

$low=0, high=4$

$\frac{high - mid}{\uparrow}$

|    |    |    |    |    |
|----|----|----|----|----|
| 0  | 1  | 2  | 3  | 4  |
| 20 | 40 | 70 | 50 | 10 |

high - (mid + 1)

mid = 2

$$\frac{2 - 0 + 1}{2} = 3$$

$$mid - low + 1$$

left  
Subarray

20, 40, 70

arr, 0, 4

arr, 0, 2

low      mid

Right Subarray

mid + 1      10, 50

C7      high

C9

arr, 3, 4

4      10

3      50

2      70

1      40

0      20

C4      C5

C3      C6

C2      C7

C1

20, 40

arr, 0, 1

high

C4

arr, 0

0

C5

arr, 1

1

arr, 2

2

arr, 3

3

arr, 4

4

i < n



arr

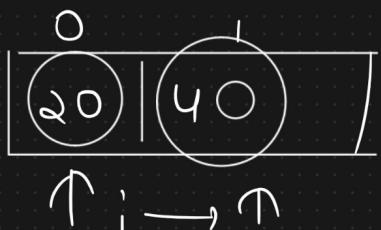
20, 40



J



T      K



$$\begin{array}{r} 20, 70 \\ 40, 70 \\ \hline \end{array}$$



|    |    |    |
|----|----|----|
| 20 | 40 | 70 |
| 0  | 1  | 2  |

0 1 2

|    |    |    |
|----|----|----|
| 20 | 40 | 70 |
| i  | i  | i  |

|    |    |
|----|----|
| 10 | 50 |
| 0  | 1  |

20, 10

20, 50

~~T<sub>j</sub>~~ T<sub>j</sub>

40, 50

|    |    |    |    |    |
|----|----|----|----|----|
| 10 | 20 | 40 | 50 | 70 |
| 0  | 1  | 2  | 3  | 4  |

→ Sorted array

$$T(n) = 2T\left(\frac{n}{2}\right) + \mathfrak{M}$$


---

|  
Conquer

merge  
Procedure

↓  
Master's Theorem/

Substitution Method

$$= \mathcal{O}(n \log n)$$


---

↳ Best, Average,

---

Worst

---

Space complexity

---

( Stack space +  
Extra space )

$$\underline{\underline{\mathcal{O}(n)}}$$


---

Extra space

---

Outplace sorting algorithm

---

Linked List

→ Dynamic data structure

↳ No contiguous storage

↖ fixed size  
Static array

→ ↗ not any fixed size  
dynamic array ↗ of an array  
ArrayList

Python → only dynamic  
array

↖ Random Access

more frequently → array

Access /

Search

)

Binary

Search

→  $O(\log n)$

Linked List ( more frequently )

→ Insertion, deletion  
( Linked List )

Skip List

{ Advance  
DSA

Load factor  $70.8$

## Working of Dynamic Array

| 1 | Size = 10 |
|---|-----------|
| 0 | 10        |
| 1 | 20        |
| 2 | 40        |
| 3 | 70        |
| 4 | 90        |
| 5 | 100       |
| 6 | 120       |
| 7 | 220       |
| 8 |           |
| 9 |           |

