

- a)
- a) Fully observable since we can see the entire map.
 - b) It is deterministic since ~~the~~ hence the agent's current state and the action selected can determine the next state / change in environment.
 - c) Sequential because the action taken will impact the next state.
 - d) Static because the map will not change.
 - e) Discrete because the number of states ~~are~~ in the environment is finite.
 - f) Single agent because there is only one agent.

b)

i) COVID-19 positive patient tracking system agent:-

- Performance: - how fast it can track the patients, ~~area~~ optimality of the search/tracking.
- Environment: - ~~area~~ 10/20 feet area surrounding the agent.
- Actuators: - notification to the agent/alarm.
- Sensors: - GPS for tracking the patient's devices, ~~other~~ other sensors.

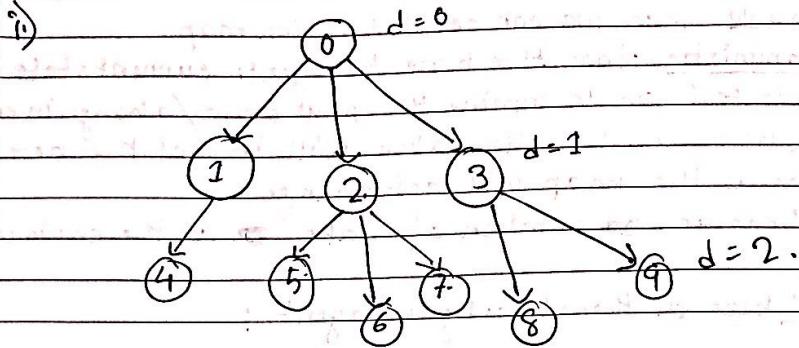
ii) Exploring the surface area of MARS.

- Performance: - mobility on rough surface, camera quality.
- Environment: - the surface of Mars / environment of Mars.
- Actuators: - accelerator, brakes, etc.
- Sensors: - IR sensor, video, sonar sensor, engine sensor, ~~GPS~~ and GPS.

Question 2

DATE

depth limit = 2.



Stack

Stack

0	1	2	3	4	5	6	7	8	9	= end
0	1	2	3	4	5	6	7	8	9	= end

Hence $d < \text{depth of goal node}$ so this is ~~not~~ incomplete!

4	1	0
---	---	---

$$b = 5, m = 10$$

ii) Total number of nodes generated = $b^m = 5^{10} = 9765625$

$$\Rightarrow 5^1 + 5^2 + 5^3 + 5^4 + 5^5 + 5^6 + 5^7 + 5^8 + 5^9 + 5^{10} = 12207030$$

Total time taken:-

10000000 $\xrightarrow{\text{node}}$ generated in 1 sec.

1 node \longrightarrow " " $\frac{1}{10000000}$

$$12207030 \longrightarrow " " \frac{12207030}{10000000} = 12.20703 \text{ secs.}$$

total space required :-

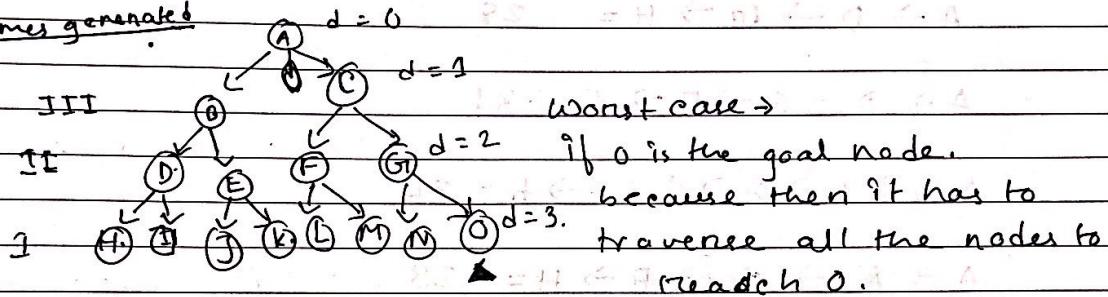
$$\begin{aligned} & O(bm + 1) \\ \Rightarrow & 0(5 \times 10) + 1 \times 1000 \text{ bytes} \\ = & 5100 \text{ bytes} \end{aligned}$$

1000 byte \rightarrow 1 kb.

$$5100 \text{ bytes} \rightarrow \frac{5100}{1000} = 5.1 \text{ kilobytes.}$$

iii) For the ~~worst~~ worst-case :-

no. of times generated



worst case \rightarrow

if O is the goal node.

because then it has to

traverse all the nodes to reach O.

