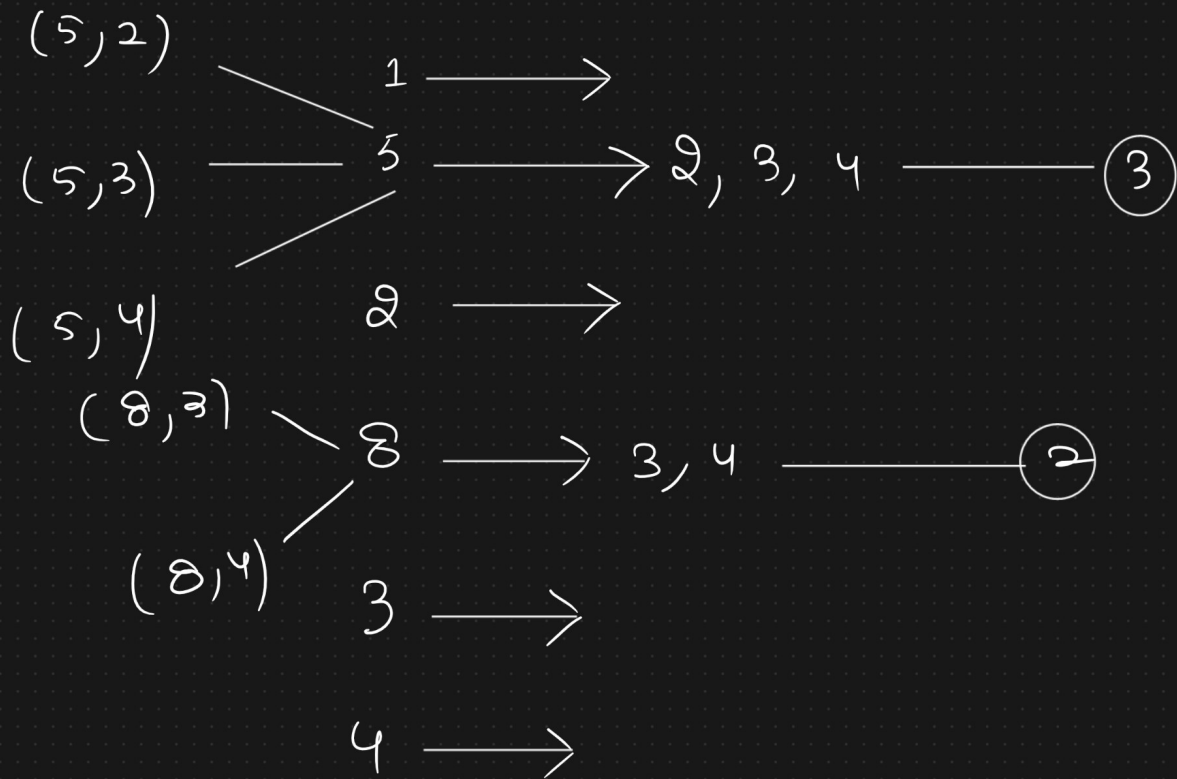


0	1	2	3	4	5
1	5	2	8	3	4

→ Inversion

$i < j$
 $arr[i] > arr[j]$



count = 5 (3 + 2)

$O(n^2)$

```

inversion(arr, n) {
  for (i = 0 to n-1) {
    for (j = i+1 to n) {
      if (arr[i] > arr[j]) {
        count += 1;
      }
    }
  }
  return count;
}
  
```

0	1	2	3	4
5	4	3	2	1

Decreasing order

$i < j$
 $arr(i) > arr(j)$
 \Downarrow
 5 ——— 4, 3, 2, 1 (4)
 4 ——— 3, 2, 1 (3)
 3 ——— 2, 1 (2)
 2 ——— 1 (1)

