

Student

# Ministry of Science and Higher Education of the Russian Federation

## National Research University Higher School of Economics

Faculty of Computer Science

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#### HOMEWORK REPORT

Practical homework Nº2

Subject: Ordered Sets for Data Analysis

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**Task.** Given the following formal context, find all formal concepts, draw the concept lattice and find all non-trivial implications.

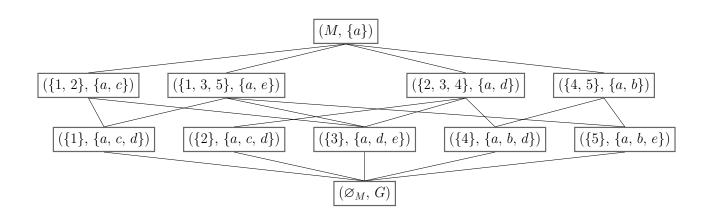
	a	b	С	d	e
1	1		1		1
2	1		1	1	
3	1			1	1
4	1	1		1	
5	1	1			1

**Solution.** Let the set of all objects be denoted as  $G = \{1, 2, 3, 4, 5\}$  and the set of all attributes as  $M = \{a, b, c, d, e\}$ . We will condense all the data about formal concepts in this context into a table, where  $A \subseteq M$  is *extent* of a formal concept and  $B \subseteq G$  is its *intent*.

A	B = A'	A"	Formal concept?
$\{\varnothing_M\}$	G	$\{\varnothing_M\}$	Yes
M	<i>{a}</i>	M	Yes
{1}	$\{a,c,e\}$	{1}	Yes
{2}	$\{a,c,d\}$	{2}	Yes
{3}	$\{a, d, e\}$	{3}	Yes
{4}	$\{a,b,d\}$	{4}	Yes
{5}	$\{a,b,e\}$	{5}	Yes
{1, 2}	$\{a, c\}$	{1, 2}	Yes
{1, 3}	$\{a\}$	M	No
{1, 4}	<i>{a}</i>	M	No
{1, 5}	$\{a, e\}$	{1, 3, 5}	No
{2, 3}	$\{a, d\}$	{2, 3, 4}	No
$\{2, 4\}$	$\{a, d\}$	$\{2, 3, 4\}$	No

A	B = A'	A''	Formal concept?
{2, 5}	<i>{a}</i>	M	No
{3, 4}	$\{a, d\}$	{2, 3, 4}	No
{3, 5}	<i>{a}</i>	M	No
$\{4, 5\}$	$\{a, b\}$	$\{4,  5\}$	Yes
{1, 2, 3}	$\{a\}$	M	No
{1, 2, 4}	$\{a\}$	M	No
{1, 2, 5}	$\{a\}$	M	No
{1, 3, 4}	<i>{a}</i>	M	No
$\{1, 3, 5\}$	$\{a, e\}$	$\{1, 3, 5\}$	Yes
$\{1, 4, 5\}$	$\{a\}$	M	No
$\{2, 3, 4\}$	$\{a, d\}$	$\{2, 3, 4\}$	Yes
$\{2, 3, 5\}$	$\{a\}$	M	No
${3, 4, 5}$	$\{a\}$	M	No
{1, 2, 3, 4}	<i>{a}</i>	M	No
{1, 2, 3, 5}	<i>{a}</i>	M	No
{1, 2, 4, 5}	<i>{a}</i>	M	No
{1, 3, 4, 5}	<i>{a}</i>	M	No
{2, 3, 4, 5}	<i>{a}</i>	M	No

Having found all the formal concepts, we are able to construct the concept lattice.



**Definition 1.**  $A \to B$ , where  $A, B \subseteq M$  holds in context (G, M, I) if  $A' \subseteq B'$ , i.e., each object having all attributes from A also has all attributes from B.

Therefore, we can construct a table of all the subsets of M, which upon application of prime operator return nonempty sets to look for all non-trivial implications.

