

# Searching & Sorting

## Searching

### Ternary Search

### Linear Search

### Binary Search

## Sorted Manner

0	1	2	3	4	5	6
2	4	8	12	20	24	35

$\text{low} = 0, \text{high} = 6$

$(\text{low} + \text{high})/2$

②

$$\text{mid} = \frac{\text{low} + (\text{high} - \text{low})}{2}$$

$$= 3$$

$x = 24$



Right

$x = 4$



Left

Iterative approach

binarySearch( $\text{arr}, \text{low}, \text{high}, \alpha$ ) {

while ( $\text{low} \leq \text{high}$ )  
if ( $\text{arr}(\text{mid}) = x$ )  
    return  $\text{mid}$  — ①

if ( $\text{arr}(\text{mid}) < x$ )  
     $\text{low} = \text{mid} + 1;$  ②

if ( $\text{arr}(\text{mid}) > x$ )

$\text{high} = \text{mid} - 1;$  ③

}

}

## Recursive approach

binarySearch (arr, low, high, x)  $\alpha$

while (low  $\neq$  high)  $\alpha$

$$mid = low + (high - low)/2;$$

C  $\left\{ \begin{array}{l} \text{if } arr(mid) = -x \\ \quad \text{return mid;} \end{array} \right.$

$T(n/2)$

Right side

if ( $arr(mid) < x$ )  $\alpha$

Recursive function

return binarySearch call  
( $arr, \underline{mid+1}, \underline{high}, x$ )

low

$T(n/2)$

Left side

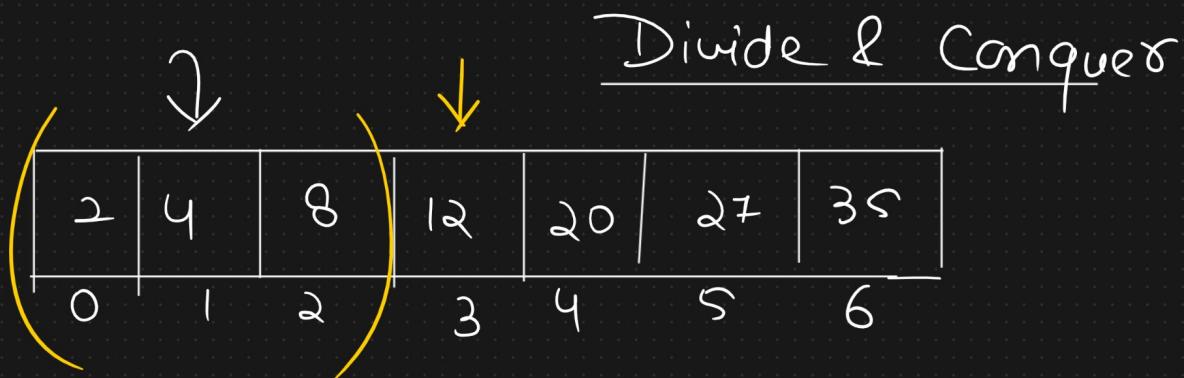
if ( $arr(mid) > x$ )  $\alpha$

return binarySearch ( $arr, \underline{low},$   
 $\underline{mid-1}, \underline{x})$ ;

