

Stack} FILO / LIFO }

insert → push()

remove → pop()

① Array

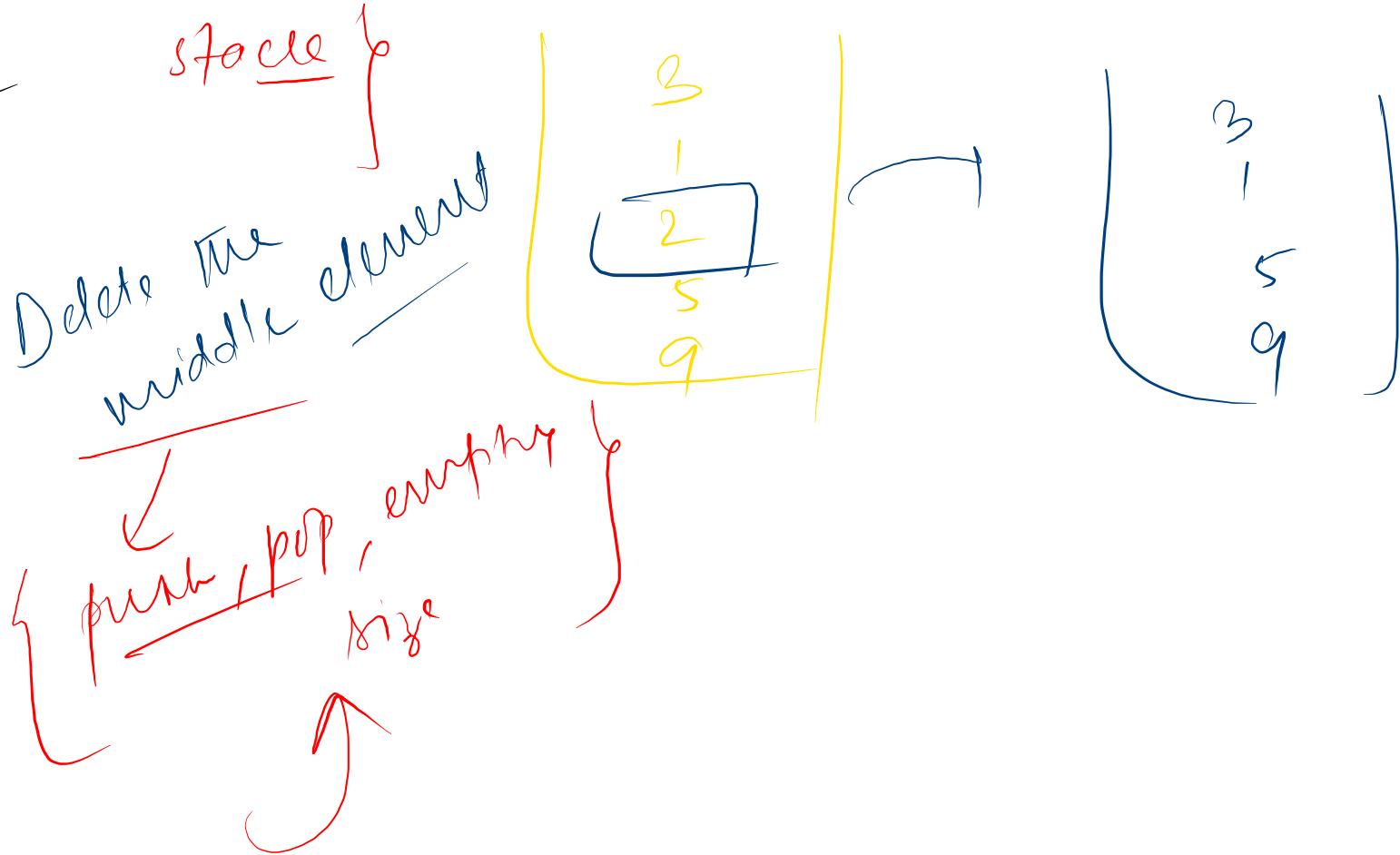
limitation :-

↓
fixed size
space

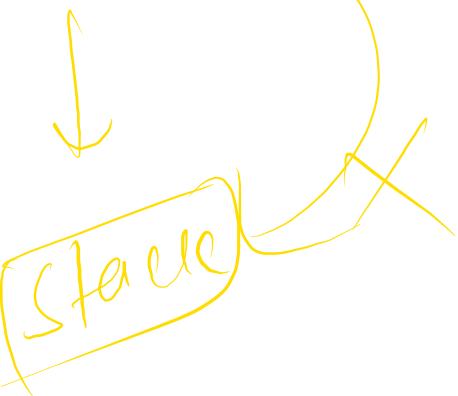
② Linked List

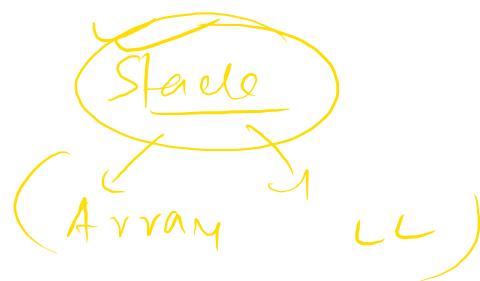
(Pr 1) → Balance the parentheses
Up ↓ Stage ←

Adriene



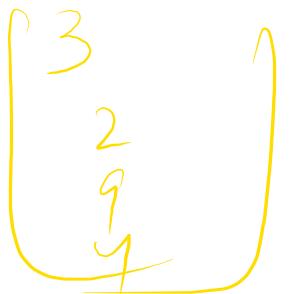
\emptyset (BS) sorted Array

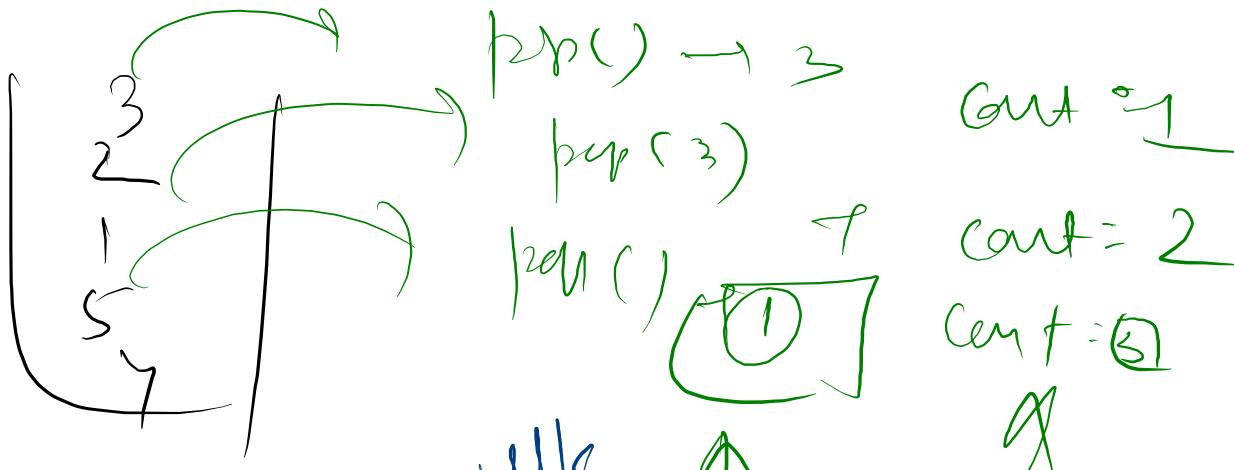




size ~ 5

~~temp~~





Size = 5

3rd element will Middle

Midd

A handwritten diagram illustrating a division problem. On the left, there is a vertical bracket enclosing three numbers: 3, 2, and 1. Above this bracket, the number 3 is written above a horizontal line. To the right of the bracket, the number 2 is written above a horizontal line, with a red arrow pointing from the top of the bracket to the number 2. Below the horizontal line, the number 1 is written. To the right of the horizontal line, the text "middle element" is written in red cursive. Below the horizontal line, the division symbol $\overline{-}$ is written, followed by the equals sign = and a circled 2.

$$\begin{array}{r} 3 \\ \overline{-} 2 \\ \quad 1 \end{array} = \textcircled{2}$$

Approach { → use 2 stacks }

```
def deleteMiddle(stack):  
    #Implement this method -> push/pop/peek  
    s2 = StackLL()
```

size = stack.size() → 5

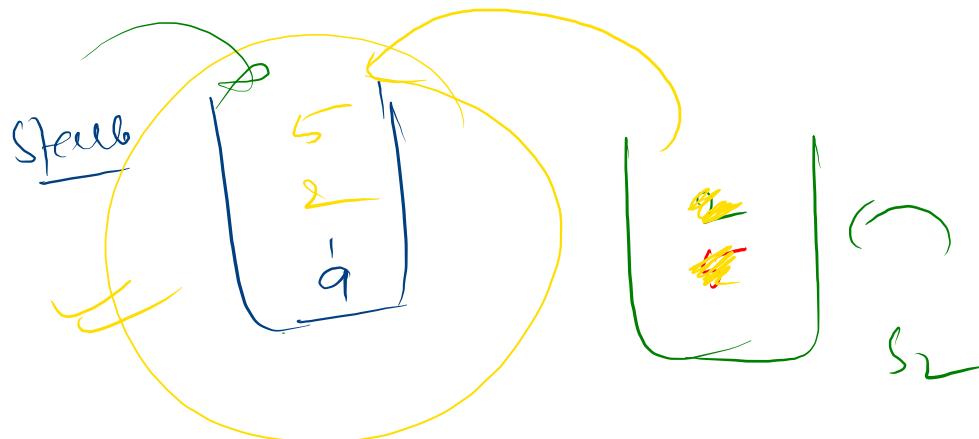
count = 0

while(count < int(size/2)):
 s2.push(stack.pop())
 count += 1

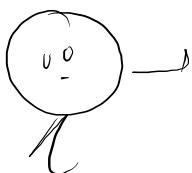
stack.pop()

while(not s2.isEmpty()):
 stack.push(s2.pop())