A	A'	$A'' \setminus A$	A	A'	$A'' \setminus A$
<i>{a}</i>	G	$\varnothing_M$	$\{a, b, d\}$	{4}	$\varnothing_M$
{b}	$\{4, 5\}$	<i>{a}</i>	$\{a, b, e\}$	{5}	$\varnothing_M$
$\{c\}$	{1, 2}	<i>{a}</i>	$\{a, c, d\}$	{2}	$\varnothing_M$
$\{d\}$	$\{2, 3, 4\}$	<i>{a}</i>	$\{a,c,e\}$	{1}	$\varnothing_M$
$\{e\}$	$\{1, 3, 5\}$	{a}	$\{a, d, e\}$	{3}	$\varnothing_M$
$\{a, b\}$	$\{4, 5\}$	$\varnothing_M$	$\{b, c, d\}$	$\varnothing_G$	$\{a, e\}$
$\{a, c\}$	{1, 2}	$\varnothing_M$	$\{b,c,e\}$	$\varnothing_G$	$\{a, d\}$
$\{a, d\}$	$\{2, 3, 4\}$	$\varnothing_M$	$\{b, d, e\}$	$\varnothing_G$	$\{a, c\}$
$\{a, e\}$	$\{1, 3, 5\}$	$\varnothing_M$	$\{c, d, e\}$	$\varnothing_G$	$\{a, b\}$
$\{b, c\}$	$\varnothing_G$	$\{a, d, e\}$	$\left\{a,b,c,d\right\}$	$\varnothing_G$	$\{e\}$
$\{b, d\}$	{4}	<i>{a}</i>	$\left\{a,b,c,e\right\}$	$\varnothing_G$	$\{d\}$
$\{b,e\}$	{5}	<i>{a}</i>	$\boxed{\{a,b,d,e\}}$	$\varnothing_G$	$\{c\}$
$\{c, d\}$	{2}	<i>{a}</i>	$\boxed{\{a,c,d,e\}}$	$\varnothing_G$	{b}
$\{c, e\}$	{1}	<i>{a}</i>	$\{b,c,d,e\}$	$\varnothing_G$	<i>{a}</i>
$\{d, e\}$	{3}	<i>{a}</i>	$\varnothing_M$	G	Trivial
$\{a,b,c\}$	$\varnothing_G$	$\{d, e\}$			

According to the table there are following non-trivial implications:  $\{b\} \to \{a\},\$   $\{c\} \to \{a\},\ \{d\} \to \{a\},\ \{e\} \to \{a\},\ \{b,c\} \to \{a,d,e\},\ \{a,b,c\} \to \{d,e\},\$   $\{b,d\} \to \{a\},\ \{b,e\} \to \{a\},\ \{c,e\} \to \{a\},\ \{c,d\} \to \{a\},\ \{d,e\} \to \{a\},\$   $\{b,c,d\} \to \{a,e\},\ \{b,c,e\} \to \{a,d\},\ \{b,d,e\} \to \{a,c\},\ \{c,d,e\} \to \{a,b\},\$   $\{a,b,c,d\} \to \{e\},\ \{a,b,c,e\} \to \{d\},\ \{a,b,d,e\} \to \{c\},\ \{a,c,d,e\} \to \{b\},\$   $\{b,c,d,e\} \to \{a\}.$ 

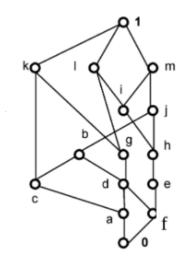
Task. Using the diagram find:

1.  $f \wedge m$ ,  $a \vee l$ ,  $i \wedge k$ ;

2.  $a \land (c \lor d), (i \land g) \lor (c \land d), \bigvee (b, c, d);$ 

3.  $\bigwedge \varnothing$ ,  $\bigvee \varnothing$ ;

and determine whether the diagram represents an upper semi-lattice, a lower semi-lattice or a lattice.



#### Solution.

1. Find  $f \wedge m$ ,  $a \vee l$ ,  $i \wedge k$ .

1.1 
$$f \wedge m = \inf\{f, m\} = f;$$

1.2 
$$a \lor l = \sup\{a, l\} = l;$$

1.3 
$$i \wedge k = f$$
.

2. Find  $a \wedge (c \vee d)$ ,  $(i \wedge g) \vee (c \wedge d)$ ,  $\bigvee (b, c, d)$ .

2.1  $a \wedge (c \vee d)$  does not exist, since there is no lowest element in set  $\{c, d\}^U = \{k, b, j, m, 1\};$ 

$$2.2 \ (i \wedge g) \vee (c \wedge d) = f \vee a = d;$$

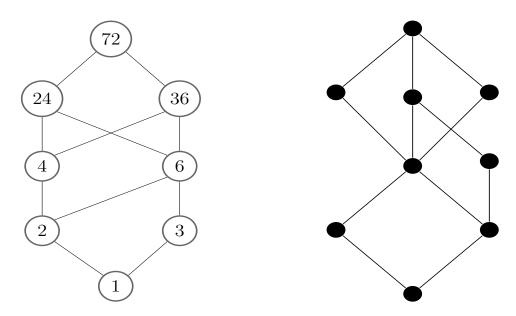
2.3 
$$\bigvee (b, c, d) = b$$
.

- 3. Find  $\bigwedge \emptyset$ ,  $\bigvee \emptyset$ .
  - 3.1 by definition,  $\bigwedge \emptyset = 1$ ;
  - 3.2 by definition,  $\bigvee \emptyset = 0$ .

Let us consider infima and suprema of all pairs of incomparable elements.

- $a \wedge f = 0$ ,  $c \wedge d = a$ ,  $c \wedge e = 0$ ,  $d \wedge e = f$ ,  $b \wedge g = d$ ,  $b \wedge h = f$ ,  $g \wedge h = f$ ,  $i \wedge j = h$ ,  $k \wedge l = g$ ,  $k \wedge m = c$ ,  $l \wedge m = i$ , therefore the diagram represents a lower-semilattice;
- $a \lor f = d$ ,  $c \lor e = j$ ,  $d \lor e = j$ ,  $b \lor g = 1$ ,  $b \lor h = j$ ,  $g \lor h = l$ ,  $i \lor j = m$ ,  $k \lor l = 1$ ,  $k \lor m = 1$ ,  $l \lor m = 1$ , however  $c \lor d$  does not exist, therefore the diagram does not represent an upper-semilattice.