## Recurrence Relation

$$T(n) \Rightarrow T\left(\frac{n}{2}\right) + c$$

$$\Rightarrow \underline{\mathcal{O}(\log n)}$$



$$x = 27$$

$$\text{low} = 4, \text{high} = 6$$

$$\text{mid} = 5$$

$$(arr(mid) == x)$$

$$\text{result} = 5$$

$$x = 4$$

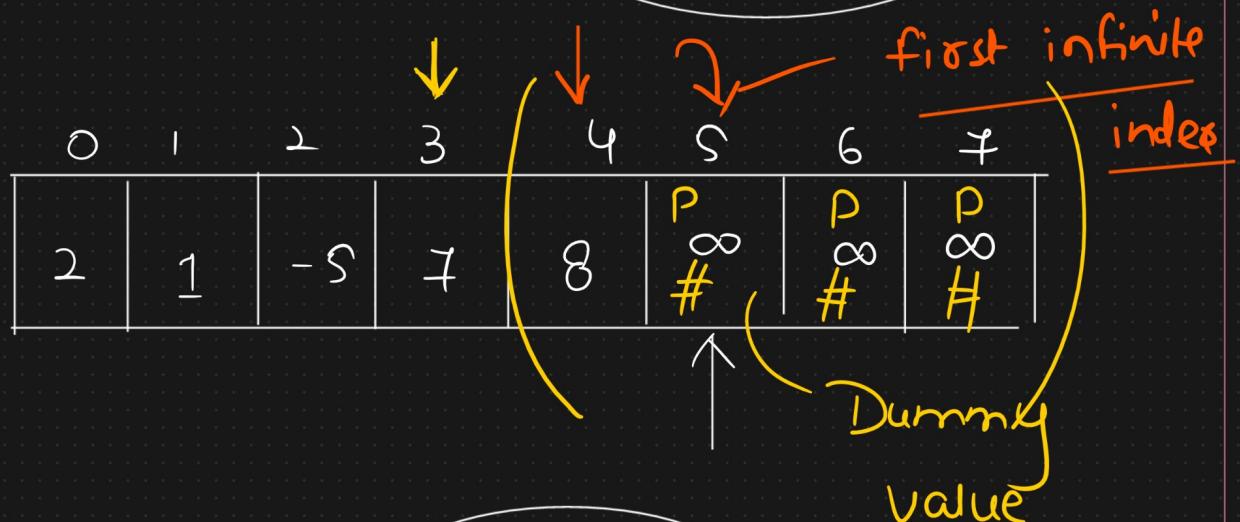
low = 0, high = 6

mid = 3

low = 0, high = 2

mid = 1

$$\text{result} = 1$$

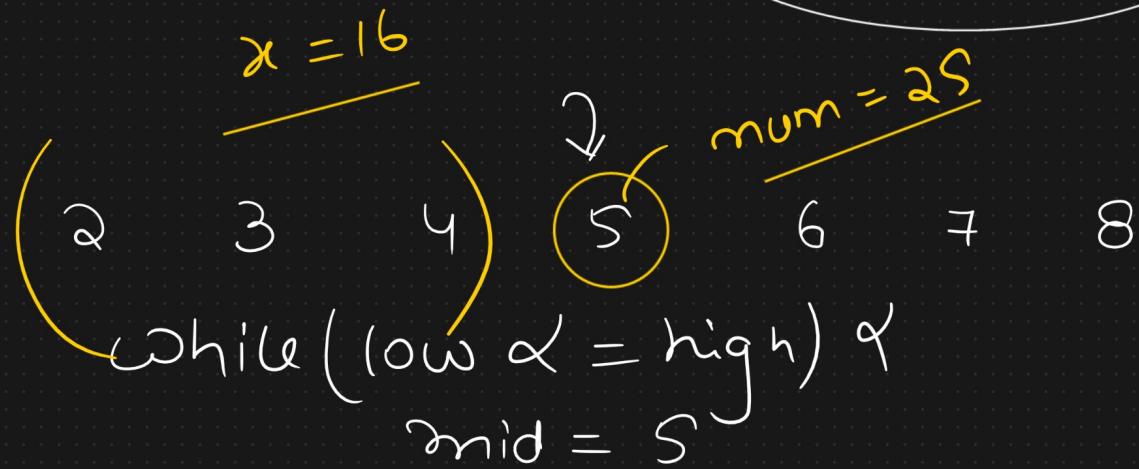
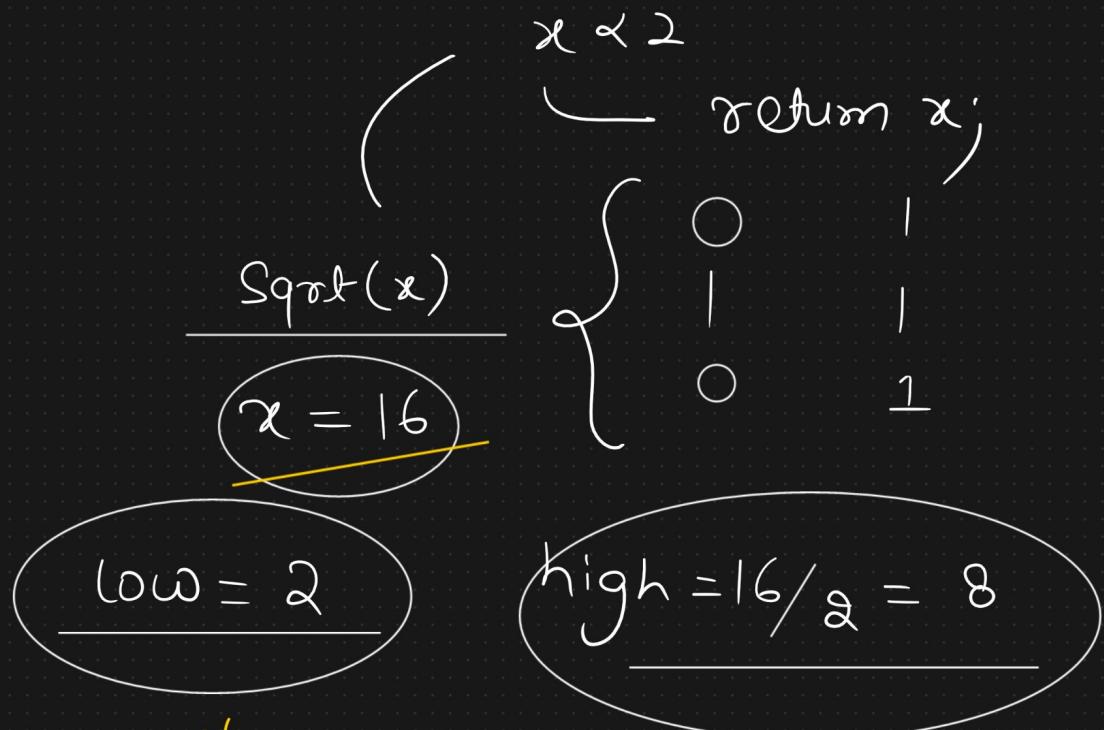


low = 0, high = 7

mid = 3

$\text{low} = 4$ ,  $\text{high} = 7$

$\text{mid} = 5$



$$\text{num} = 5 * 5 = 25$$

if  $\text{num} \geq x$

$\rightarrow \text{left}$

$\rightarrow \text{high} = \text{mid} - 1$

if  $\text{num} < x$

    → Right

$\text{low} = \text{mid} + 1$

if  $\text{num} == x$

}

return  $\text{mid}$ ;

return high,

    → Lower bound

$x = 8$

Expected Result  
= 2

$\text{low} = 2$

$\text{high} = 4$

↓

2      3      4

$\text{mid} = 3$

$\text{num} = 9$

$\text{num} \geq x :$

$$\underline{\underline{low = 2}} \quad \underline{\underline{right = 2}}$$

$$\underline{\underline{mid = 2}}$$

$$\text{num} = 4$$

$$\frac{\text{num} < x}{\curvearrowleft \text{right}} \quad \begin{array}{l} low = mid + \\ \quad \quad \quad \downarrow \\ \text{right} \end{array}$$

