

# Optimization #1

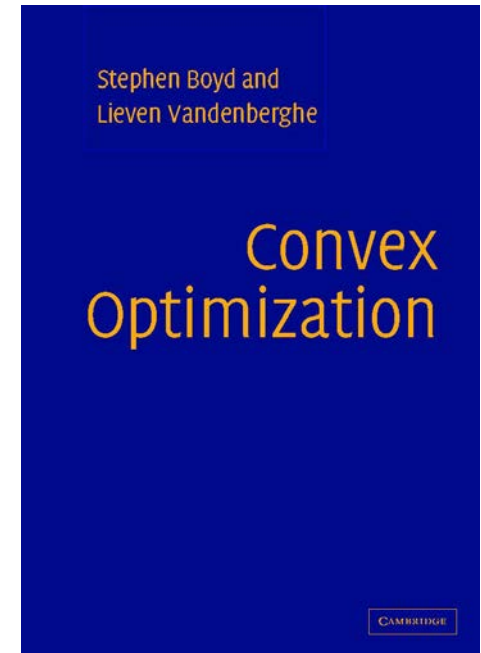
Dep. EE, KPU

2020년 2학기

# Book & Homepage

## ★ Convex Optimization

- Stephen Boyd & Lieven Vandenberghe
  - Cambridge University Press (2009)
- Authors' homepage
  - <https://web.stanford.edu/~boyd/cvxbook/>
- Boyd's Lecture Videos:
  - <https://web.stanford.edu/class/ee364a/videos.html>
- 교재의 pdf 파일
  - <http://www.stanford.edu/~boyd/cvxbook/> 에서 다운로드 가능



# Basic Concept & Definition

# Optimization Problems

$$\begin{array}{llll} \min_{\mathbf{x} \in \mathbb{R}^n} f(\mathbf{x}) & & \text{Objective function} \\ \text{subject to} & g_i(\mathbf{x}) \leq 0, & i = 1, \dots, r & \text{inequality constraints} \\ & h_k(\mathbf{x}) = 0, & k = 1, \dots, m & \text{equality constraints} \end{array}$$

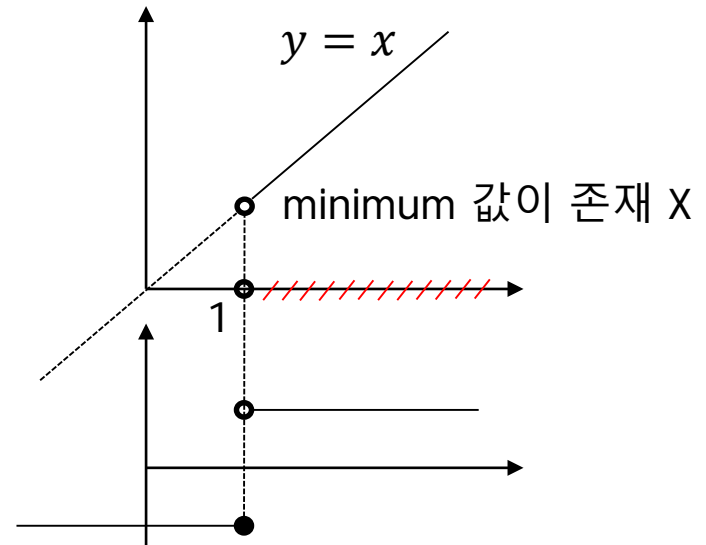
$n \times 1$  vector  $\rightarrow$   $\mathbf{x}$

- Feasible Solution
  - Constraints를 만족시키는 solution
- Feasible Set
  - Feasible solution 들의 모든 집합
- Optimal Solution
  - $\mathbf{x}^*$ 는 모든 constraints를 만족시킴
  - constraints를 만족시키는 모든  $\bar{\mathbf{x}}$ 에 대해서  $f(\mathbf{x}^*) \leq f(\bar{\mathbf{x}})$

# Optimization Problems

- 최적화 문제는 항상 optimal solution을 가지는가?

$$\begin{array}{ll}\min & x \\ \text{s. t.} & g(x) \geq 0 \\ \text{where} & \\ g(x) = & \begin{cases} 1, & \text{for } x > 1 \\ -1, & \text{for } x \leq 1 \end{cases}\end{array}$$



This problem does not have an optimal solution

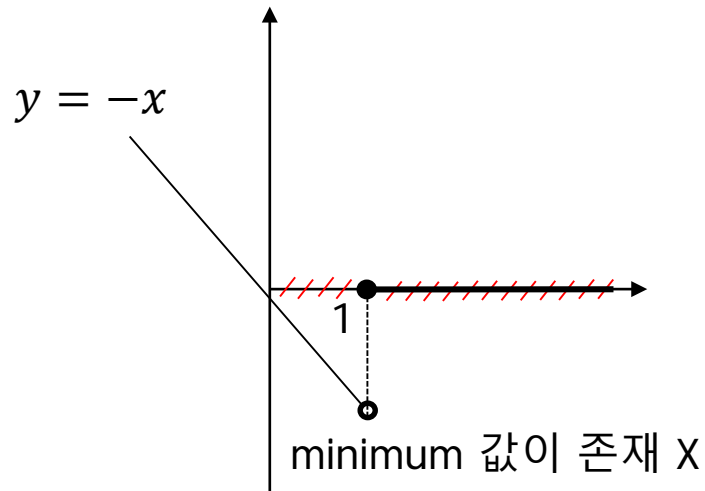
# Optimization Problems

$$\min f(x)$$

$$\text{s. t. } x \geq 0$$

where

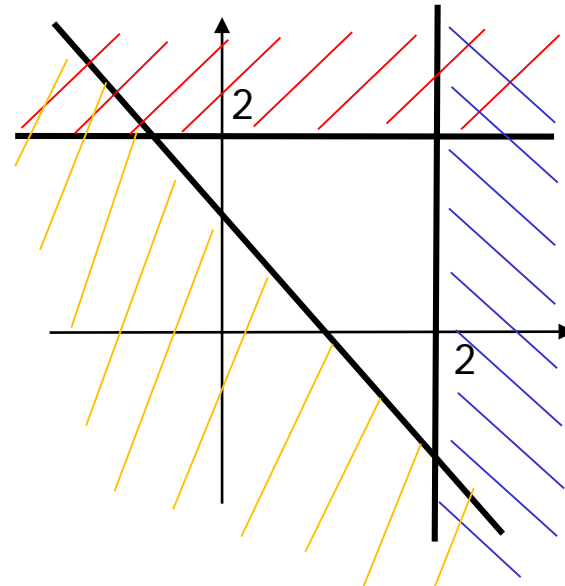
$$f(x) = \begin{cases} -x, & \text{for } x < 1 \\ 0, & \text{otherwise} \end{cases}$$



