

SCLD LAB ASSIGNMENT 2

GROUP NO : 31

GROUP MEMBERS:

- 1.) SHREYAS JENA : 20CS30049
- 2.) PRANAV NYATI : 20CS30037
- 3.) UMANG SINGLA : 20CS10068
- 4.) UTSAV VINAY MEHTA : 20CS10069

1 Question 1 (BCD)

1. Develop circuits to convert from 4-bit binary to 2-digit BCD

Solution:

Table 1: Binary to BCD Converter

Binary Code (Input)				BCD Code (Output)				
A	B	C	D	V	W	X	Y	Z
0	0	0	0	0	0	0	0	0
0	0	0	1	0	0	0	0	1
0	0	1	0	0	0	0	1	0
0	0	1	1	0	0	0	1	1
0	1	0	0	0	0	1	0	0
0	1	0	1	0	0	1	0	1
0	1	1	0	0	0	1	1	0
0	1	1	1	0	0	1	1	1
1	0	0	0	0	1	0	0	0
1	0	0	1	0	1	0	0	1
1	0	1	0	1	0	0	0	0
1	0	1	1	1	0	0	0	1
1	1	0	0	1	0	0	1	0
1	1	0	1	1	0	0	1	1
1	1	1	0	1	0	1	0	0
1	1	1	1	1	0	1	0	1

From above table we can construct logical expressions for V, W, X, Y, Z as follow:

For V:

$$\begin{aligned}
 V &= \overline{A}\overline{B}\overline{C}\overline{D} + \overline{A}\overline{B}C\overline{D} + \overline{A}B\overline{C}\overline{D} + \overline{A}B\overline{C}D + \overline{A}BC\overline{D} + \overline{A}BCD \\
 V &= \overline{A}\overline{B}C(\overline{D} + D) + \overline{A}B\overline{C}(\overline{D} + D) + \overline{A}BC(\overline{D} + D) \\
 V &= \overline{A}\overline{B}C + \overline{A}B\overline{C} + \overline{A}BC \\
 V &= \overline{A}\overline{B}C + \overline{A}B(\overline{C} + C) \\
 V &= \overline{A}\overline{B}C + \overline{A}B \\
 V &= \overline{A}(\overline{B}C + B) \\
 V &= \overline{A}((B + C)(B + \overline{B})) \\
 \mathbf{V} &= \mathbf{\overline{A}B + \overline{A}C}
 \end{aligned}$$

For W:

$$\begin{aligned}
 W &= \overline{A}\overline{B}\overline{C}\overline{D} + \overline{A}\overline{B}\overline{C}D \\
 W &= \overline{A}\overline{B}\overline{C}(\overline{D} + D) \\
 \mathbf{W} &= \mathbf{\overline{A}\overline{B}\overline{C}}
 \end{aligned}$$

For X:

$$\begin{aligned}
X &= \overline{A}B\overline{C}.\overline{D} + \overline{A}B\overline{C}D + \overline{A}BC\overline{D} + \overline{A}BCD + ABC\overline{D} + ABCD \\
X &= \overline{A}B\overline{C}(\overline{D} + D) + \overline{A}BC(\overline{D} + D) + ABC(\overline{D} + D) \\
X &= \overline{A}B\overline{C} + \overline{A}BC + ABC \\
X &= \overline{A}B(\overline{C} + C) + ABC \\
X &= \overline{A}B + ABC \\
X &= B(\overline{A} + AC) \\
X &= B((\overline{A} + A)(\overline{A} + C)) \\
\mathbf{X} &= \mathbf{\overline{A}B + BC}
\end{aligned}$$