1 copy

2 Insertion

of new

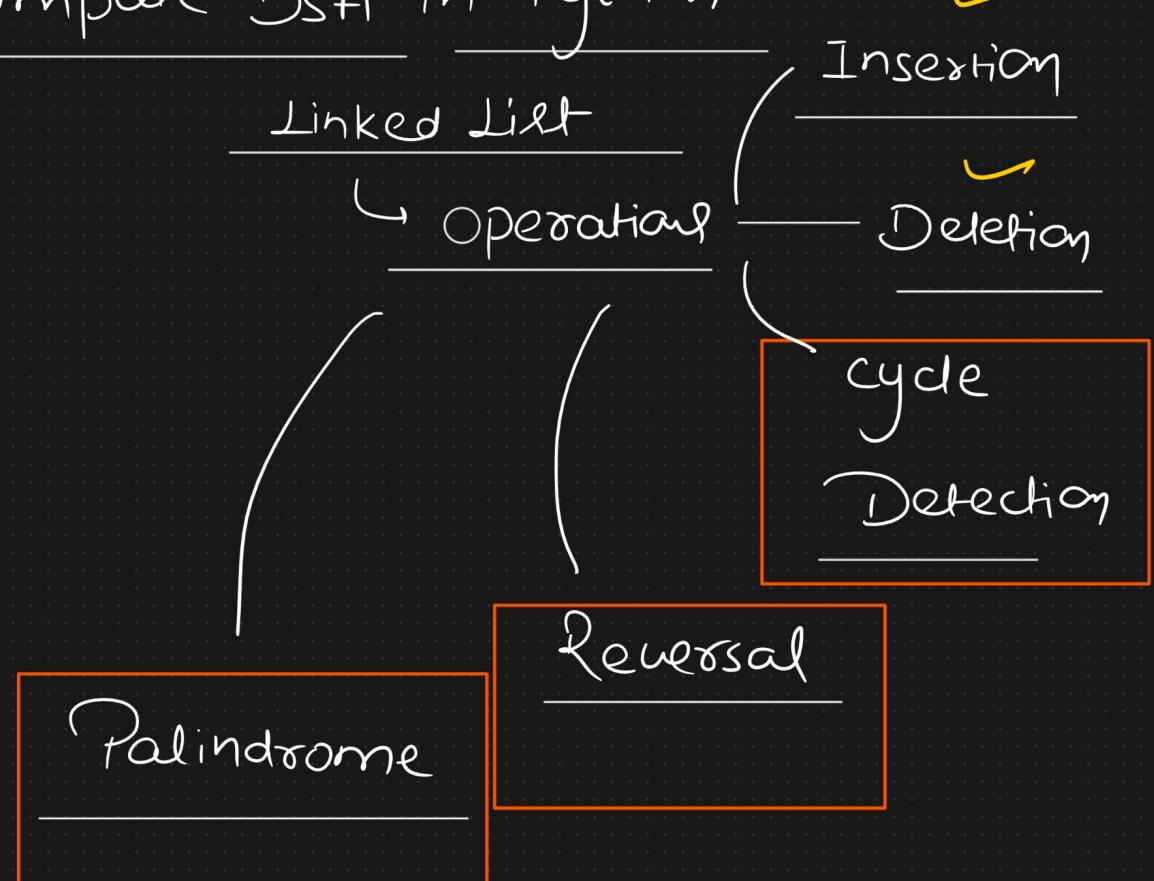
element

270

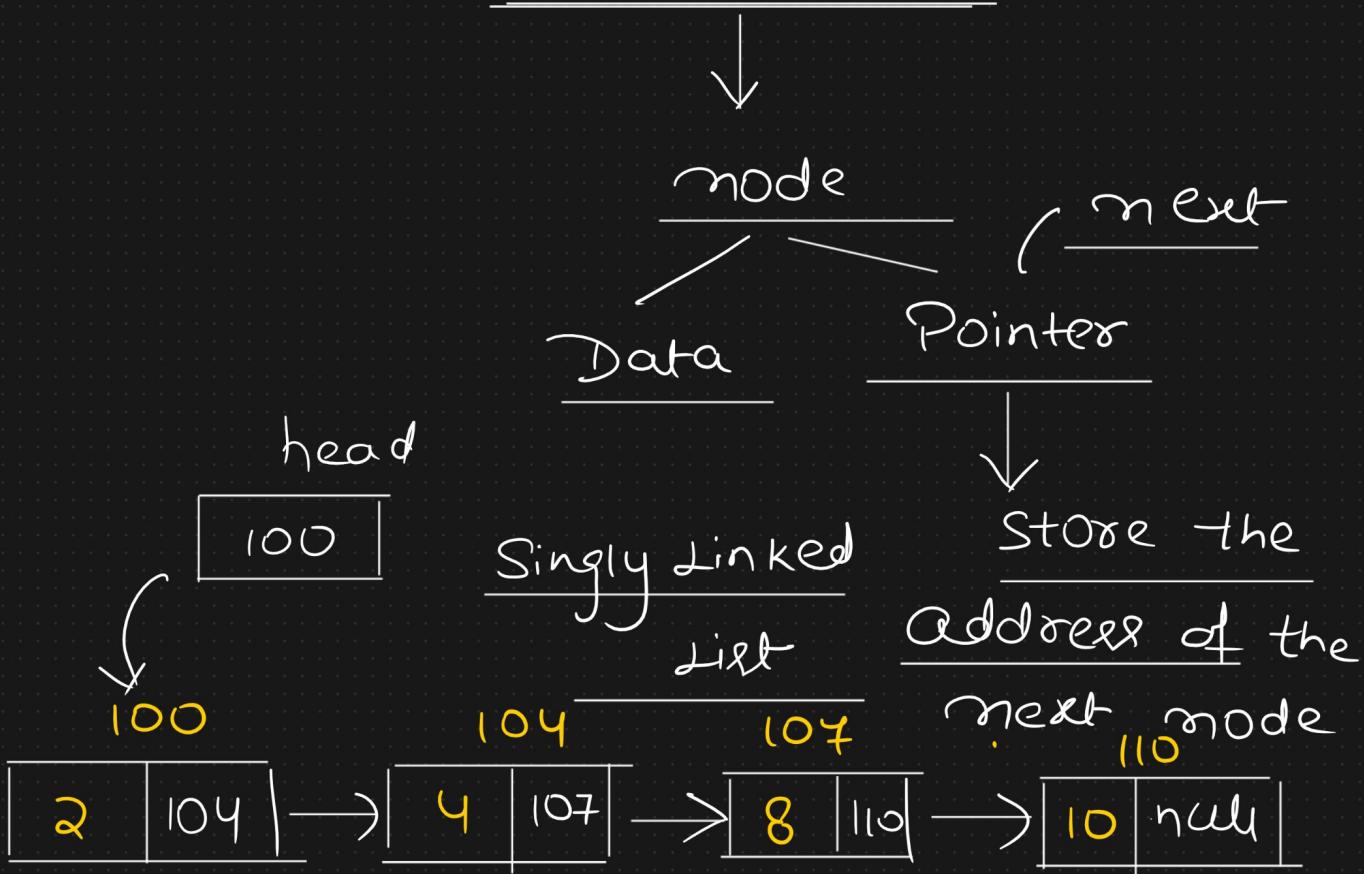
$$10 \times 2 = 20 \quad 2$$

|     |
|-----|
| 10  |
| 20  |
| 40  |
| 70  |
| 90  |
| 100 |
| 120 |
| 220 |
| 270 |
| 19  |