This problem does not have an optimal solution

# Infeasible Problem

$$\begin{array}{ll}\min & x_1 - x_2 \\ \text{s. t.} & \\ & x_1 + x_2 \leq 1 \\ & x_1 \geq 2 \\ & x_2 \geq 2\end{array}$$

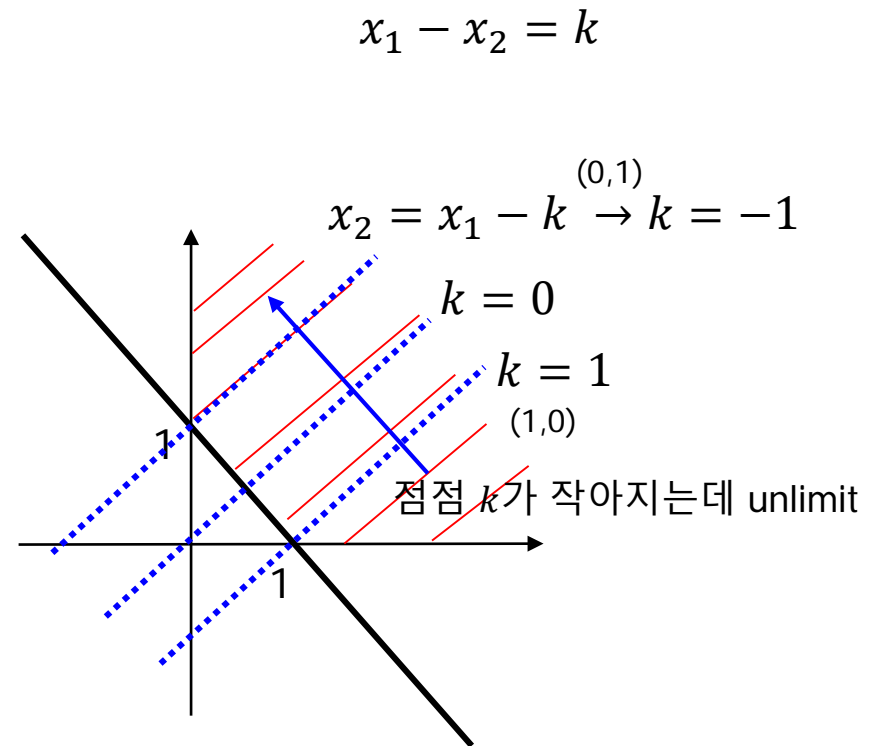


3개의 constraints를 동시에 만족시키는 해가 없음

→ Infeasible problem

# Unbounded Problem

$$\begin{array}{ll}\min & x_1 - x_2 \\ \text{s. t.} & x_1 + x_2 \geq 1 \\ & x_1, x_2 \geq 0\end{array}$$



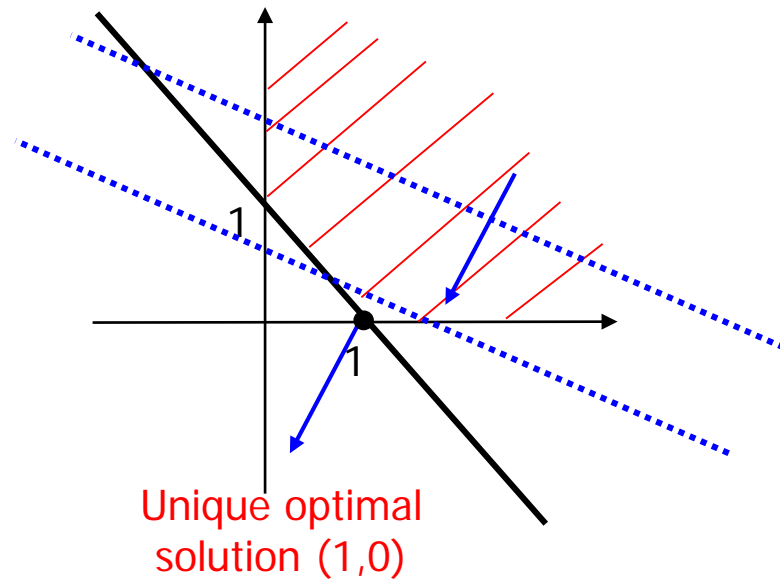
Objective function decreases w/o any bound if the solution moves along vector  $(1,0)$  from point  $(0,1)$

Feasible problem이냐 unbounded로 solution이 없음



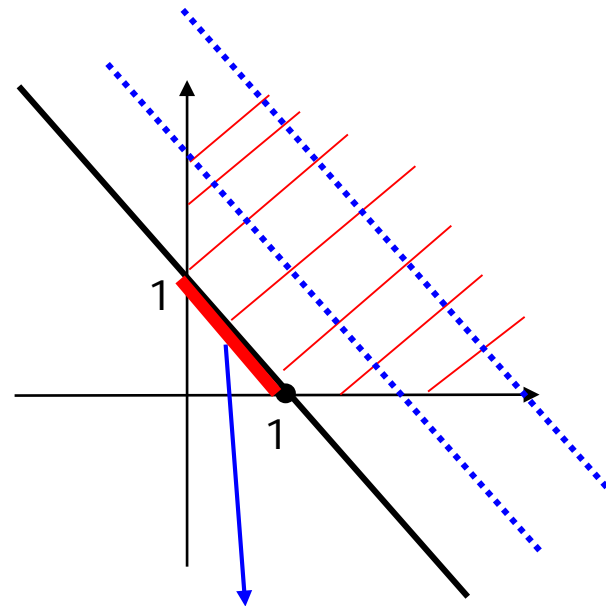
# Unique optimal solution

$$\begin{array}{ll}\min & x_1 + 2x_2 \\ \text{s. t.} & x_1 + x_2 \geq 1 \\ & x_1, x_2 \geq 0\end{array}$$