$Tc \sim O(n)$
 $Sc \sim O(1)$



1/2

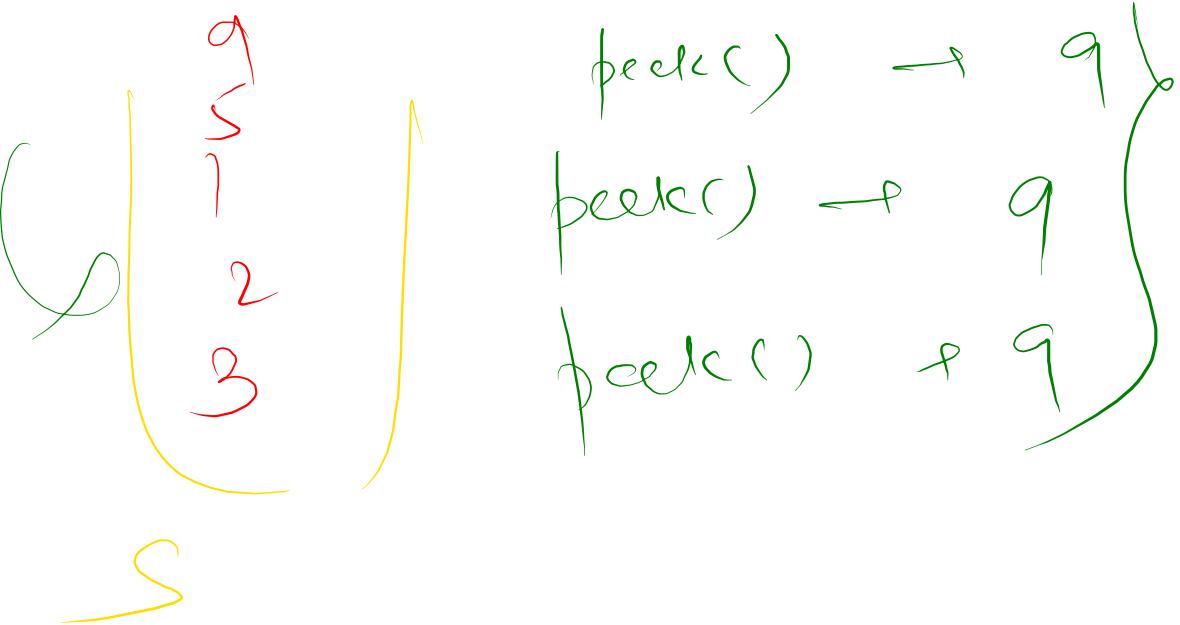


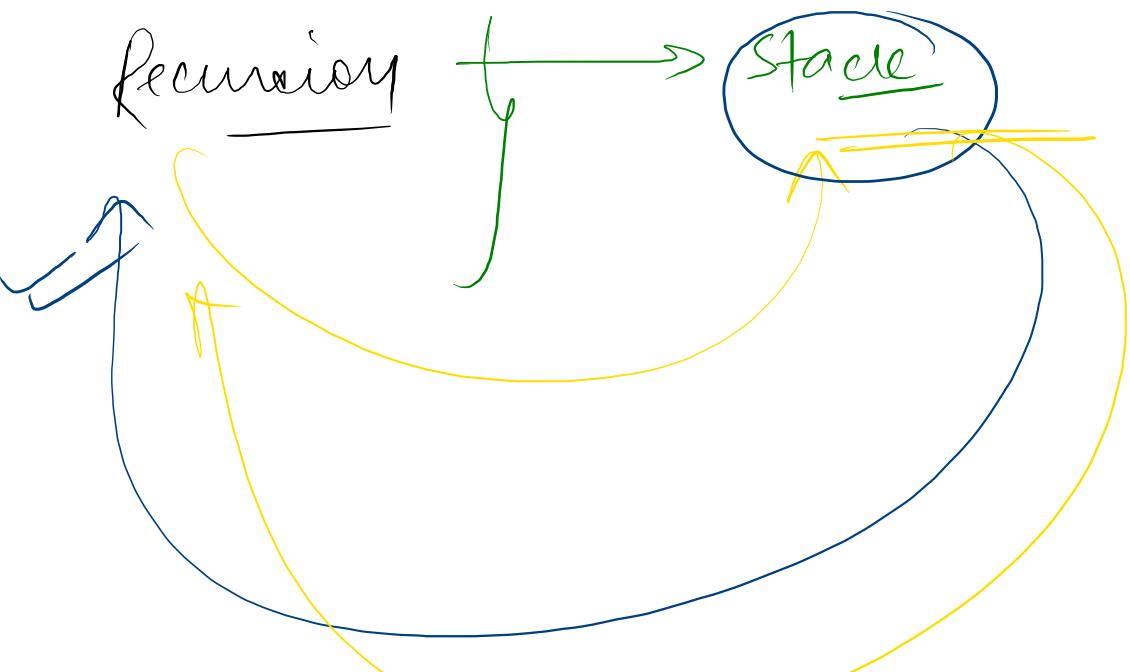
Additional Stack

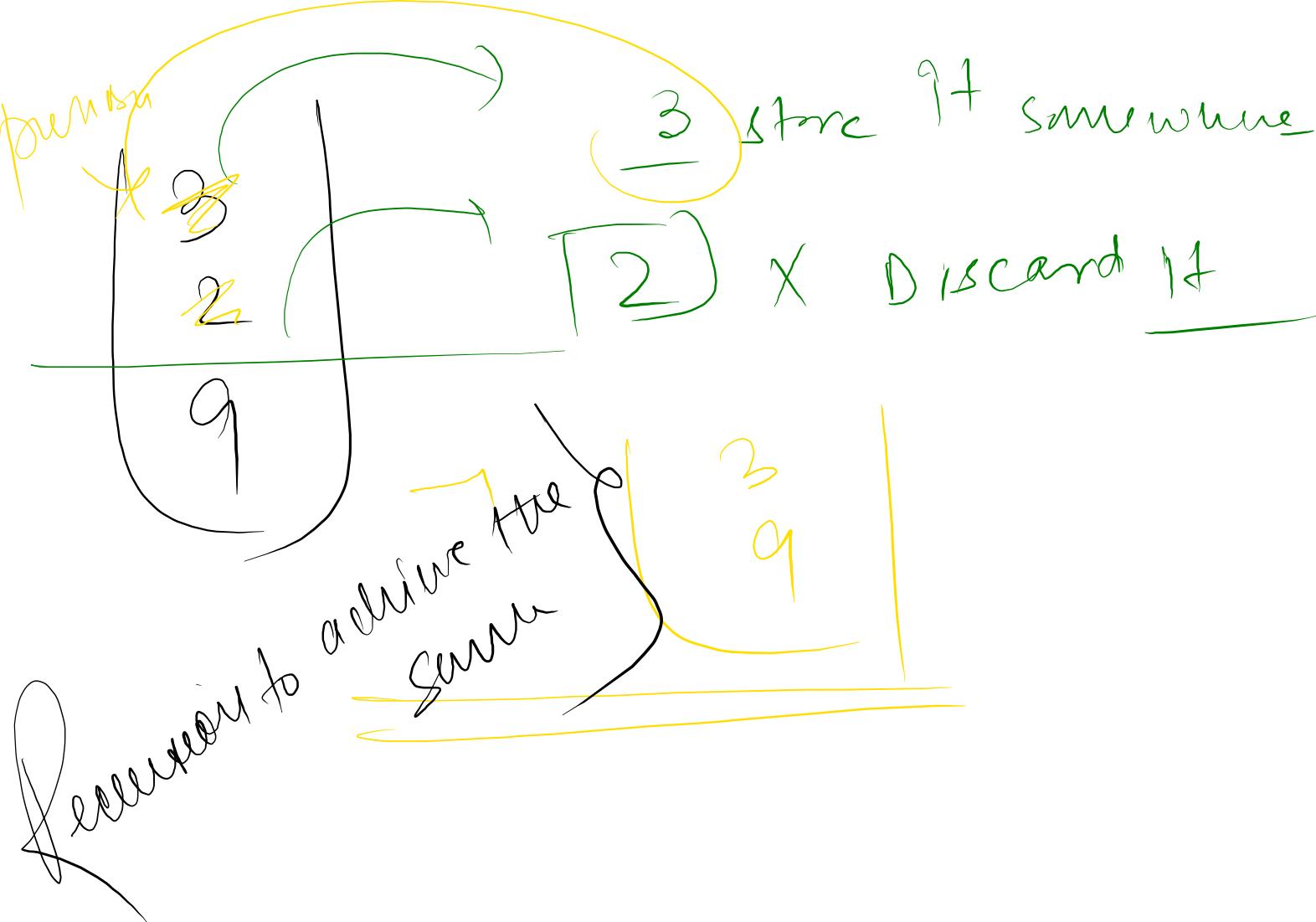
No stack allowed

Stack

push (pop)







```

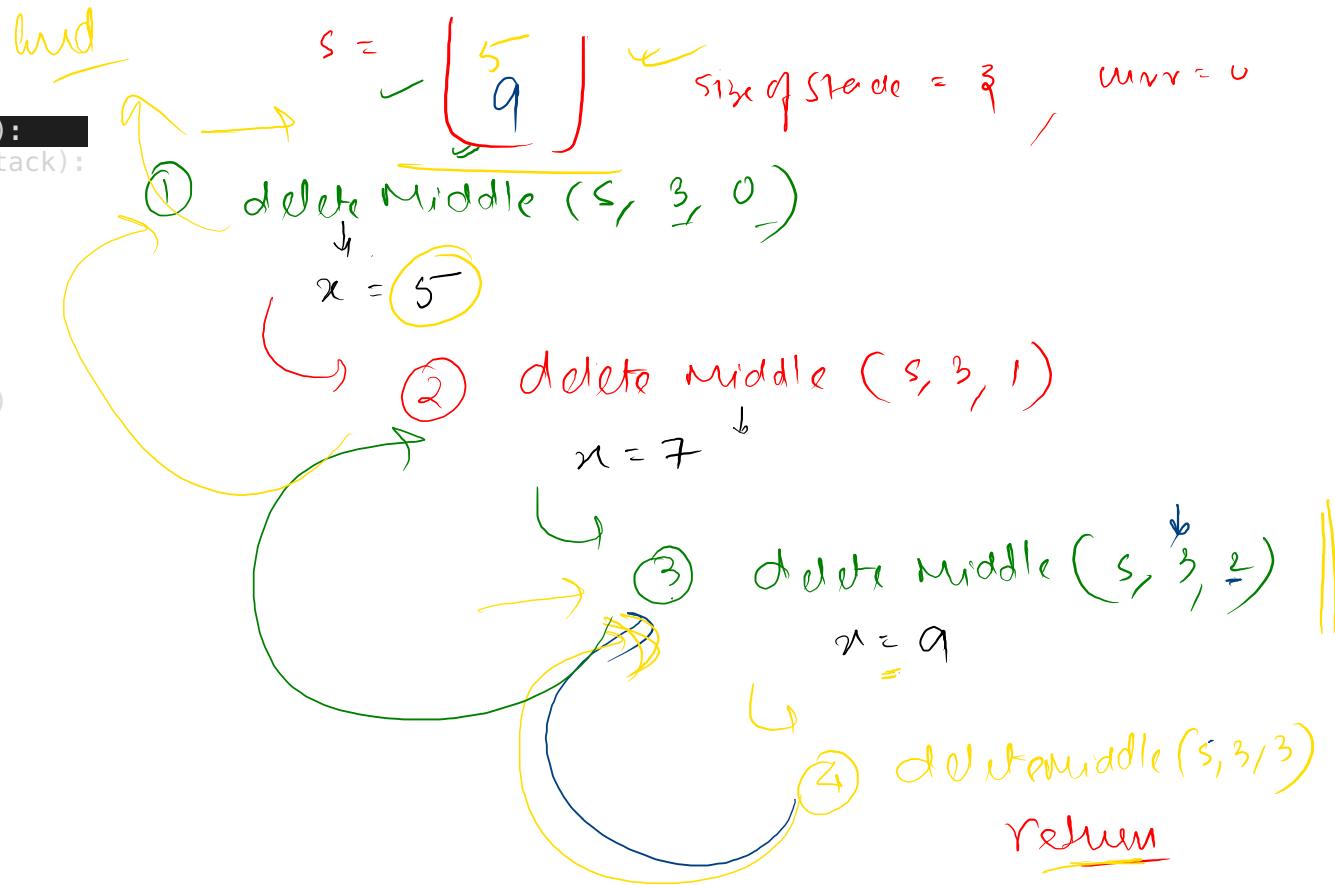
def deleteMiddle(s, sizeOfStack, curr):
    if (s.isEmpty() or curr == sizeOfStack):
        return

    # Remove the current item
    x = s.pop()

    # Recursively call the fn
    deleteMiddle(s, sizeOfStack, curr+1)

    # Put the item back
    if(curr != int(sizeOfStack/1)):
        s.push(x)

```



$f(x)$!

(x) \rightarrow Local variable }

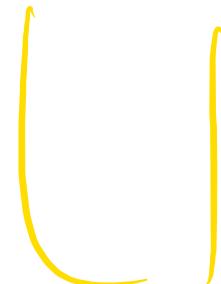


3
Y



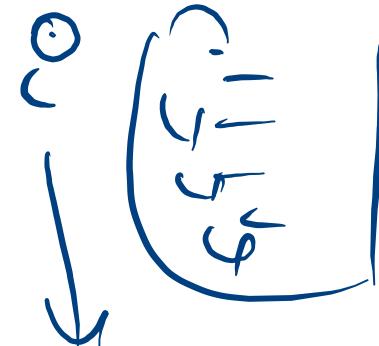
Practise

empty



stacks

+



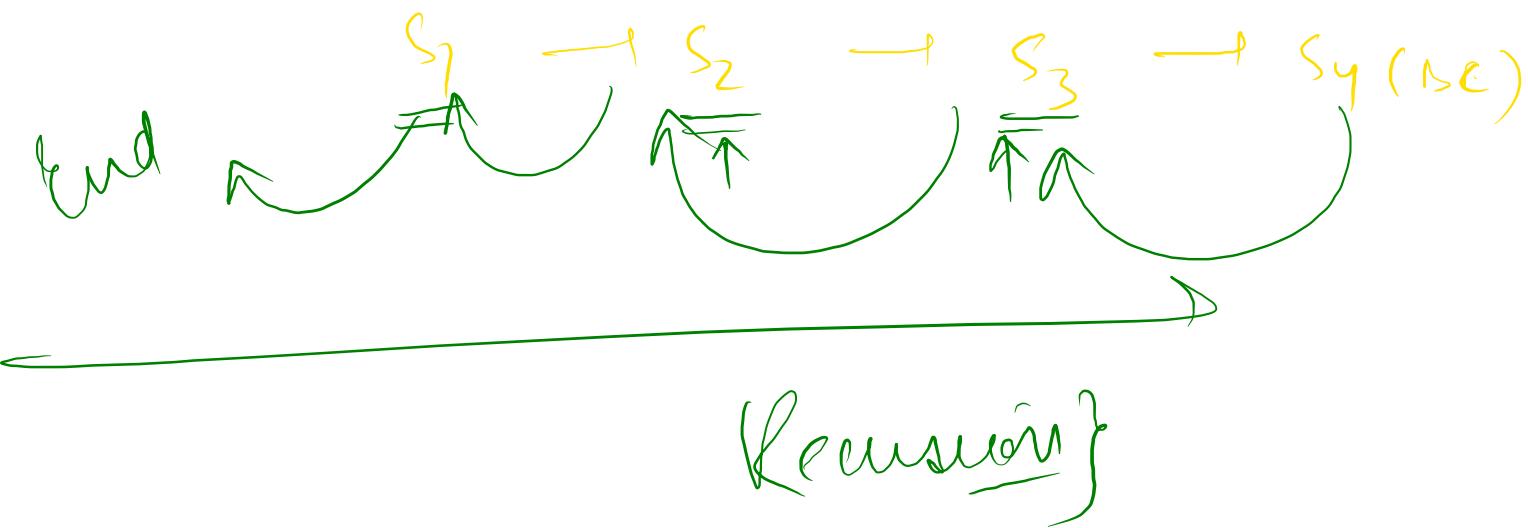
stack

Revised

Reunions

Tree Graph

→ Based at
Revision



TMUC

SC

Chana 18

Stack } → Array
 (pop)

10:50 AM

No Breaks

10:55 AM

(Queue)