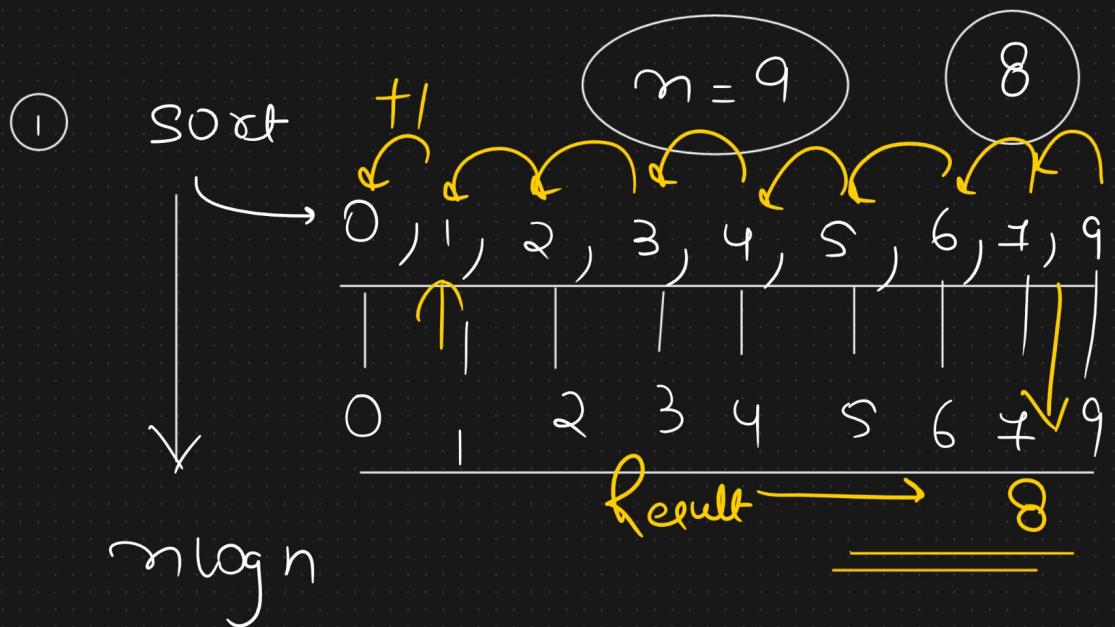
$$\underline{\underline{low = 3}} \quad \underline{\underline{high = 2}}$$

Missing number

$$n = 9$$

9, 7, 6, 4, 2, 0, 1, 3, 5
0    1    2    3    4    5    6    7    8

Output = 8



Space →  $\mathcal{O}(n)$

② Gauss' formula

Sum of  $n$  natural numbers

$$\frac{n(n+1)}{2}$$

$$\text{Expected Sum} = \frac{5 \times 10 * 11}{2} = 55$$

c ←

$\left\{ \begin{array}{l} \text{for } (\text{num}: \text{num}) \times \\ \quad \text{ActualSum} += \text{num} \\ \quad \quad \quad \downarrow 47 \end{array} \right.$

$$55 - 47 = 8$$

$$\text{ExpectedSum} - \text{actualSum} = 8$$


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$\left\{ \begin{array}{l} \text{Time complexity} = O(n) \\ \text{Space Complexity} = O(1) \end{array} \right.$