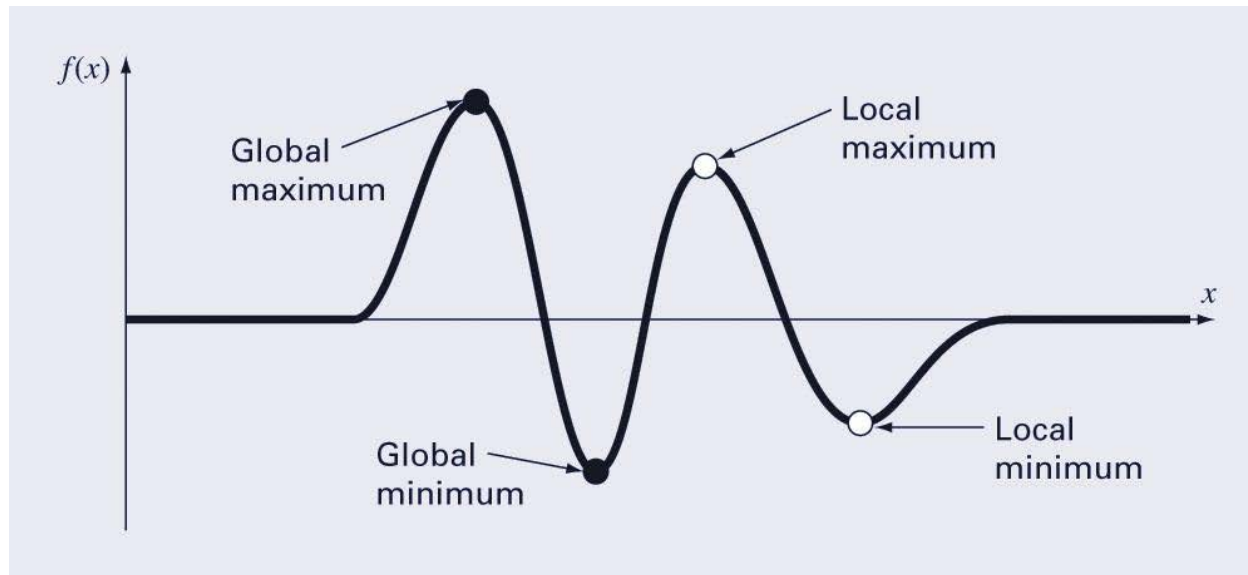
# Multiple optimal solutions

$$\begin{array}{ll}\min & x_1 + x_2 \\ \text{s. t.} & x_1 + x_2 \geq 1 \\ & x_1, x_2 \geq 0\end{array}$$



Multiple optimal  
solutions

# Global vs Local Optimum



# Unconstrained Optimization Problem

- 1D Optimization Problem

# Unconstrained Optimization Problem

$$\min_{x \in R} f(x)$$

- (0<sup>th</sup> order info.) only objective value  $f$  is available
- (1<sup>st</sup> order info.)  $f'$  is available, but not  $f''$
- (2<sup>nd</sup> order info.) both  $f'$  and  $f''$  are available
- Higher-order information tends to give more powerful algorithms.
- These methods are also used in multi-dimensional optimization as line search for determine how far to move along a given direction

Inequality constraints just limit the searching range! (Omit)

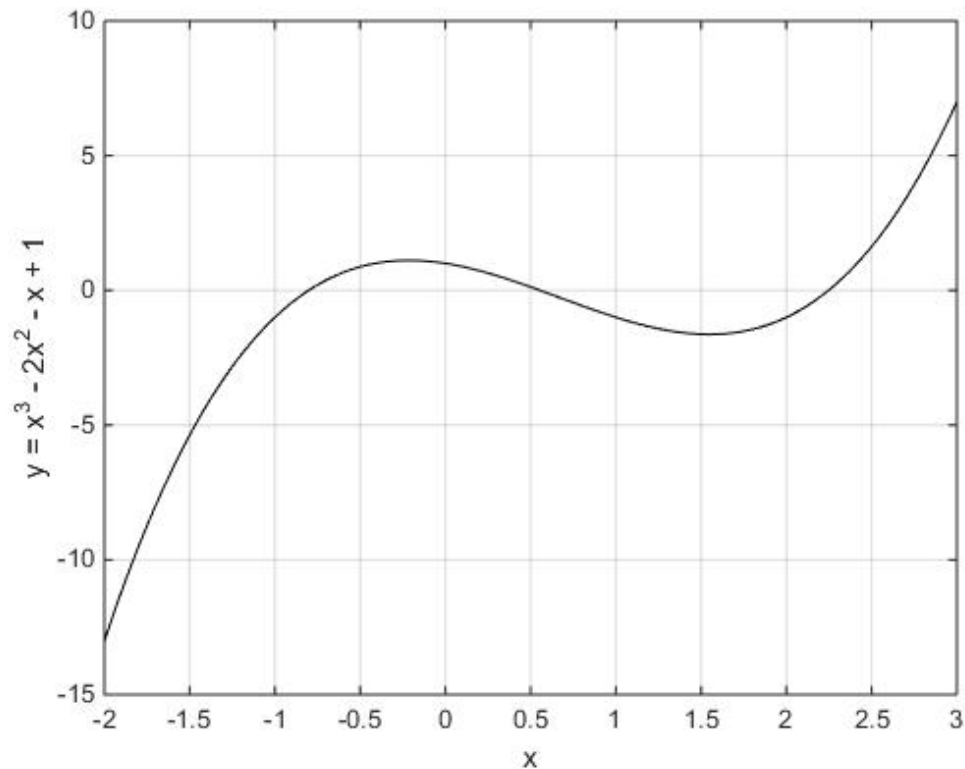
# Graphical Solution

## ■ Procedures

- ▶ 1. Plot each constraint as an equation and then decide which side of the line is feasible
- ▶ 2. Find the feasible region
- ▶ 3. Find the coordinations(좌표) of the corner points of the feasible region
- ▶ 4. Substitute the corner point coordinates in the objective function
- ▶ 5. Choose the optimal solution

# Graphical Solution

$$\min_{x \in \mathbb{R}} x^3 - 2x^2 - x + 1$$



고차식으로 가면? 그림을 그리기 용이하지 않음!!

# Iterative Algorithm

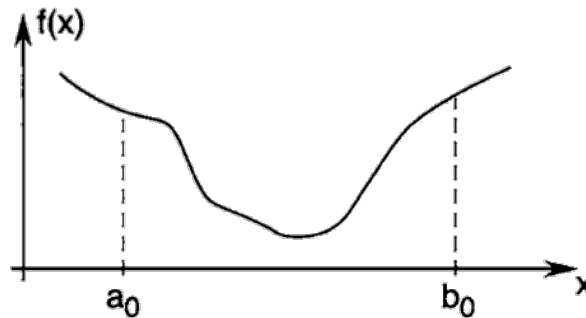
- Most optimization problems cannot be solved in a closed form (a single step).
- For them, we develop iterative algorithms:
  - Start from an initial candidate solution :  $x^{(0)}$
  - Generate a sequence of candidate solutions(iterates):  $x^{(1)}, x^{(2)}, \dots$
  - Stop when a certain condition is met; return the candidate solution
- In a large number of algorithms,  $x^{(k+1)}$  is generate from  $x^{(k)}$ , that is, using the information of  $f$  at  $x^{(k)}$ .
- In some algorithms,  $x^{(k+1)}$  is generate from  $x^{(k)}, x^{(k-1)}, \dots$ . But, for time and memory consideration, most history iterates are not kept in memory.



