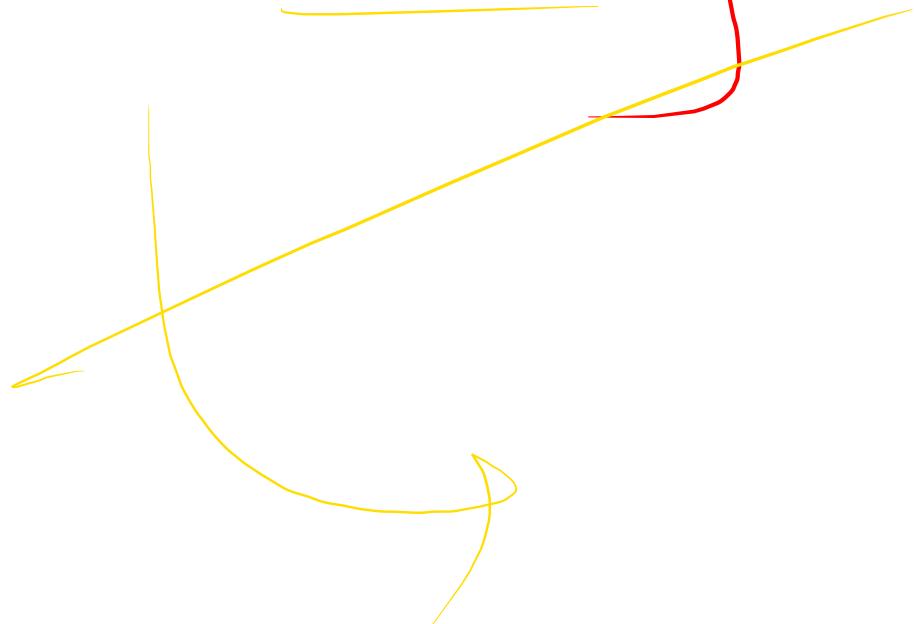


Sorting

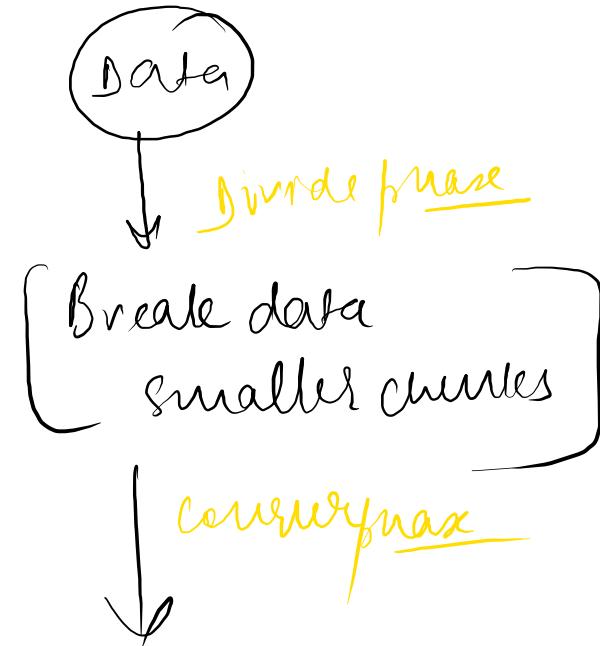
① Bubble $O(n^2)$ [frequency]

