 2 students. [5 Points]
- 9. Find the sid, sname, and salary of each student who has at least two friends such that these friends have a common major but provided that it is not the 'Mathematics' major. [5 Points]

# 4 Translating TRC queries to SQL

Consider the following queries formulated in TRC. Translate each of these queries to an equivalent SQL query. This underscores the close correspondence between TRC and SQL. The SQL queries should be included in the assignment1.sql file.

```
10.
                                                              \{(s_1.sid, s_1.sname) \mid Student(s_1) \land \exists d \in Department, \ w \in employedBy (
                                                                                                                  d.deptName = w.deptName \land s_1.sid = w.sid \land d.mainOffice = LuddyHall \land d.deptName = w.deptName \land s_1.sid = w.sid \land d.mainOffice = LuddyHall \land d.deptName \land s_1.sid = w.sid \land d.mainOffice = LuddyHall \land d.deptName \land s_1.sid = w.sid \land d.mainOffice = LuddyHall \land d.deptName \land s_1.sid = w.sid \land d.mainOffice = LuddyHall \land d.deptName \land s_1.sid = w.sid \land s_1.sid = w.sid \land s_1.sid = w.sid \land s_2.sid = w.sid \land s_3.sid = w.sid \land s_4.sid = w.
                                                                                                                  \exists s_2 \in Student \ (hasFriend(s_1.sid, s_2.sid) \land s_2.homeCity \neq Bloomington)) \}.
                                                          [5 Points]
11.
                                                              \{s_1.sid \mid Student(s_1) \land \forall s_2 \in Student(hasFriend(s_1.sid, s_2.sid) \rightarrow \}
                                                                                                                  \exists sm_1 \in studentMajor, sm_2 \in studentMajor(sm_1.sid = s_1.sid \land sm_2.sid = s_2.sid \land
                                                                                                                  sm_1.major = sm_2.major \land sm_1.sid \neq sm_2.sid).
                                                          [5 Points]
12.
                                                              \{(s_1.sid, s_2.sid) \mid Student(s_1) \land Student(s_2) \land s_1.sid \leq s_2.sid \land s_1.sid \leq s_2.sid \land s_
                                                                                                                  \forall f_1 \in hasFriend(s_1.sid = f_1.sid_1 \rightarrow \exists f_2 \in hasFriend(f_2.sid_1 = s_2.sid \land f_1.sid_2 = f_2.sid_2))\}.
                                                   [5 Points]
  13.
                                                              \{m.major \mid Major(m) \land \neg (\exists s \in Student \exists sm \in student Major(s.sid = sm.sid \land substitution = sm.sid \land
                                                                                                                  sm.major = m.major \land s.homeCity = 'Bloomington').
                                                          [5 Points]
```

# 5 Formulating Queries in the Tuple Relational Calculus

The TRC solutions of these problems should be included in the assignment1.pdf file. You need not write SQL queries for this section.

14.

(a) Find each pair (d, m) where d is the name of a department and m is a major of a student who is employed by that department and who earns a salary of at least 20000. [2.5 Points]

#### Answer:

```
 \{ (d.deptName,\ m.major) \mid Department(d) \land Major(m) \land \exists s \in Student(studentMajor(m.major, s.sid) \land \exists w \in employedBy(s.sid = w.sid \land d.deptName = w.deptName \land w.salary \geq 20000)) \}
```

(b) Find each pair  $(s_1, s_2)$  of sids of different students who have the same (set of) friends who work for the CS department. [2.5 Points]

#### Answer:

```
 \{(s1.sid, s2.sid) \mid Student(s1) \land Student(s2) \land s1.sid \neq s2.sid \land \forall w \in employedBy(w.deptName = CS \rightarrow ((hasFriend(s1.sid, w.sid) \rightarrow hasFriend(s2.sid, w.sid)) \land hasFriend(s2.sid, w.sid) \rightarrow hasFriend(s1.sid, w.sid))))\}
```

15.