# Iterative Algorithm

## Golden Section Search

Given a closed interval  $[a_0, b_0]$ , a unimodal function (that is, having one and only one local minimizer in  $[a_0, b_0]$ ), and only objective value information, find a point that is no more than  $\epsilon$  away from that local minimizer.



- Mid-point : evaluate the mid-point, that is compute  $f\left(\frac{a_0+b_0}{2}\right)$ .
- But, it cannot determine which half contains the local minimizer. Does not work!!!
- Then? Two-points Approach!

# Golden Section Search

- ▶ (Step 0) Let  $a_0 = a$  and  $b_0 = b$ , Let  $k = 0$

- $b_k \triangleq$  upperlimit,  $a_k \triangleq$  lowerlimit

- ▶ (Step 1) If  $b_k - a_k < \epsilon$ , Stop.

Otherwise, go to Step 2.

- ▶ (Step 2) Let  $\hat{a}_k = a_k + \tau(b_k - a_k)$

$$\hat{b}_k = b_k - \tau(b_k - a_k)$$

where  $\tau = \frac{3-\sqrt{5}}{2} \approx 0.382$  (Golden Ratio)

1. When  $f(\hat{a}_k) < f(\hat{b}_k)$ ,

- (a) If  $f(a_k) > f(\hat{a}_k)$ , then let  $a_{k+1} = a_k$ ,  $b_{k+1} = \hat{b}_k$

- (b) If  $f(a_k) \leq f(\hat{a}_k)$ , then let  $a_{k+1} = a_k$ ,  $b_{k+1} = \hat{a}_k$

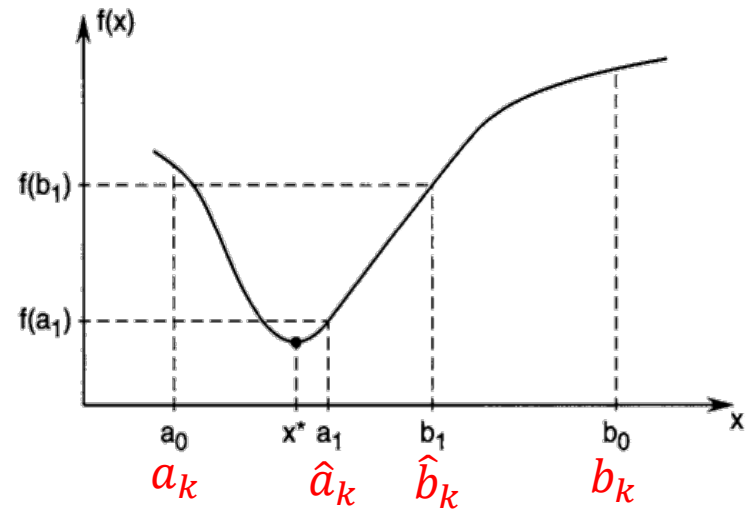
2. When  $f(\hat{a}_k) > f(\hat{b}_k)$ ,

- (a) If  $f(\hat{b}_k) < f(b_k)$ , then let  $a_{k+1} = \hat{a}_k$ ,  $b_{k+1} = b_k$

- (b) If  $f(\hat{b}_k) \geq f(b_k)$ , then let  $a_{k+1} = \hat{b}_k$ ,  $b_{k+1} = b_k$

3. When  $f(\hat{a}_k) = f(\hat{b}_k)$ , let  $a_{k+1} = \hat{a}_k$ ,  $b_{k+1} = \hat{b}_k$

- ▶ Let  $k = k + 1$ . Go to Step 1.



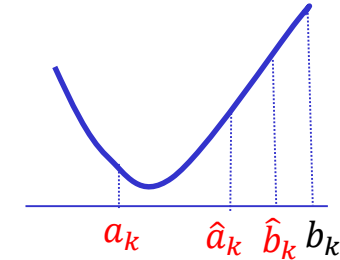
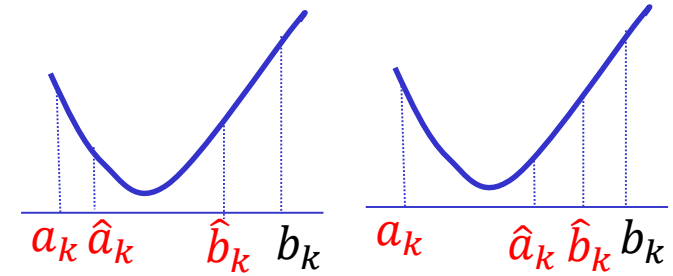
뒷장에 상세히

# Golden Section Search

1. When  $f(\hat{a}_k) < f(\hat{b}_k)$ ,

(a) If  $f(a_k) > f(\hat{a}_k)$ , then let  $a_{k+1} = a_k$ ,  $b_{k+1} = \hat{b}_k$

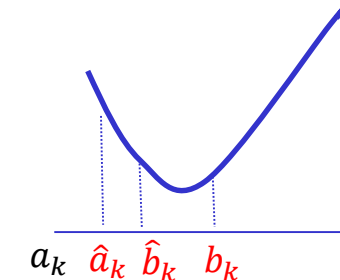
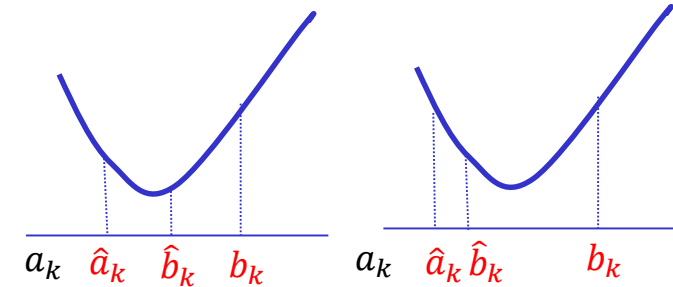
(b) If  $f(a_k) \leq f(\hat{a}_k)$ , then let  $a_{k+1} = a_k$ ,  $b_{k+1} = \hat{a}_k$