② let b be the branching factor, & d is the depth.

so,

$$N(IDS) = db + (d-1)b^2 + (d-2)b^3 + \dots + (1)b^d.$$

$\therefore O(b^d)$ \leftarrow is the time complexity.

$db \rightarrow d=1$ is generated

~~at depth 1~~ \rightarrow generated

at depth = 1, it is generated d number of times.

$(d-1)b^2 \rightarrow$ at depth = 2, it is generated ^{fed} $(d-1)$ number of times.

$(1)b^d \rightarrow$ at depth = d, it is generated 1 time only.

Typo

Question 3

DATE

?) All the path cost from A to H

$$A \rightarrow B \rightarrow E \rightarrow H = 23$$

$$A \rightarrow B \rightarrow F \rightarrow H = 25$$

$$A \rightarrow C \rightarrow F \rightarrow H = 22$$

$$\checkmark A \rightarrow C \rightarrow G \rightarrow H = 20$$

$$A \rightarrow D \rightarrow G \rightarrow H = 28$$

$$A \rightarrow D \rightarrow G \rightarrow F \rightarrow H = 31$$

$$A \rightarrow D \rightarrow G \rightarrow C \rightarrow F \rightarrow H = 36$$

$$A \rightarrow \cancel{F} \rightarrow G \rightarrow F \rightarrow H = 23$$

$$A \rightarrow C \rightarrow F \rightarrow B \rightarrow E \rightarrow H = 30$$

$$A \rightarrow B \rightarrow F \rightarrow G \rightarrow H = 26$$

$$A \rightarrow B \rightarrow F \rightarrow C \rightarrow G \rightarrow H = 31$$

$$A \rightarrow C \rightarrow F \rightarrow G \rightarrow H = 23$$

~~A → D → G → F → E → H~~

$$A \rightarrow D \rightarrow G \rightarrow C \rightarrow F \rightarrow B \rightarrow E \rightarrow H = 44$$

$$A \rightarrow D \rightarrow G \rightarrow F \rightarrow B \rightarrow E \rightarrow H = 39$$