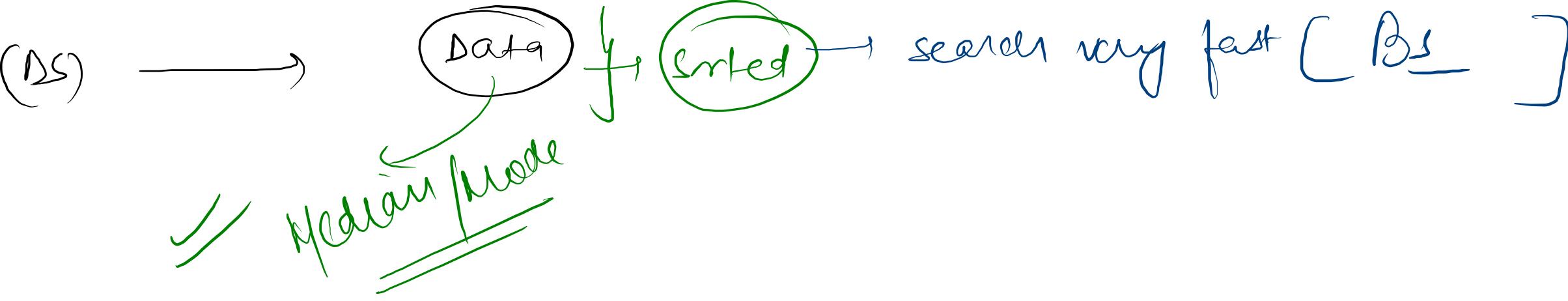
② Selection

③ Insertion

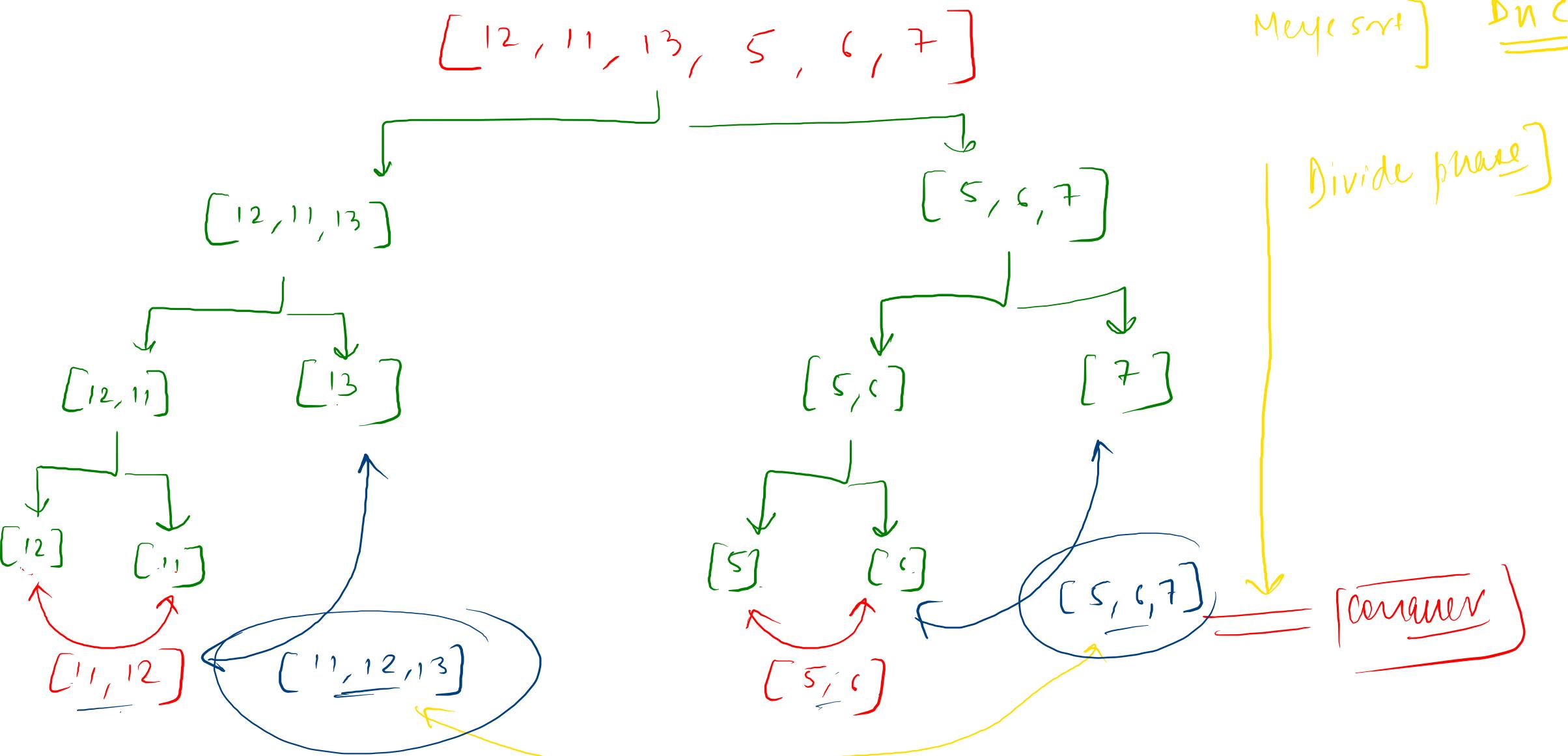


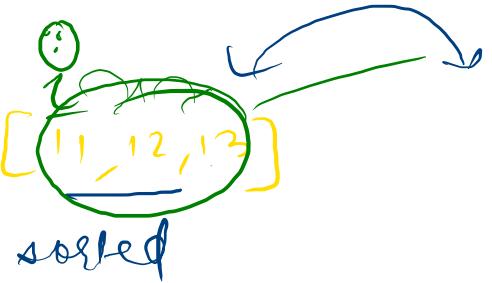
Merge Sort → Dn C | Divide and Conquer Algorithm