2. When  $f(\hat{a}_k) > f(\hat{b}_k)$ ,

(a) If  $f(\hat{b}_k) < f(b_k)$ , then let  $a_{k+1} = \hat{a}_k$ ,  $b_{k+1} = b_k$

(b) If  $f(\hat{b}_k) \geq f(b_k)$ , then let  $a_{k+1} = \hat{b}_k$ ,  $b_{k+1} = b_k$

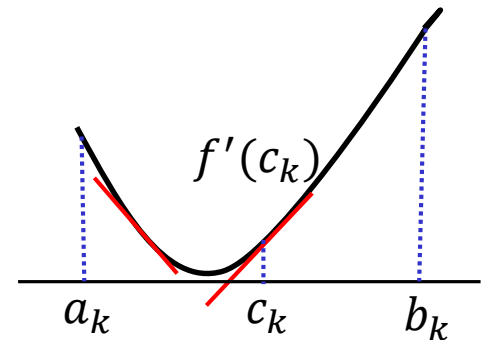


3. When  $f(\hat{a}_k) = f(\hat{b}_k)$ , let  $a_{k+1} = \hat{a}_k$ ,  $b_{k+1} = \hat{b}_k$

# Bisection Search

## Bisection Search

- Given a closed interval  $[a_0, b_0]$ , a continuously differentiable, a unimodal function, and derivative information, find a point that is no more than  $\epsilon$  away from that local minimizer.
  - Mid-point works with derivative!
- 
- ▶ (Step 0) Let  $a_0 = a$  and  $b_0 = b$ , Let  $k = 0$ 
    - $b_k \triangleq$  upperlimit,  $a_k \triangleq$  lowerlimit
  - ▶ (Step 1) If  $b_k - a_k < \epsilon$ , Stop.  
Otherwise, go to Step 2.
  - ▶ (Step 2) Let  $c_k = \frac{1}{2}(a_k + b_k)$ 
    1. If  $f'(c_k) = 0$ , then  $c_k$  is the local minimizer.
    2. If  $f'(c_k) > 0$ , then  $a_{k+1} = a_k$ ,  $b_{k+1} = c_k$
    3. If  $f'(c_k) < 0$ , then  $a_{k+1} = c_k$ ,  $b_{k+1} = b_k$
  - ▶ Let  $k = k + 1$ . Go to Step 1.



# Newton's Method

## Newton's Method

- Given a twice continuously differentiable function and objective, derivative, and 2<sup>nd</sup> derivative information, find an approximate minimizer.
- **Newton's method does not need intervals but must start sufficiently close to  $x^*$  !!!**
- Iteration: minimize the quadratic approximation

$$x_{k+1} \leftarrow \arg \min q(x) := f(x_k) + f'(x_k)(x - x_k) + \frac{1}{2}f''(x_k)(x - x_k)^2$$

*k*번째 iteration point인  $x_k$ 에서 2nd order approximation 최소화하는  $x$ 를 찾음  
(해당 함수 미분해서 0되는 지점 찾음)

Taylor series expansion

$$f(x) = \sum_{n=0}^{\infty} \frac{f^{(n)}(a)}{n!} (x - a)^n = f(a) + f'(a)(x - a) + \frac{1}{2}f''(a)(x - a)^2 + \dots$$

2<sup>nd</sup> order Taylor series expansion at  $x_k$

$$f(x) \approx f(x_k) + f'(x_k)(x - x_k) + \frac{1}{2}f''(x_k)(x - x_k)^2$$

# Newton's Method

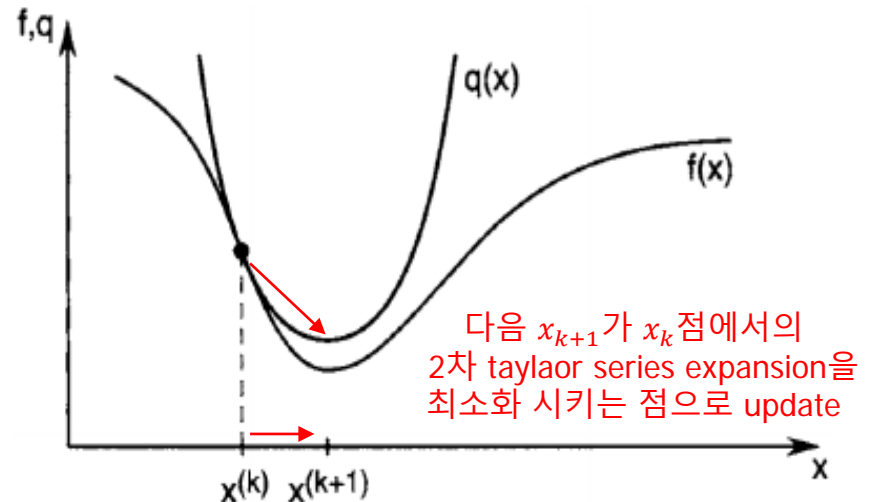
$$\begin{aligned} f(x) &\approx f(x_k) + f'(x_k)(x - x_k) + \frac{1}{2}f''(x_k)(x - x_k)^2 \\ &= f(x_k) + f'(x_k)\Delta x + \frac{1}{2}f''(x_k)(\Delta x)^2 \end{aligned}$$

- Derivative with respect to  $\Delta x$

$$0 = \frac{d}{d\Delta x} \left( f(x_k) + f'(x_k)\Delta x + \frac{1}{2}f''(x_k)(\Delta x)^2 \right) = f'(x_k) + f''(x_k)\Delta x$$