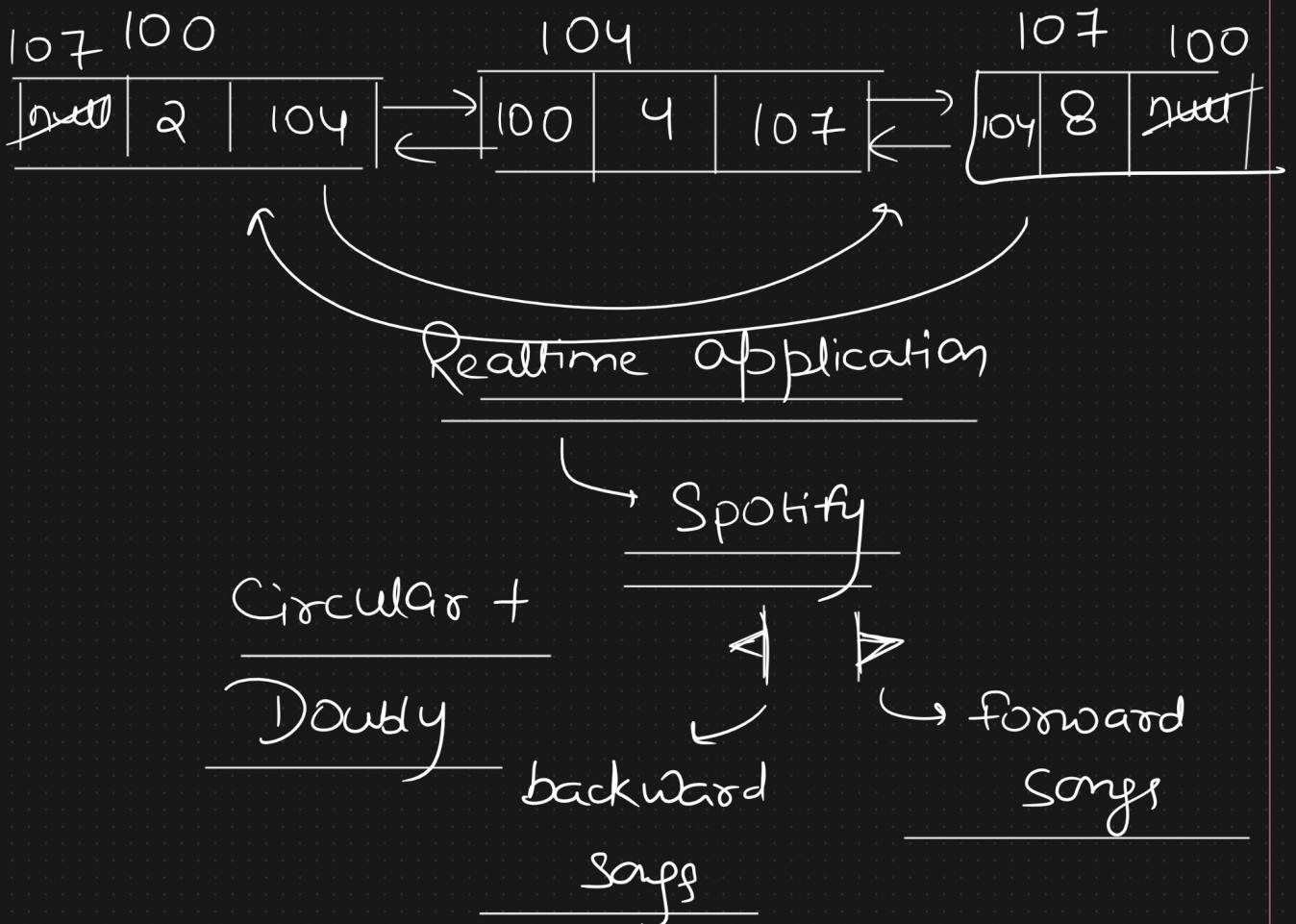
# Complete DSA in Python



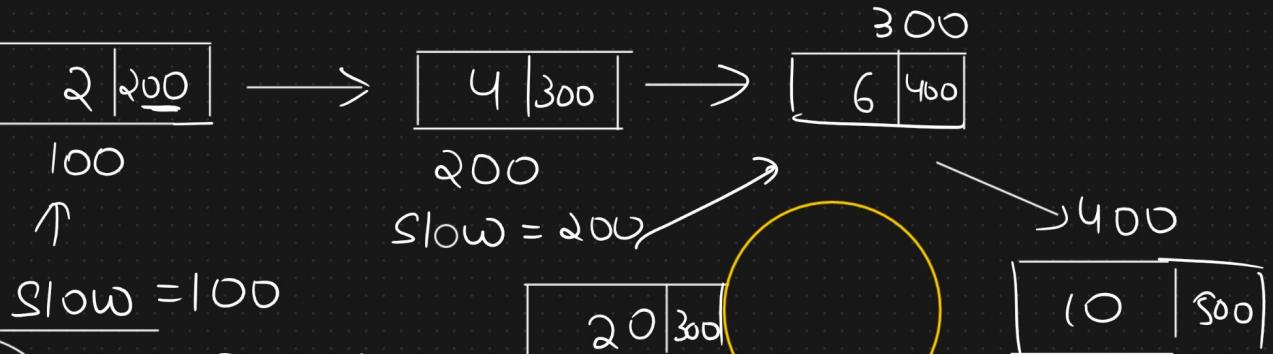
## Linked List



# Doubly Linked List

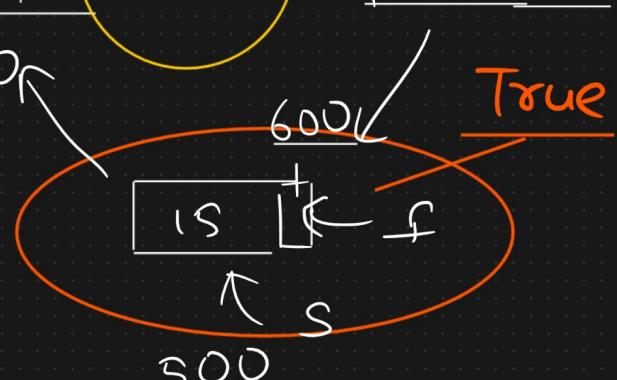


Slow → One step at a time,  
fast → Two step at a time



datatype  
Boolean value

True / False



if ( $\text{Slow} == \text{fast}$ )