Merge sort] DnC



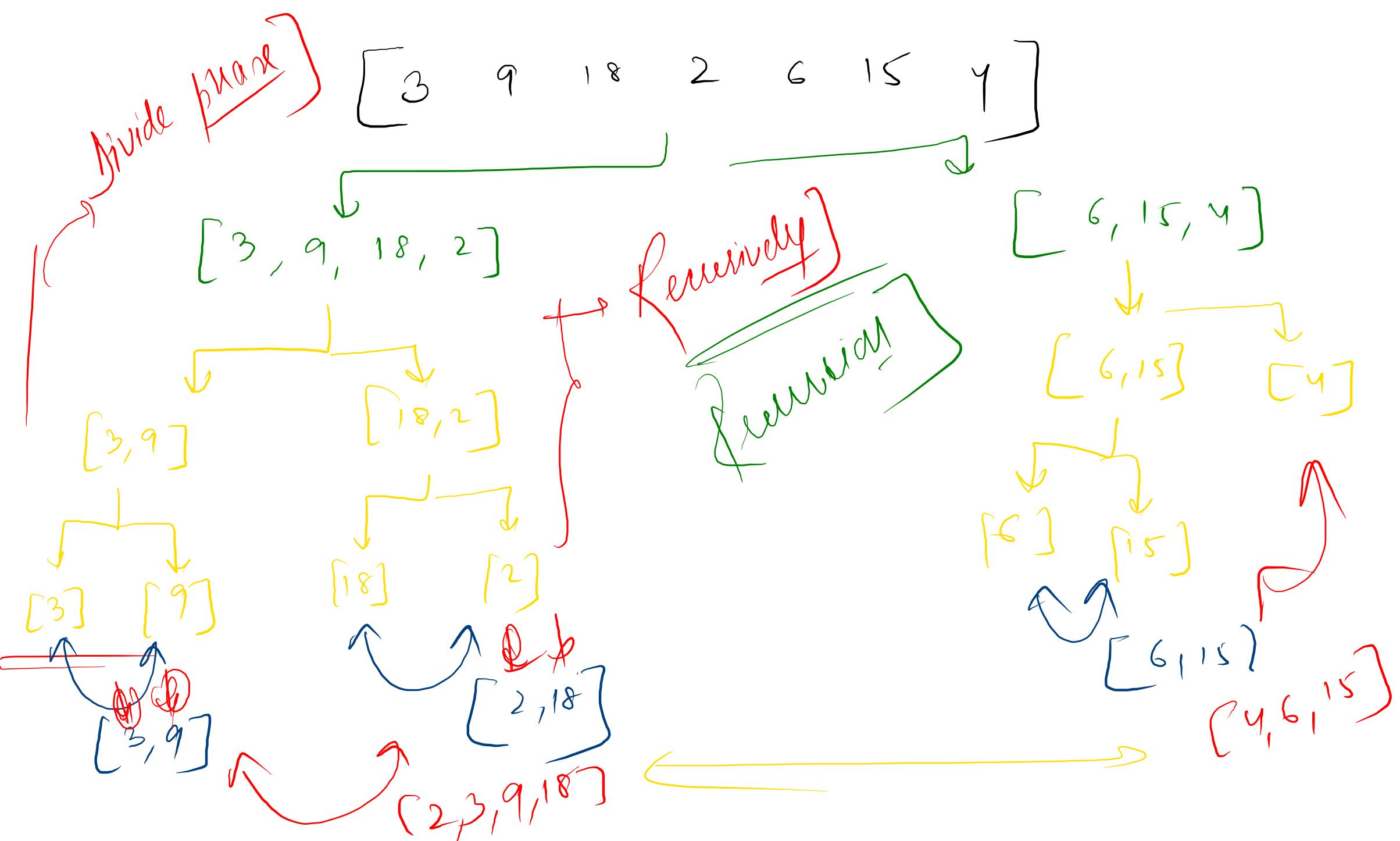

[11, 12, 13]
sorted

[5, 6, 7]
sorted

combine so that after
combination we get sorted
array

[5, 6, 7, 11, 12, 13] sorted array

~~Push SMT~~



~~2, 3, 4, 6, 9, 15~~

[2, 3, 9, 18]

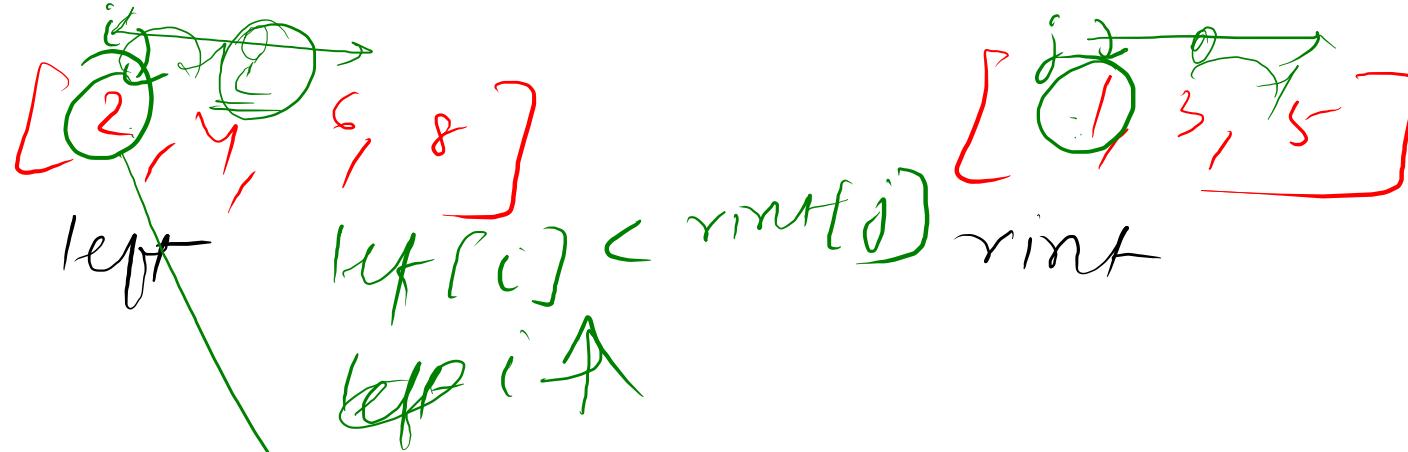
left

[4, 2, 15]

right

meep

[2, 3, 4, 6, 9, 15, k]



arr



```
def merge(left, right, arr):  
    i=0 #index of left array  
    j=0 #index of right array  
    k=0 # index of the result arr
```

```
        while i < len(left) and j < len(right):  
            if(left[i] <= right[j]):  
                arr[k] = left[i]  
                i +=1  
            else:  
                arr[k] = right[j]  
                j+=1  
            k +=1  
  
            #Check if any element is left in left array  
            while i < len(left):  
                arr[k] = left[i]  
                i +=1  
                k +=1  
  
            #Check if any element is left in right array  
            while j < len(right):  
                arr[k] = right[j]  
                j +=1  
                k +=1
```

[, , ,]

[2, 1, 6]
i ↑ j ↑ k ↑
i=0 j=1 k=2

[1, 2, 6]
i ↑ j ↑ k ↑
i=0 j=1 k=2

Terminates

Terminates

[1, 2, 3, 4, 6]
i ↑ j ↑ k ↑
i=0 j=1 k=2

Sorted

```
def mergeSort(arr): # this fn sorts the arr passed as argument
```

```
if (len(arr)<=1):  
    return  
# Divide into two parts  
mid = len(arr)//2
```

```
#Left half of the arr  
left = arr[:mid]
```

```
#Right half of the arr  
right = arr[mid:]
```

```
#sort the left part  
mergeSort(left)
```

```
#sort the right part  
mergeSort(right)
```

```
#logic to merge both individually sorted left and right  
merge(left, right, arr)
```

```
def merge(left, right, arr):  
    i=0 #index of left array  
    j=0 #index of right array  
    k=0 # index of the result arr
```

```
while i<len(left) and j<len(right):  
    if(left[i]<= right[j]):  
        arr[k] = left[i]  
        i +=1  
    else :  
        arr[k] = right[j]  
        j+=1  
    k +=1
```