$$\Delta x = -\frac{f'(x_k)}{f''(x_k)}$$

$$x_{k+1} \leftarrow x_k + \Delta x = x_k - \frac{f'(x_k)}{f''(x_k)}$$



Second derivative → 계산량 많은 문제

Quadratic approximation with  $f''(x) > 0$

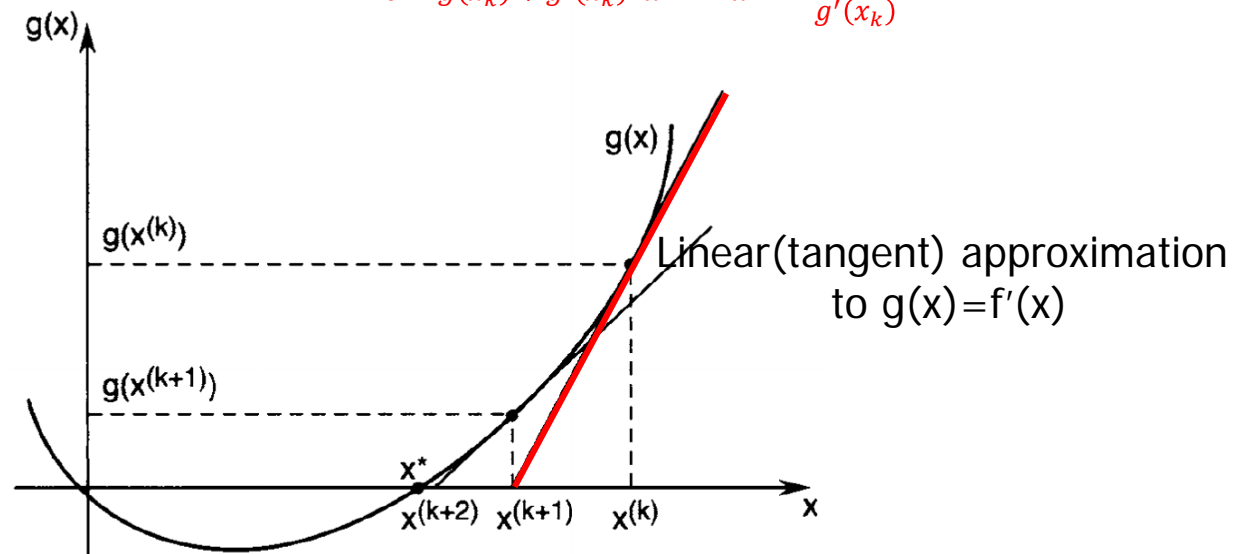
$$\arg \min q(x) := f(x_k) + f'(x_k)(x - x_k) + \frac{1}{2}f''(x_k)(x - x_k)^2$$

# Newton's Method

- Let us define  $g(x) = f'(x) \rightarrow g'(x) = f''(x)$
- Newton's method is a way to solve  $g(x) = f'(x) = 0$

$$x_{k+1} \leftarrow x_k - \frac{f'(x_k)}{f''(x_k)} = x_k - \frac{g(x_k)}{g'(x_k)}$$

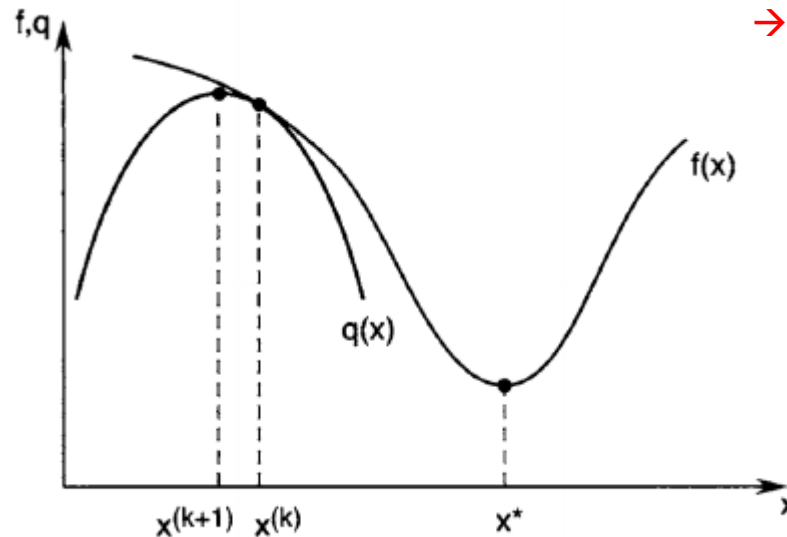
$$0 = g(x_k) + g'(x_k)\Delta x \rightarrow \Delta x = -\frac{g(x_k)}{g'(x_k)}$$



# Newton's Method

- ▶ What if  $f''(x_k) < 0$ ?

$q(x)$  (2차 approximation)가 위로 볼록 함수 → 미분해서 0 되는 지점을 찾는데 이 지점은 최소값 찾는 방향에서 더욱 멀어짐  
→ 처음에 초기값을 솔루션 근처로 잘 잡아야함



즉, Newton's iteration 은 발산함



# Secant Method

- Recall Newton's method uses  $f''(x_k)$  for minimizing  $f(x)$ :

$$x_{k+1} := x_k - \frac{f'(x_k)}{f''(x_k)}$$

- If  $f''$  is not available or is expensive to compute, we can approximate

$$f''(x_k) \approx \frac{f'(x_k) - f'(x_{k-1})}{x_k - x_{k-1}}$$

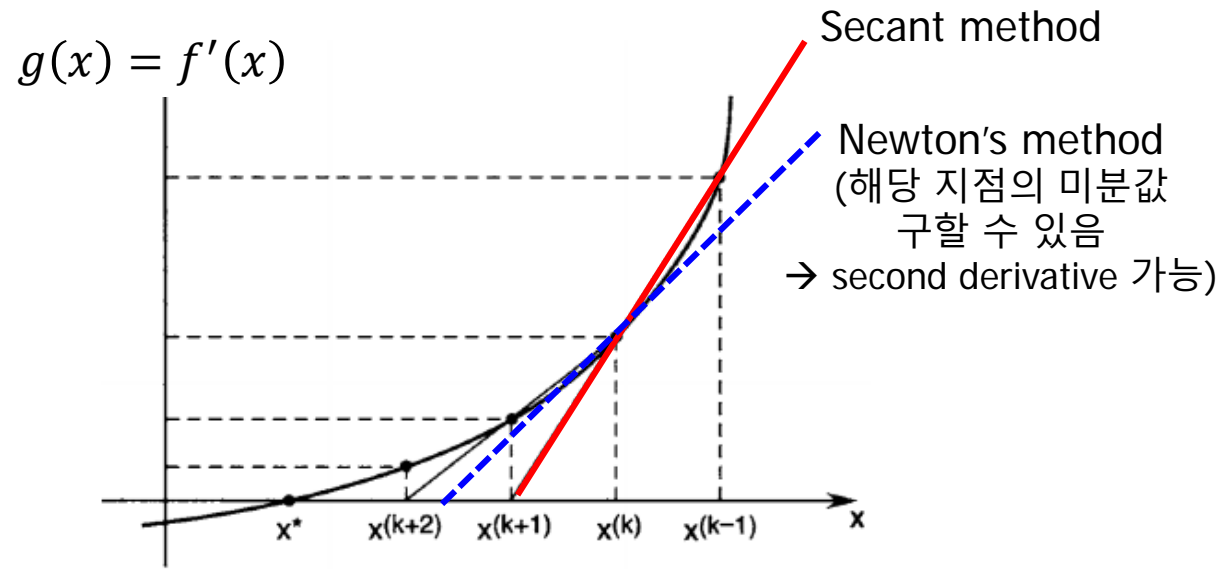
- This gives the iteration of the secant method:

$$x_{k+1} := x_k - \frac{(x_k - x_{k-1})f'(x_k)}{f'(x_k) - f'(x_{k-1})} = \frac{x_{k-1}f'(x_k) - x_kf'(x_{k-1})}{f'(x_k) - f'(x_{k-1})}$$

Note: the method needs two initial points. (미분 해야 하므로)

# Secant Method

$$f''(x_k) = \frac{f'(x_k) - f'(x_{k-1})}{x_k - x_{k-1}}$$



The secant method is slightly slower than  
Newton's method but is cheaper!

