COL 100: 11 notes

· Computer Science fundamentally involves computing!

Note: Computing can be performed even 10/0 Computers. ____ mechanism to perform Computation This Explicitly o Computation Tollow A sequence of written down set of rules unambiguous ruiles _ is called an Eg: Addition, GCD [Mathematics] algorithm Bisecting an angle, [Geometry, Art...]

Chemical reaction & neation [chem.]
Homeostasis?... [Bio, Med. ...]

. Let us take some examples - Computing tools: abacus, sticks & stones, paper & sen,... · Straight Edge & Compass How is it a computing tool? Specify lines; W: Given a square can't measure Construct another square lengtes of twice the area of the Original square. Justi I cation Assume DABCD of side a 70 AC= VZa · Draw the diagonal AC

· Draw a square with AC

as an Edge recipe : Area of ACEF= 202

A. Bisect a given angle using straight-edge & Assume a point B is given as the vertex of the angle Vitte · Place compass on B, stretch it to any keyth that evill stay on the angle · Draw an arc; two intersection pts. are obtained · Place the compass on

A & C and draw ares
in the interior of the angle 3 c

· Connect the vertex D with B Justification B.t. < ABD = < CBD · Conquest As > "equal areas rinitations of Computing tools => equal angles Eg: Nith

Straight edge + compass

we cannot trisect an Not a primitive

Step

that 2xample Revist Jeu square with Ac Something can't - Draw a square as an edge down farther

Observations

1. Every computing tool has a "model" of computation

Q: weat is - ??

3. Methods of abstraction -- how compound Exprs can be named & manipulated as

· Perpendicular at a pt on a line was used to construct a square a vog. construction alg to construct

A: Model of computation gives essa framework to organize our idean about produses i.e: 3 mechanisms 1 Princitive operation and expressions -> simplest entities of the Eg: Nat, + 7. Method of combination of building complex IL nicher Expressions

Use of abstraction -- Deparates logical subproblems - avoiding repetitions of copies of smilar solutions 2. Based on teuse princitives, the computational model's power or Expressivity is defined Turing's Result [With just 6 primitive Expressions]

Can compute any "computably enumerable"

problem => As a result power of (c) = power (lython) = power (sml)....

But based on particular combinations & abstraction thuis Duitability for a task may differ for instance: MATLAB: Excellent for vectors, matrices C: for data n/ws LISP/SML: for arbitrary structured data

Our Computing Tool . Definite i/P and O/P · Unambiguous Control
Storage Logic · Specifier solution as a finite process => # of steps in ten computation is finite inst. Programmiz Lanjuage_ Language via which communication

with m/c is possible

Program: Algorithm written in a proje language

[well defined grammars in the communication] Computing models mathematical Expressions · functional (declarative): of commands · Josepherative: sequence · Take a gens g · Compute g² Compute y = Vx · check if g2 NX s.t. y2 = x · If yes - return of Mathematical Dej". Imperative = (g + g/x) reprat.

Mu model of computation (Functional model)
Our model of computation (Functional model) Le Closenes to mathematics
Principles Expressions La Chosenes to makenatics makes it convenient to me
Primitive Expressions - Constants, Variables and functions Eg: N, Z, R x, y, z MXN -> N Therifiers N x P -> N to data objects
- Constants, Variables and
Eg: N, Z, R x, y, z MXN -> N
TB Identifiers NXP - N to data
to data
objects
- Basic Operations, addition, subtraction, division, multiplication, boolean operations, etc.
booleon oberations etc.
multiplication, Court of
. On rowing Entites
- Marring mechanism for various Entities 1 12/2 Repeating definitions
to be used ω/o repeating definitions
Eg: Square (n) = n*n

Methods of Combinations
 Composition of Junctions
 Eg: Sum-sq(xy) = Square(x) + square(y)
 Industrie definitions
 Eg: factorial

Methods of abstraction

- Naming & using groups of objects

and Expressions as a single unit

Eg: Factorial ij n<1 $n! = \begin{cases} 1 \\ 1 \times 2 \times \dots \times n \end{cases}$ Math notation if n<1 Inductive Variant $ui = \begin{cases} u \times (u-T)j \\ T \end{cases}$ otherwise n* fact (n-1); fun fact n= if n<1 then 1 Else val fact = fn: int → int - fact 4; Evaluation has talken Val it = 24 : int 1. vars to values
for names with definitions
args with actual arg. operators tell offerwise 2. Order left bright 3. Priority 01 operators

Computation

$$= 3 \times (2 \times (2-1))$$

$$= 3 \times (2 \times 1)$$

$$= 3 \times (2 \times 1) = 3 \times 2 = 6$$

Q. How about $n_0 = \begin{cases} 1 & \text{if } n < 1 \\ (n+1)! & \text{otherwise} \end{cases}$ but computationally incorrect Mathematically Correct Observation 31 = (3+1)[/(3+1) guarantee that all = 414 1 its computation = ((4+1)!/4+L)/4 terminati = (51/5)/4 · Not all mathematical functions are algorithms

Another Eg:

Sqvt(x) = { y if y * y = x }

Sqvt(x) = { 0 if # y: y * y = x}

No Description
of how to Evaluate" Another Eg: { y if y = y = x} ->
Sqv+(x) = { 0 if #y: y = y = x} SML Jano - 5/j ____ terminate the i/p wal it = 5: int - [val] x= 5; Les Keyword: defining a new name