[12, 11, 13, 5, 6, 7] n

[5, 5, 7] $n/2$

$n/4$

$n/8$

$Tc = O(n \log n)$

Left arranic

[12, 11, 13]

[12, 11]

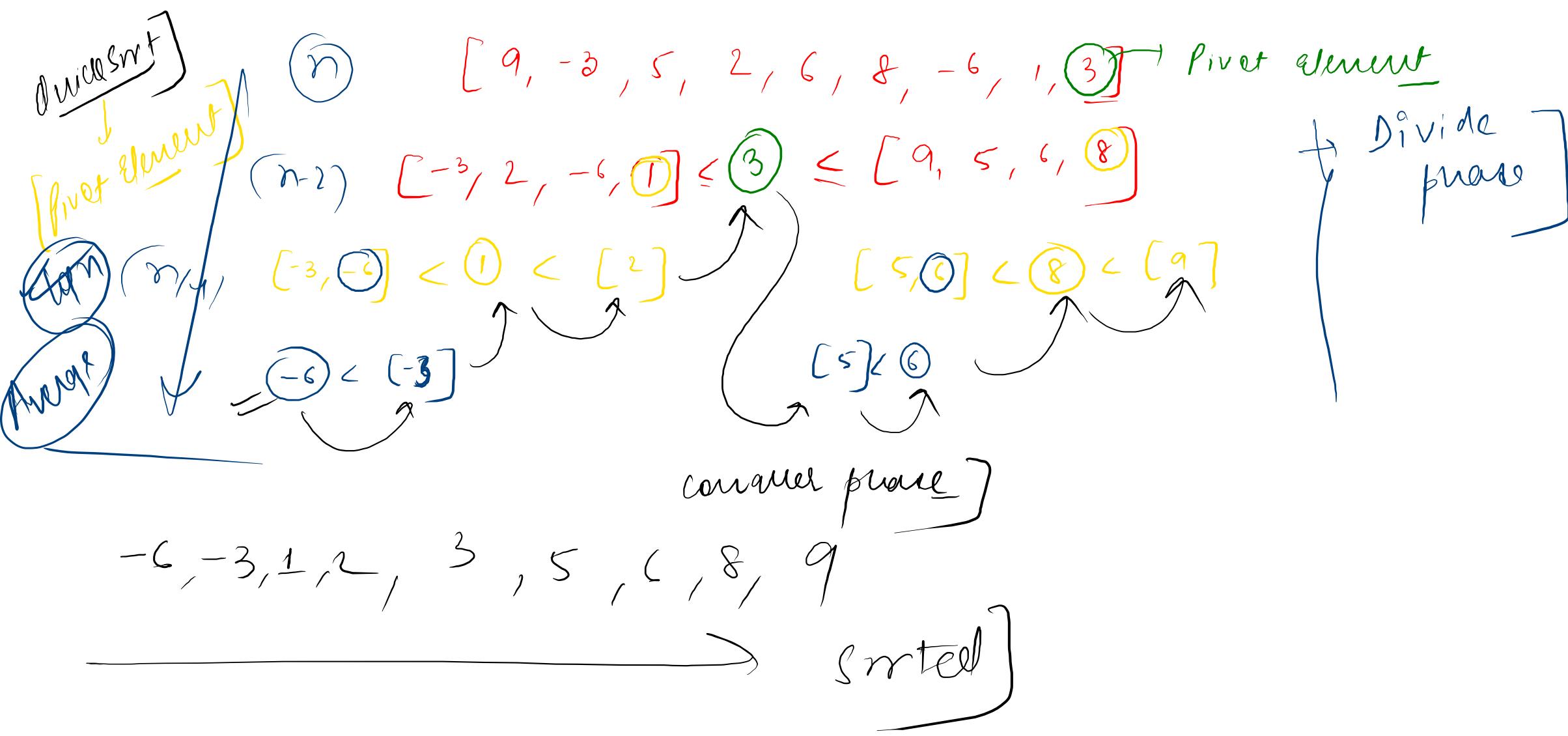
13

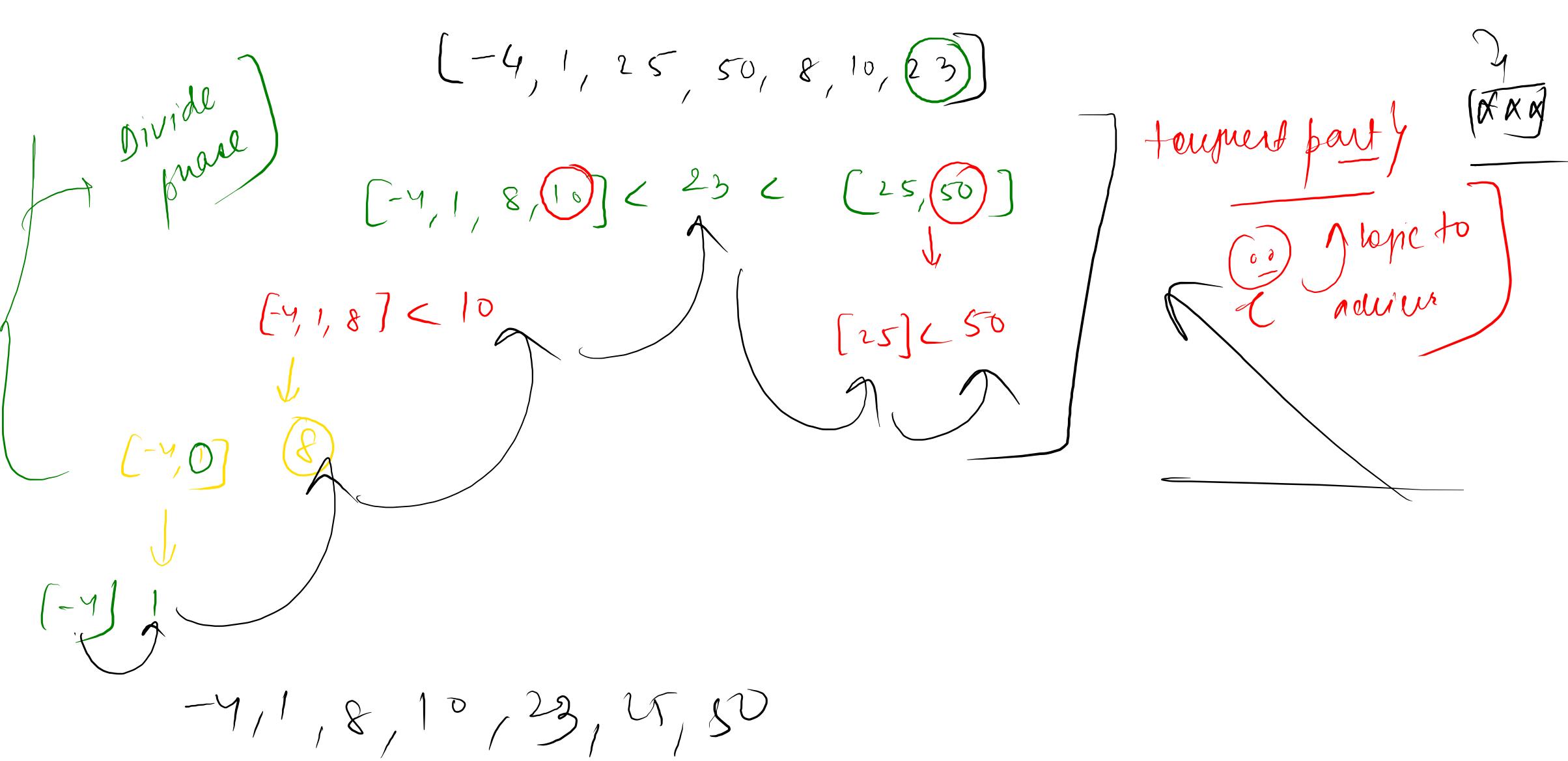
Linear $O(n)$

For $(\log n)$ ~~\rightarrow~~ n steps to merge

$O(m \log n)$

Quick Sort } Divide and Conquer Algorithm





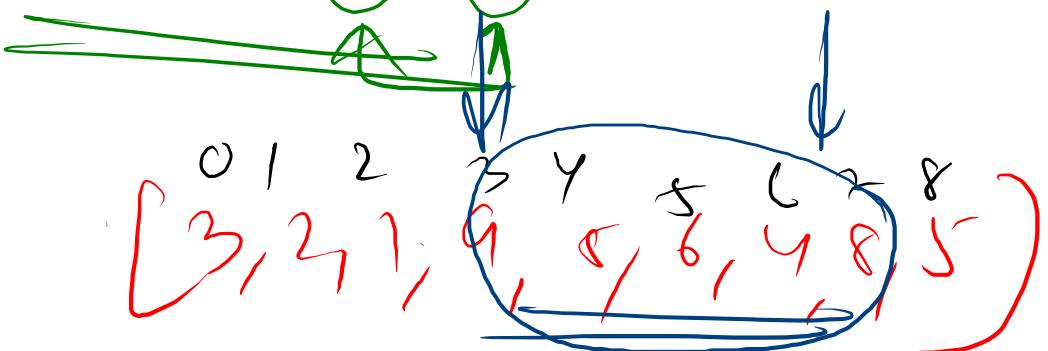
$$O(\log n) + O(\log n) \rightarrow \underline{O(n)}$$

[map]

for arr, start, end

↳ whether arr → [low, high]

```
def quickSort(arr, low, high):
```



arr[3:7] ↳

3 ↳ ↳
2 ↳ ↳
1 ↳ ↳
4, 5, 6, 8, 9, ↳
5 ↳ ↳
sorted

```
def quickSort(arr, low, high):
```

arr: [5, 4, 1, 2, 3, 9, 1, 8, 6]

quickSort(arr, 2, 6) ↗ [5, 4, 1, 2, 3, 8, 9, 1, 6]
sorted

Am I clear!

