

# Intermediate Code Generation

## Mini-C Compiler

*CS4031 - Compiler Construction | Fall 2025*

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## 1. Overview

Intermediate Code Generation is the fourth phase of the Mini-C compiler. It transforms the AST into TACKY IR (Three-Address Code), a low-level representation that is machine-independent but close to assembly.

Location: src/backend/ir/

### 1.1 Design Principles

- Three-Address Code: Each instruction has at most 3 operands
- Flat Structure: No nested expressions; complex expressions are broken down
- Explicit Control Flow: Jumps and labels instead of structured statements
- Typed Operations: Each operation knows its operand types
- SSA-like Temporaries: New temporaries for intermediate results

## 2. TACKY IR Instructions

### 2.1 Value Operations

Instruction	Format	Description
Copy	dst = src	Copy value
Unary	dst = op src	Unary operation
Binary	dst = left op right	Binary operation
Load	dst = *src	Load from pointer
Store	*dst = src	Store to pointer
GetAddress	dst = &src	Get address

### 2.2 Control Flow Instructions

Instruction	Format	Description
Jump	goto label	Unconditional jump
JumpIfZero	if src == 0 goto label	Conditional jump
JumpIfNotZero	if src != 0 goto label	Conditional jump
Label	label:	Jump target
Return	return src	Return from function
FunCall	dst = call func(args)	Function call

## 3. Translation Rules

### 3.1 Expression Translation

#### Binary Expression: $a + b * c$

```
tmp.1 = b * c
tmp.2 = a + tmp.1
result = tmp.2
```

#### Short-Circuit AND: $a \&\& b$

```
tmp.1 = a
    if tmp.1 == 0 goto false_label
    tmp.2 = b
    if tmp.2 == 0 goto false_label
    tmp.3 = 1
    goto end_label
false_label:
    tmp.3 = 0
end_label:
    result = tmp.3
```

#### Ternary: $a ? b : c$

```
tmp.1 = a
    if tmp.1 == 0 goto else_label
    tmp.2 = b
    goto end_label
else_label:
    tmp.2 = c
end_label:
    result = tmp.2
```

### 3.2 Statement Translation

#### If-Else Statement:

```
tmp.1 = condition
    if tmp.1 == 0 goto else_label
    ; then_body instructions
    goto end_label
else_label:
    ; else_body instructions
end_label:
```

#### While Loop:

```
continue_label:
    tmp.1 = condition
    if tmp.1 == 0 goto break_label
    ; body instructions
```

```
    goto continue_label
break_label:
```

### For Loop:

```
; init instructions
start_label:
    tmp.1 = condition
    if tmp.1 == 0 goto break_label
    ; body instructions
continue_label:
    ; post instructions
    goto start_label
break_label:
```

## 4. Complete Translation Example

### Source Code:

```
int factorial(int n) {  
    if (n <= 1)  
        return 1;  
    return n * factorial(n - 1);  
}
```

### TACKY IR Output:

```
function factorial(n):  
    ; if (n <= 1)  
    tmp.1 = n <= 1  
    if tmp.1 == 0 goto else.1  
  
    ; return 1  
    return 1  
  
else.1:  
    ; n - 1  
    tmp.2 = n - 1  
  
    ; factorial(n - 1)  
    tmp.3 = call factorial(tmp.2)  
  
    ; n * factorial(n - 1)  
    tmp.4 = n * tmp.3  
  
    ; return result  
    return tmp.4
```

## Array Sum Example:

```
int sum_array(int *arr, int size) {
    int total = 0;
    for (int i = 0; i < size; i = i + 1) {
        total = total + arr[i];
    }
    return total;
}
```

## TACKY IR Output:

```
function sum_array(arr, size):
    ; int total = 0
    total = 0
    ; int i = 0
    i = 0

for.start.1:
    ; i < size
    tmp.1 = i < size
    if tmp.1 == 0 goto for.break.1

    ; arr[i] - compute address
    tmp.2 = i * 4           ; sizeof(int) = 4
    tmp.3 = arr + tmp.2     ; pointer arithmetic
    tmp.4 = *tmp.3          ; load arr[i]

    ; total = total + arr[i]
    tmp.5 = total + tmp.4
    total = tmp.5

for.continue.1:
    ; i = i + 1
    tmp.6 = i + 1
    i = tmp.6
    goto for.start.1

for.break.1:
    ; return total
    return total
```

## 5. IR Properties

Property	Description
Machine Independence	No register references, no specific instruction set
Type Preservation	Operations maintain type information
Explicit Addressing	All memory accesses made explicit

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Linear Representation	Flat sequence of instructions
SSA-like	Temporaries typically assigned once