# Comparisons

- Comparisons of different 1D search methods
  - Golden section search (and Fibonacci search):
    - One objective evaluation at each iteration
    - Narrows search interval by less than half each time
  - Bisection search:
    - One derivative evaluation at each iteration
    - Narrows search interval to exactly half each time
  - Secant method:
    - Two points to start with; then one derivative evaluation at each iteration
    - Must start near  $x^*$ , has superlinear convergence
  - Newton's method:
    - One derivative and one second derivative evaluations at each iteration
    - Must start near  $x^*$ , has quadratic convergence, fastest among the four

# Comparisons

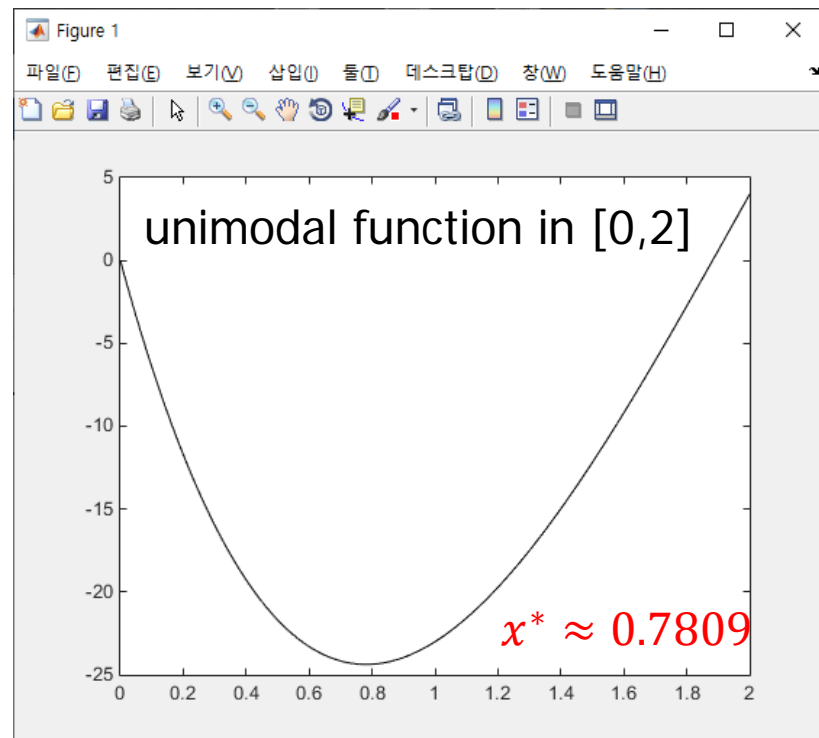
- 이외에도 많은 알고리즘
  - Gradient descent 알고리즘 등 (다음 주 강의)

# 실습

- MATLAB으로 Golden Section Search 알고리즘을 구현하여 해를 구하시오.

$$\min f(x) = x^4 - 14x^3 + 60x^2 - 70x$$

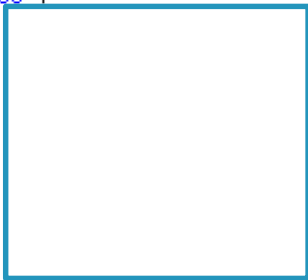


- Initial Interval:  $[0, 2]$
- Required Accuracy  $\varepsilon = 0.00001$
- 각 iteration마다  $a_k$  값을 확인하시오.



# 실습

```
1  %%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%
2  %%% Sample Code of Golden Section Method %%%
3  %%% Copyright) Prof. Seung-ho Chae %%%%%%%%%
4  %%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%
5  - clc; clear all; close all;
6
7  - epsilon = 0.00001;
8  - stepsize = (3 - sqrt(5))/2;
9  - a_k = 0;
10 - b_k = 2;
11 - dx = b_k - a_k;
12 - f = b_k^4 - 14*b_k^3 + 60*b_k^2 - 70*b_k;
13 - iteration = 0;
14
15 - fprintf('iter      a_k      b_k      dx      f(b_k)\n')
16 - fprintf('-----  -----  -----  -----  -----#n')
17 - fprintf('%3i %12.6f %12.6f %12.6f %12.6f\n', iteration, a_k, b_k, dx, f)
18
19 - while ( (b_k - a_k) > epsilon )
20
21 -     iteration = iteration + 1;
22
23 -     hat_a_k = a_k + stepsize*(b_k - a_k);
24 -     hat_b_k = b_k - stepsize*(b_k - a_k);
25
26 -     f_a_k = a_k^4 - 14*a_k^3 + 60*a_k^2 - 70*a_k; % objective function at a_k
27 -     f_b_k = b_k^4 - 14*b_k^3 + 60*b_k^2 - 70*b_k;
28
29 -     f_hat_a_k = hat_a_k^4 - 14*hat_a_k^3 + 60*hat_a_k^2 - 70*hat_a_k;
30 -     f_hat_b_k = hat_b_k^4 - 14*hat_b_k^3 + 60*hat_b_k^2 - 70*hat_b_k;
31
32 -     if (f_hat_a_k < f_hat_b_k)
33 -         check = 1;
34 -     elseif (f_hat_a_k > f_hat_b_k)
35 -         check = 2;
36 -     else
37 -         check = 3;
38 -     end
39
```