∴ ~~the~~ The smallest path cost from A → H is 20

~~heuristic~~ value of A is, $h(A) = 12$

∴ pathcost > heuristic value

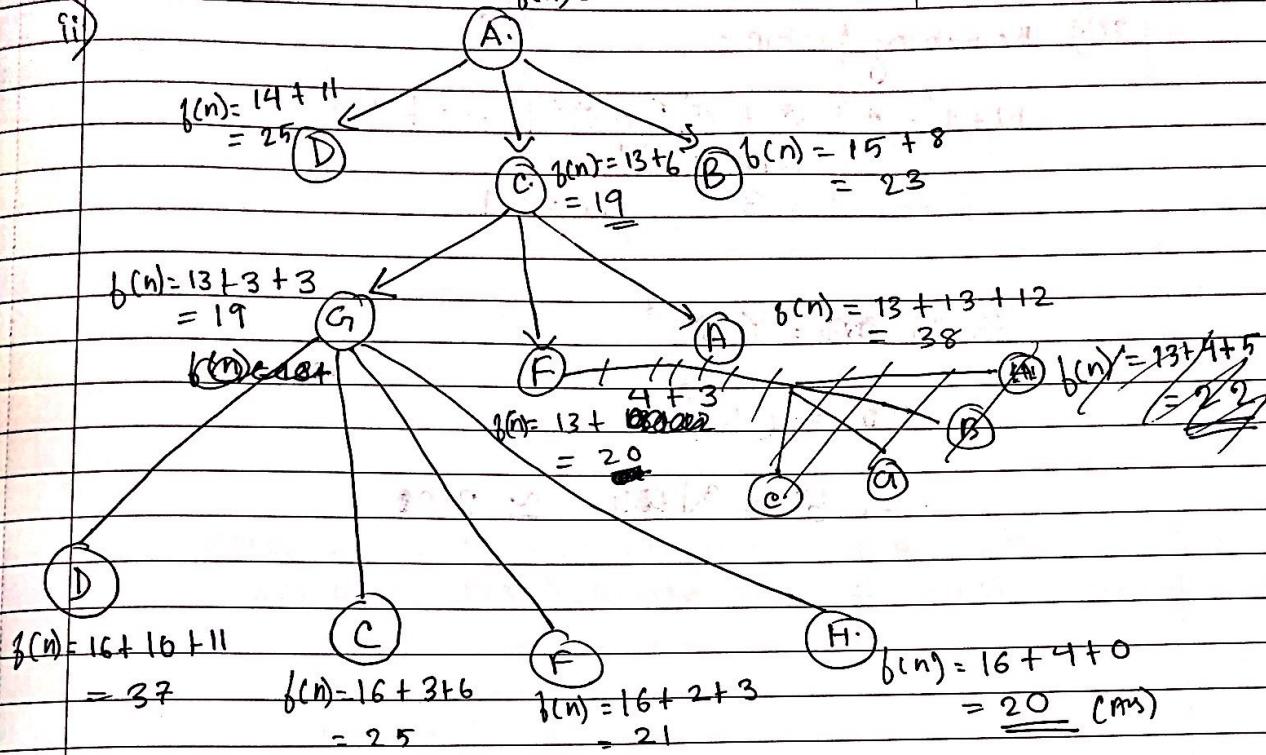
$$20 > 12$$

(shown.)

$$f(n) = 12 + 0 = 12$$

DATE _____

ii)



let the optimal path cost be C^*

~~Let the cost of optimal solution be C^*~~

Hence, even after finding the solution, since there was a node ~~with~~ with the same path cost as C^* , it expanded the other node just to check if the ~~value~~ path cost to goal via that node is lesser than C^* .
if there was a

let the suboptimal goal node be G_1 ,

$$\therefore f(G_1) \text{ } \cancel{\text{=}} \text{ } g(G_1) + h(G_1)$$

~~$$\therefore C^* \neq f(G_1) > C^*$$

 2. $\cancel{\text{if }} \cancel{\text{shown}}$~~

If there was a suboptimal path to the goal node, the A* search will not stop there, it will expand other nodes to find the most optimal solution.

Typo

iii) Branching factor:-

$$N+1 = 1 + b^* + (b^*)^2 + (b^*)^3 + \dots + (b^*)^d.$$

~~$N+1 \approx (b^*)^d$~~

here, $d = 3$, $N = 11$.

$$\Rightarrow 11 = (b^*)^3$$

$$\Rightarrow 11 = (b^*)^3$$

$$\therefore b^* = \sqrt[3]{11} \approx 2.11$$

Question 4

DATE

String	n	r	F
010110	22	88	7744π
100001	33	132	17424π
010001	17	68	4624π
101010	42	168	28224π