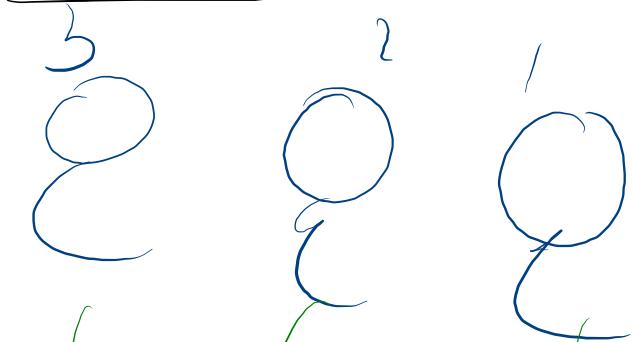
Queue DS

FIFO / LIFO

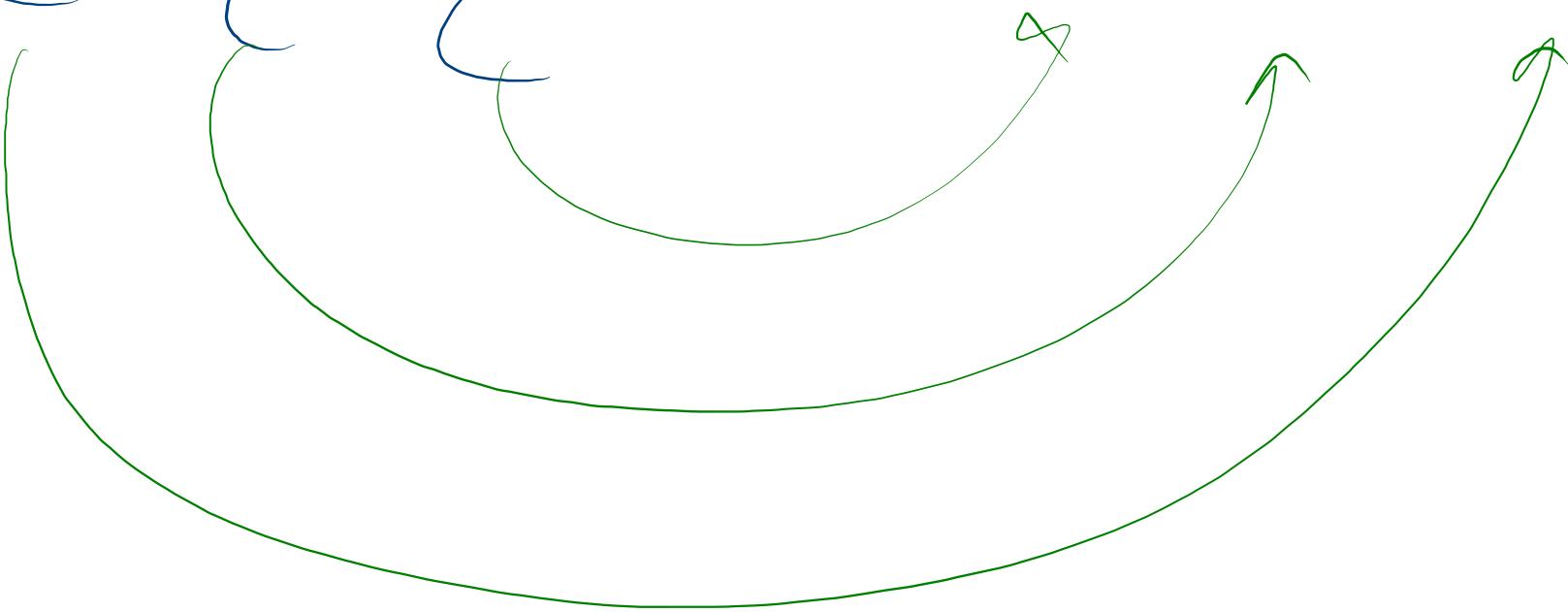


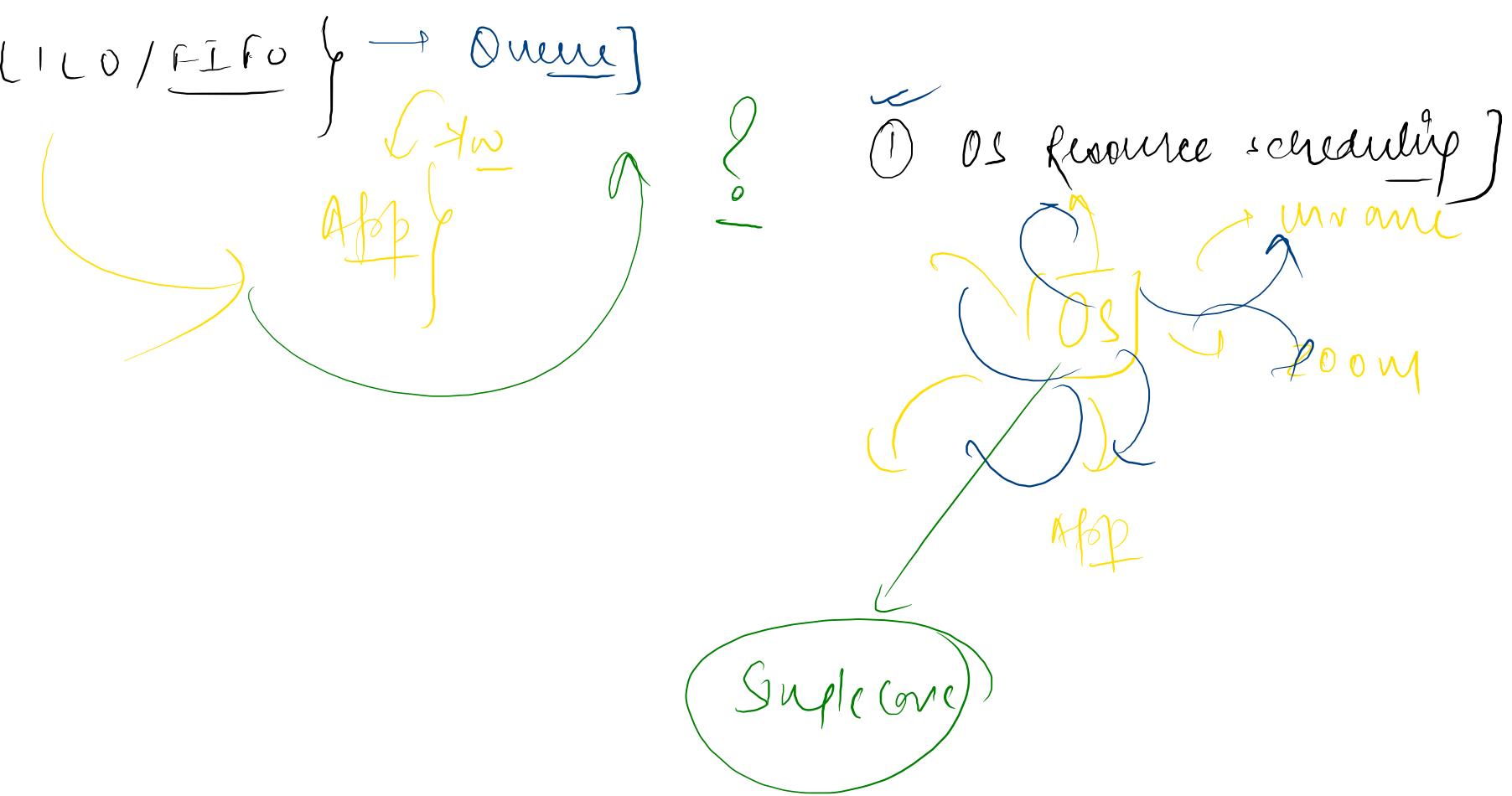
Queue

Queue of FIFO



Bank
Deposit money

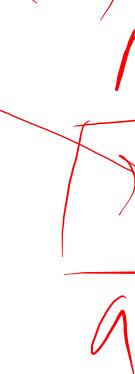




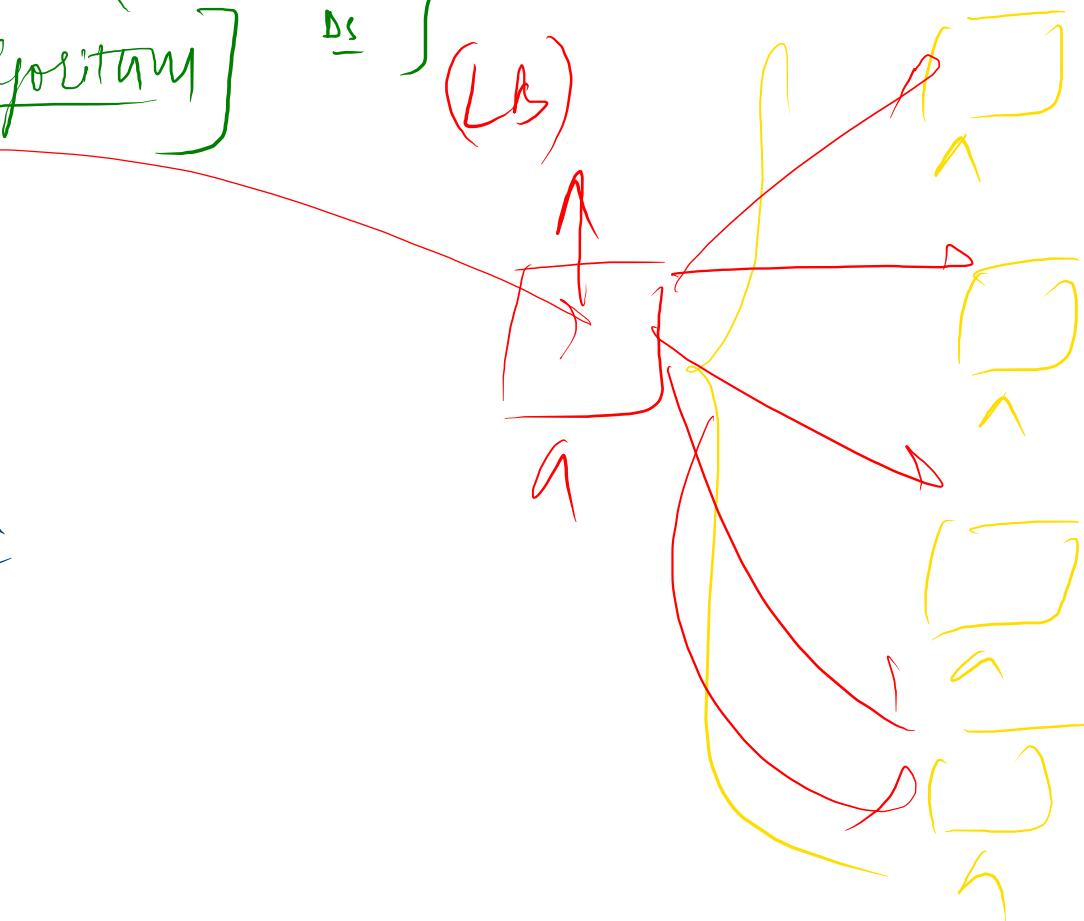
②

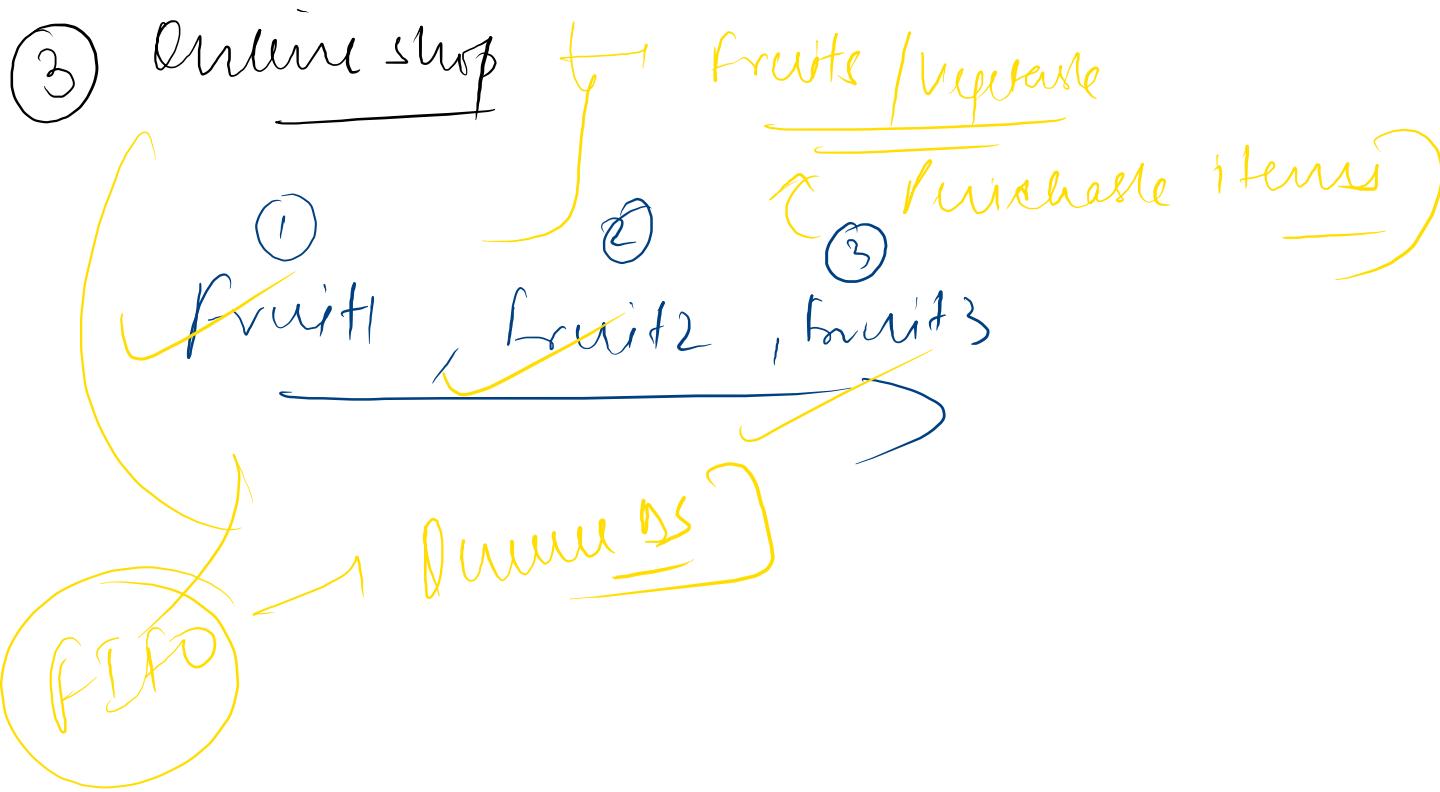
Load balancers {
↳ KR algorithm}

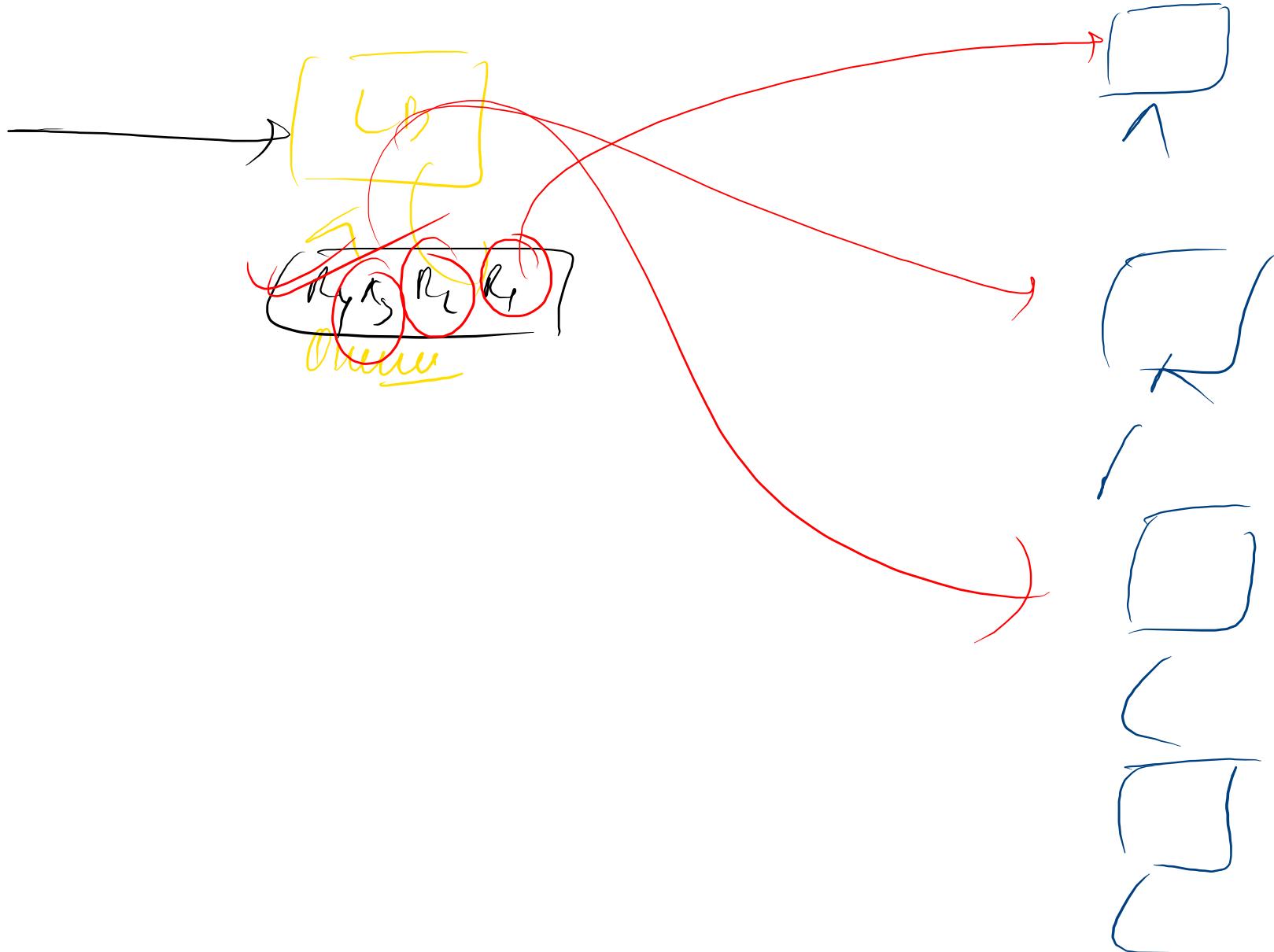
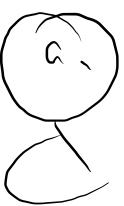
(Lb)