For Y:

$$\begin{aligned}
Y &= \overline{A}.\overline{B}C\overline{D} + \overline{A}.\overline{B}CD + \overline{A}BC\overline{D} + \overline{A}BCD + AB\overline{B}.\overline{D} + AB\overline{C}D \\
Y &= \overline{A}.\overline{B}C(\overline{D} + D) + \overline{A}BC(\overline{D} + D) + AB\overline{C}(\overline{D} + D) \\
Y &= \overline{A}.\overline{B}C + \overline{A}BC + AB\overline{C} \\
Y &= \overline{A}C(\overline{B} + B) + AB\overline{C} \\
\mathbf{Y} &= \mathbf{A\overline{B}C + \overline{A}C}
\end{aligned}$$

For Z:

For Z, it can be easily seen from the table that Z is 1 whenever D is one. So, we can conclude that:
 $\mathbf{Z = D}$

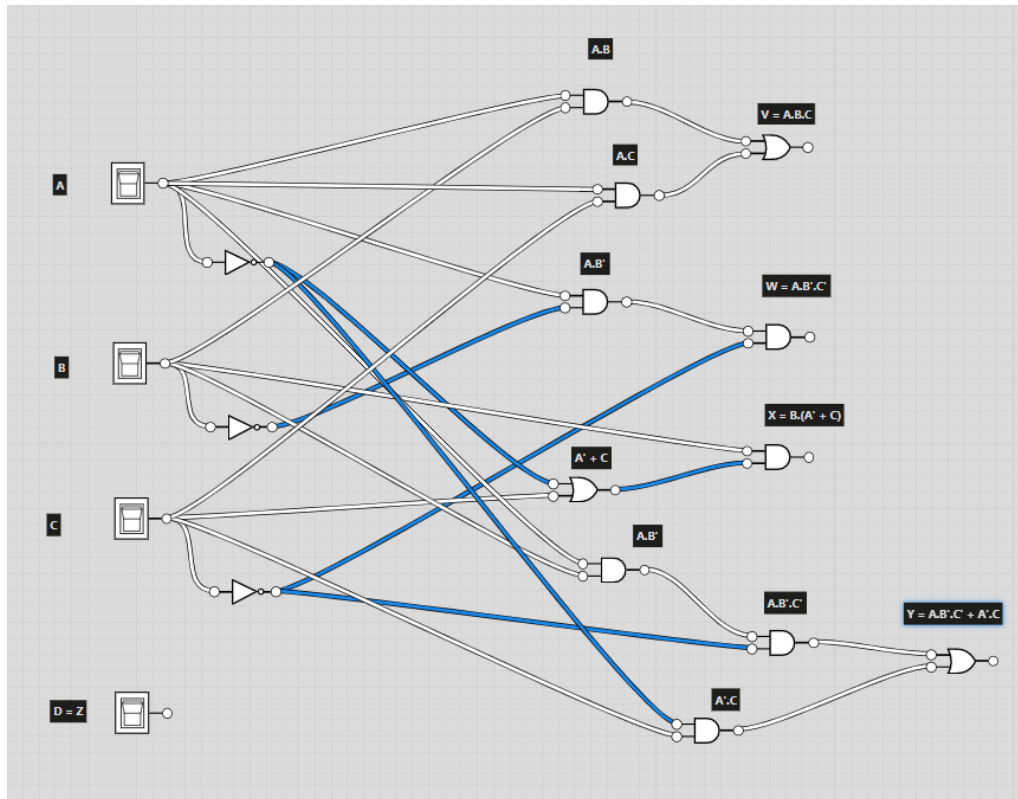


Fig 1: Circuit diagram to convert from 4-bit binary to 2-digit BCD

2 Question 2 (Gray)

1. Develop circuits to convert from 4-bit Gray to 4-bit binary and vice-versa

Solution:

Table 2: Binary to Gray Converter

Binary Code (Input)				Gray Code (Output)			
B1	B2	B3	B4	G1	G2	G3	G4
0	0	0	0	0	0	0	0
0	0	0	1	0	0	0	1
0	0	1	0	0	0	1	1
0	0	1	1	0	0	1	0
0	1	0	0	0	1	1	0
0	1	0	1	0	1	1	1
0	1	1	0	0	1	0	1
0	1	1	1	0	1	0	0
1	0	0	0	1	1	0	0
1	0	0	1	1	1	0	1
1	0	1	0	1	1	1	1
1	0	1	1	1	1	1	0
1	1	0	0	1	0	1	0
1	1	0	1	1	0	1	1
1	1	1	0	1	0	0	1
1	1	1	1	1	0	0	0

Conversion from binary code to gray code:

1. Most Significant Bit (MSB) of Gray code is same as that of Binary code.
2. For subsequent bits, the below mentioned rule is followed:

$$G_i = B_i \oplus B_{i+1}$$

3. Thus, for 4-bit binary to 4-bit gray conversion, we have the following:

$$\begin{aligned}G1 &= B1 \\G2 &= B1 \oplus B2 \\G3 &= B2 \oplus B3 \\G4 &= B3 \oplus B4\end{aligned}$$

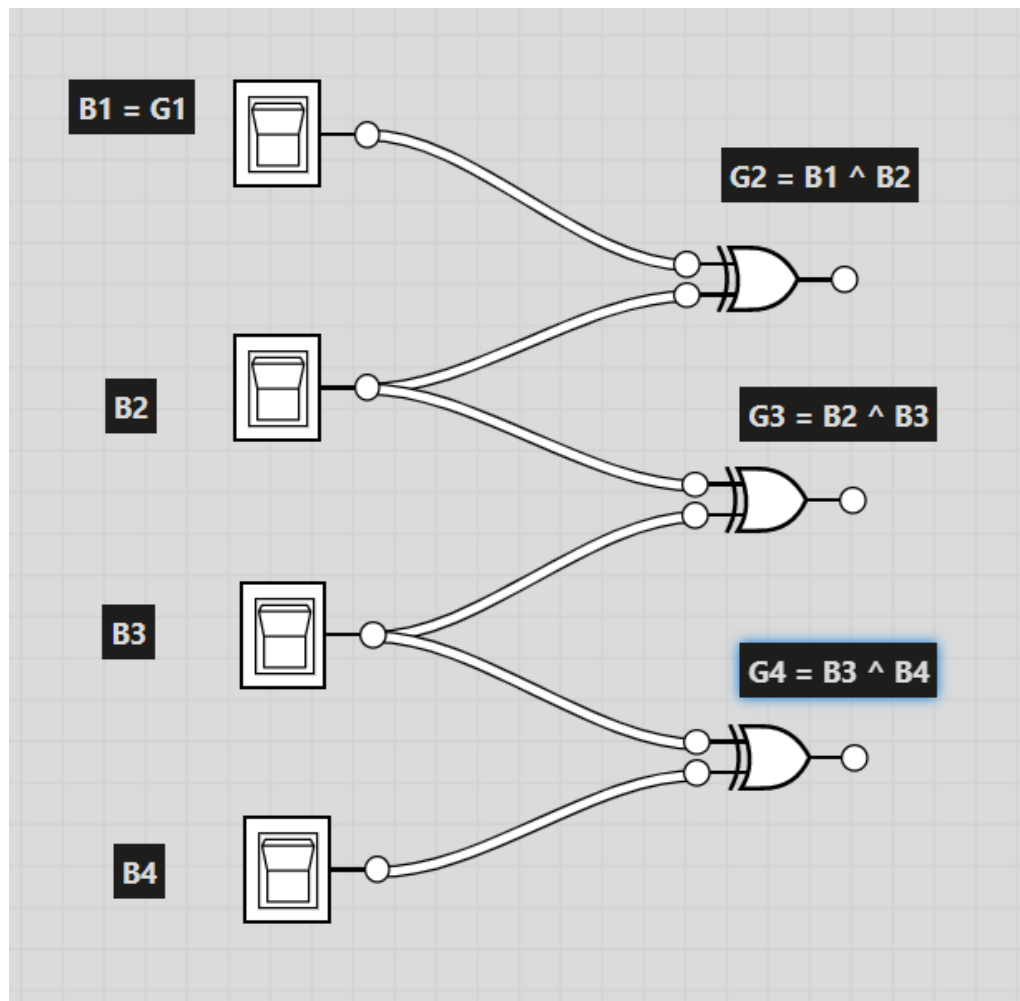


Fig 2: Circuit diagram to convert from 4-bit binary to 4-bit Gray Code.