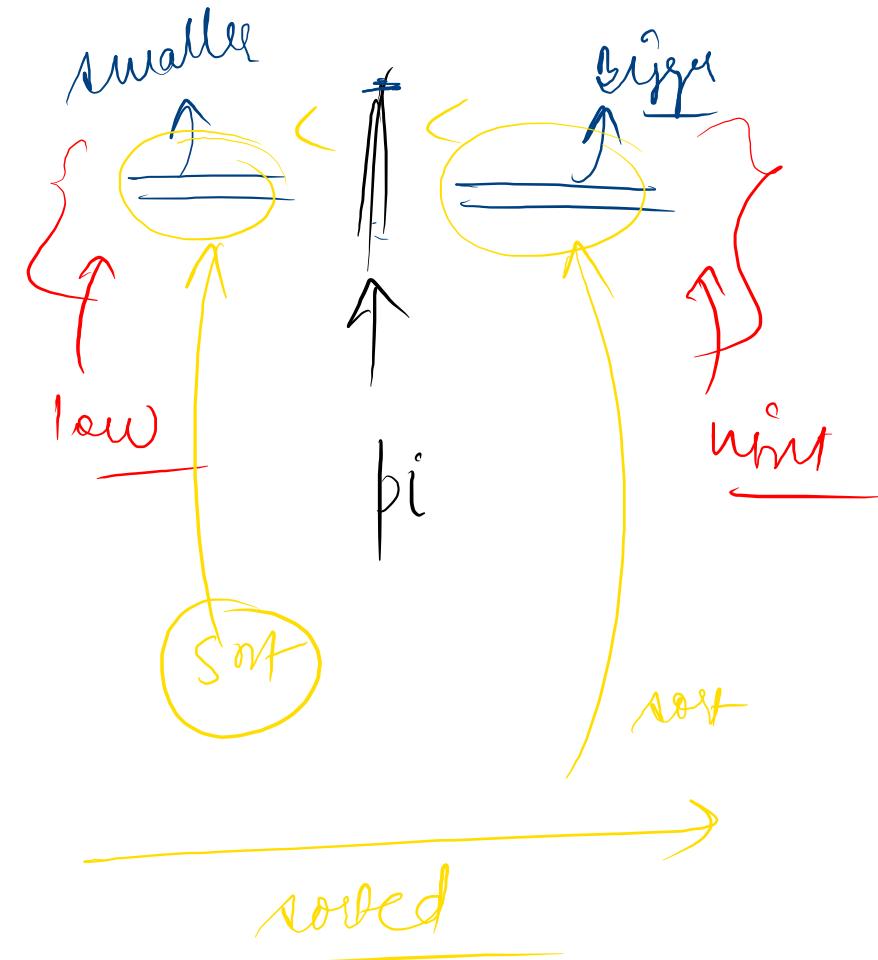
```

def quickSort(arr, low, high):
    # Base condition
    if low >= high :
        return

    #find the pivot point and re-arrange the array around it
    pi = partition(arr, low, high)

    #left of pi will be all small
    # right of pi will be all big

```



```

def quickSort(arr, low, high):
    # Base condition
    if low >= high :
        return

    #find the pivot point and re-arrange the array around it
    pi = partition(arr, low, high)

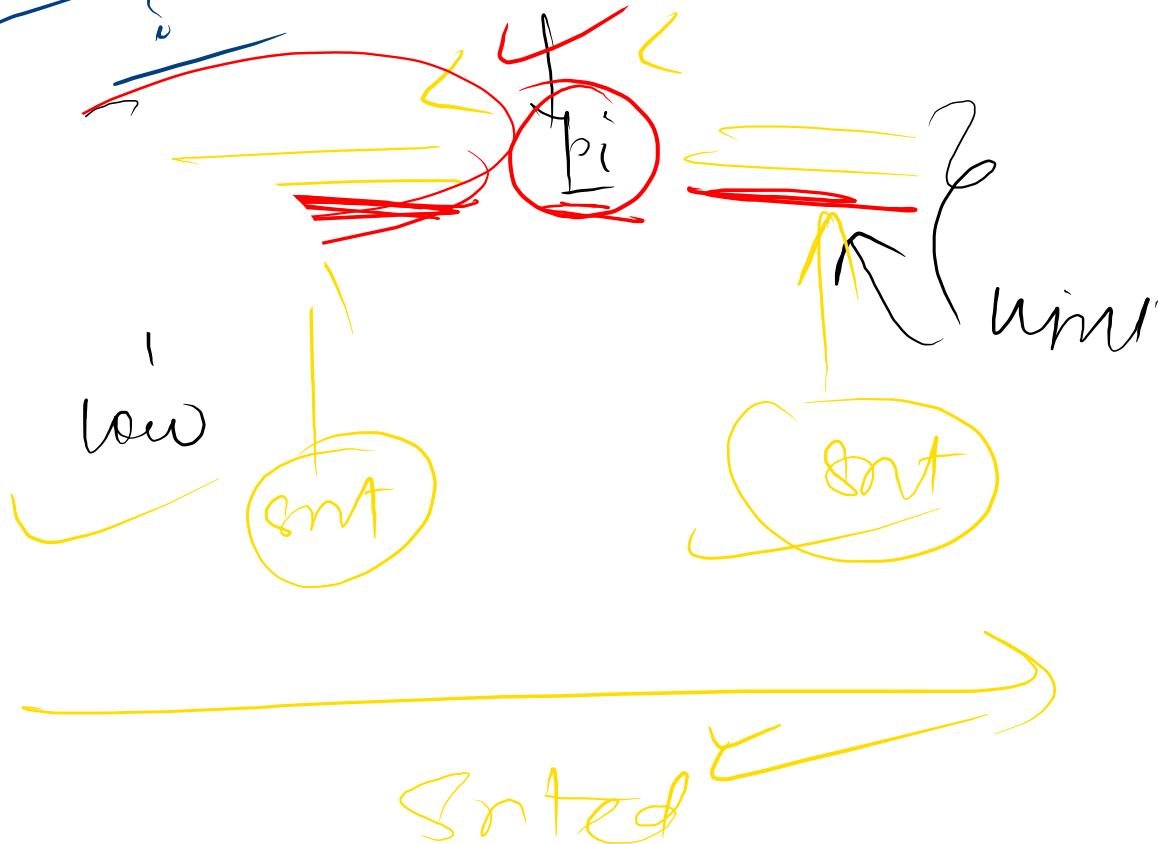
    #left of pi will be all small
    # right of pi will be all big

    #sort the left part
    quickSort(arr, low, pi-1)

    #sort the right part
    quickSort(arr, pi+1, high)

```

How we partition the array?



```

def partition(arr, low, high):
    # We choose right most element as the pivot element
    pivot = arr[high] = 1