India | Bandwidth



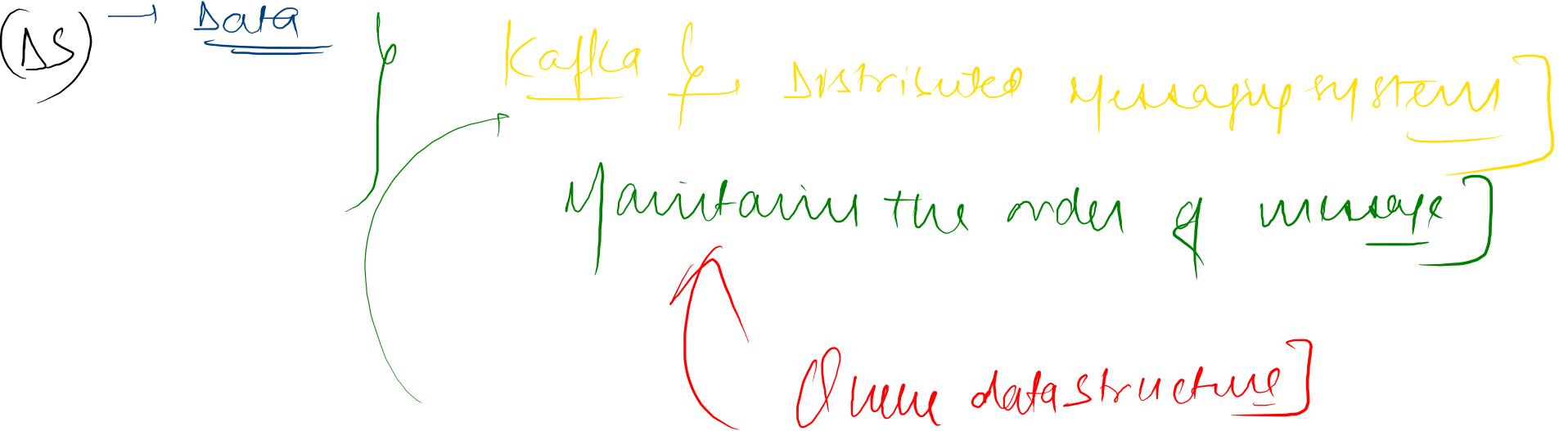


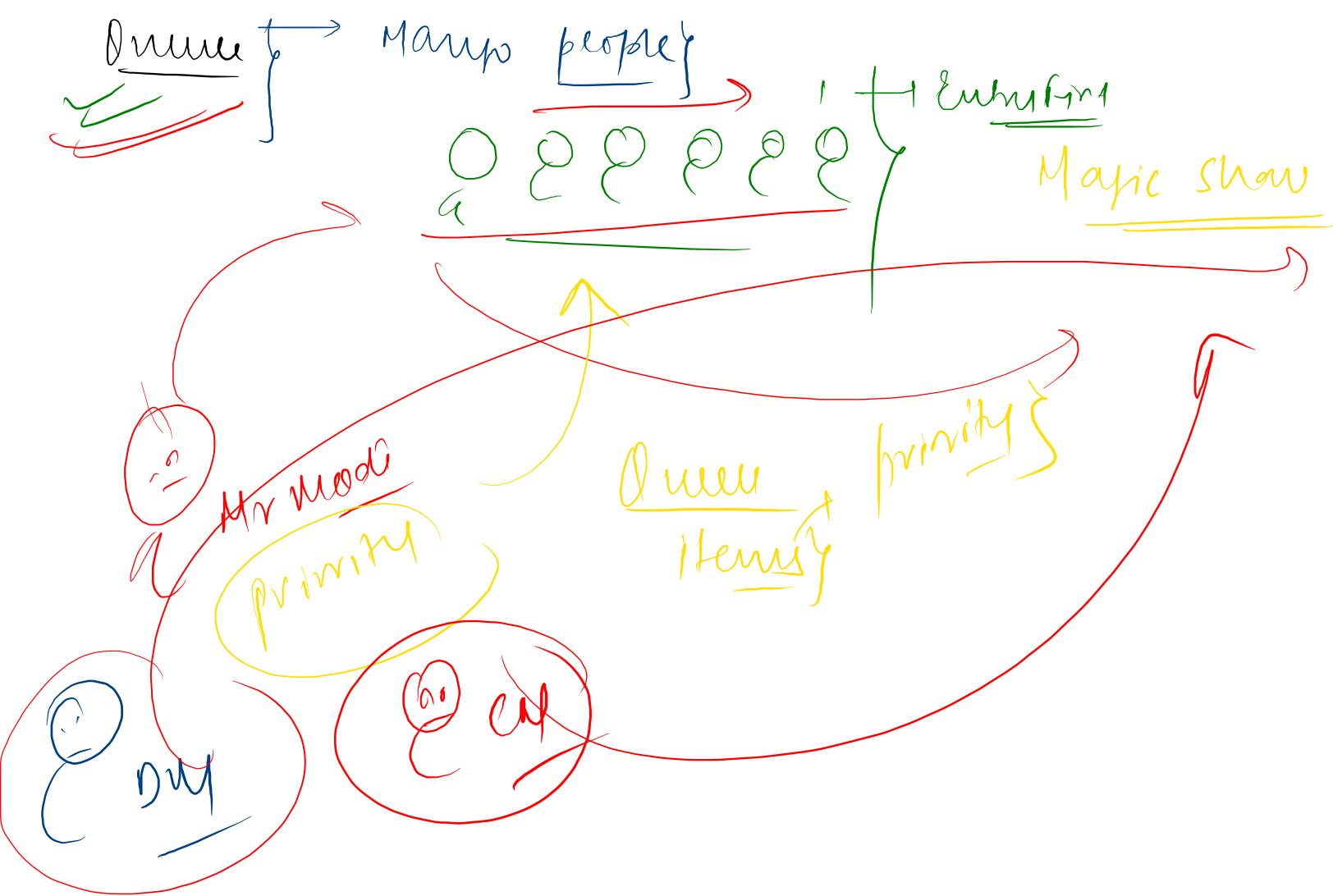


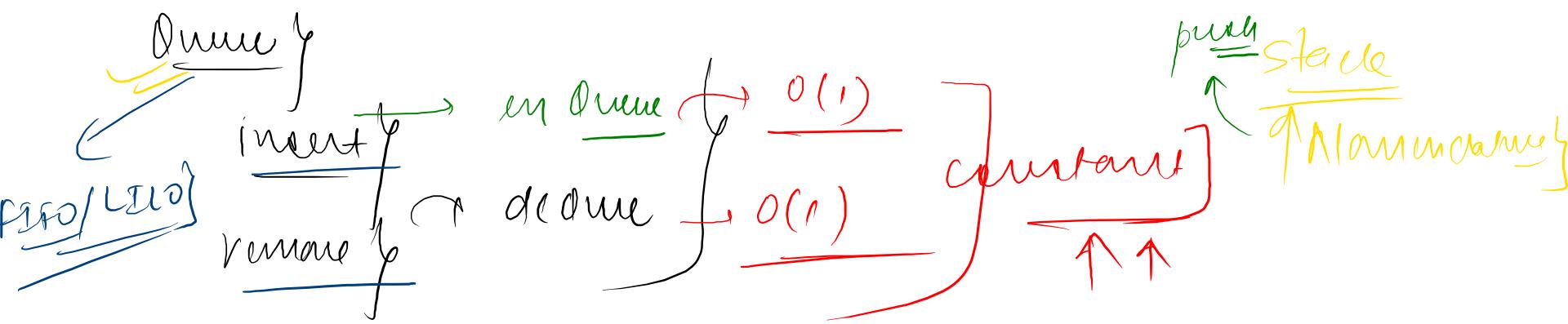


$\tau \propto n$

$\tau \propto \eta_h$







12c²

→ Smartly use Array [] solve this

Q. enqueue(5)

Q. enqueue(13)

Q. enqueue(73)

Q. enqueue(33)

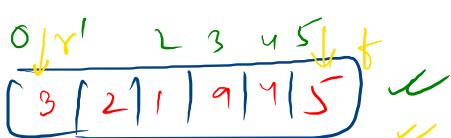
Q. deQueue() → 5

Q. deQueue() → 13

Q. enqueue(23)

Q. enqueue(72)

Q. enqueue(25)



$$\text{Condition } (f+1) \% \text{size} = -\infty$$

$$\Rightarrow (5+1)\%6 = -0$$

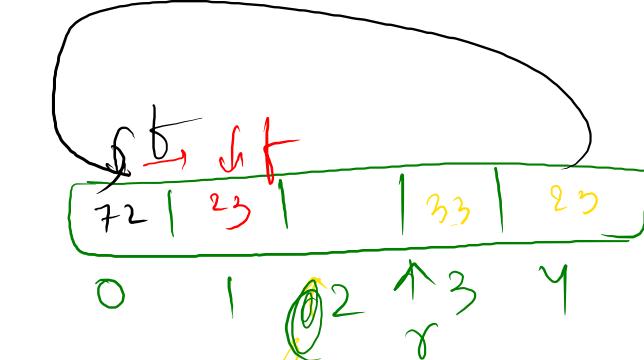
$$\Rightarrow 0 = -0$$

Q. deQueue() → 73

if it's full?

full X

Q. enqueue(51)



$$(f+1) \% 5$$

$$= (4+1)\%5 = 0$$

$$r = 2$$

$$\text{front} = -1 / 0 / 1 / 2 / 3 / 3 / 3 / 4 / (4+1)\%5 = 0$$

$$\text{rear} = -1 / 0 / 0 / 0 / 1 / 2 / 3$$

{ No element }

$$(f+1) \% \text{size} = -\infty$$

$$\Rightarrow (1+1)\%5 = 0 \text{ or } 2$$

- Queue
- ① $f \rightarrow \text{remove}$
 - $r + \text{dequeue}$
 - \rightarrow facts
 - ② empty $\rightarrow f = r = -1$
 - ③ pull $\rightarrow (f+1) \% \text{size} = r \% \text{size}$
- $(r = v \% \text{size})$
- $y r (\frac{\text{size}}{\text{size}})$
- so that index is always in
correct range*

front = 3

capacity = 5

f[0]

f[1] + f[2] = 3
3 + 5 = 8

$$f = 4$$

$$f = f \cdot c$$

$$c = 5$$

$$y = 4 \cdot 5 = 4$$

$$f = 5$$

$$c = 5$$

more
ways
last index

$$f \cdot c = 5 \cdot 5 = 0$$

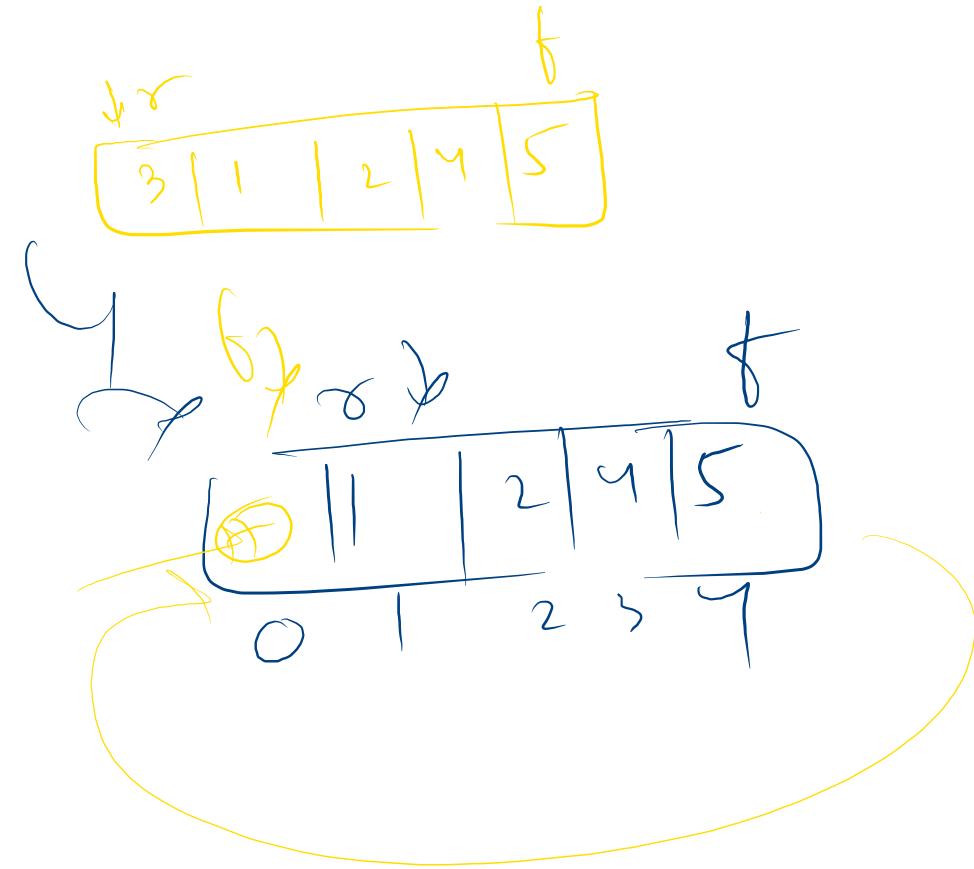
0th index

```

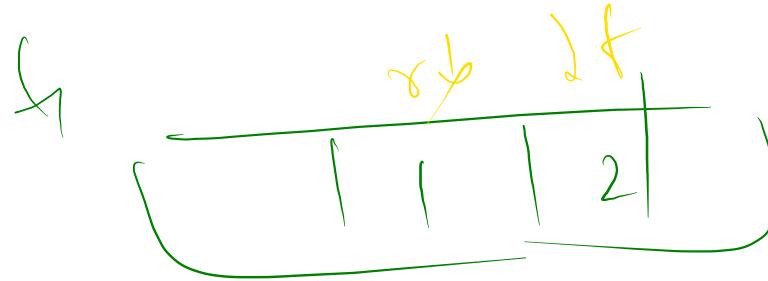
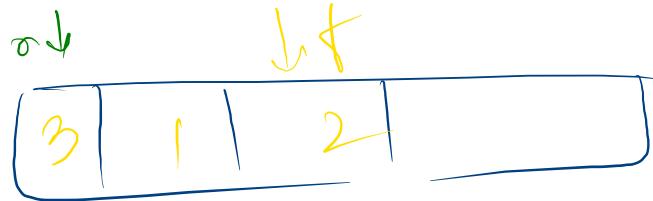
def enqueue(self, data):
    if self.isFull():
        print("Queue overflow")
        return
    # Move the front to the right one position
    self.front = (self.front + 1) % self.capacity # ensure index is correct
    self.arr[self.front]=data

```

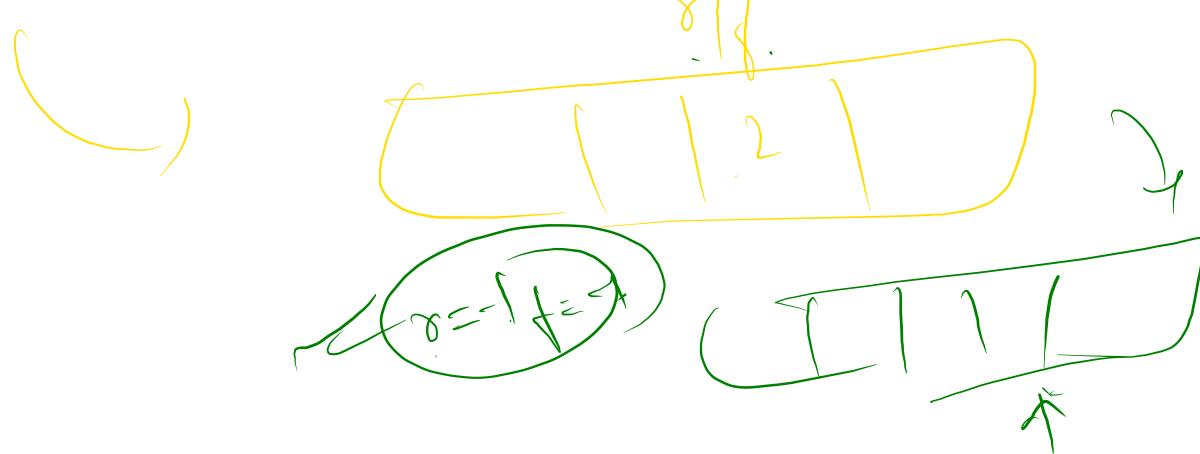
(4+1) / 5
= 5

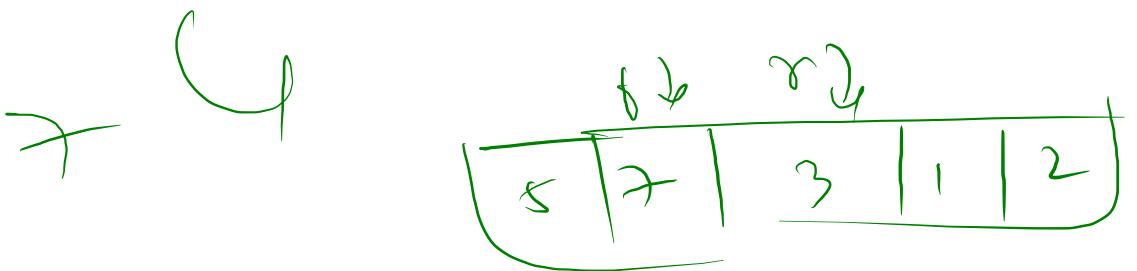
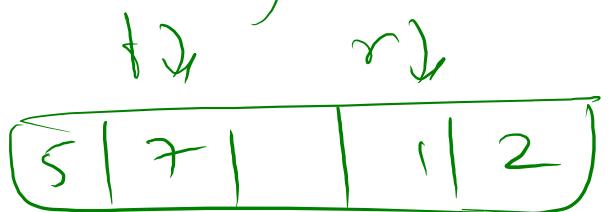
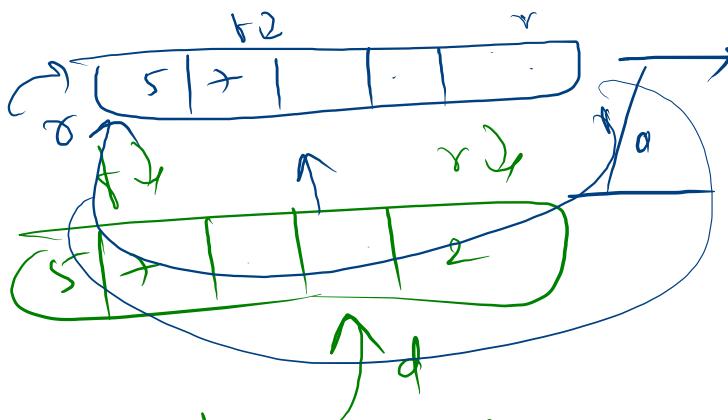
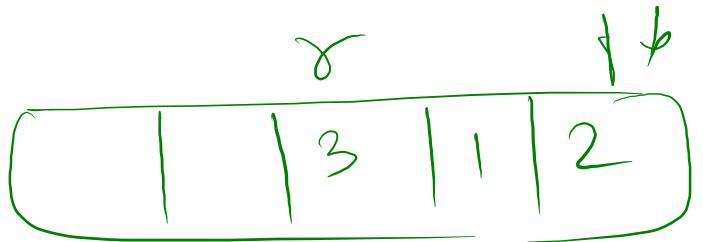


