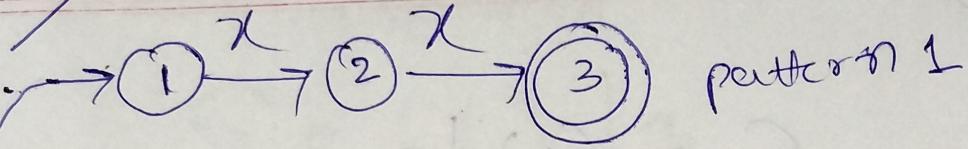
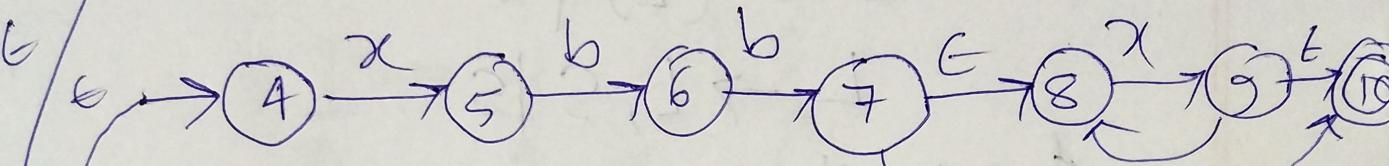