```

```

    # Traverse through the entire array
    # Compare the elements with the pivot element
    i = low-1 # i indicate the available index for the smaller ele
    for j in range(low, high):
        if(arr[j] <= pivot):
            # We need to set it in the left part of the array
            i = i+1 # position where arr[j] should be

```

```

    # swap arr[j] and i index elements
    arr[i], arr[j] = arr[j], arr[i]

```

```

    # Swap the pivot element with the next greater element of i
    arr[i+1], arr[high] = arr[high], arr[i+1]

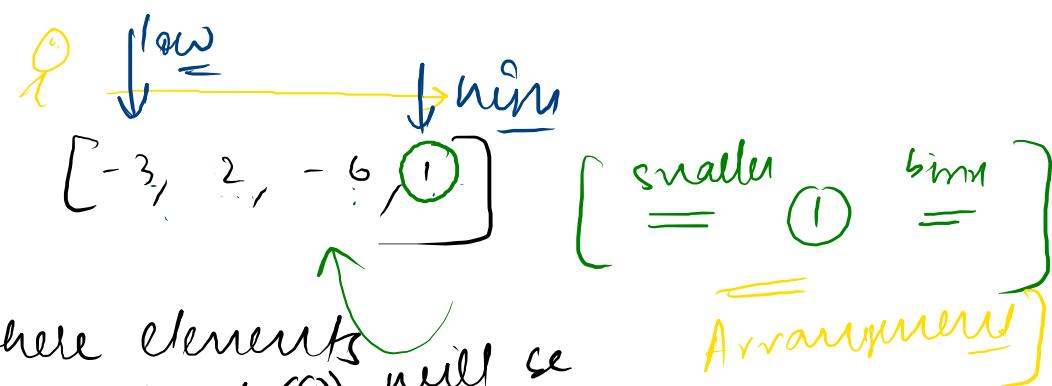
```

```

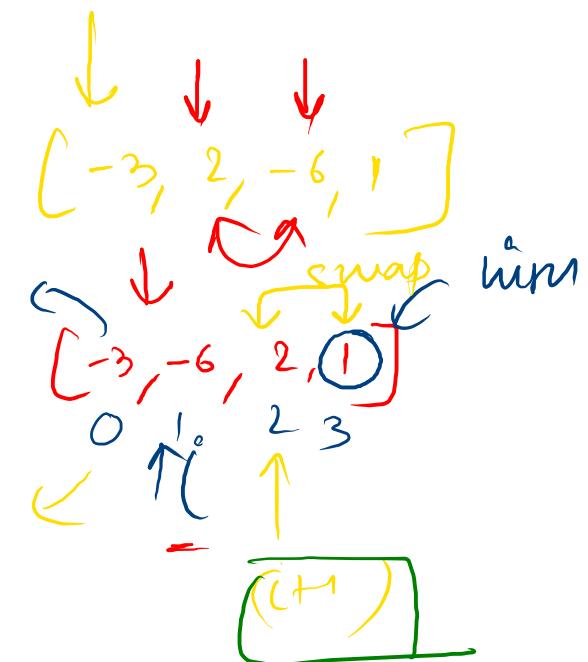
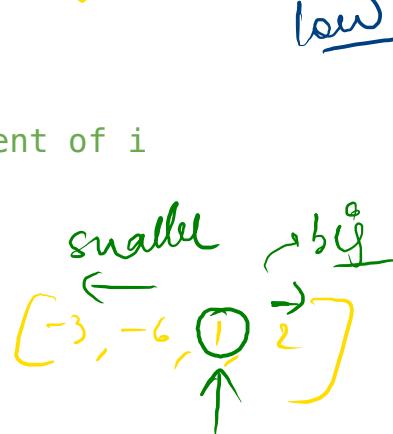
return i+1

```

index of pivot element



i → index where elements smaller than pivot (1) will be stored
 $i = 0-1 = -1 = H$ → 1
 $j = i+1$



Arrangement

= ① =
 = ② =



```
def partition(arr, low, high):
    # We choose right most element as the pivot element
    pivot = arr[high] = 1
```

low high i

```
# Traverse through the entire array
# Compare the elements with the pivot element
i = low - 1 # i indicate the available index for the smaller
elements than pivot
for j in range(low, high):
    if(arr[j] <= pivot):
        # We need to set it in the left part of the array
        i = i + 1 # position where arr[j] should be

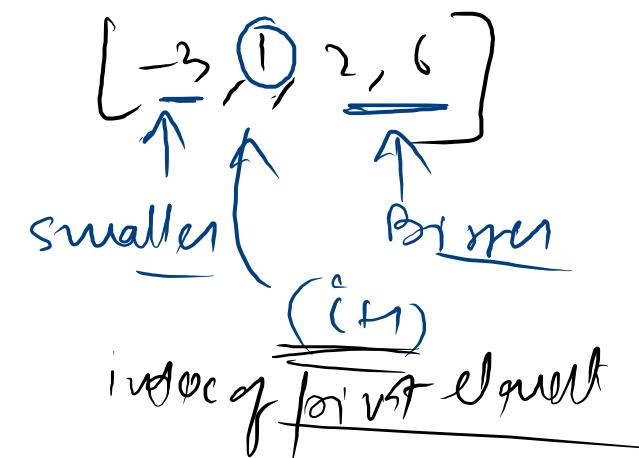
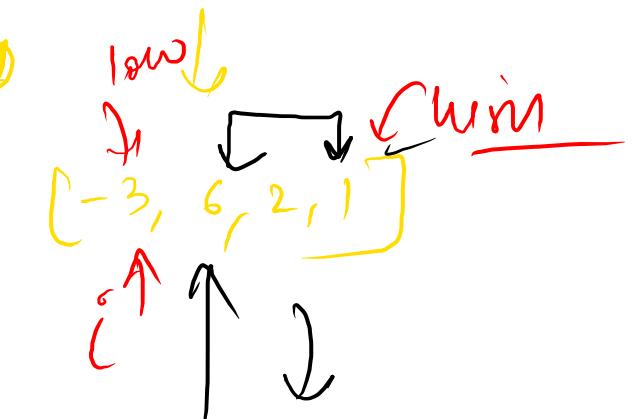
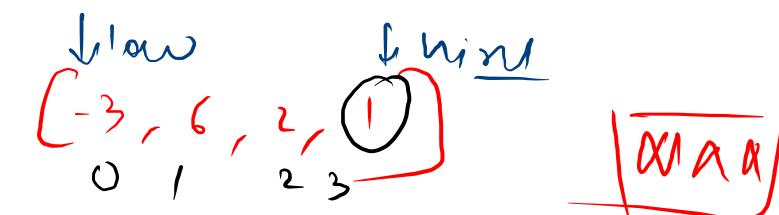
        # swap arr[j] and i index elements
        arr[i], arr[j] = arr[j], arr[i]

    # Swap the pivot element with the next greater element of i
    arr[i+1], arr[high] = arr[high], arr[i+1]
```