Conversion from gray code to binary code:

1. Most Significant Bit (MSB) of Binary code is same as that of Gray code.
2. For subsequent bits, the below mentioned rule is followed:

$$B_i = G_i \oplus B_{i-1}$$

3. Thus, for 4-bit gray to 4-bit binary conversion, we have the following:

$$\begin{aligned} B1 &= G1 \\ B2 &= G2 \oplus B1 \\ B3 &= G3 \oplus B2 \\ B4 &= G4 \oplus B3 \end{aligned}$$

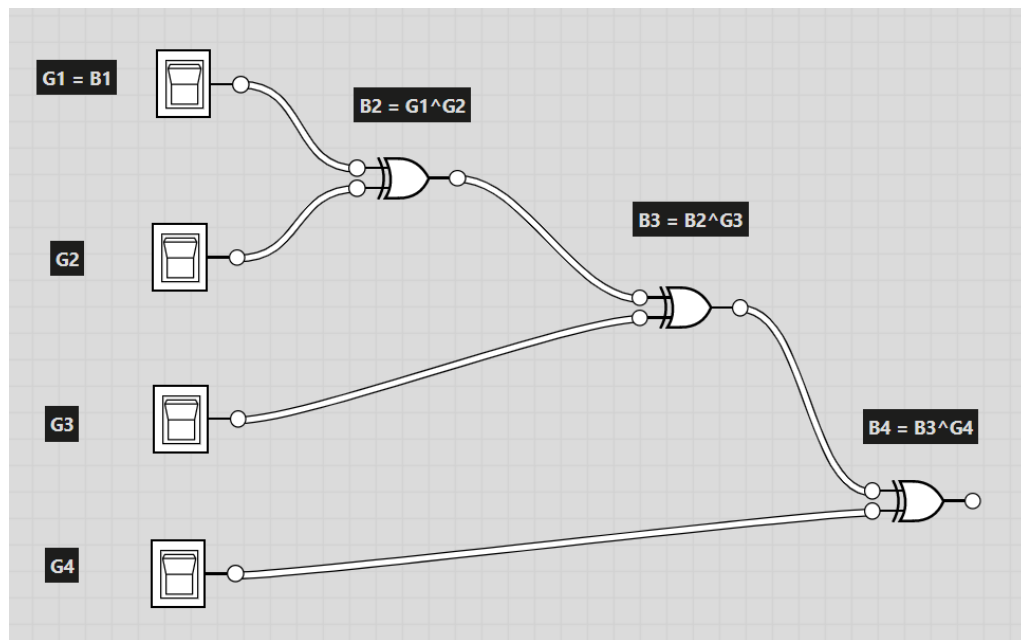


Fig 3: Circuit diagram to convert from 4-bit Gray Code to 4-bit binary.

3 Question 3 (Excess-3)

1.
 - Develop a half adder for handling two bits
 - Develop a full adder using half adders and any additional logic
 - Develop a ripple carry adder needed for this assignment using full adders
 - Develop circuits to convert from excess-3 to 4-bit binary and vice-versa

