

Design of FSM to detect the sequence 1011 using Verilog

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Abstract—Tis project goal is to implement and validate the design of a sequence detector in Verilog. It helps us get a better understanding of the FSM level of abstraction

I. INTRODUCTION

A finite state machine is based on a fictitious sequential circuit having one or more states. There can only ever be one active state of this machine. It suggests that the system must switch between states in order to carry out various activities. The device is in charge of recording a situation's state at a specific moment. The adjustments that are put into place immediately affect the states that generate the desired results. In short, Finite State Machines (FSMs) are a functional abstraction for sequential circuits with centralised "states" of operation. Combinatorial logic computes outputs and the following state as a function of inputs and the current state at each clock edge. The block diagram of a finite state machine is shown in the following figure (Fig. 1)

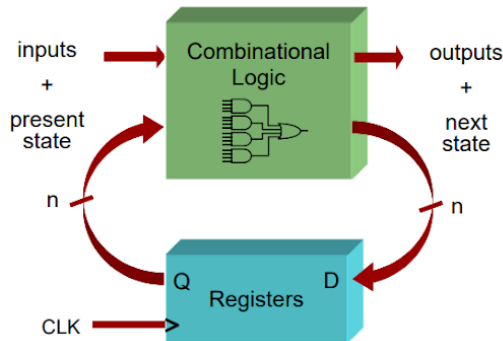


Fig. 1. Block diagram of FSM

A sequence detector, which is a sequential state machine, creates an output 1 whenever the target sequence has been detected after accepting string of bits as input. The current state, as well as any outside input or incoming bit sequences, determine the output. There are two categories of sequence detectors: overlapping and non-overlapping. In contrast to a non-overlapping sequence detector, where the end bit of one sequence does not become the first bit of the following sequence, an overlapping sequence detector assumes the last bit of one sequence as the first bit of the following sequence.

In this experiment, the FSM continually analyses a binary sequence coming from a digital input and only turns on the output when it detects a "1011" sequence.

We compare the output of the FSM level of abstraction, which is here referred to as the DUT, and the output of the structural description or the reference model. A structural model is defined by a system employing basic parts like digital gates and adders. By integrating different modules and gates, structural modelling explains how a hardware module's components are connected and how the hardware is designed.

II. FSM LEVEL OF ABSTRACTION

A. Implementation details

Fig. 2 shows the architectural design of the FSM, which consists of a state register for storing the current state and a combinational network for creating the output and the subsequent state.

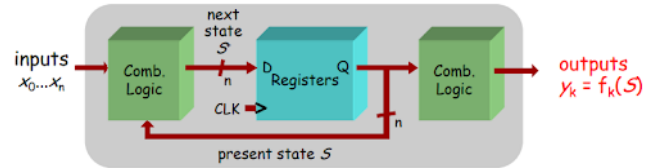


Fig. 2. Architectural layout of FSM

Fig. 3 shows the state transition diagram of the built "1011" sequence detector. It is considered that the sequence detector in question is non-overlapping in nature. Two three-bit registers, one holding the current state and the other the next state, make up the design. Any one of the five states S0 (sequence: 0), S1 (sequence: 1), S2 (sequence: 10), S3 (sequence: 101), and S4 can be applied to a register (sequence: 1011). The output variable changes to logic 1 when the process reaches the final state S4.

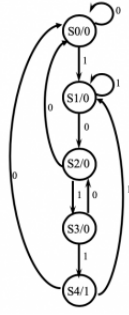


Fig. 3. State transition diagram of 1011 sequence detector

State transition takes place in the always block because at the clock's positive edge, the register holding the current state is changed. In the event of a reset, the register is assigned to the initial state S0; otherwise, it is allocated to the following state, which is determined based on the value of the input sequence bit.