**Task.** Determine if each of the following two lattice diagrams is: distributive, modular. Answer in detail.



**Definition 2.** A lattice elements of which satisfy equations

$$x \wedge (y \vee z) = (x \wedge y) \vee (x \wedge z),$$
  
$$x \vee (y \wedge z) = (x \vee y) \wedge (x \vee z)$$

is called distributive.

**Theorem 1.** A lattice is distributive iff it contains (a sublattice which is) neither a pentagon, nor a diamond.

**Definition 3.** A lattice such that if  $x \leq z$ , then  $x \vee (y \wedge z) = (x \vee y) \wedge z$  is called *modular*.

**Theorem 2.** A lattice is modular iff it does not contain (a sublattice which is) a pentagon.

Solution. Let us consider the first diagram. It represents poset

$$(\{1, 2, 3, 4, 5, 6, 24, 36, 72\}, |).$$

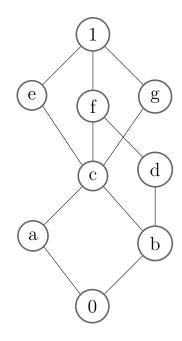
First, we have to find out whether this diagram indeed represents a lattice by finding suprema and infima of all pairs of incomparable elements.

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- $24 \wedge 36$  does not exist,  $4 \wedge 6 = 2$ ,  $2 \wedge 3 = 1$ ;
- $24 \lor 36 = 72$ ,  $4 \lor 6$  does not exist,  $2 \lor 3 = 6$ .

Since there is a pair of elements which does not have infimum and a pair which does not have supremum, this object cannot be considered lattice, therefore distributivity and modularity do not apply to the object represented by the first diagram.

Let us consider the second diagram. For the sake of clarity, we will assign a one-letter name for each node while naming top and bottom nodes 0 and 1 respectively. Resulting diagram:



Again, we have to check whether this diagram represents a lattice.

- Infima.  $a \wedge b = 0$ ,  $c \wedge d = b$ ,  $e \wedge f = c$ ,  $e \wedge g = c$ ,  $f \wedge g = c$ .
- Suprema.  $a \lor b = c, c \lor d = f, e \lor f = 1, e \lor g = 1, f \lor g = 1.$

Since each pair of incomparable elements has infinum and supremum, the diagram represents a lattice.

This diagram contains a sublattice  $\{c, e, f, g, 1\}$  which is a diamond. Therefore, according to Theorem 1, it cannot be distributive. However, this diagram does not contain any pentagons, hence, according to Theorem 2, it is modular.

Task. Show that distributivity implies absorption.

**Solution.** Suppose that operators  $\vee$  and  $\wedge$ , defined on a bounded lattice, are distributive. Since the lattice is bounded,  $\vee$  and  $\wedge$  have identity elements  $1_{\vee}$  and  $1_{\wedge}$  respectively. Then, for any x, y we have

$$x \wedge (x \vee y) = (x \wedge x) \vee (x \wedge y) = x \vee (x \wedge y).$$

Also from distributivity follows that

$$x \lor (x \land y) = (x \land 1_{\land}) \lor (x \land y) = x \land (1_{\land} \lor b).$$

Since identity element of  $\wedge$  is a zero for  $\vee$  and vice versa, we get

$$x \wedge (1_{\wedge} \vee b) = x \wedge 1_{\wedge} = x,$$

hence

$$x = x \land (x \lor y) = x \lor (x \land y)$$
 QED.

**Task.** Prove the following two rules using only Armstrong rules without using the properties of the operation  $(\cdot)'$ .

1. 
$$\frac{X \to Y \cup Z}{X \to Y; X \to Z},$$

$$2. \ \frac{X \to Y \setminus X}{X \to Y}.$$

#### Solution.

1. Using second Armstrong rule twice we get at first  $\frac{Y \to Y}{Y \cup Z \to Y}$  and then  $\frac{Y \cup Z \to Y}{(Y \cup Z) \cup X \to Y}$ . After that, using third Armstrong rule we get

$$\frac{X \to Y \cup Z, (Y \cup Z) \cup X \to Y}{X \to Y},$$

dually, after applying second Armstrong rule twice and third Armstrong rule once we get  $\frac{X \to Y \cup Z, \ (Z \cup Y) \cup X \to Z}{X \to Z}$ ; QED.

2.  $\frac{X \to Y \setminus X}{X \to Y}$ . Let  $Y = X \cup X_1$ . Using first Armstrong rule, we get  $\frac{X \to Y \setminus X}{X_1 \cup X \to X \cup X_1}$ , then using third rule we get

$$\frac{X \to X_1, X_1 \cup X \to X \cup X_1}{X \to X \cup X_1}; \text{QED}.$$