Solution:

a) The circuit that adds two bits is called half adder

A	B	Sum	Carry
0	0	0	0
0	1	1	0
1	0	1	1
1	1	0	1

Table 3: Truth Table for Half Adder

From the above truth table it can be seen that:

$$\text{Sum} = \overline{A}B + A\overline{B} = A \text{ XOR } B$$

$$\text{Carry} = AB$$

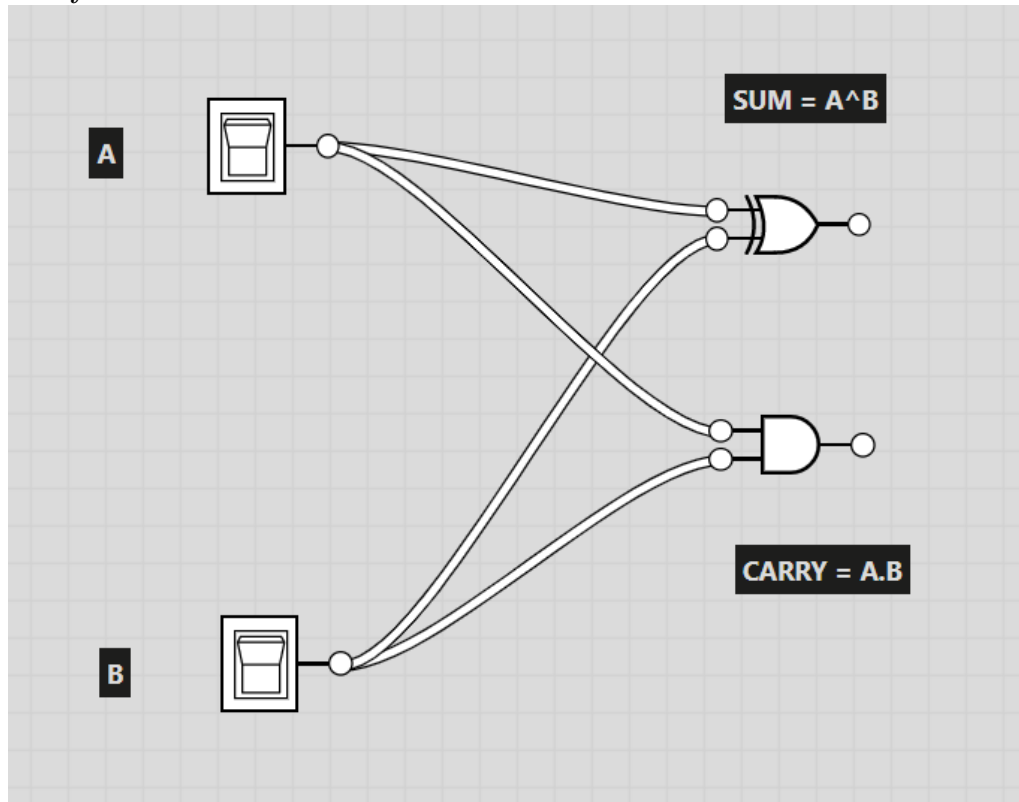


Fig 2: Circuit diagram for Half Adder

- b) Full Adder is the adder which adds three inputs and produces two outputs. The first two inputs are A and B and the third input is an input carry as C-IN. The output carry is designated as C-OUT and the normal output is designated as S which is SUM.

Input			Output	
P	Q	C-IN	Sum	C-Out
0	0	0	0	0
0	0	1	1	0
0	1	0	1	0
0	1	1	0	1
1	0	0	1	0
1	0	1	0	1
1	1	0	0	1
1	1	1	1	1

Table 4: Truth Table of Full Adder

From the above truth table it can be seen that:

For Sum:

$$\begin{aligned}\text{Sum} &= \overline{P}.\overline{Q}C + \overline{P}Q\overline{C} + P\overline{Q}.\overline{C} + PQC \\ \text{Sum} &= C(\overline{P}.\overline{Q} + P\overline{Q}) + \overline{C}(\overline{P}Q + P\overline{Q}) \\ \text{Sum} &= C \text{ XOR } (P \text{ XOR } Q)\end{aligned}$$