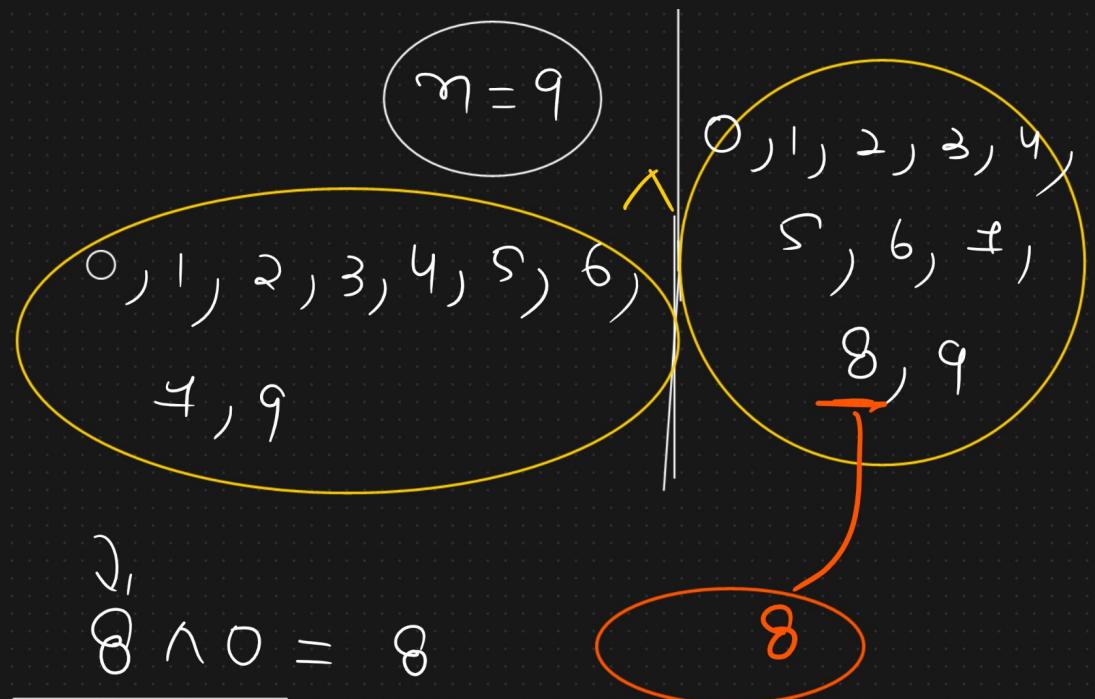
### 3 XOR Operation

$\hookrightarrow$  Bitwise manipulation

# Truth Table

A	B	O/P
✓ 0	✓ 0	✓ 0
0	1	1
1	0	1
✓ 1	✓ 1	✓ 0

$$\neg A \wedge 0 = A$$



time complexity =  $\mathcal{O}(n)$

Space complexity =  $\mathcal{O}(1)$

$$\begin{array}{c}
 \text{missingEle} = 3 \\
 \boxed{\text{missingEle}} = 3 \\
 \boxed{[0, 1, 3]} \rightarrow 4 \\
 \underline{6 - 4 = 2}
 \end{array}$$

$$\begin{array}{c}
 i = 0 \quad | \quad i = 1 \quad | \quad i = 2 \\
 0 \wedge 0 = \underline{0} \quad | \quad 1 \wedge 1 = \underline{0} \quad | \quad 2 \wedge 3 \\
 \text{missingEle} = 3 \wedge 2 \wedge \cancel{3} \\
 \underline{\underline{\Rightarrow 0 \wedge 2 = 2}}
 \end{array}$$

$$\text{missingEle} = \text{missingEle} \wedge i \wedge \text{numsl[i]}$$

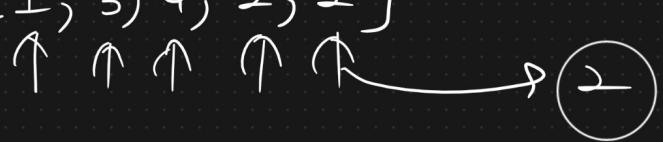
$$\left\{ \begin{array}{c}
 \begin{array}{c}
 i \\
 \boxed{0 \ 1 \ 2} \\
 3, 0, 1) \rightarrow \boxed{\text{missingEle}} \\
 \text{numsl[i]} \\
 \downarrow \quad \downarrow \quad \downarrow \quad \downarrow \quad \downarrow \quad \downarrow \\
 3 \wedge \cancel{0} \wedge 3 \wedge 1 \wedge 0 \wedge 2 \wedge 1 \\
 \underline{\underline{0 \wedge 2 = 2}}
 \end{array}
 \end{array} \right.$$

Approach 1

Set

Not any duplicates

0	1	2	3	4
[1, 3, 4, 2, 2]				