The RTL code for system design at the FSM level of abstraction is shown in Fig. 4. Fig. 5 shows the testbench codes created for verifying the DUT.

```

1 // Code your design here
2 module FSM_1011_det (input clk, x, rst, output detect);
3
4     parameter S0 = 3'b000;
5     parameter S1 = 3'b001;
6     parameter S10 = 3'b010;
7     parameter S101 = 3'b011;
8     parameter S1011 = 3'b100;
9
10    reg [2:0] curr_state, next_state;
11    assign detect = curr_state == S1011 ? 1 : 0;
12    always @ (posedge clk)
13    begin
14        if (rst)
15            curr_state <= S0;
16        else
17            curr_state <= next_state;
18        end
19    always @ (*)
20    begin
21        case (curr_state)
22
23            S0: begin
24                next_state = x ? S1 : S0;
25            end
26
27            S1: begin
28                next_state = x ? S10 : S1;
29            end
30
31            S10: begin
32                next_state = x ? S101 : S0;
33            end
34
35            S101: begin
36                next_state = x ? S1011 : S10;
37            end
38
39            S1011: begin
40                next_state = x ? S1 : S0;
41            end
42
43        endcase
44    end
45 endmodule
46

```

Fig. 4. RTL code

```

1 // Code your testbench here
2 // or browse Examples
3 module FSM_1011_tb;
4     reg clk, x, rst;
5     wire detect;
6
7     FSM_1011_det u0 ( .clk(clk), .rst(rst), .x(x), .detect(detect) );
8
9     initial begin
10        forever #5 clk = ~clk;
11    end
12
13    initial begin
14        $monitor ("At time = %t, Out = %d", $time, detect);
15        clk = 0;
16        rst = 1;
17        x = 0;
18
19        #20 rst = 0;
20
21        #10 x = 1;
22        #10 x = 0;
23        #10 x = 1;
24        #10 x = 1;
25        #10 x = 0;
26        #10 x = 1;
27        #10 x = 0;
28
29        #20 $finish;
30    end
31
32    initial begin
33        $dumpfile("dump.vcd");
34        $dumpvars(0);
35    end
36
37 endmodule

```

Fig. 5. Testbench of DUT

B. Experimental results

Using Cadence software, Verilog architecture of the FSM level of abstraction with a clock, reset, and two registers is simulated and synthesised. Figs. 6 and 7 depict the synthesised circuit and observed waveform, respectively.

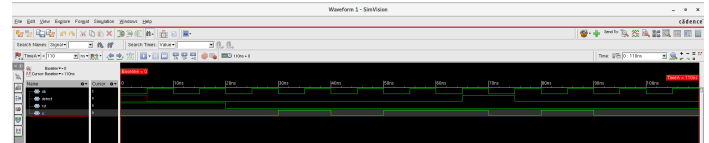


Fig. 6. Output Waveform

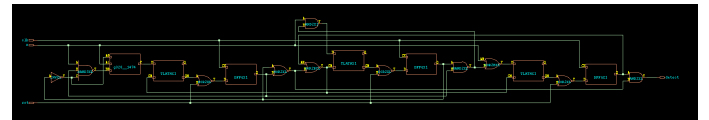


Fig. 7. Synthesized Circuit Diagram

III. STRUCTURAL DESCRIPTION

A. Implementation details:

The state transition diagram in Fig. 3 is used to create the state table of the sequence detector, depicted in Fig. 8.

| Present states | | | Input x | Next states | | | Output y |
|----------------|---|---|------------|----------------|----------------|----------------|-------------|
| A | B | C | | A ⁺ | B ⁺ | C ⁺ | |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 0 | 1 | 0 | 0 | 1 | 0 |
| 0 | 0 | 1 | 0 | 0 | 1 | 0 | 0 |
| 0 | 0 | 1 | 1 | 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 1 | 0 | 1 | 0 | 1 | 1 | 0 |
| 0 | 1 | 1 | 0 | 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | 1 | 1 | 0 | 0 | 0 |
| 1 | 0 | 0 | 0 | 0 | 0 | 0 | 1 |
| 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 |
| 1 | 0 | 1 | 0 | X | X | X | X |
| 1 | 0 | 1 | 1 | X | X | X | X |
| 1 | 1 | 0 | 0 | X | X | X | X |
| 1 | 1 | 0 | 1 | X | X | X | X |
| 1 | 1 | 1 | 0 | X | X | X | X |
| 1 | 1 | 1 | 1 | X | X | X | X |

Fig. 8. State Table

The three-bit state parameter of the FSM level of abstraction has distinct registers assigned to each of its three bits. Three one-bit registers are needed to handle this, and an additional one-bit register is needed for the output. The state table, which is applicable to the development of the combinational logic, serves as the source of the Karnaugh Maps for each of these registers.