~~move()~~



~~pop()~~





delete

$r = r\% \text{ capacity}$
when $r < \text{capacity}$

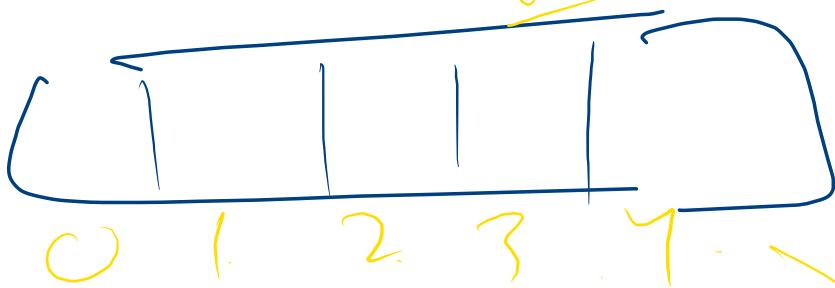
$\gamma \leq \gamma_{cap}$

$(\gamma+1) \leq (\gamma+1)^{cap}$

γ *valid index*

γ *capacity*

$\gamma \leq \gamma_{cap}$



$r=4$

γ
 $(\gamma+1)$

$(\gamma+1)$
 γ

```

class Queue :
    def __init__(self, capacity):
        self.arr = [None]*capacity
        self.capacity = capacity
        self.front = -1
        self.rear = -1 #in the begining both front and rear should be -1
        self.size = 0

    def isEmpty(self):
        return self.front == -1

    def isFull(self):
        return (self.front +1)%self.capacity == self.rear

    def enqueue(self, data):
        if self.isFull():
            print("Queue overflow")
            return
        #Move the front to the right one position
        self.front = (self.front + 1) % self.capacity # ensure index is correct
        self.arr[self.front]=data

        if self.rear == -1: #when the first element is inserted
            self.rear=0
        self.size +=1

    def dequeue(self):
        if self.isEmpty():
            print("Queue underflow")
            return
        data = self.arr[self.front]
        #Check if this was the only element
        if self.rear == self.front:
            self.rear = -1
            self.front = -1
            #Because it was the only element, and after removal it's empty
        else:
            self.rear = (self.rear+1)%self.capacity #ensure we are in correct index range
            self.size -=1
        return data

    def size(self):
        return self.size

    def traverse(self):
        for i in range(self.size):
            print(self.arr[self.rear+i]%self.capacity)

```

$q = \text{Queue}(5)$
 $f = 0$ $r = 0$
 $\text{size} = 0$
 $\text{capacity} = 5$
 $\boxed{5 \mid \text{None/None/None/None/None}}$

$\text{isEmpty}() \rightarrow \text{True}$

$\text{isFull}() \rightarrow \text{False}$

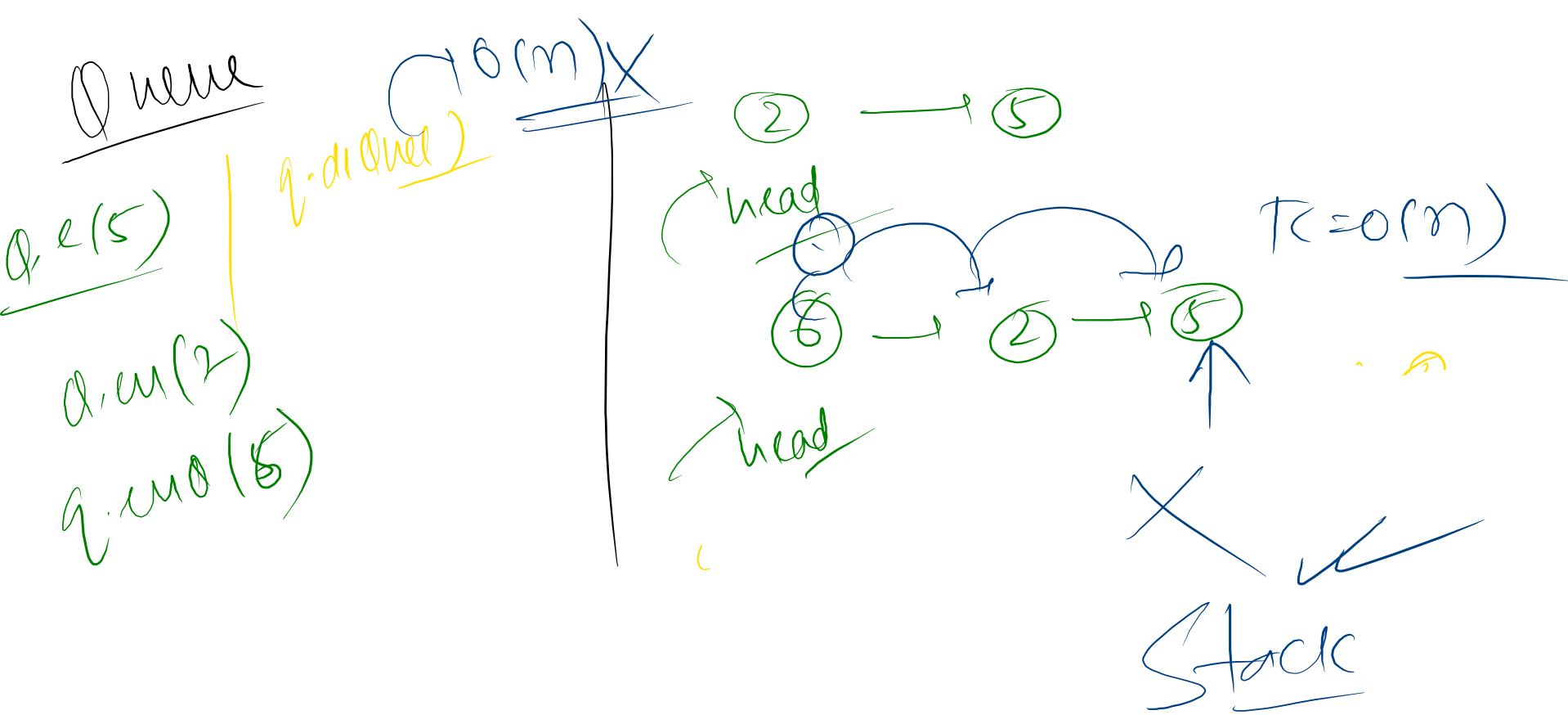
$\text{enqueue}(5)$
 $\text{dequeue} \rightarrow 5$

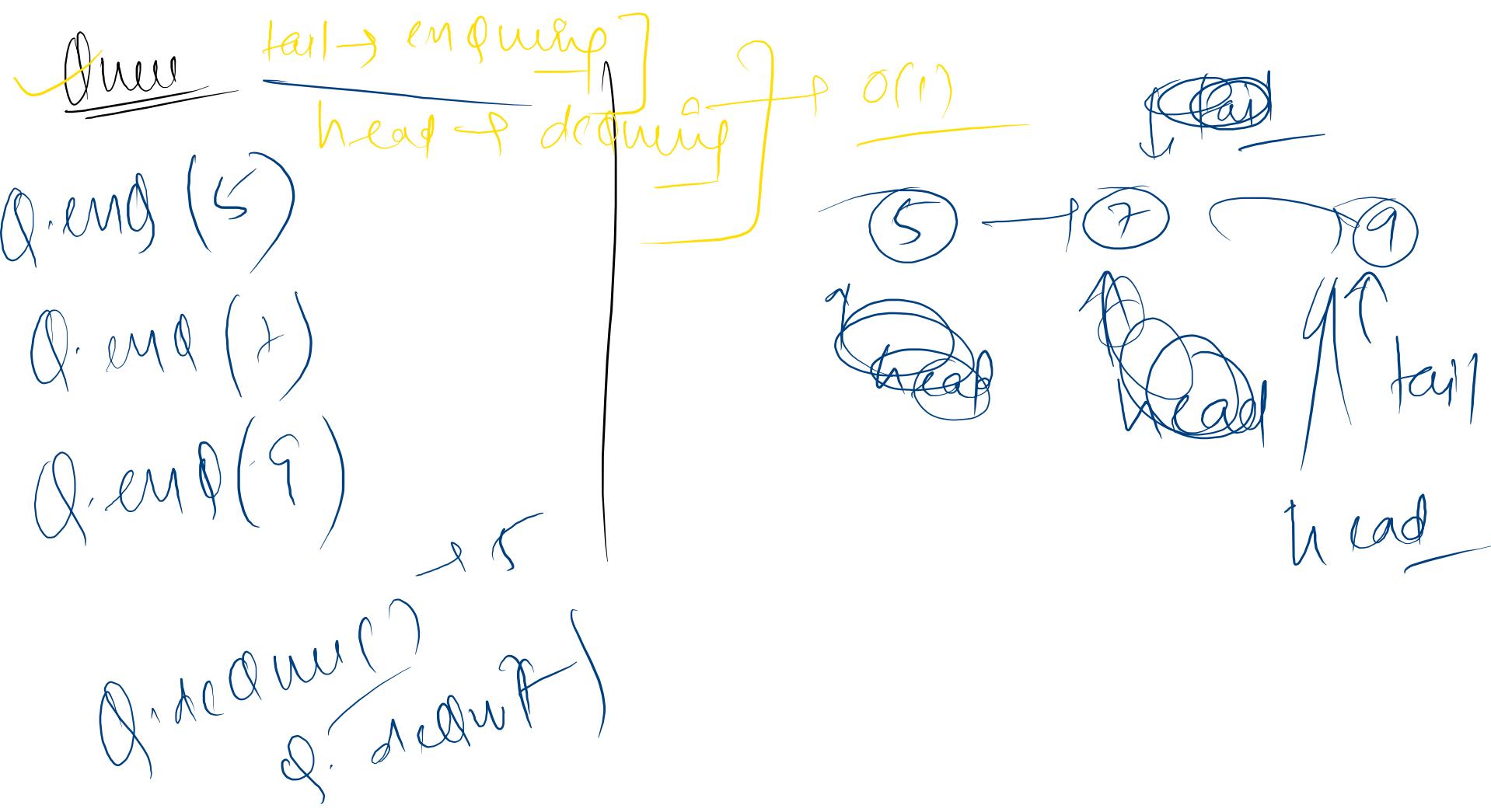
$\text{isEmpty}() \rightarrow \text{True}$

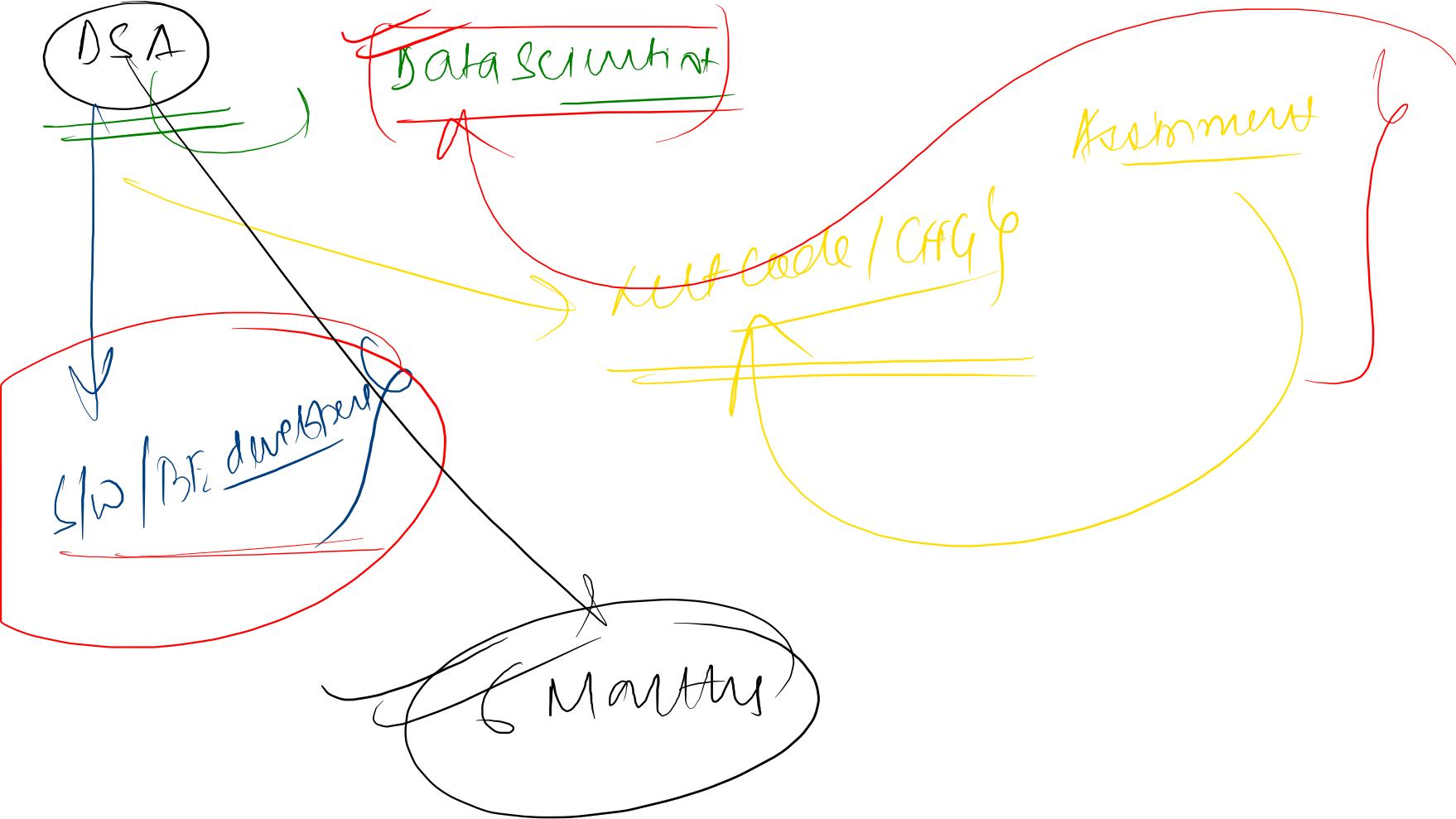
Implement Queue with C

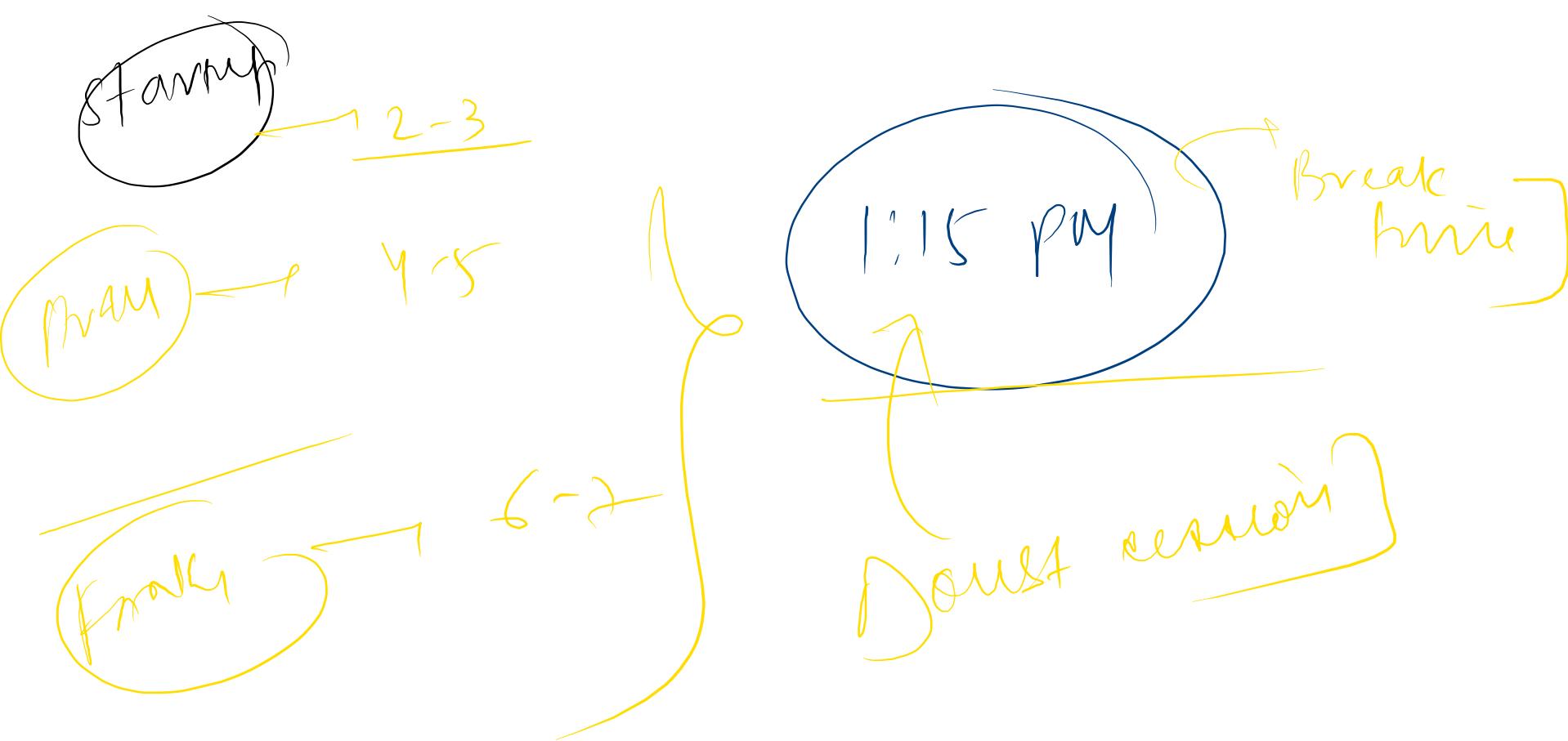
o—
o—

80
→ Treat me :)