return i+1

$$i = 0 - 1 = -1 \rightarrow 0$$

i = 0



Reversed } 9 8 7 5 4 3 1
by worst case $T = O(n^L)$

Worst }

1 9 8 7 5 4 3 (n)

n times/steps

(r) corner

$O(n^2)$

3

4 1 9 8 7 (n-1)

5

6

7

8

9

10

11

12

13

14

15

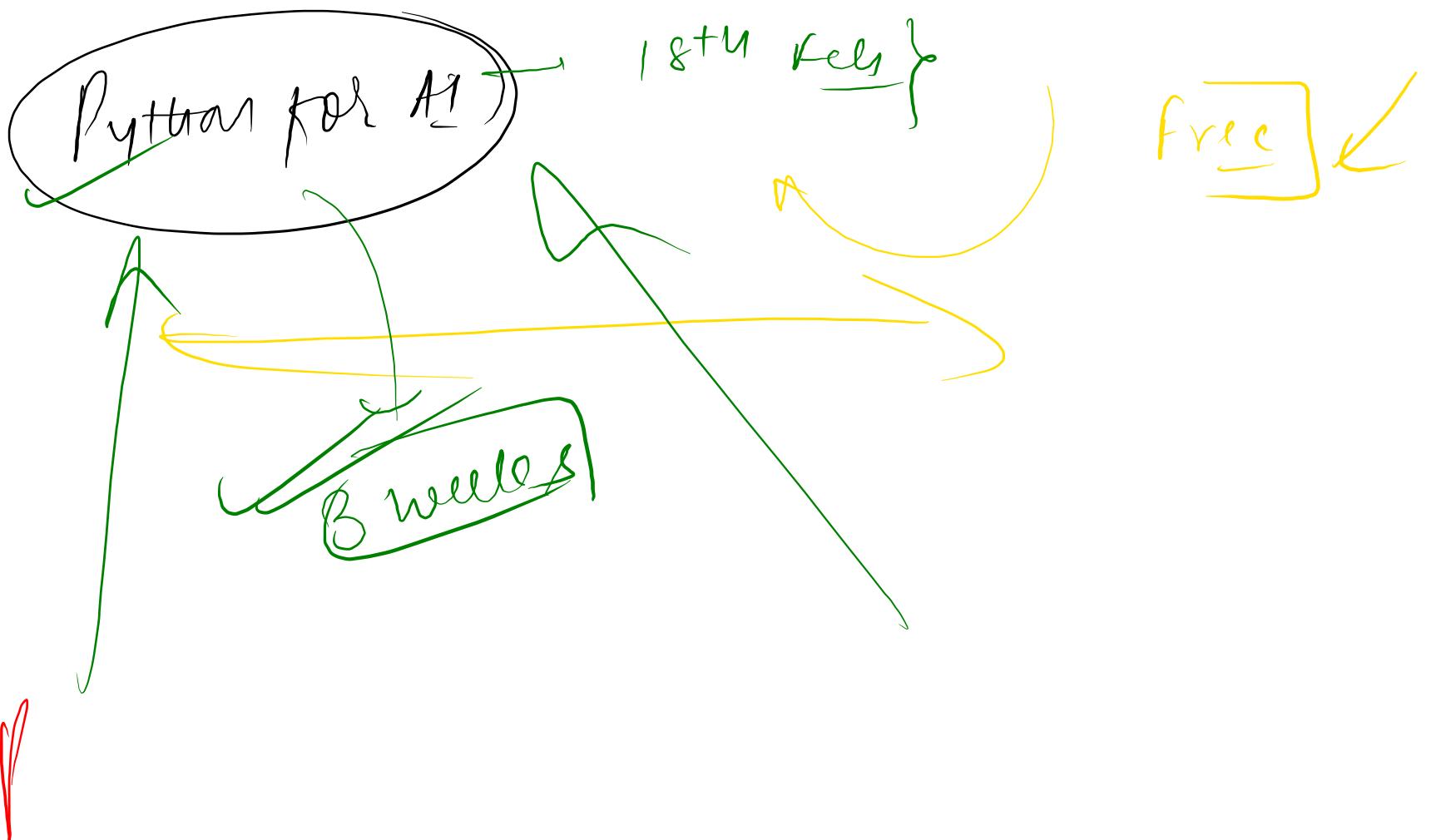
16

17

18

19

Interview } obviously } $O(n^2)$ } Reversed array)



MergeSort([3, 4, 5, 1, 2, 9]) → [1, 3, 4, 5, 6, 9]

MergeSort([3, 4, 5]) → [3, 4, 5]