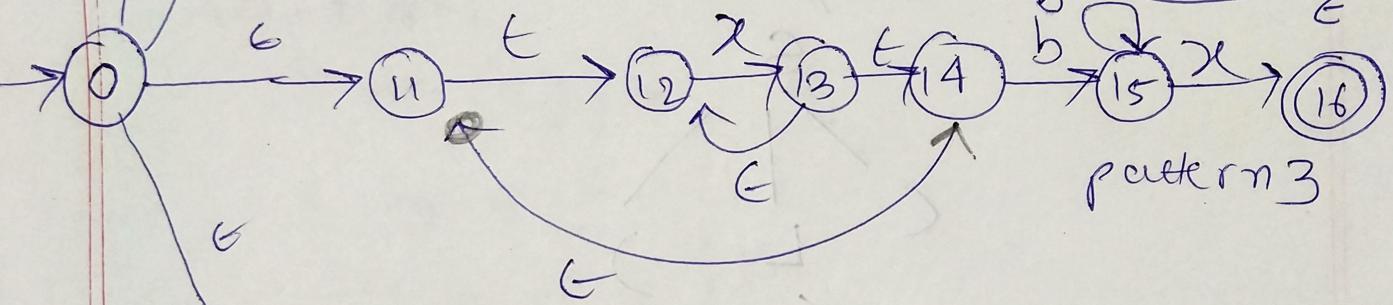
1



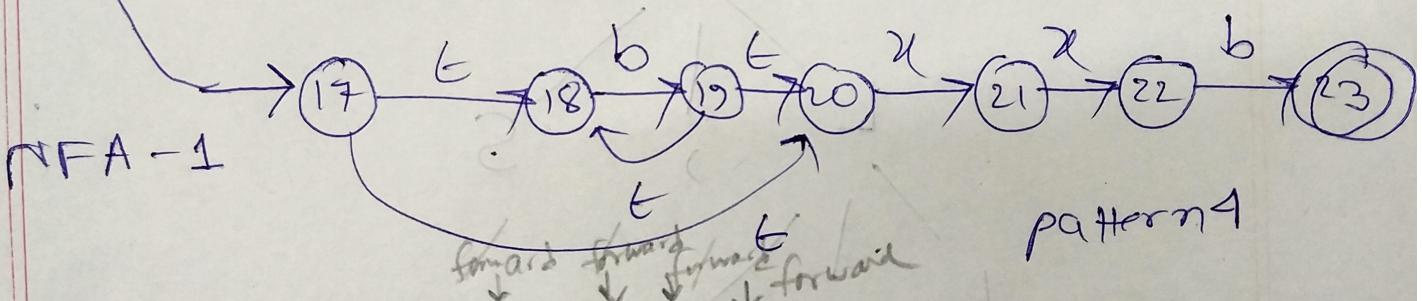
pattern 1



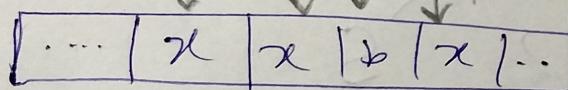
pattern 2



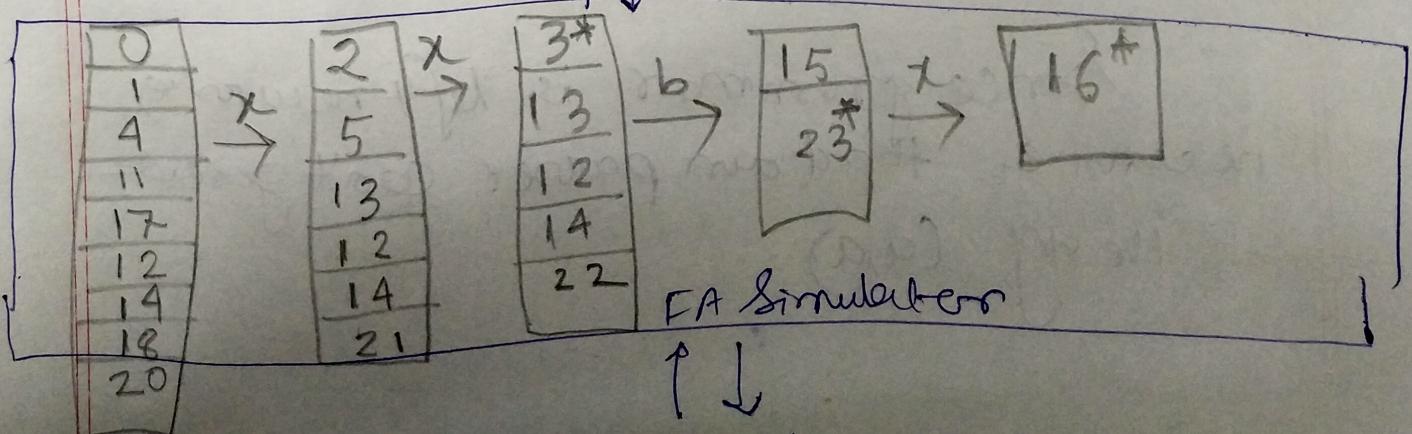
pattern 3



pattern 4



input buffer



So, $\lambda x b \lambda$ matches ^{NFA - 1} pattern 3 $\lambda^* (b + \lambda) \lambda$

(5)

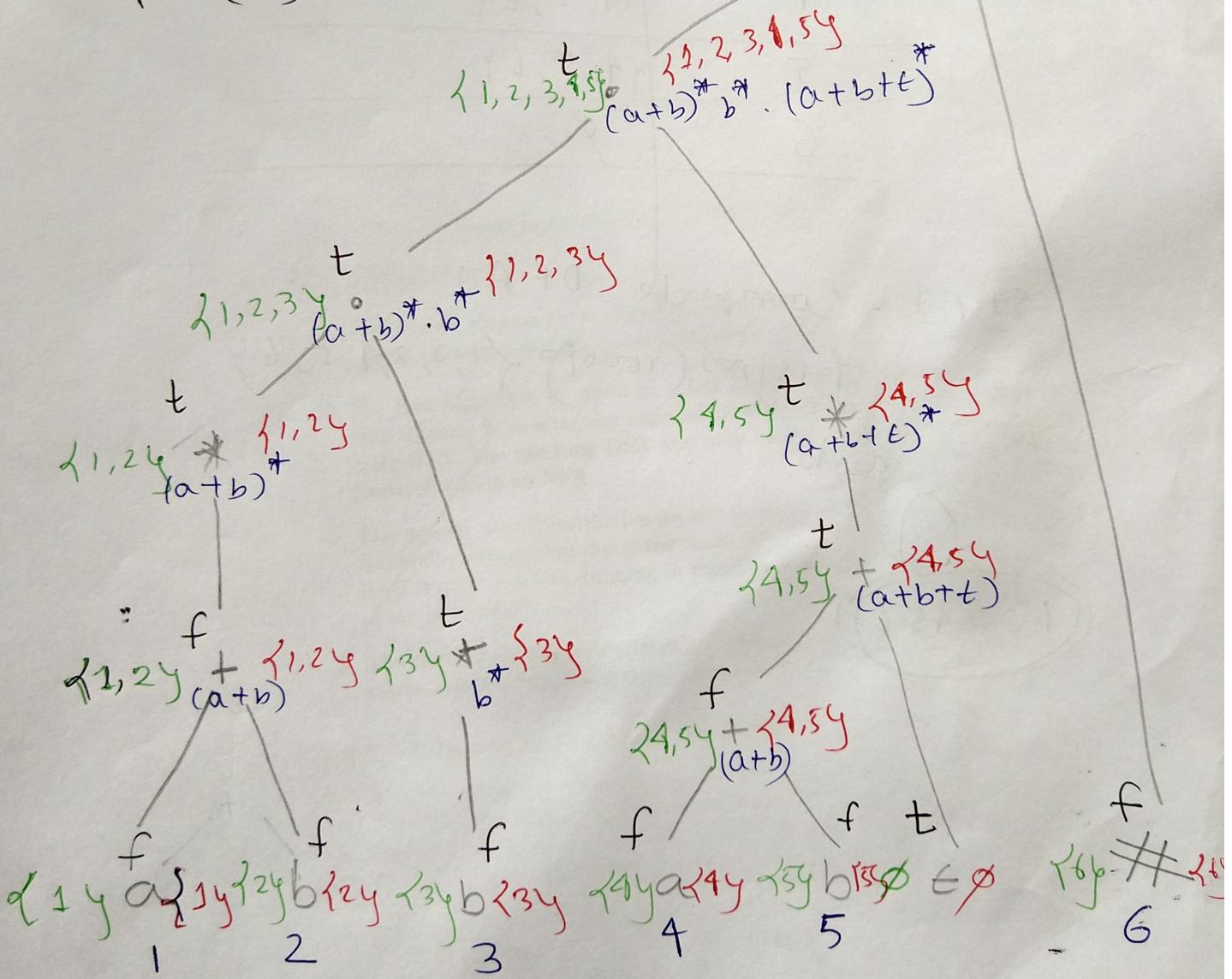
Convert the RE directly to DFA:

$$(a+b)^* b^* (a+b+t)^*$$

Solution:-

Step 1: Augmented RE: $(a+b)^* b^* (a+b+t)^* \#$

Step 2: Computing nullable (n), firstpos(n), lastpos(n):

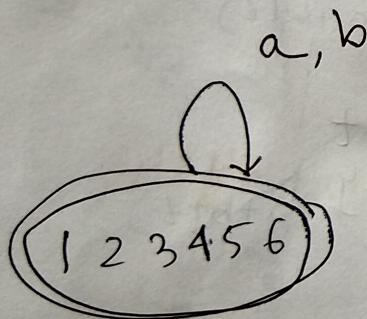


Step 3: Compute $\text{followpos}(n)$:

Node (n)	$\text{followpos}(n)$
1	$\{1, 2, 3, 4, 5, 6\}$
2	$\{1, 2, 3, 4, 5, 6\}$
3	$\{3, 4, 5, 6\}$
4	$\{4, 5, 6\}$
5	$\{4, 5, 6\}$
6	\emptyset

Step 4: Compute DFA:

$\text{firstpos}(\text{root}) = \{1, 2, 3, 4, 5, 6\}$

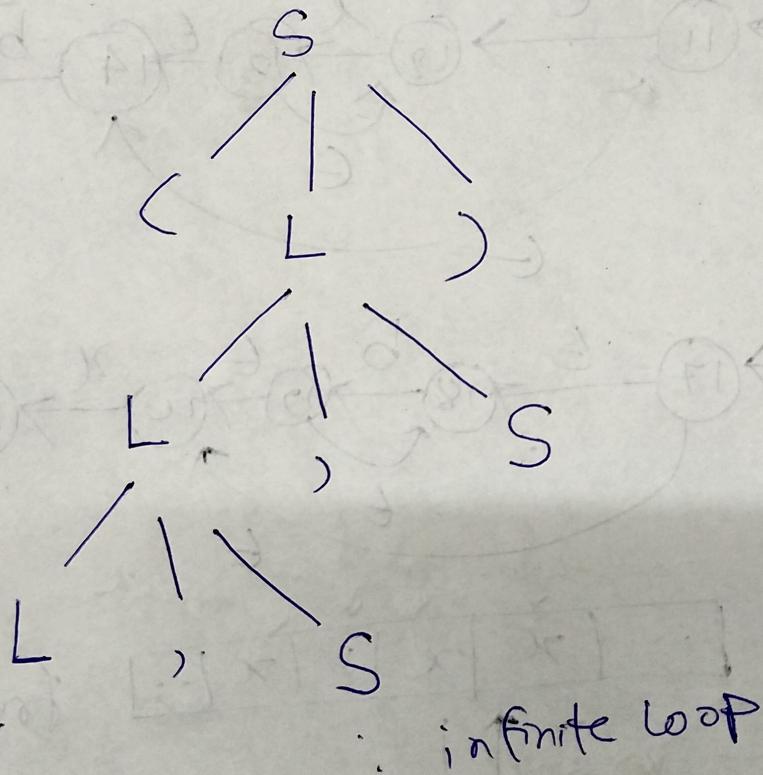


3/

$$S \rightarrow (L) | a$$

$$L \rightarrow L, S | S$$

Input "(a,a)"



Since grammar is left recursive,
recursive top down parser can't generate
the string (a,a)

The solution is to eliminate left recursion.

$$\begin{array}{l} S \rightarrow (L) | a \\ \frac{L}{A} \rightarrow \frac{L, S}{A^* \alpha} | \frac{S}{\beta} \end{array}$$

Solution:

$$\begin{array}{l} S \rightarrow (L) | a \\ L \rightarrow S L' \\ L' \rightarrow , S L' | \epsilon \end{array}$$