For Carry:

$$\begin{aligned}\text{Carry} &= \overline{P}BC + P\overline{Q}C + PQ\overline{C} + PQC \\ \text{Carry} &= (\overline{P}Q + P\overline{Q})C + PQ(\overline{C} + C) \\ \text{Carry} &= (P \text{ XOR } Q)C + PQ\end{aligned}$$

From the above question we can construct the circuit for Full adder as follows:

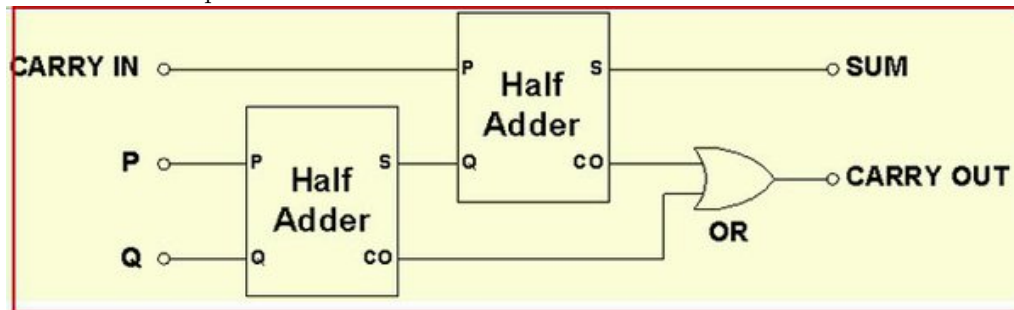


Fig 4: Circuit diagram for full adder

- c) A ripple carry adder is a logic circuit in which the carry-out of each full adder is the carry in of the succeeding next most significant full adder. It is called a ripple carry adder because each carry bit gets rippled into the next stage.

A1	A2	A3	A4	B4	B3	B2	B1	S4	S3	S2	S1	Carry
0	0	0	0	0	0	0	0	0	0	0	0	0
0	1	0	0	0	1	0	0	1	0	0	0	0
1	0	0	0	1	0	0	0	0	0	0	0	1
1	0	1	0	1	0	1	0	0	1	0	0	1
1	1	0	0	1	1	0	0	1	0	0	0	1
1	1	1	0	1	1	1	0	1	1	0	0	1
1	1	1	1	1	1	1	1	1	1	1	0	1

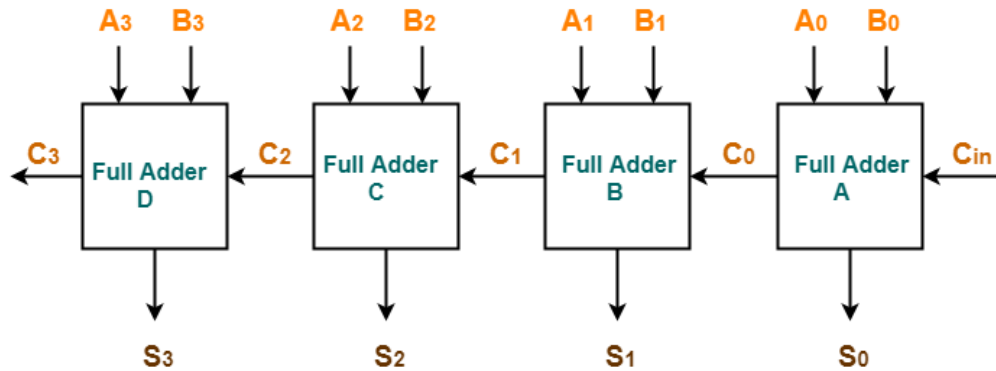
The general equation for n-bit ripple carry adders is:

$$SUM_i = FULL - ADDER(A_i, B_i, CARRY_{i-1})$$

$$CARRY_i = FULL - ADDER(A_i, B_i, CARRY_{i-1})$$

$$CARRY_0 = 0$$

From the above question we can construct the circuit for Ripple Adder as follows:



4-bit Ripple Carry Adder

Fig 5: Circuit diagram for Ripple Adder

d) Excess-3: