

- (a) Design an equivalent left-recursion free and left factored grammar of the following grammar G.

$$P \rightarrow (E) \mid a$$

$$E \rightarrow E, P \mid P$$

Applying left recursion removal,

$$P \rightarrow (E) \mid a$$

$$E \rightarrow P E'$$

$$E' \rightarrow , P E' \mid \epsilon$$

- (b) Construct FIRST and FOLLOW sets for non-terminals of the resulting grammar.

$$\text{FIRST}(P) = \{ (, a \}$$

$$\text{FIRST}(E) = \{ (, a \}$$

$$\text{FIRST}(E') = \{ , , \epsilon \}$$

$$\text{FOLLOW}(P) = \{ \$, , , ) \}$$

$$\text{FOLLOW}(E) = \{ ) \}$$

$$\text{FOLLOW}(E') = \{ ) \}$$



(c) Construct LL(1) Parse Table from the resulting grammar.

Ans:

	\$	,	(	)	a
P			$P \rightarrow (E)$		$P \rightarrow a$
E			$E \rightarrow PE'$		$E \rightarrow PE'$
E'		$E' \rightarrow , PE'$		$E' \rightarrow E$	

(d) Validate the following Syntax through Predictive Parsing using the constructed Parse Table.  
 $(a, (a, (a, a)))$

$( \quad a \quad , \quad ( \quad a \quad , \quad ( \quad a \quad , \quad a \quad ) \quad ) \quad ) \quad )$   
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Output

$P \rightarrow (E)$	✓	✗	
$E \rightarrow PE'$	✗	✓	
$P \rightarrow a$	✓	✗	
$E' \rightarrow , PE'$	✗	✗	
$P \rightarrow (E)$	✗	✓	
$E \rightarrow PE'$	✗	✗	
$P \rightarrow a$	✗	✓	
$E' \rightarrow , PE'$	✗	✓	
$P \rightarrow (E)$	✓	✗	
$E \rightarrow PE'$	✗	✗	
$P \rightarrow a$	✓	✗	
$E' \rightarrow , PE'$	✗	✗	✓
$P \rightarrow a$	\$	E'	✓

$E' \rightarrow E$   
 $E' \rightarrow E$   
 $E' \rightarrow E$

Stack