time complexity  
 $= O(n)$

Set

1
3
4
2

Space  
 $\hookrightarrow O(n)$

Approach 2

Sort

$O(n \log n)$

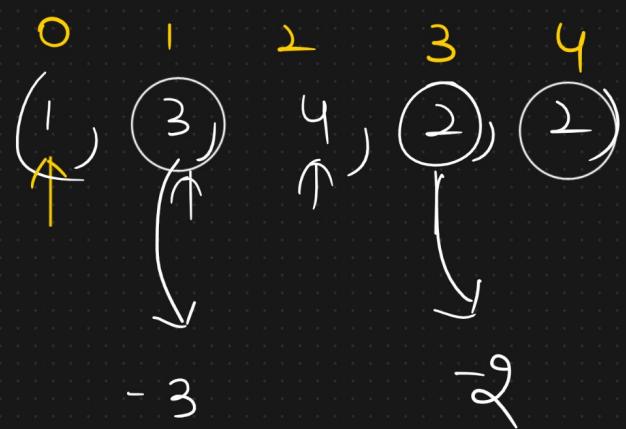
$n = s$

0	1	2	3	4
(1, 2, 2, 3, 4)				



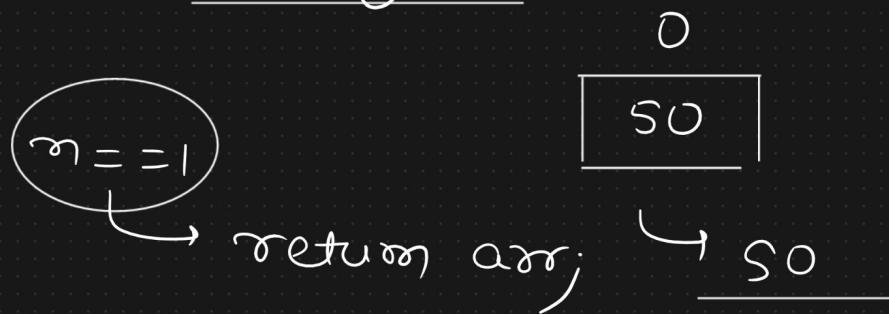