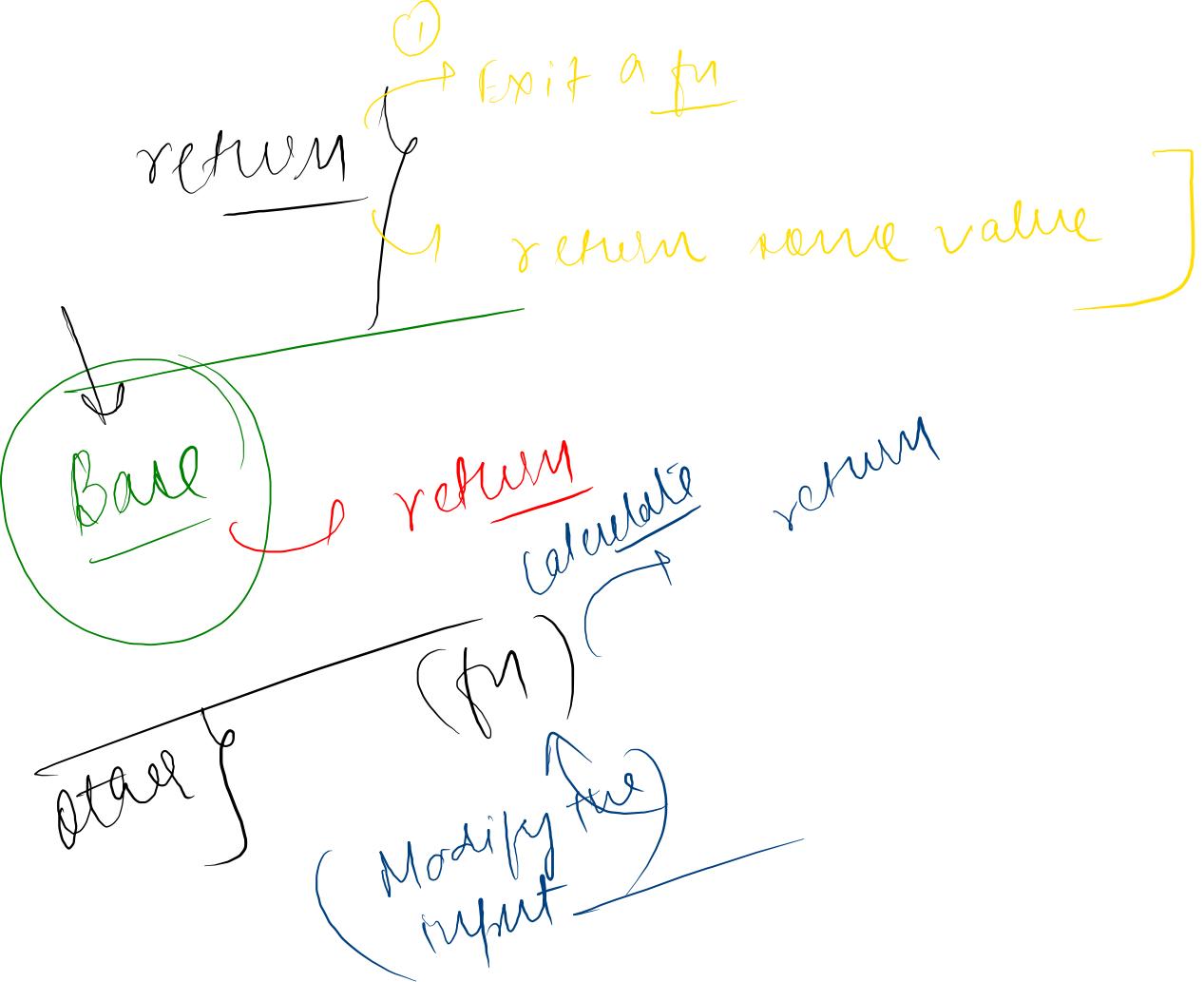









```
def revv_rec(arr,si,ei):      #Why this is not working
    if si>ei:
        return
    else:
        return revv_rec(arr,si+1,ei-1)
        arr[si],arr[ei]=arr[ei],arr[si]
```



```
def findrotate_ii(arr):
    start=0
    end=len(arr)-1

    if arr[start]<arr[end]:
        return 'array is not rotated'
    while start<=end:
        mid=(start+(end-start)//2)

        if arr[mid]<arr[mid+1] and arr[mid]<arr[mid-1]:
            return mid
        elif arr[mid]<arr[end]:
            #start=mid+1
            end=mid-1
        else:
            arr[mid]>arr[start]
            #end=mid-1
            start=mid+1
    return
```

```
def iterative_function():
    i = n
    while(i>2):
        i = i^(1/25)
```

$$\left(\left(a^{\frac{1}{25}} \right)^c \right) \in a^{\frac{c}{25}}$$

$T_C = \text{No of operations}$

$$\left\{ \frac{1}{25} \right\}$$

What is the time complexity of the 'iterative_function' in terms of Big O notation?

- $O(\log_2(\log_2 n))$
- $O(\log_2(\log_2 n))$
- $O(\log_2(\log_2 n))$
- $O(\log n)$

\downarrow
Shows $\log_2(\log_2 n)$

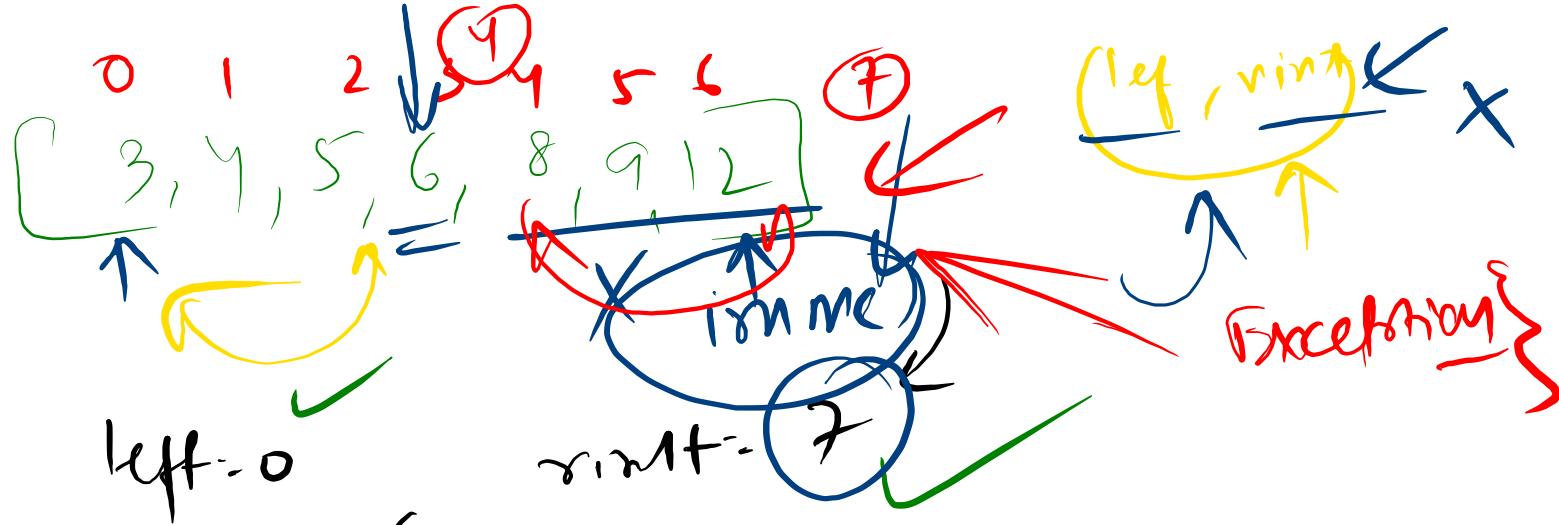
$$\begin{aligned} (2^5)^2 &= (8)^2 = 64 \\ (2)^c &\in 2 \times 2 \times 2 \times 2 \times 2 \times 2 \end{aligned}$$

$i \leftarrow 1$
1st
 $i = n$
2nd
 $i = (i)^{\frac{1}{25}}$
3rd
 $i = [(i)^{\frac{1}{25}}]^{\frac{1}{25}}$
nth
 $i = ((i^{\frac{1}{25}})^{\frac{1}{25}})^{\frac{1}{25}}$

$$k^{\text{th}} \equiv i^{\frac{1}{25^{k-1}}} \quad | \quad (k+1)^{\text{th}}$$

$$\Rightarrow i^{\frac{1}{25^k}} \propto n$$

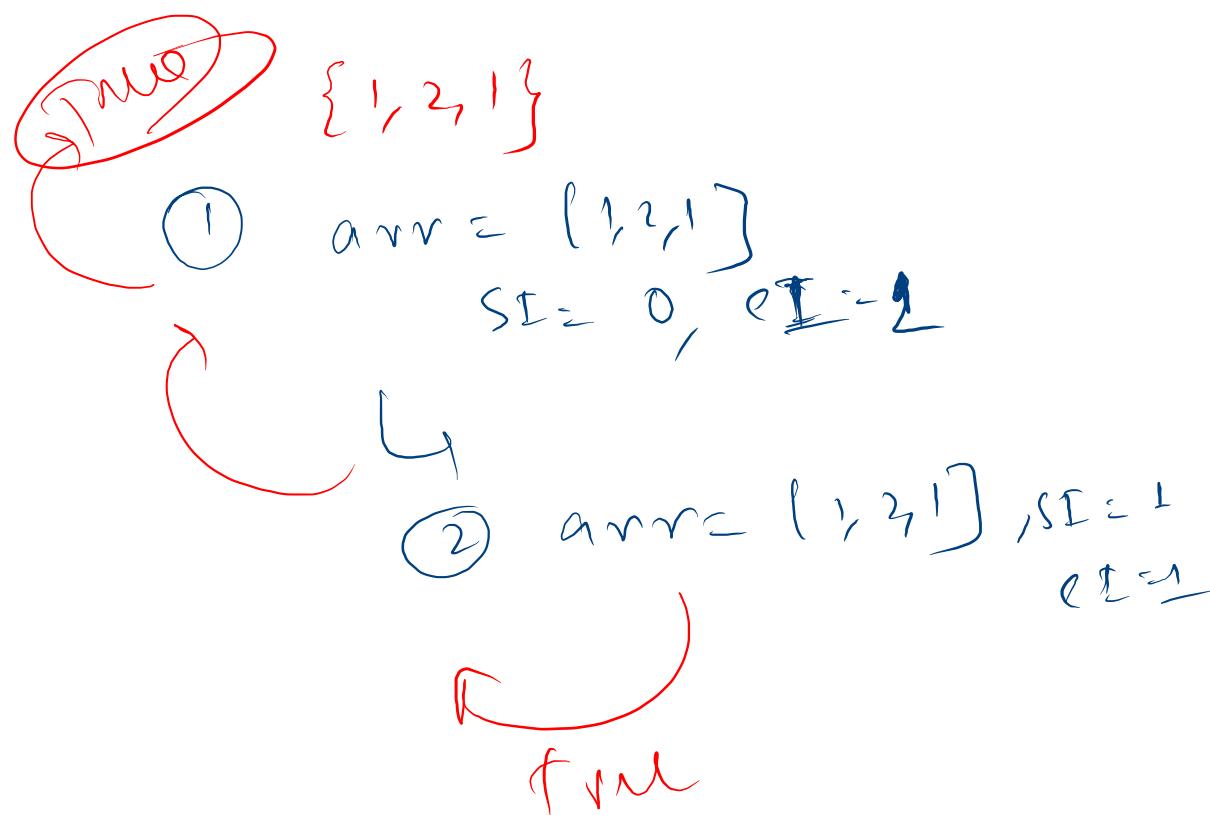
$$\begin{aligned} \text{WPS}_n &= \text{WPS}_{25} \\ \text{WPS}_{25} &= \text{WPS}_{25} (\text{WPS}_n) \end{aligned}$$



$$\frac{6+10}{2} = \textcircled{3}$$

~~fn(left, right)~~

```
def isPalindrome_rec(arr,sI,eI):
    # Base Case
    if sI>eI:
        return False
    if sI==eI:
        return True X
    # Logic
    if arr[sI]!=arr[eI]:
        return False
    # Recursive Call
    return
isPalindrome_rec(arr,sI+1,eI-1)
```



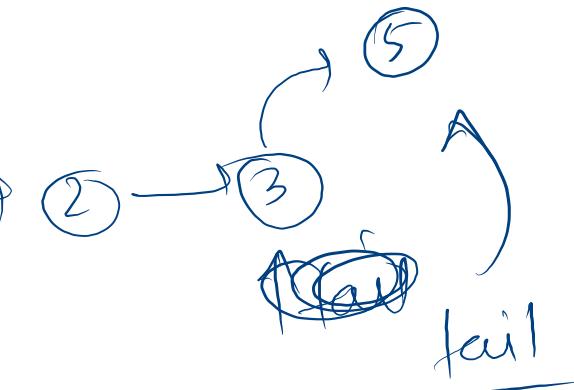
```
else :  
    self.tail.setNext(node)  
    self.tail = self.tail.getNext()  
    self.size += 1  
    can you explain this sir once , misse yestedays class , just this  
much rest I got
```

increasing the size by 1

queue

5

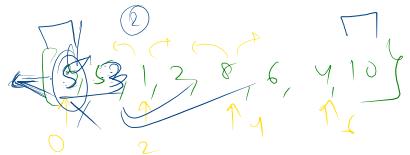
head



```

arr=[3,5,1,2,8,6,4,10]
def sortinwave(arr):
    for i in range (0,len(arr)-1,2): index 1 6 #Rather than mentioning can we
    start range from 2
        if i>0 and arr[i]>arr[i-1] index 1 #even if do not add condition i>0 and
        i<len(arr)-1 am not getting
            arr[i],arr[i-1]=arr[i-1],arr[i] #out of index error why? see below
    code and output
    elif i<len(arr)-1 and arr[i]>arr[i+1]: #code is not working see 15 and
26
        arr[i],arr[i+1]=arr[i+1],arr[i]
    return arr

```



Node

node = Node(5)



node.set data(6)



