# Operating Systems - Homework 1

# Github repo

# Subata Naveen Khan – 18119

1. File: subata\_q1.c:

```
#include "kernel/types.h"
#include "kernel/stat.h"
#include "user/user.h"
int main (void) {
   int p_ptc[2], p_ctp[2]; // ptc = parent-to-child, ctp = child-to-parent
   char byte[1];
                            // char datatype has size of 1 byte
                            // no. of times byte gets sent from each side
   int n = 3;
   int i;
   pipe(p_ptc);//child will read from p_ptc[0], parent will write to p_ptc[1]
   pipe(p_ctp);//parent will read from p_ctp[0], child will write to p_ctp[1]
       if (fork() == 0) {
       close (p_ctp[0]);  // won't read from ctp pipe
       for (i = 0; i < n; i++)
           read(p_ptc[0], byte, 1);
           printf("child receives: %c\n", byte[0]);
           byte[0] = 'c';
           printf("child sends: %c\n", byte[0]);
          write(p_ctp[1], byte, 1);
       }
       close(p_ctp[1]);
       close(p_ptc[0]);
       sleep(1); //making sure everything is finished printing before exiting
       exit(0);
   } else {
                             // parent
                           // won't read from ptc pipe
       close(p_ptc[0]);
       close(p_ctp[1]);  // won't write to ctp pipe
       for (i = 0; i < n; i++)
           byte[0] = 'p';
           printf("parent sends: %c\n", byte[0]);
           write(p_ptc[1], byte, 1);  // byte sent to child
```

```
read(p_ctp[0], byte, 1);
    printf("parent receives: %c\n", byte[0]);
}

close(p_ptc[1]);  // closing remaining pipe ends
    close(p_ctp[0]);

wait(0);  // waiting for child
}

exit(0);
}
```

#### Output:

```
parent sends: p
child receives: p
child sends: c
parent receives: c
parent sends: p
child receives: p
child sends: c
parent receives: c
parent sends: p
child receives: p
child receives: c
```

# 2. File: subata\_q2.c:

```
#include "kernel/types.h"
#include "kernel/stat.h"
#include "user/user.h"
#include "kernel/fcntl.h"
int main(int argc, char *argv[]) {
   int p[2];
   if (pipe(p) < 0) {
       printf("pipe creation failed\n");
       exit(0);
   }
   int pid = fork();
   if (pid == 0) {
                                 // child
                                  // won't read from pipe
       close(p[0]);
       countsyscalls(p[1]);
                                 // passes p[1] to kernel to write to
       exec(argv[1], argv + 1); // execute the target program
       printf("exec failed\n");
```

# File: syscall.c:

```
void
syscall(void)
 int num;
 struct proc *p = myproc();
 num = p->trapframe->a7;
 if(num > 0 && num < NELEM(syscalls) && syscalls[num]) {</pre>
// if calling program called countsyscalls
   if(p->syscall_pipe >= 0) {
       struct file *f = p->ofile[p->syscall_pipe];
            uint64 counter_signal = 1;
            filewrite(f, counter_signal, sizeof(counter_signal));
   // Use num to lookup the system call function for num, call it,
   // and store its return value in p->trapframe->a0
   p->trapframe->a0 = syscalls[num]();
 } else {
   printf("%d %s: unknown sys call %d\n",
            p->pid, p->name, num);
   p->trapframe->a0 = -1;
 }
```

#### Output:

```
$ subata_q2 Myprog
Hello World!
number of system calls made: 15
$ subata_q2 subata_q1
parent sends: p
child receives: p
child sends: c
parent receives: c
parent sends: p
child receives: p
child sends: c
parent receives: c
parent sends: p
child receives: p
child sends: c
parent receives: c
number of system calls made: 121
$ subata_q2 wc
this is a test
for question 2
2 7 30
number of system calls made: 13
```

- 3. a. Pipes automatically clean themselves up, as opposed to file redirection, where the shell would need to remove the temorary file that is created.
  - Pipes can handle arbitrarily long streams of data. With file redirection, the disk must have enough free space to store all the data.
  - Pipes allow for parallel execution of the different pipeline stages. With file redirection, the first program needs to finish before the next can start.
  - Pipes provide efficient inter-process communication. The reading and writing processes can communicate directly and without delay. With file redirection, there is no real-time communication so it is much less efficient.
  - b. The main reason this code cannot be implemented in xv6 is that xv6 is a bare bones operating system with limited functionality, so it does not support the operations required.
    - xv6 has an 1s command, but it doesn't have options like -a1
    - xv6 doesn't have a tr command