Space  $\rightarrow O(n)$

## Approach 3    Negative Masking

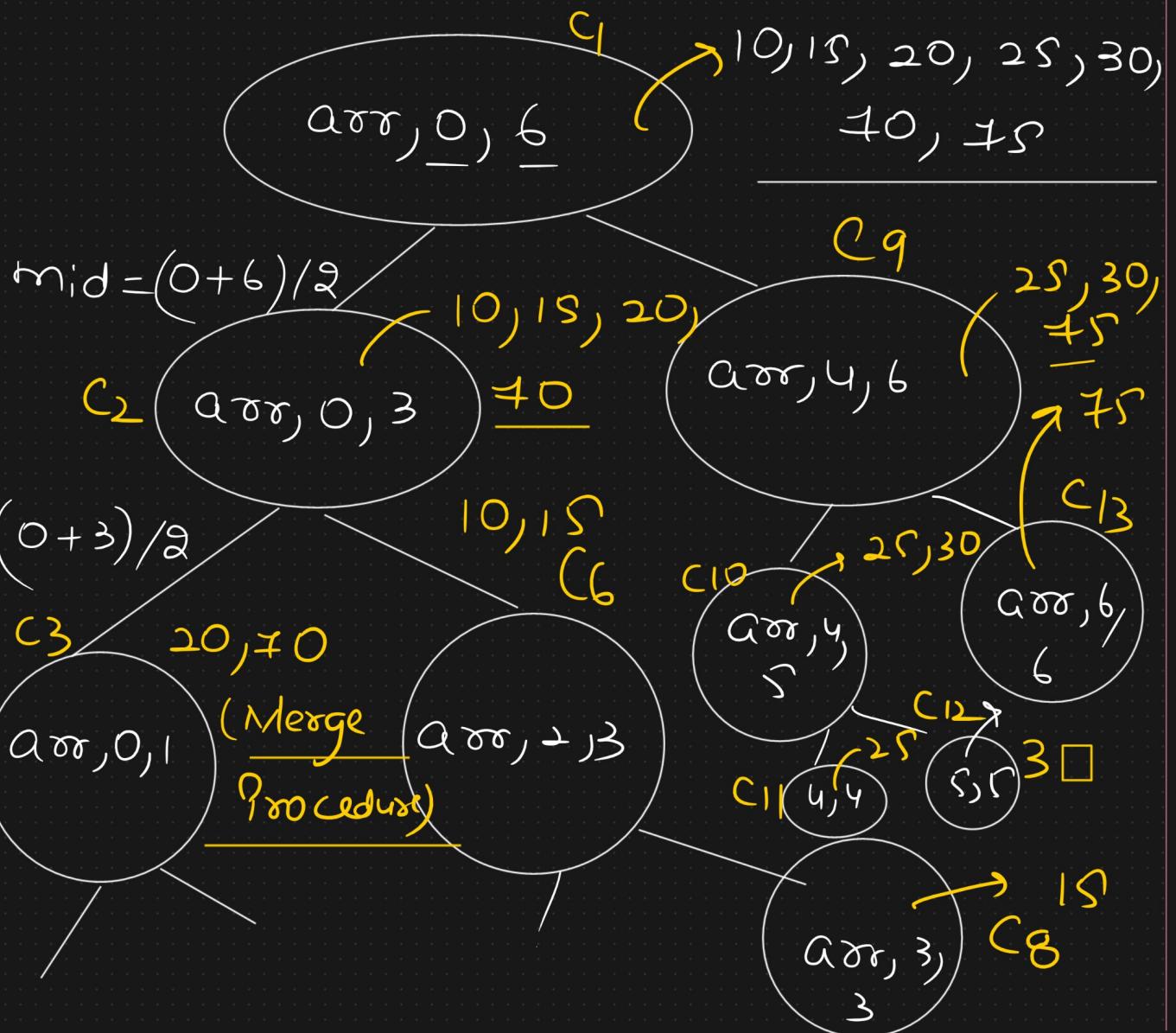


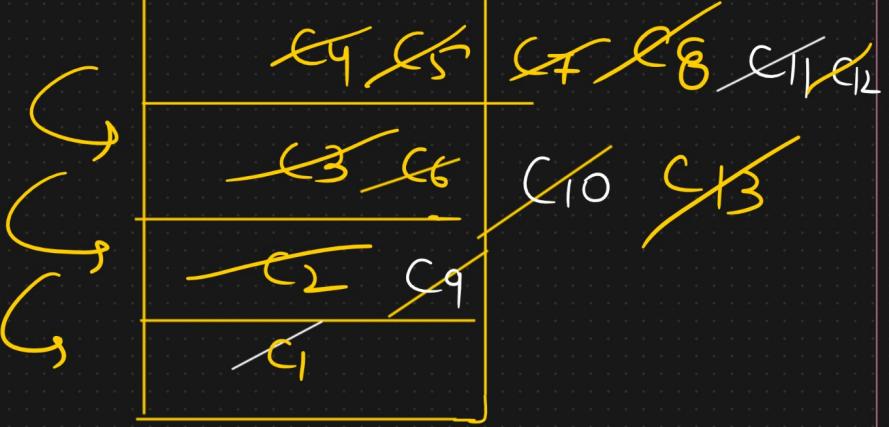
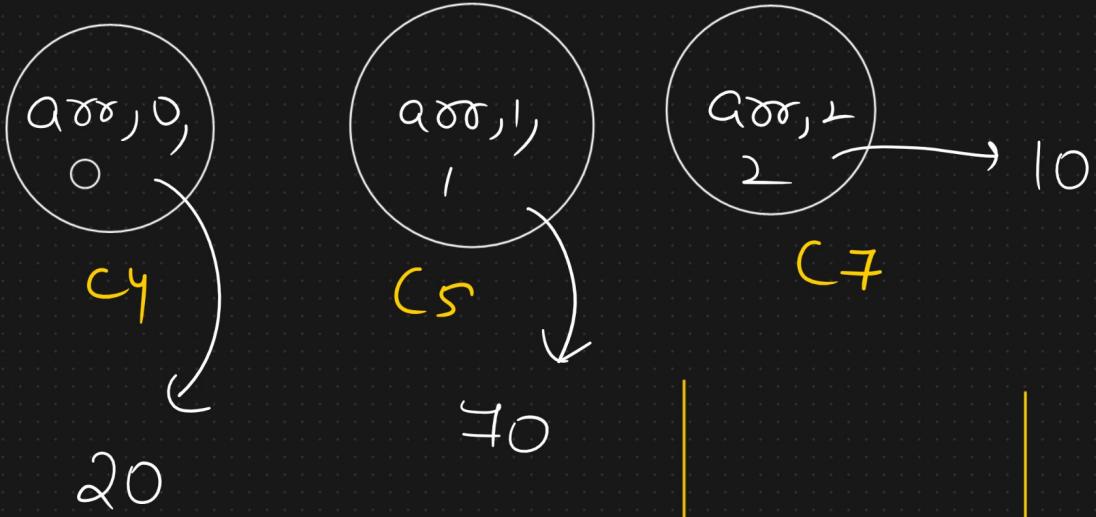
# Divide & conquer