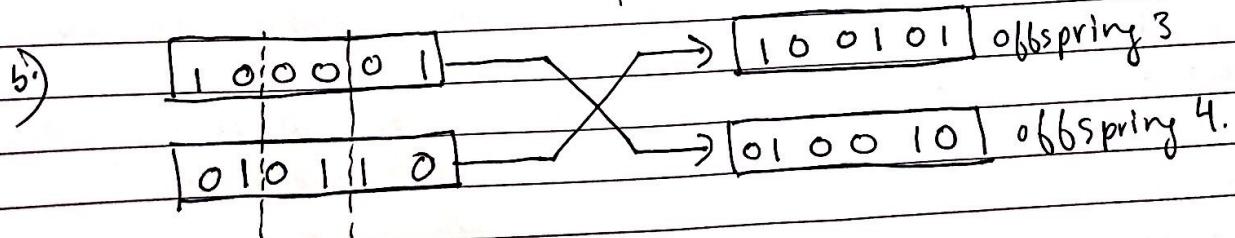
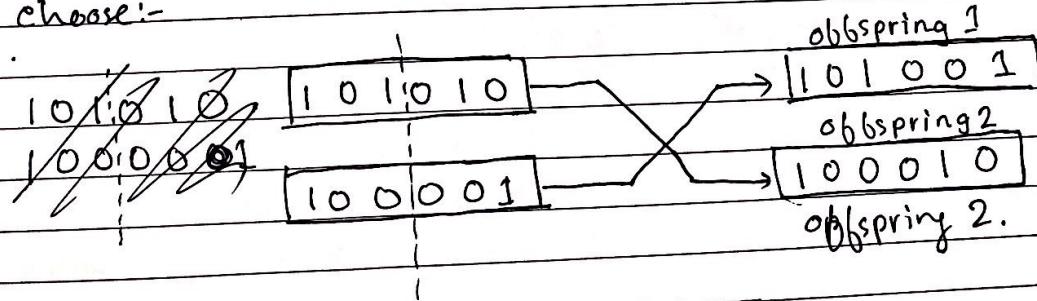
i) Here, we are always keeping $k=4$ states in the memory where all of these works as a candidate for crossover. :-

010110
100001
010001
101010

iii) The chromosomes chosen according to the fitness are:-

010110
100001
101010

iv.) a.) we choose:-



Typo

iv) Mutation:-

Final Offspring
 $10100\downarrow \rightarrow 101011$

~~100010~~ \rightarrow ~~100000~~ 110010

$100101 \rightarrow 100101$

$0100\downarrow 10 \rightarrow 0100110$

Diagram of mutation is as follows:

Initial

Second

Third

Fourth

Final Offspring

Initial

Second

Third

Fourth

Fifth

Sixth

Seventh

Eighth

Ninth

Tenth

Eleventh

Twelfth

Thirteenth

Fourteenth

Fifteenth

Sixteenth

Seventeenth

Eighteenth

Nineteenth

Twentieth

Twenty-first

Twenty-second

Twenty-third

Twenty-fourth

Twenty-fifth

Twenty-sixth

Twenty-seventh

Twenty-eighth

Twenty-ninth

Thirty

Thirty-one

Thirty-two

Thirty-three

Thirty-four

Thirty-five

Thirty-six

Thirty-seven

Thirty-eight

Thirty-nine

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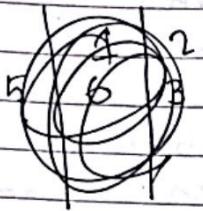
Forty-nine

Forty

Question 5

DATE

Goal State.



1	2	3	
4	5	6	$f = 0$
7	8		

Let's consider the Manhattan distance as the heuristic value (since it dominates misplaced tiles)

$$\text{let, } f = -(\text{Manhattan distance})$$

start.

	1	2
5	6	3
4	7	8

$$f = - (1 + 1 + 1 + 1 + 1 + 1 + 1) = -8$$

$$= -(8)$$

move up

move left.

5	1	2
6	3	
4	7	8

$$f = -9$$

Ridge.

1	2	
5	6	3
4	7	8

$$f = -7$$

Local Minima left

move up

1	2	2
5	.	3
4	7	8

$$f = -8$$

Ridges

1	2	.
5	6	3
4	7	8

$$f = -6$$



Two problems that can occur are :-

(1) Local Minima:- Hill climbing will get stuck at a local minima, which is not a satisfactory solution.

(2) Ridges:- It will not overcome ridges. Search will only oscillate slowly.

Typo

To overcome these problems, we can modify the Hill climbing search escape search:-

(1) We can modify the choice of neighbor.

for example neighbors can be chosen randomly which will help Hill climbing search escape the local maxima ridges.

(2) We can modify the criteria for accepting neighbor for the current state.

for example:- Accept the neighbor if it is better, if it is not better, then accept it with some fixed probability.

This helps HIC to escape local maxima.