/

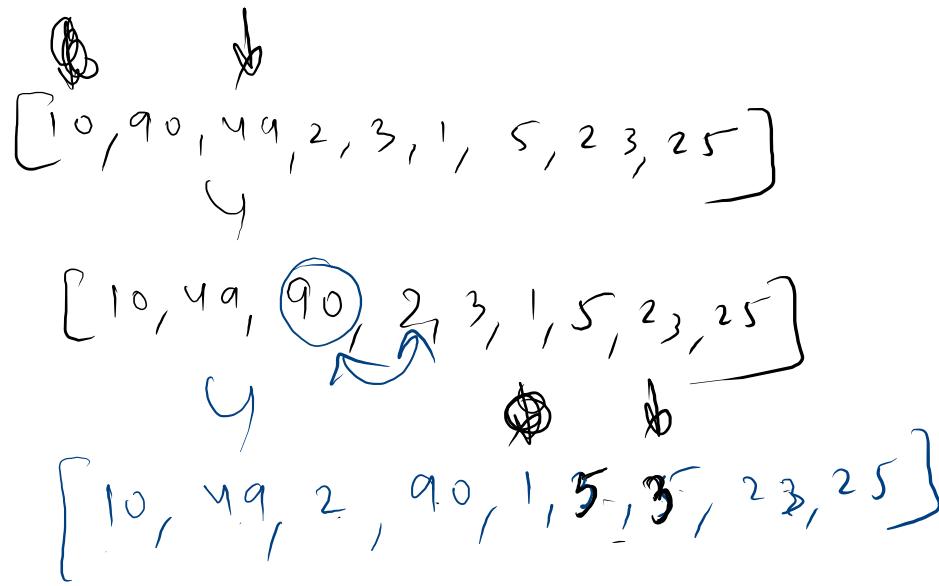
```

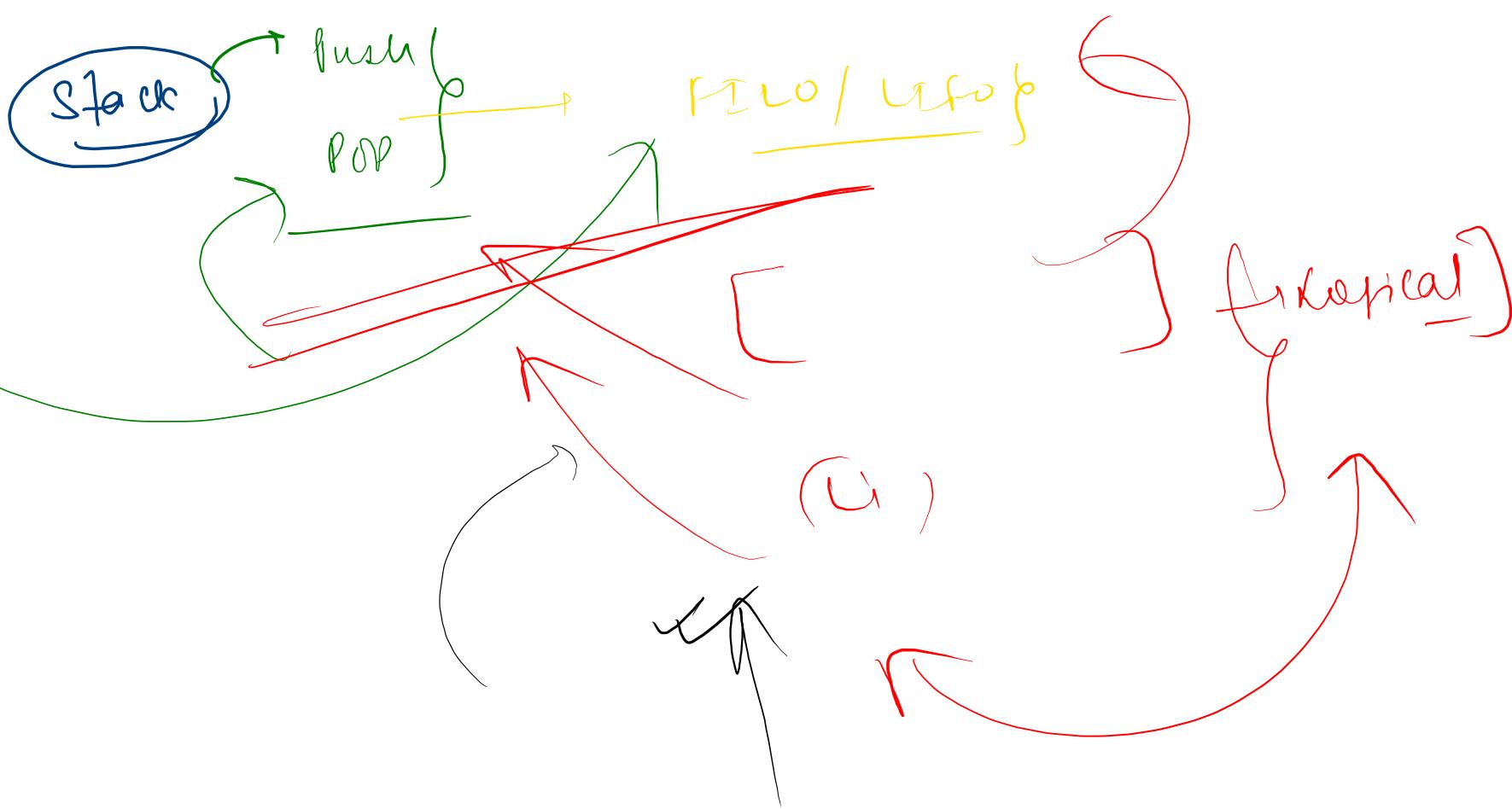
def sortInWave(arr):
    ← 6-8
    for i in range(0, len(arr) ← 2):
        #Comparing with left element
        if i>0 and arr[i] > arr[i-1]:
            arr[i], arr[i-1] = arr[i-1], arr[i]

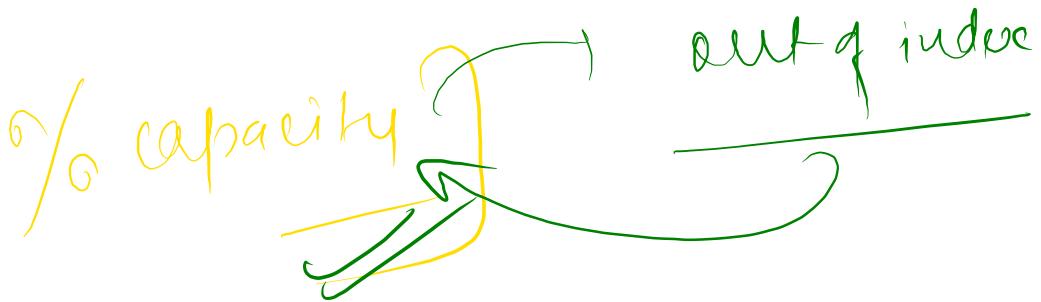
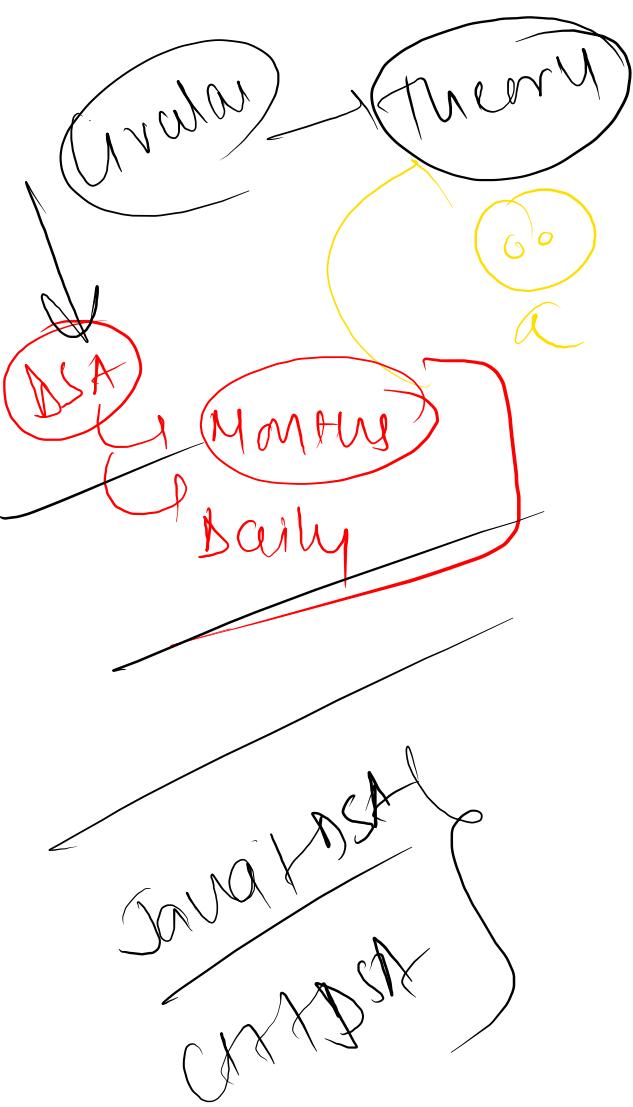
        #Compare with the right element
        if i<=len(arr)-1 and arr[i] > arr[i+1]:
            arr[i], arr[i+1] = arr[i+1], arr[i]

arr = [10, 90, 49, 2, 3, 1, 5, 23, 25]

```

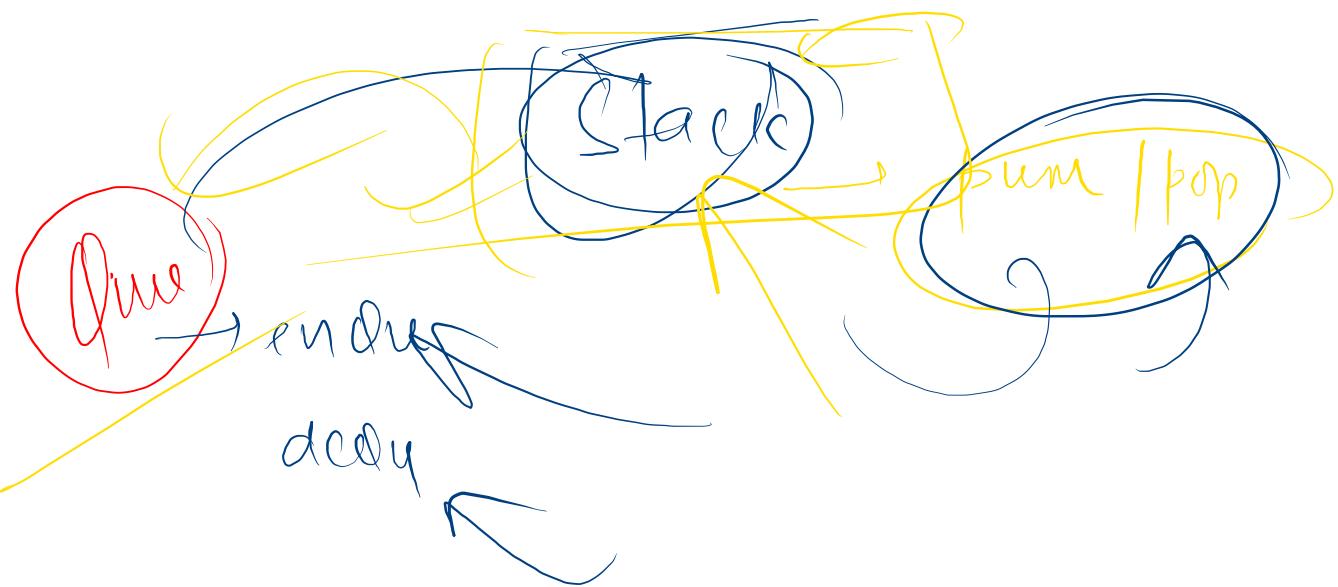


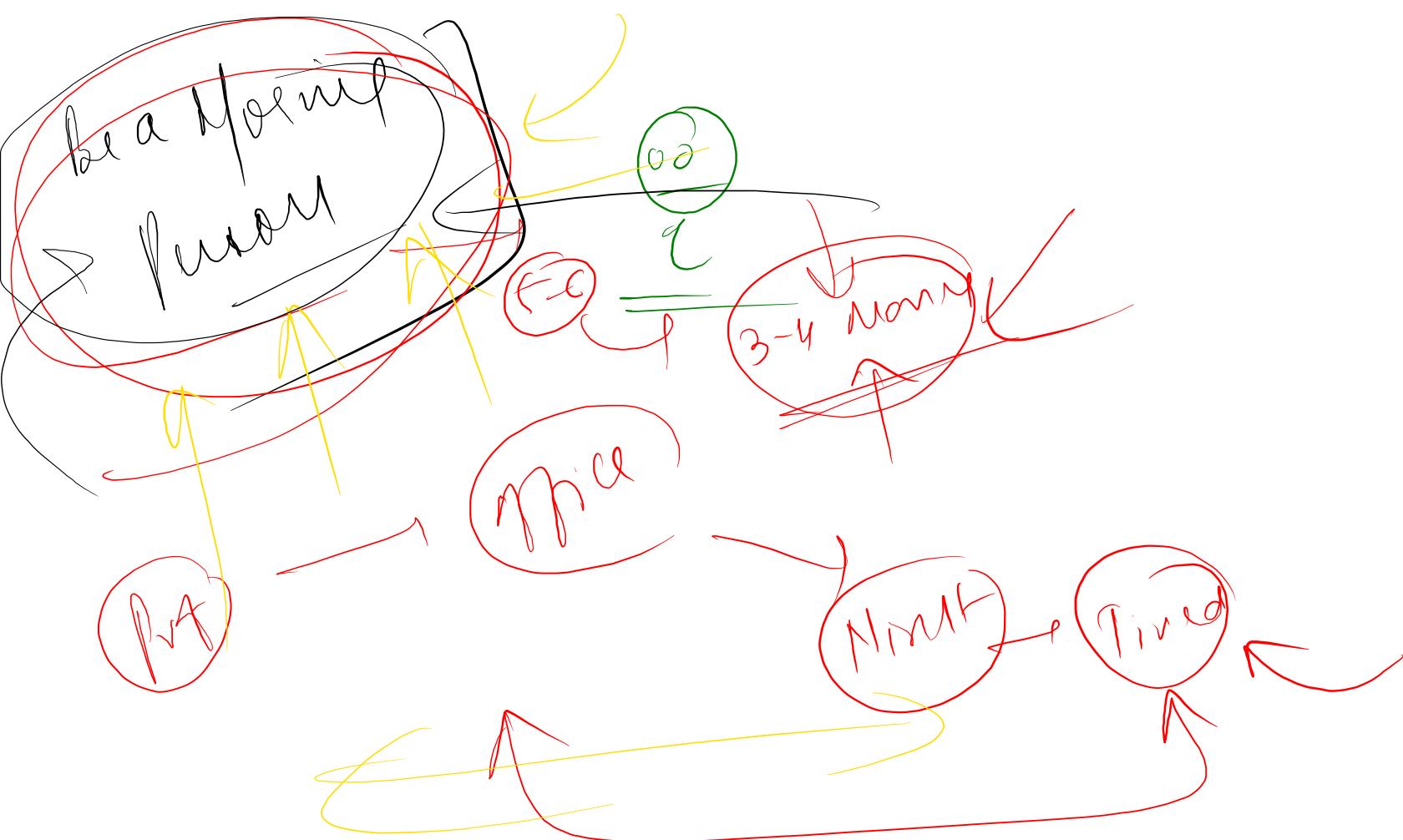




~~Python + DSA~~

Big product Bonded ↗





DSA

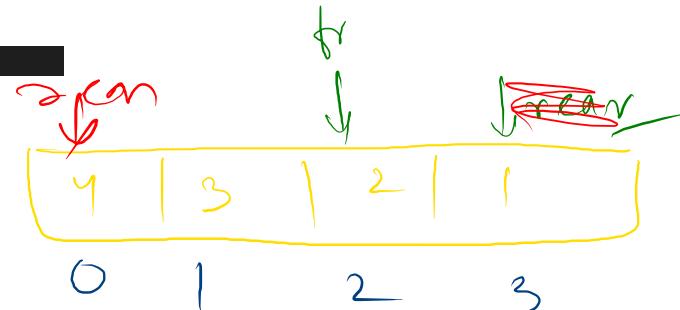
everydays

2 days/ week


```

def deQueue(self):
    if self.isEmpty():
        print("Queue underflow")
        return
    data = self.arr[self.rear]
    #Check if this was the only element
    if self.rear == self.front:
        self.rear = -1
        self.front = -1
        #Because it was the only element, and after removal it's empty
    else:
        self.rear = (self.rear+1)%self.capacity #ensure we are in correct
        index range
        self.size -=1
    return data