Worst

case

scenario

10

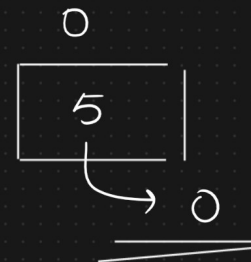
Increasing order

0	1	2	3	4
1	2	3	4	5

Best case
scenario

1 ———> 0
 2 ———> 0
 3 ———> 0
 4 ———> 0
 5 ———> 0

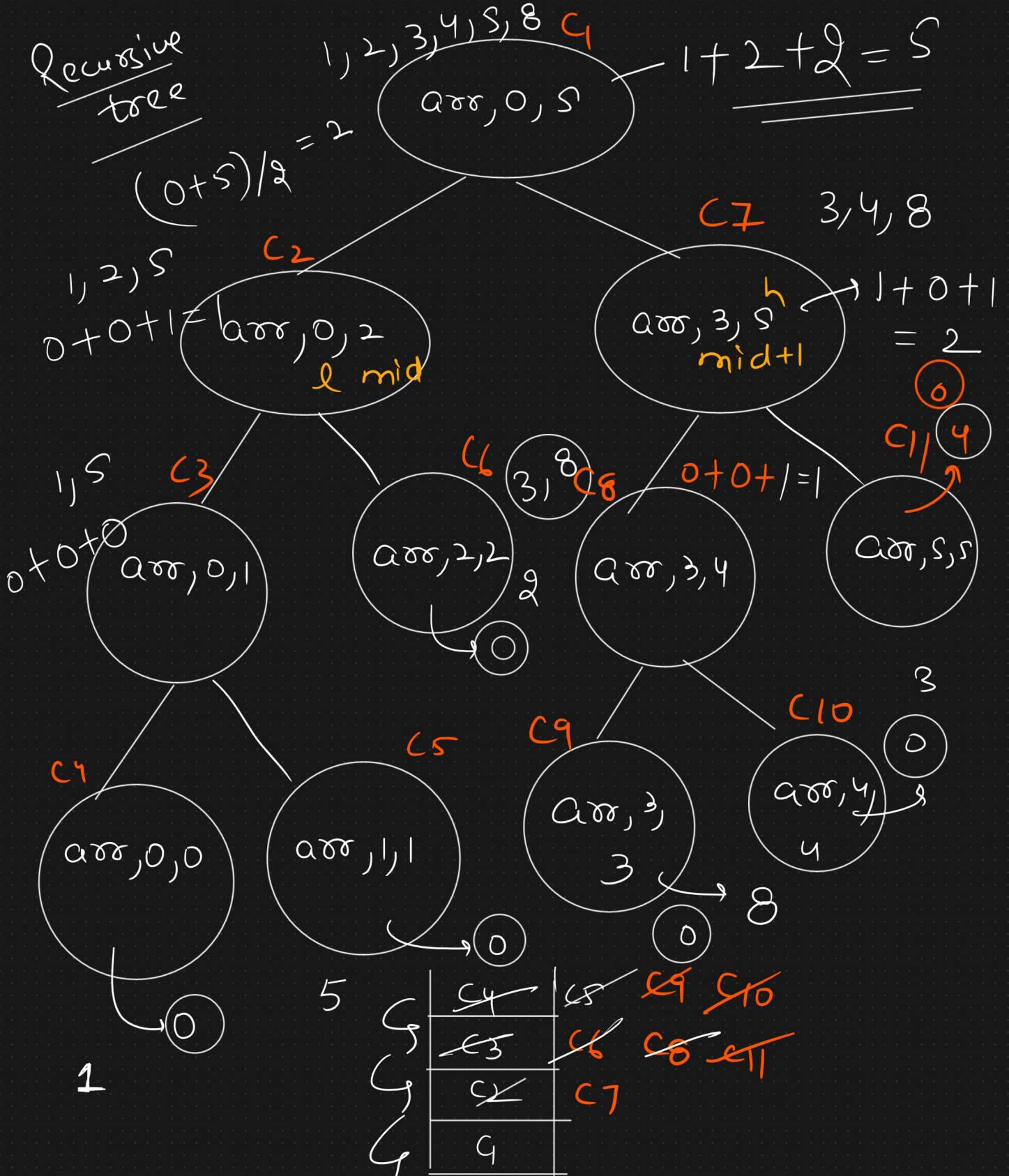
inversion = 0



Divide & Conquer

0	1	2	3	4	5
1	5	2	8	3	4

Recursive tree



inversionCount(arr, l, h) {

if (l < h) {

Divide ——— mid = l + (h - l) / 2

conquer ——— {
2T(n/2) { count += inversionCount(arr, l, mid);
count += inversionCount(arr, mid + 1, h);

combine ——— {
count += mergeProcedure
→ n (arr, l, mid, h);

return count;

Recurrence Relation

$$T(n) = 2T(n/2) + n$$

$$= \underline{\underline{O(n \log n)}}$$