

\* Note: I will go over this in class again if I see necessary.

## Cyclic Redundancy Check (CRC)

CRC can detect both single and burst errors

Commonly used CRCs CRC-1 (parity), CRC-3-GSM, CRC-4-ITU, CRC-5-EPC, CRC-5-USB, CRC-7-MVB (Train communication), CRC-8 (satellite DVB-S2), CRC-8/16-CCITT (wireless, Bluetooth), CRC-40 (GSM control channel-slower/reliable).

You can check hardware implementation of CRC and parity check [here](#).

In CRC, you should know about the input message, the generator and the codeword.

When you divide message+padding by generator, you should get a remainder. Add that remainder to your message to get the codeword.

Lets go through the CRC process then,

### Transmit Process

1. Your input message bit pattern
2. Your generator
3. Add padding (0s) after the message. The size of the padding will be  $\text{number\_of\_bits}(\text{generator}) - 1$
4. Divide the message+padding with the generator
5. Store the remainder of size =  $\text{number\_of\_bits}(\text{generator}) - 1$
6. Add that to your message and send.

### Receive Process

1. Divide the received message by the generator
2. Did you get a remainder of 0's?
3. If not, there is/are error(s)
4. If yes, your recieved bits are okay.

Let's go through an example,

## Implementation (simple long division)- Transmitter side

```
----- TRANSMISSION PROCESS (longer division)
-----

Your polynomial (input message) = 1101011011 (10
bits)
Your generator = 10011 (5 bits)
number of padding bits = number_of_bits(10011) - 1
                        = 4

Your input message after padding = 11010110110000
(14 bits)

The division process is a XOR at each step you see
below.

Multiply generator with 1 when MSB is 1.
```

```

    _1100001010_____
10011| 11010110110000
      10011| | | | | |
      ----| | | | | |
          10011| | | | | |
          10011| | | | | |
          ----| | | | | |
              00001| | | | |
              00000| | | | |
              ----| | | | |
                  00010| | | | |
                  00000| | | | |
                  ----| | | | |
                      00101| | | | |
                      00000| | | | |
                      ----| | | | |
                          01011| | | | |
                          00000| | | | |
                          ----| | | | |
                              10110| | | | |
                              10011| | | | |
```

```

-----|||
      01010||
      00000||
-----||
      10100|
      10011|
-----|
      01110
      00000
-----
      1110

```

So your transmitted message is 11010110111110 where 1110 is the remainder

----- shorter division -----

this division took a long while to solve. How about making it shorter (optional read)

```

      _1100001010_
10011| 11010110110000
      10011|||||||
      ----|||||||
      10011|||||||
      10011|||||||
      ----|||||||
      000010110      ; we need 5 bits to make MSB
1
      10011      ; we add 4 zeros in the
quotient
      -----|||
      0010100      ; we need 2 bits to make MSB
1
      10011      ; we add 1 zero in the
quotient
      -----|
      001110      ; we are still adding bits
after the division is complete, add a zero to in the
quotient.

```

----- arithmetic example -----

Check an arithmetic exmple below,

```

      _20040_____
32 | 641281
    64
    ---
    128      ; adding 3 numbers to make it greater
than 32
    128      ; added 2 zeros in the quotient
    ---
    0001     ; after all that it still has 1 left in
dividend
              ; add zero in quotient

```

## Impelementation (simple long division)- Receiver side (no error)

Receive side should receive -  
 message + crc\_calculated(remainder) = 11010110111110  
 1101011011(1110)  
 Now we go through same process of division with  
 generator --  
 10011  
 But this time with dividend,

```

      _1100001010_____
10011 | 11010110111110
      10011 |||||
      ----- |||||
      10011 |||||
      10011 |||||
      ----- |||||
      00001 |||||
      00000 |||||
      ----- |||||
      00010 |||||
      00000 |||||
      ----- |||||

```

```

      00101| | | |
      00000| | | |
-----| | | |
      01011| | | |
      00000| | | |
-----| | | |
      10111| | |
      10011| | |
-----| | |
      01001| |
      00000| |
-----| |
      10011|
      10011|
-----|
      00000
      00000
-----
      0000

```

## Implementation- Receiver Side (With error)

Receive side should receive -  
 message + crc\_calculated(remainder) = 11010110011110  
 But unfortunately received = 11010110011110  
 Now we go through same process of division with  
 generator --  
 10011  
 But this time with dividend,

```

      _1100001000_
10011| 11010110011110
      10011x|x|x|x
      ----x|x|x|x
      10011|x|x|x|x
      10011|x|x|x|x
      -----|x|x|x|x
      00001x|x|x|x
      00000x|x|x|x
      -----x|x|x|x
      00010|x|x|x
      00000|x|x|x

```

```

-----|x|x|x
    00100x|x|x
      00000x|x|x
-----x|x|x
    01001|x|x
      00000|x|x
-----|x|x
    10011x|x
      10011x|x
-----x|x
    00001|x
      00000|x
-----|x
    00011x
      00000x
-----
    00110
      00000
-----
    0110

```

remainder is non-zero  
 so there is a receive error