(a) Find each major for which there exists a student with that major and who does not only have friends who also have that major. [2.5 Points]

#### Answer:

```
 \{m.major \mid Major(m) \land \exists s \in Student(studentMajor(m.major, s.sid) \land \neg (\forall s1 \in Student(hasFriend(s1.sid, s.sid) \rightarrow studentMajor(m.major, s1.sid)))\}
```

(b) Find the sid and sname of each student whose home city is different than those of his or her friends. [2.5 Points]

#### Answer:

```
\{(s.sid, s.sname) \mid Student(s) \land \neg \exists f(hasFriend(s.sid, f.sid2) \land \exists s2(Student(s2) \land s2.sid = f.sid2 \land s.homeCity = s2.homeCity))\}
```

16.

(a) Find the sid, sname, and salary of each student who has at least two friends such that these friends have a common major but provided that it is not the 'Mathematics' major. [2.5 Points]

#### Answer:

```
 \{s.sid, s.sname, e.salary \mid Student(s) \land \exists e \in employedBy(e.sid = s.sid) \land \exists f1 \exists f2(hasFriend(f1) \land hasFriend(f2) \land s.sid = f1.sid1 \land s.sid = f2.sid1 \land f1.sid2 \neq f2.sid2 \land \exists sm1 \exists sm2(studentMajor(sm1) \land studentMajor(sm2) \land sm1.sid = f1.mid \land sm2.sid = f2.mid \land sm1.major = sm2.major \land sm1.major \neq' Mathematics')) \}
```

(b) For each department, list its name along with the highest salary made by students who are employed by it. [2.5 Points]

#### Answers

```
 \{ d.dname, w.salary \mid Department(d) \land employed By(w) \land d.dname = w.dname \land \neg \exists w1 (employed By(w1) \land d.dname = w1.dname \land w.salary < w1.salary) \}
```

17.

Find the sid, sname of each student who (a) has home city Bloomington, (b) works for a department where he or she earns a salary that is higher than 20000, and (c) has at least one friend. [5 Points]

#### Answer:

```
\{ s.sid, s.sname \mid Student(s) \land s.city =' Bloomington' \land \exists w (employedBy(w) \land w.sid = s.sid \land w.salary > 20000) \land \exists f (hasFriend(f) \land s.sid = f.sid1) \}
```

# 6 Formulating constraints in the safe Tuple Relational Calculus and as boolean SQL queries

Formulate the following constraints in safe TRC and as boolean SQL queries.

The TRC solutions of these problems should be included in the assignment1.pdf file and the SQL solutions should be included in the assignment1-SQL.sql file.

Here is an example of what is expected for your answers.

**Example 1** Consider the constraint "Each major is the major of a student." In safe TRC, this constraint can be formulated as follows:

```
\forall m \ Major(n) \rightarrow \exists sm \ (studentMajor(sm) \land sm.major = m.major)
```

Alternatively, the constraint can be formulated in safe TRC as

```
\neg \exists m(Major(m) \land \neg \exists sm(studentMajor(sm) \land sm.major = m.major)).
```

This constraint can be specified using the following boolean SQL query:

- 18. Consider the constraint "Some major has fewer than 2 students with that major."
  - (a) Formulate this constraint in safe TRC. [2.5 points]

#### Answer:

```
\{\exists mMajor \land \neg \exists s1 \in Student \exists s2 \in Student(s1.sid \neq s2.sid \land studentMajor(s1.sid, m.major) \land studentMajor(s2.sid, m.major))\}
```

- (b) Formulate this constraint as a boolean SQL query. [2.5 points]
- 19. Consider the constraint "Each student is employed by a department and has at least two majors."

(a) Formulate this constraint in safe TRC. [2.5 points]

#### Answer:

```
 \{ \forall sStudent(s) \rightarrow (\exists w(employedBy(w) \land s.sid = w.sid) \land \exists sm1 \exists sm2(studentMajor(sm1) \land studentMajor(sm2) \land sm1.sid = s.sid \land sm2.sid = s.sid \land sm1.major \neq sm2.major)) \}
```

- (b) Formulate this constraint as a boolean SQL query. [2.5 points]
- 20. Consider the constraint "Each student and his or her friends work for the same department."
  - (a) Formulate this constraint in safe TRC. [2.5 points]

#### Answer:

```
 \{ \ \forall f \forall w 1 \forall w 2 ((hasFriend(f) \land employedBy(w1) \land employedBy(w2) \land f.sid1 = w1.sid \land f.sid2 = w2.sid) \rightarrow w1.dname = w2.dname) \}
```

(b) Formulate this constraint as a boolean SQL query. [2.5 points]