$\hookrightarrow$  Cycle in linked

Floyd's cycle

Detection

head

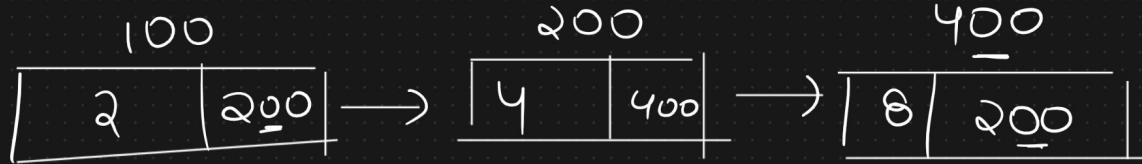
100

Algorithm

List

$\hookrightarrow$  return

true



while ( $\text{Slow} != \text{null}$   $\&$   $\text{Fast} != \text{null}$   $\&$   $\text{Slow} \neq \text{Fast}$ )

$\text{Slow} = \cancel{\text{Slow}}$ ,  $\text{Fast} = \cancel{\text{Fast}}$

Fast.next != null



Same mode

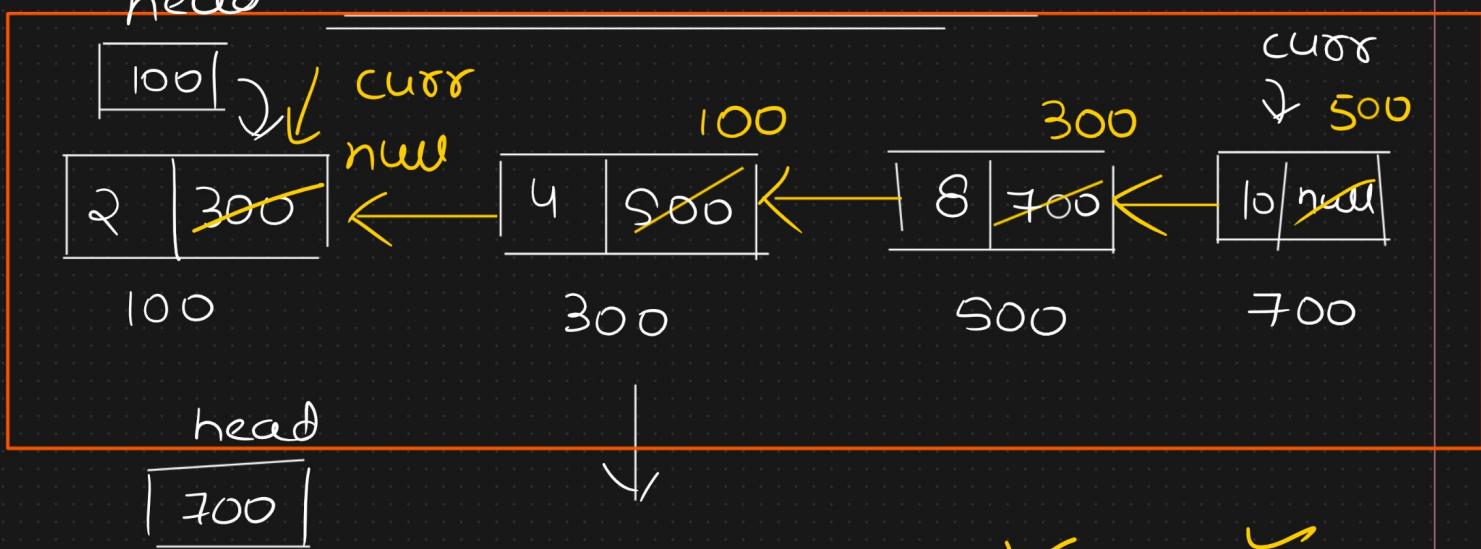
Point

Loop/cycle in  
a given Linked List

time complexity  
 $\hookrightarrow \Theta(n)$

Space Complexity  
 $\hookrightarrow \Theta(1)$

## Reversal of Linked List



$$\text{curr} = \text{head} = 100 \quad 300 \quad 500 \quad 700 \quad \text{null}$$

$$\text{nextPtr} = \text{null} \quad 300 \quad 500 \quad 700$$

$$\text{prev} = \text{null} \quad 100 \quad 300 \quad 500$$

$O(n)$  while ( $\text{curr} \neq \text{null}$ )

nextPtr = curr.next;

