

Greedy Approach

constraint
Optimization Problem
(Max, Min)

1) Fractional Knapsack

constraint

total weight $\leq M$

Max Profit

Complexity - $n \log n$

Approach

for (int i = 0 to n-1) {

P/w (Ratio)

①

}

②

Sort P/w Desc order

③ for (i = 0 to n-1) {
M decrease — till the value becomes 0
totalProfit increase

}

Weight

10

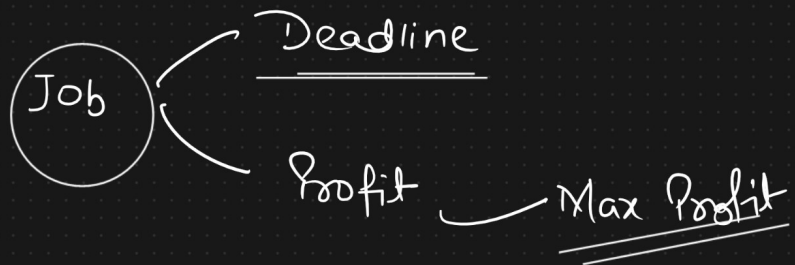
M =

9

9/10

Fractional Knapsack

Job Scheduling



$$\min(1, 6) = 1$$

Job	J ₁	J ₂	J ₃	J ₄	J ₅	J ₆	J ₇	J ₈	J ₉
Profit	55	65	75	60	40	50	85	68	45
Deadline	5	2	7	3	2	1	4	5	3

	1	2	3	4	5	6	7
	J ₂	J ₅		J ₇	J ₈		J ₃
	T	F	F	F	F	F	F
	0	1	2	3	4	5	0

Job Id →	1	2	3	4	5	6	7	8	9
Deadline →	5	2	7	3	2	1	4	5	3
Profit →	55	65	75	60	70	50	85	68	45
	↑	↑	↑	↑	↑	↑	↑	↑	↑

Constraint → Need to complete a job within given deadline.

Optimization → Max Profit

1) Sort the jobs → Profit → Decreasing order

Job Id	7	3	5	8	2	4	1	6	9
Deadline	4	7	2	5	2	3	5	1	3
Profit	85	75	70	68	65	60	55	50	45
	↑	↑	↑	↑	↑	↑	↑	↑	↑

min(6,6) = 6, min(1,6) = 1, min(1,6) = 1, min(4,4) = 4, min(0,6) = 0

min(3,6) = 3

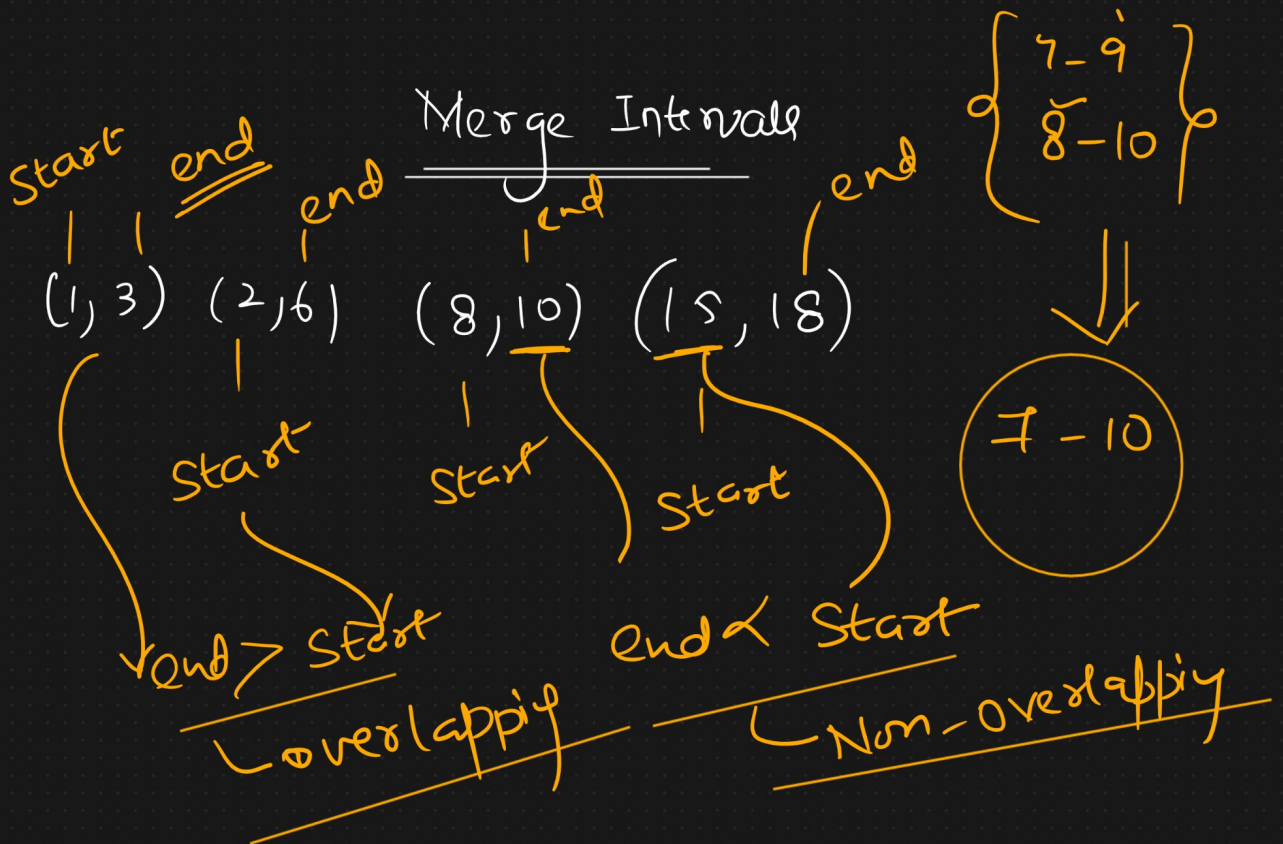
maxDeadline = 7

2)

Job →	2	5	4	7	8	—	3
result →	T	F	T	T	T	F	T

$\min(\text{maxDeadline} - 1,$

$\text{arr}(i).\text{get}.\text{deadline}(i) - 1)$





interval

(1, 4)

4 > 2 true

interval

(1, -

(8, 10) (15, 16)

10 < 15

```

for(int[] interval: intervals){
    // No overlapping
    // 10 < 15 -yes
    if(merged.isEmpty() || merged.getLast()[1] < interval[0]){
        merged.add(interval);
    }
    else{
        // overlapping
        // max(LastEnd, end)
        // [1, 6]
        // [1, 4] [2, 3] - [1, 4]
        merged.getLast()[1] = Math.max(merged.getLast()[1], interval[1]);
    }
}

```

Maximum Value return

Data
compression

Huffman coding
└─ Huffman tree

Max frequency — Lower
bit

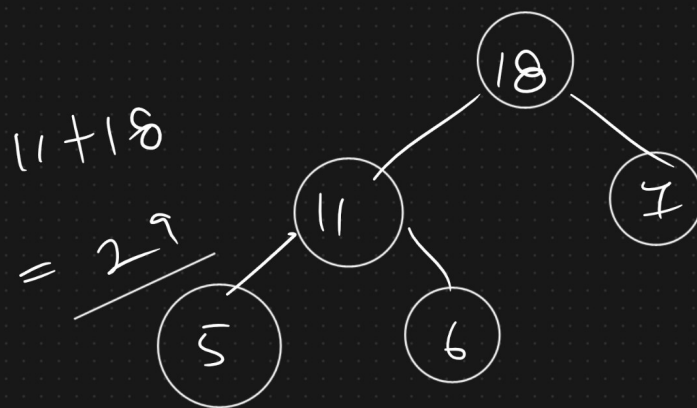
Low frequency — Higher
bit

Optimization

Optimal Merge Pattern

{
A = 5 ✓
F2 = 7
F3 = 6 ✓

Minheap



Priority
Queue

Assignment Problem

<https://leetcode.com/problems/task-scheduler/description/>

QuickSort Algorithm

Live session