```



$$3+1$$

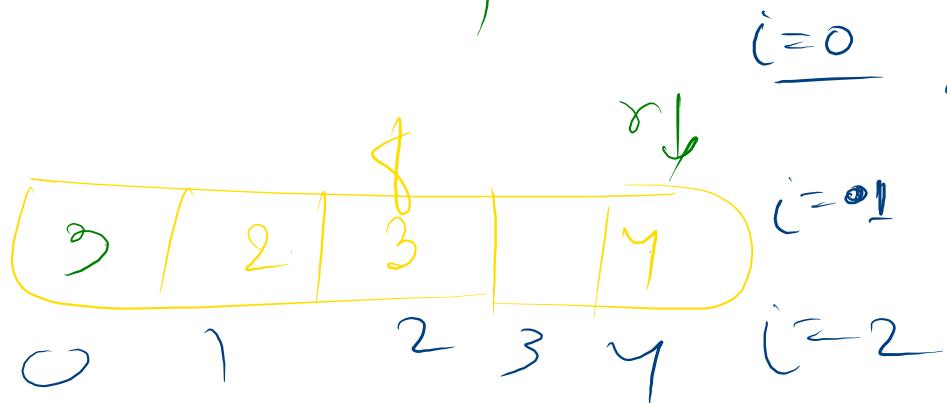
$$= 4$$

Beyond valid index

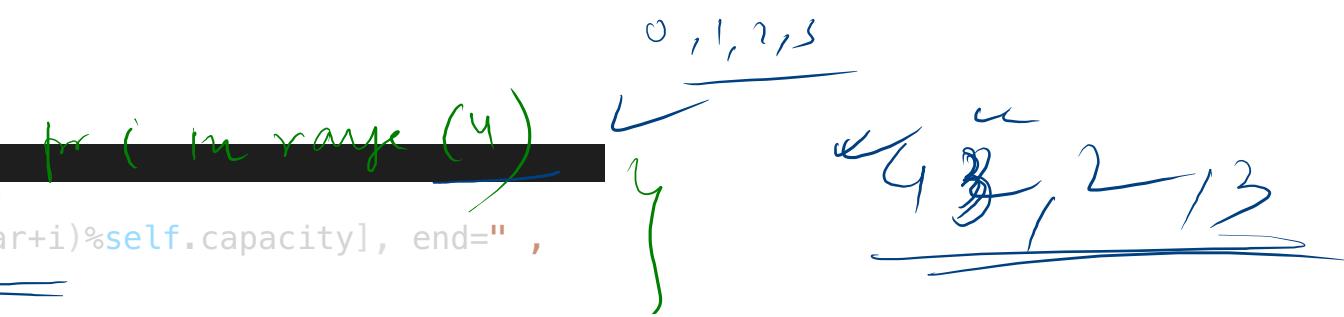
$$4 \% 4$$

$$= \underline{0}$$

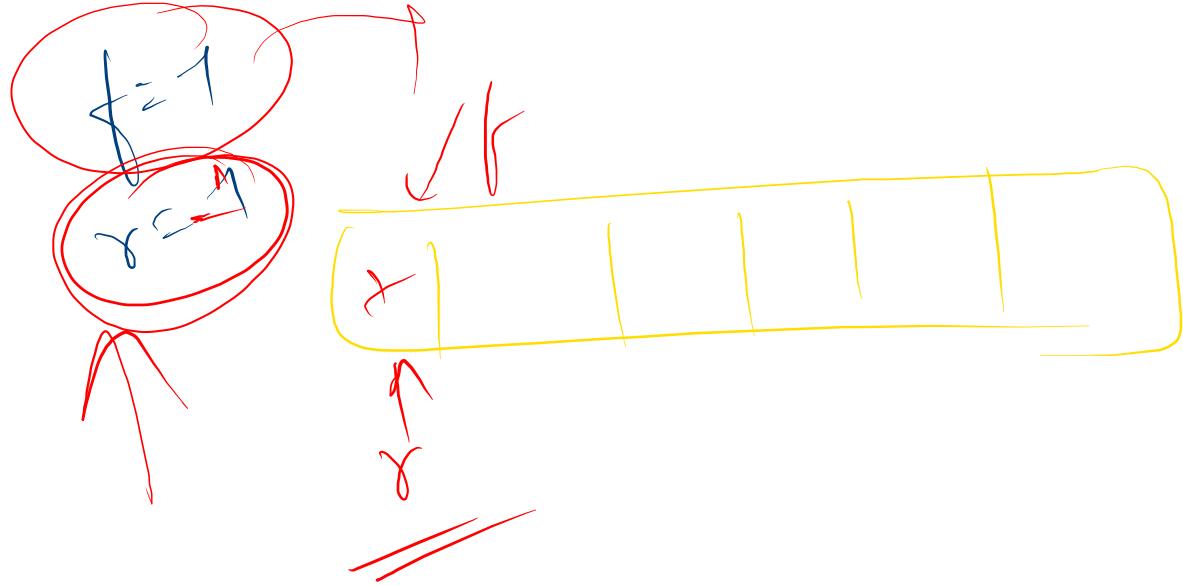
```
def traverse(self):
    for i in range(self.size):
        print(self.arr[(self.rear+i)%self.capacity], end=" ",
```



$$\begin{aligned}
 & \text{arr}[0+0/5] = \text{arr}[0] \\
 & \text{arr}[0+1/5] = \text{arr}[1] \\
 & \text{arr}[0+2/5] = \text{arr}[2] \\
 & \text{arr}[0+3/5] = \text{arr}[3]
 \end{aligned}$$

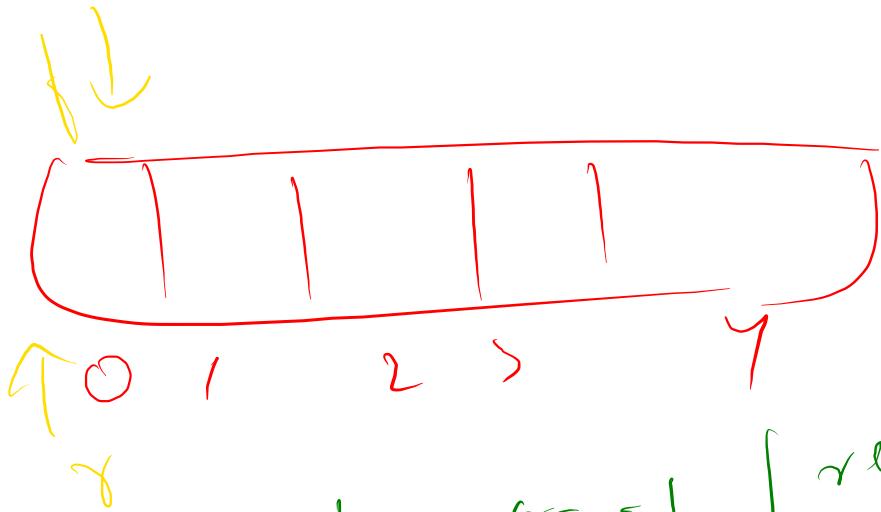


$\partial_{\gamma} \psi(\gamma)$

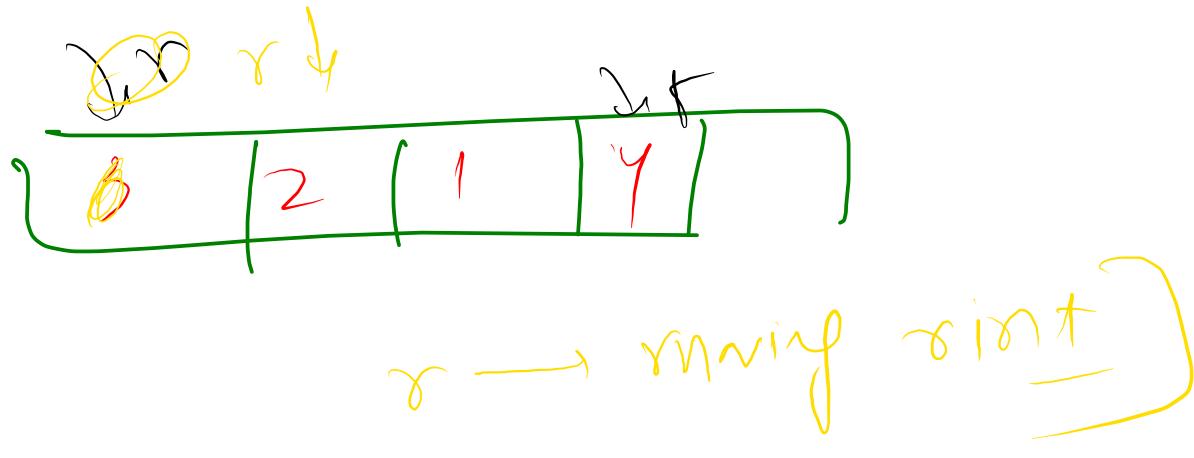


drawn

$f^{-1} = -x$



$f = -x$ | $f \rightarrow$ next to +
| $f = -x$ | $f \rightarrow$ next to +
why ! Just present



Circular array
implementation

(ans)

0	1	2	3	4	5	6
5	7	11	10	21	33	

↑
 γ

⑥ ↑ artq
index
value

$q \cdot F_2(5)$

$q \cdot F_2(7)$

$q \cdot F_2(11)$

$q \cdot F_2(10)$

$q \cdot F_2(21)$

$q \cdot F_2(33)$

$$f+1 = \gamma$$

$$\textcircled{6} = 0$$

$$(f+1)_{\gamma, \textcircled{6}} = \gamma \text{ y. size}$$

$$\Rightarrow 5_{\gamma, 6} = 0 \text{ y. 6}$$

$$= 0 = \underline{0} \quad \uparrow \text{full}$$

$q' E(7)$

$q' E(11)$

$q' E(13)$

$\cancel{q' E(51)}$ $\cancel{n'}$

1	2	3
73	23	31

pr

$q' E(73)$

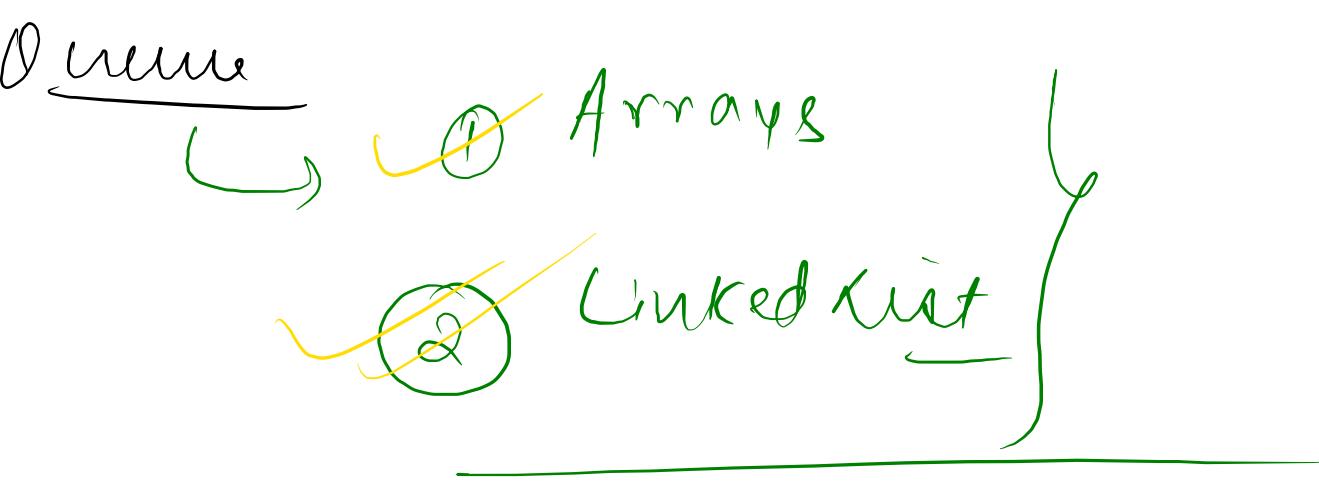
$q' E(23)$ $q' E(31)$

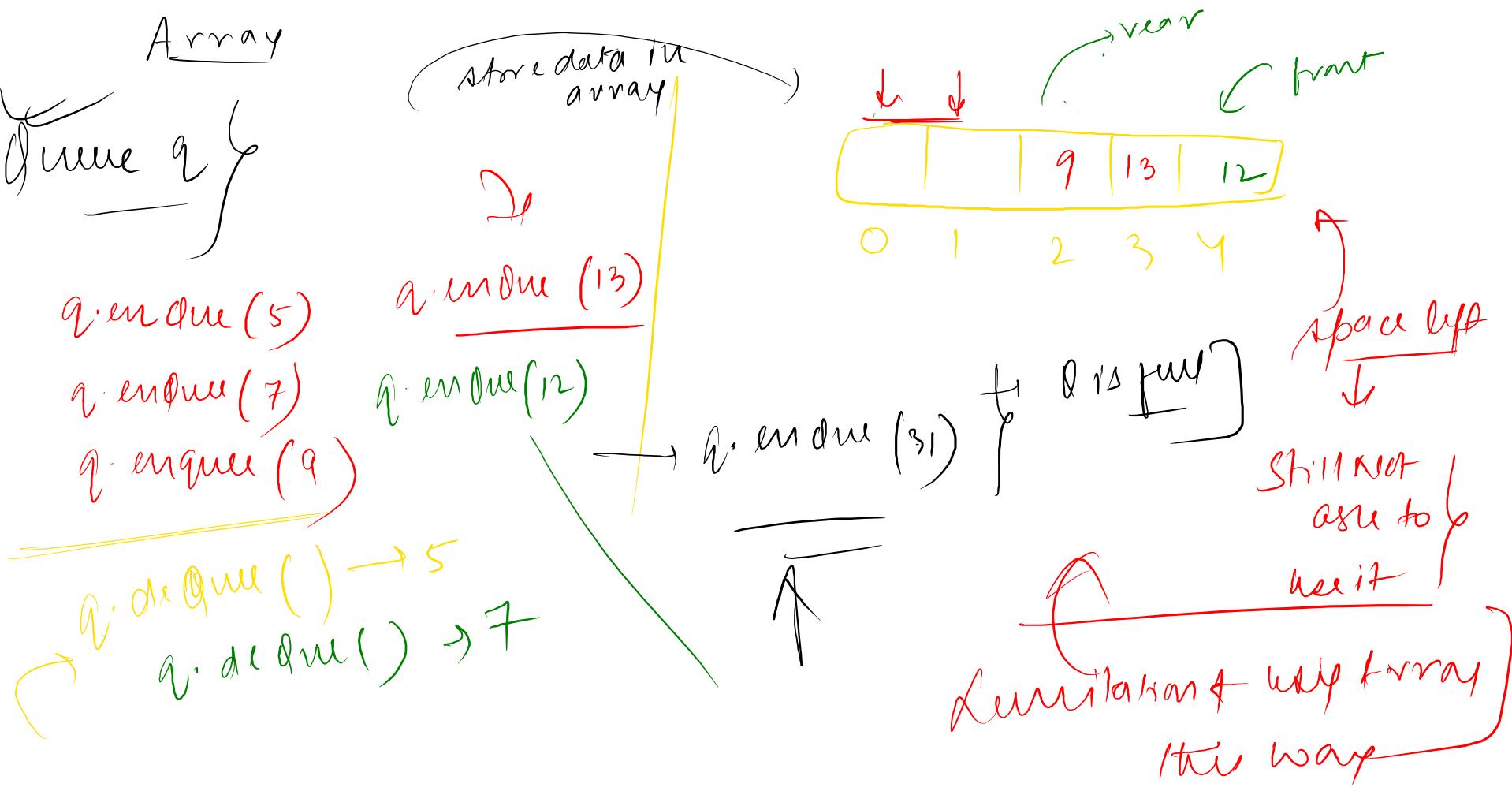
$f+1^{th}$ index
 $f+1 = 4^{th}$ invalid
 4% size

$f+1 = 8$ $4 \times 4 = 0$
 $i+1 = 2 \{ \cancel{4^{th}}$

Conditions to check

$$\text{full} \rightarrow \frac{(f+1)\% \text{ capacity}}{\text{---}} = \gamma \quad \left. \begin{array}{l} \text{---} \\ \text{---} \end{array} \right\} \begin{array}{l} \text{Ons} \\ \text{full} \end{array}$$





Queue front → index where we
 enqueue

q.enqueue(5) rear → index where we
 enqueue

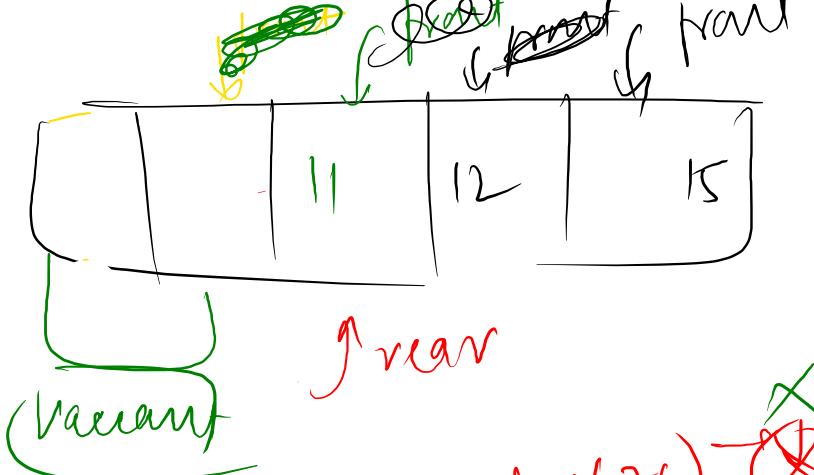
q.enqueue(2)

q.enqueue(11)

q.enqueue(7) → 5
q.enqueue(7) → 7

q.enqueue(12)

q.enqueue(15)



q.enqueue(7) → ~~7~~

value → O(1)
dequeue →

TC ↑