# 실습

```
40 - switch check
41 - case 1
42 - 
43 -
44 -
45 -
46 -
47 -
48 -
49 -
50 - case 2
51 - 
52 -
53 -
54 -
55 -
56 -
57 -
58 -
59 - case 3
60 - 
61 -
62 - end
63
64 - f = b_k^4 - 14*b_k^3 + 60*b_k^2 - 70*b_k;
65 - dx = b_k-a_k;
66
67 - fprintf('%3i %12.6f %12.6f %12.6f %12.6f\n', iteration, a_k, b_k, dx, f)
68
69 - end
```

a\_k, b\_k를 어떻게 update를  
해야할지 잘 고민해 보기

# 실습

iter	a_k	b_k	dx	f(b_k)
0	0.000000	2.000000	2.000000	4.000000
1	0.000000	1.236068	1.236068	-18.958161
2	0.472136	1.236068	0.763932	-18.958161
3	0.472136	0.944272	0.472136	-23.592462
4	0.652476	0.944272	0.291796	-23.592462
5	0.652476	0.832816	0.180340	-24.287887
6	0.721360	0.832816	0.111456	-24.287887
7	0.763932	0.832816	0.068884	-24.287887
8	0.763932	0.806504	0.042572	-24.349526
9	0.763932	0.790243	0.026311	-24.366907
10	0.773982	0.790243	0.016261	-24.366907
11	0.773982	0.784032	0.010050	-24.369296
12	0.777821	0.784032	0.006211	-24.369296
13	0.777821	0.781660	0.003839	-24.369583
14	0.779287	0.781660	0.002372	-24.369583
15	0.780193	0.781660	0.001466	-24.369583
16	0.780193	0.781099	0.000906	-24.369600
17	0.780539	0.781099	0.000560	-24.369600
18	0.780753	0.781099	0.000346	-24.369600
19	0.780753	0.780967	0.000214	-24.369601
20	0.780835	0.780967	0.000132	-24.369601
21	0.780835	0.780917	0.000082	-24.369602
22	0.780866	0.780917	0.000051	-24.369602
23	0.780866	0.780897	0.000031	-24.369602
24	0.780878	0.780897	0.000019	-24.369602
25	0.780878	0.780890	0.000012	-24.369602
26	0.780878	0.780886	0.000007	-24.369602

$$x^* \approx 0.7809$$

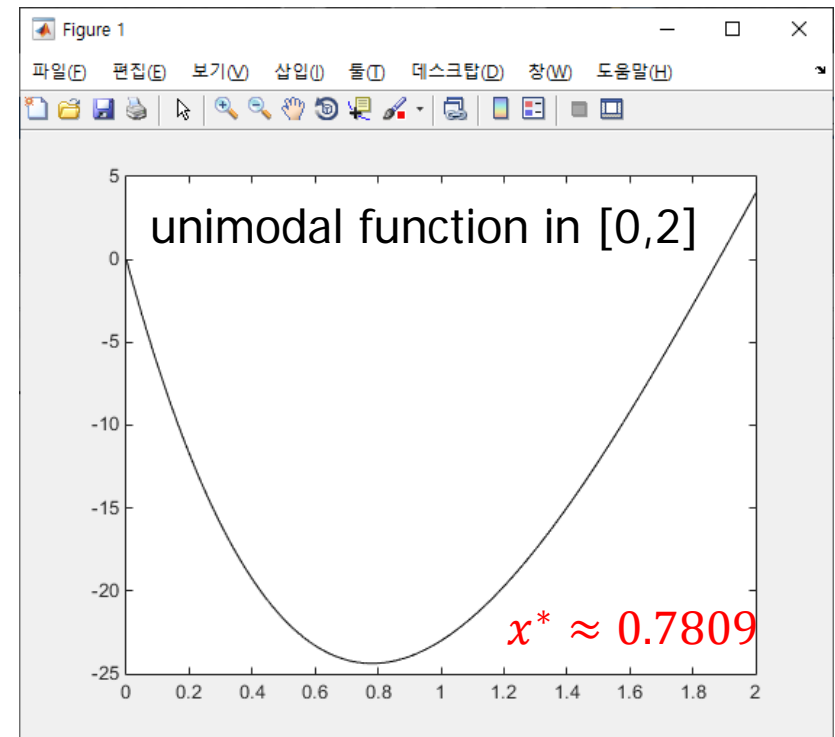


# Homework

- MATLAB으로 bisection Search 알고리즘을 구현하여 해를 구하시오.

$$\min f(x) = x^4 - 14x^3 + 60x^2 - 70x$$

- Initial Interval:  $[0,2]$
- Required Accuracy  $\varepsilon = 0.00001$
- 각 iteration마다  $c_k$  값을 확인하시오.



# Homework

iter	a_k	b_k	dx	f(b_k)
0	0.000000	2.000000	2.000000	4.000000
1	0.000000	1.000000	1.000000	-23.000000
2	0.500000	1.000000	0.500000	-21.687500
3	0.750000	1.000000	0.250000	-24.339844
4	0.750000	0.875000	0.125000	-24.105225
5	0.750000	0.812500	0.062500	-24.339096
6	0.750000	0.781250	0.031250	-24.369597
7	0.765625	0.781250	0.015625	-24.362377
8	0.773438	0.781250	0.007813	-24.367886
9	0.777344	0.781250	0.003906	-24.369214
10	0.779297	0.781250	0.001953	-24.369524
11	0.780273	0.781250	0.000977	-24.369590
12	0.780762	0.781250	0.000488	-24.369601
13	0.780762	0.781006	0.000244	-24.369601
14	0.780884	0.781006	0.000122	-24.369602
15	0.780884	0.780945	0.000061	-24.369601
16	0.780884	0.780914	0.000031	-24.369602
17	0.780884	0.780899	0.000015	-24.369602
18	0.780884	0.780891	0.000008	-24.369602

$$x^* \approx 0.7809$$

# Homework

- 다음 문제를 MATLAB으로 Newtons' method를 구현하여 해를 구하시오.

$$\min f(x) = x^4 - 14x^3 + 60x^2 - 70x \quad x^* \approx 0.7809$$

- ▶ Initial point:  $x_0 = 0$
  - ▶  $x_1 = 0.5833$
  - ▶  $x_2 = 0.7631$
  - ▶  $x_3 = 0.7807$
  - ▶  $x_4 = 0.7809$
- Produce highly accurate solutions in
- Just a few steps.

# Homework

iter	x	dx	f(x)
0	0.00000000	1.00000000	0.00000000
1	0.58333333	0.58333333	-23.07981289
2	0.76310273	0.17976939	-24.35978265
3	0.78071929	0.01761656	-24.36960073
4	0.78088404	0.00016475	-24.36960157
5	0.78088405	0.00000001	-24.36960157

$$x^* \approx 0.7809$$

Only 5 steps!