

## QuickSort

Merge sort

↳ combine

↓

Merge procedure

$$\text{mid} = l + (h - l) / 2$$

$l, m-1$   
 $QS(0, 1)$

$n = 6$

$m$



$l = 0$

$h = 5$

$x = \text{Pivot} = \text{arr}[l] = 40$

$i \rightarrow$  to take smaller elements than pivot

$j \rightarrow$  to take larger element than pivot

$QS(3, 5)$

↳ ~~60~~<sup>3</sup>, ~~70~~<sup>4</sup>, ~~60~~<sup>5</sup>, 90<sup>5</sup>  
 $x$      $j$      $j$

QuickSort

↳ Partition

↓

Divide the array

into two subarrays

↳ No combine required in quickSort

$QS(m+1, h)$

Partition (arr, l, h) {

i = l

pivot = arr[l];

for (j = l + 1; j <= h; j++) {

if (arr[j] < pivot) {

i = i + 1;

swap(arr[i], arr[j]);

}

}

correct  
index  
of pivot  
element

swap(arr[l], arr[i])

return i;

Note

1. Pivot Element reaches to its correct position
2. Left side of pivot — smaller element  
Right side of pivot — bigger element

# QuickSort Algorithm

function  
name

quickSort(arr, l, h)

if ( $l < h$ )

1. Divide the array

$m = \text{partition}(\text{arr}, l, h);$

2. Conquer the subproblems  
via recursion

$m-l+1$

$T(m-l)$

quickSort(arr, l, m-1);

Recursion

$T(h-m)$

quickSort(arr, m+1, h);

✓

↓

$h-(m+1)+1$

$h-m-1+1$

## Recurrence Relation

$$T(n) = \begin{cases} 1 & ; n=1 \\ \underbrace{T(m-l)}_{\downarrow} + \underbrace{T(h-m)}_{\downarrow} + n & ; n \geq 1 \end{cases}$$

partition



QS(0,2)

0	1	2	3	4	5	6
10	30	40	50	70	60	90

Left Subarray  
Right Subarray  
QS(4,6)

Best case

Average case

$$T(n) = T(n/2) + T(n/2) + n$$

$$= 2T(n/2) + n$$

$$= O(n \log n)$$

algorithm

0	1	2	3	4	5
10	50	40	60	65	90

QS(0,0) QS(2,5)

Worst case

$$T(n) = T(n-1) + n$$

$$= O(n^2)$$

Insertion sort  
Sorted / Almost sorted  
 $O(n)$

Nothing

0	1	2	3	4
10	20	30	40	50

$i$

QuickSort → Geniune scenario

↳ Highly unsorted

Space Complexity

↓

Inplace sorting

↓

$O(1)$

→ Not a stable algorithm

$10^1, 10^2, 10^3, 10^4, 10^5$   
0 1 2 3 4

→ Randomized QuickSort

→ Pivot element

(Randomly chosen)

0	1	2	3	4	5
<del>7</del>	4	6	<del>5</del>	1	3