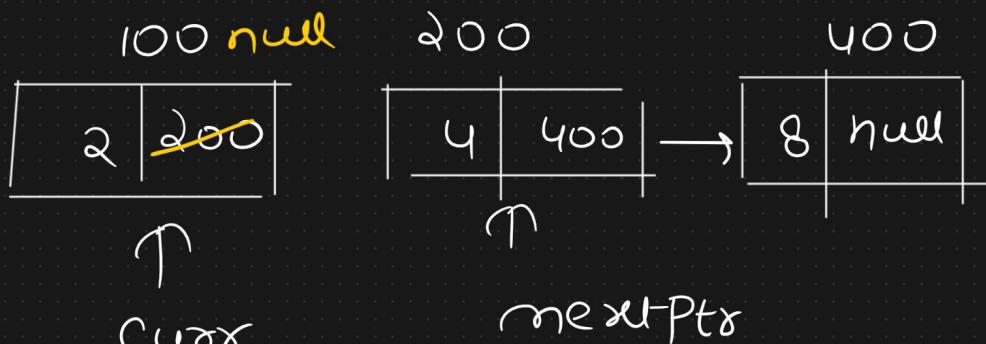
curr.next = prev;

prev = curr;

curr = nextPtr;

head = prev;

return head;



while ( $\text{curr} \neq \text{null}$ )  
   $\xrightarrow{\text{Complete LL}}$

  nextPtr = curr.next;

  curr.next = prev;

  prev = curr;

  curr = nextPtr;

# Palindrome

## Linked List



## ArrayList ↪ 2 Pointers



## Approach

low = 0 /

high = 3 - 2

time complexity =  $O(n)$       low      high

while (low < high) {

if ( $\text{arr}(\text{low}) == \text{arr}(\text{high})$ ) {

$\text{low} = \text{low} + 1$ ;

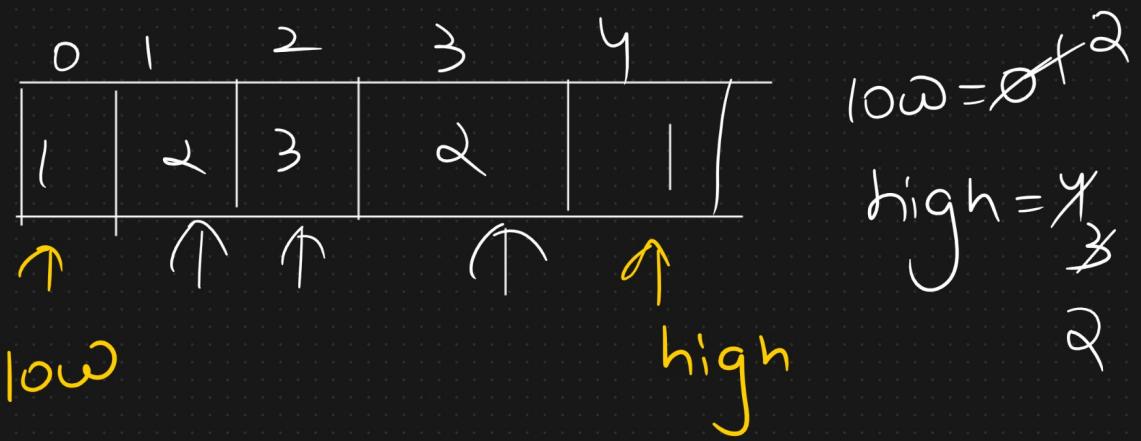
$\text{high} = \text{high} - 1$ ;

}

return true

$\text{arr}(\text{low}) != \text{arr}(\text{high})$  {

return false {



Space complexity =  $\mathcal{O}(n)$

ArrayList

Dynamic

Array

Total  $K$  time complexity =  $\mathcal{O}(n)$   
Space complexity =  $\mathcal{O}(1)$