## Mergesort



0	1	2	3	4	5	6
20	70	10	15	25	30	75





MergeSort

- ① Divide - the bigger problem into smaller subproblems

- ② Conquer → Recursion

- ③ Combine

→ merge procedure

## Merge Procedure

0	1	2	3	4	5	6	7				
10	20	30	40	11	21	31	41				

Left Subarray      Right Subarray

$\underline{40, 31 = 31}$

$\underline{40, 41 = 40} \quad \checkmark$

$10 | 20, 11 \rightarrow 11 | 20, 21 = 20$

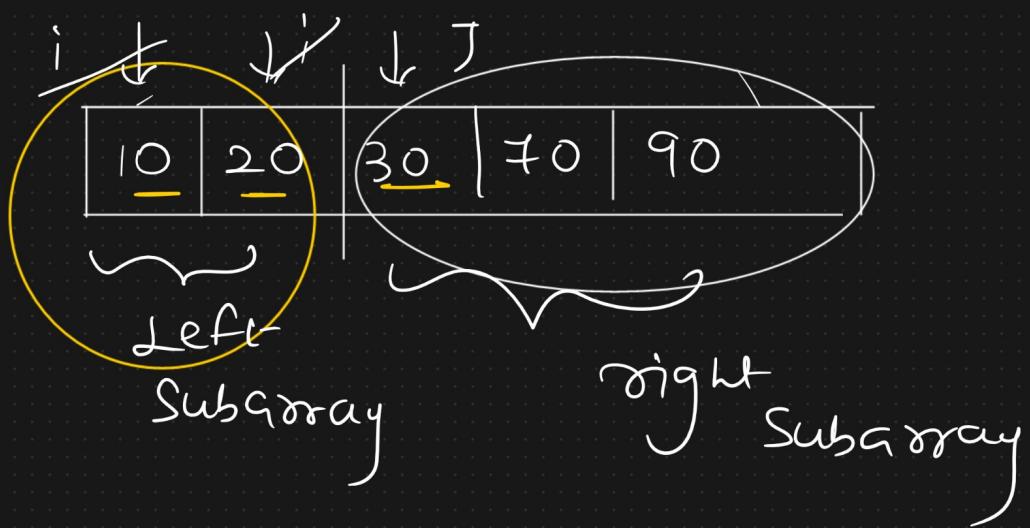
0	1	2	3	4	5	6	7				
10	11	20	21	30	31	40	41				

$$m = 4 \quad | \quad n = 4$$

Worst Case

$$\text{Comparisons} = m + n - 1$$

$$4 + 4 - 1 = 7$$



$$\frac{2 \text{ comparison}}{\left\{ \begin{array}{l} 10 < 30 \\ 20 < 30 \end{array} \right\}} = \frac{10}{20} = 10$$



$$\frac{\text{Best case}}{\text{num of comparison}} = \frac{\min(m, n)}{\min(2, 3)} = \frac{2}{2}$$

mergesort( arr, low, high )

if ( low < high )

Divide       $mid = \lfloor low + (high - low) / 2 \rfloor$

Conquer      {      mergesort( arr, low, mid )  
                        mergesort( arr, mid+1, high )

mergeProcedure( arr, low,

mid, high )

Combine

Bottom to up

merge Sorted

array