Figures 9, 10, 11, and 12 show the Karnaugh Maps for the updated values of the state parameter's first bit (A), second bit (B), last bit (C), and output, which would be assigned to the registers, accordingly.

| DA | | C,x | | | | | |
|-----|----|-----|---|----|----|----|----|
| | | | | 00 | 01 | 11 | 10 |
| A,B | 00 | 0 | 0 | 0 | 0 | 0 | 0 |
| | 01 | 0 | 0 | 0 | 1 | 0 | 0 |
| | 11 | - | - | - | - | - | - |
| | 10 | 0 | 0 | - | - | - | - |

Fig. 9. Karnaugh Map of first bit A

| DB | | C,x | | | | | |
|-----|----|-----|---|----|----|----|----|
| | | | | 00 | 01 | 11 | 10 |
| A,B | 00 | 0 | 0 | 0 | 0 | 0 | 1 |
| | 01 | 0 | 1 | 0 | 0 | 1 | 1 |
| | 11 | - | - | 0 | - | - | - |
| | 10 | 0 | 0 | - | - | - | - |

Fig. 10. Karnaugh Map of second bit B

| DC | | C,x | | | | | |
|-----|----|-----|---|----|----|----|----|
| | | | | 00 | 01 | 11 | 10 |
| A,B | 00 | 0 | 1 | 1 | 0 | 0 | 0 |
| | 01 | 0 | 1 | 0 | 0 | 0 | 0 |
| | 11 | - | - | - | - | - | - |
| | 10 | 0 | 1 | - | - | - | - |

Fig. 11. Karnaugh Map of last bit C

| DC | | C,x | | | | | |
|-----|----|-----|---|----|----|----|----|
| | | | | 00 | 01 | 11 | 10 |
| A,B | 00 | 0 | 1 | 1 | 0 | 0 | 0 |
| | 01 | 0 | 1 | 0 | 0 | 0 | 0 |
| | 11 | - | - | - | - | - | - |
| | 10 | 0 | 1 | - | - | - | - |

Fig. 12. Karnaugh Map of output Y

The following formulations represent the logical expressions generated from the Karnaugh Maps:

$$DA(A, B, C, x) = BCx$$

$$DB(A, B, C, x) = Cx' + BC'x$$

$$DC(A, B, C, x) = B'x + C'x$$

$$Y(A, B, C, x) = A$$

The RTL code for the structural design of the "1011" sequence detector is shown in Fig. 13. Fig. 14 shows the testbench codes created for verifying the reference model.

```

1 // Code your design here
2 module FSM_1011_det_model (input clk, x, reset, output detect_model);
3
4   wire D_A, D_B, D_C;
5   reg A, B, C;
6   assign D_A = B & C & x;
7   assign D_B = (C & ~x) | (B & ~C & x);
8   assign D_C = (~B & x) | (~C & x);
9   assign detect_model = A;
10
11   always @(posedge clk)
12   begin
13     A <= reset? 1'b0 : D_A;
14     B <= reset? 1'b0 : D_B;
15     C <= reset? 1'b0 : D_C;
16   end
17 endmodule

```

Fig. 13. RTL code

```

1 // Code your testbench here
2 // or browse Examples
3 module FSM_1011_model_tb;
4     reg clk, x, reset;
5     wire detect_model;
6
7     FSM_1011_det_model u0 ( .clk(clk), .reset(reset), .x(x),
8     .detect_model(detect_model) );
9
10    initial begin
11        forever #5 clk = ~clk;
12    end
13
14    initial begin
15        $monitor ("At time = %t, Out = %d", $time, detect_model);
16        clk = 0;
17        reset = 1;
18        x = 0;
19
20        #20 reset = 0;
21
22        #10 x = 1;
23        #10 x = 0;
24        #10 x = 1;
25        #10 x = 0;
26        #10 x = 1;
27        #10 x = 0;
28
29        #20 $finish;
30    end
31
32    initial begin
33        $dumpfile("dump.vcd");
34        $dumpvars(0);
35    end
36 endmodule

```

Fig. 14. Testbench code

```

1 // Code your testbench here
2 // or browse Examples
3 module FSM_1011_model_tb;
4     reg clk, x, reset;
5     wire detect_model_fsm, detect_model_struct, comp;
6     assign comp = detect_model_fsm == detect_model_struct;
7
8     FSM_1011_det_model dut ( .clk(clk), .rst(reset), .x(x), .detect_model_fsm );
9     FSM_1011_det_model m ( .clk(clk), .reset(reset), .x(x), .detect_model_struct );
10
11    always @ (posedge clk) begin
12        if (comp) begin
13            $display("Output are matched");
14        end
15        else begin
16            $display("Output are Not matched");
17        end
18    end
19
20    initial begin
21        forever #5 clk = ~clk;
22    end
23
24    initial begin
25        // $monitor ("At time = %t, Out = %d", $time, detect_model);
26        clk = 0;
27        reset = 1;
28        x = 0;
29        #20 reset = 0;
30
31        #10 x = 1;
32        #10 x = 0;
33        #10 x = 1;
34        #10 x = 0;
35        #10 x = 1;
36        #10 x = 0;
37
38        #20 $finish;
39    end
40
41    initial begin
42        $dumpfile("dump.vcd");
43        $dumpvars(0);
44    end
45 endmodule

```

Fig. 17. Testbench performing the comparison

B. Experimental results

Cadence software is used to simulate and synthesis the Verilog design of the structural description, which is the golden reference model. The circuit has a clock and four registers. In Figures 15 and 16, respectively, the observed waveform and the synthesised circuit are displayed.

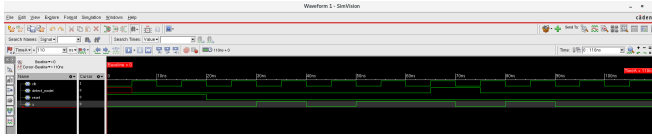


Fig. 15. Output Waveform

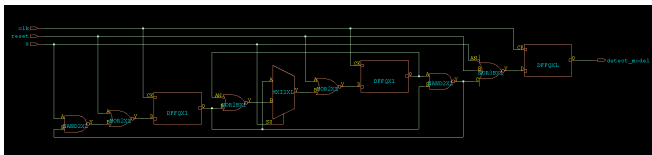


Fig. 16. Synthesized Circuit Diagram

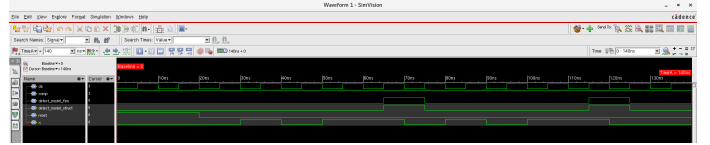


Fig. 18. Output Waveform

B. Experimental results

Utilizing Cadence software, the Verilog design that compares the structural description modules and FSM degree of abstraction is simulated. In Fig.18, the observed waveform is displayed.

V. ANALYSIS

The following table (Fig. 19) shows the comparison of the DUT and reference model with regards to power, area and time.

| | Power (nW) | Area (nm ²) | Timing (ps) |
|--------------------------|------------|-------------------------|-------------|
| FSM level of abstraction | 7634.430 | 135.485 | 1265 |
| Structural Description | 5692.098 | 76.447 | 788 |

Fig. 19. Comparative analysis

IV. COMPARISON OF THE TWO MODELS

A. Implementation details:

The FSM level of abstraction and structural description modules are both instantiated and their outputs are compared. The testbench which executes this task is shown below in Fig. 17.

VI. CONCLUSION

The "1011" sequence detector's FSM level of abstraction is effectively applied in Verilog design and is synthesised and simulated using Cadence software. The DUT is also validated using the structural description, which functioned as the golden reference model. The reference model exhibits minimal time, area and power parameters when compared to the FSM level of abstraction. We consider the FSM level of abstraction to be the best way to build the sequence detector since constructing the schematic diagram from scratch in structural description prevents the compiler from optimising the design to meet requirements that are not merely concerned with exactness.