

ASSIGNMENT

Course Code CSC310A

Course Name COMPILERS

Programme B.TECH

Department CSE

Faculty FET

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Semester/Year 6SEM/2017

Course Leader/s MS SUVIDHA

Declaration Sheet					
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Programme	B.TECH			Semester/Year	6
Course Code	CSC310A				
Course Title	COMPILERS				
Course Date		to			
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Declaration

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Signature of the Student			Date	
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Signature of the Cours	e Leader and date	Signature of the F	Reviewe	er and date

Faculty of Engineering and Technology				
Ramaiah University of Applied Sciences				
Department	Computer Science and Engineering	Programme	B. Tech in Computer Science and Engineering	
Semester/Batch	06 th /2017			
Course Code	CSC310A	Course Title	Compilers	
Course Leader	Mr. Hari Krishna S. M. & Ms. Suvidha			

	Assignment					
Regi	ster No.	17ETCS002213	Name of the student	Sudhanshu shekhar		
Se				Marks		
cti on		Marking S	M ax M ark s	First Examiner Marks	Moderator	
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rt ^	A 1.1	Identification and grouping of Tokens		04		
	A 1.2	Implementation in <i>Lex</i>		02		
	A 1.3	Design of Context Free Grammar		04		
	A 1.4	Implementation in <i>Yacc</i>		06		
	A 1.5	Results and Comments		04		
		Part-A 1 Max Marks		20		
	Total Assignment Marks			20		

Course Marks Tabulation					
Component- CET B Assignment	First Examiner	Remarks	Second Examiner	Remarks	
A.1					
Marks (out of 20)					

PART - A

A 1.1 Identification and grouping of Tokens

Tokens are sequences of characters that have a certain meaning or represent a particular group. Lexemes are a sequence of characters in the source code that is matched by the pattern for a token. In the given problem, the set of tokens are:

Reserved words:

```
Enter the exp: int main () {
  int a;
  int b;
  scanf("%d",&b);
  int c[100];
  int s=0;
  if(a>b)
  s=a+b;
  else
  s=a-b;
  int i;
  for(i=0;i<3;i++)
  s=s+i;
  printf("s = ",s);
  }
  Input accepted.</pre>
```

LEXEMES	TOKENS
Int	Data Type
main	Reserved
	Keywords
(Separators
)	Separators
{	Separators
Int	Data Type
a	Identifier
;	Separators
Int	Data Type
b	Identifier
;	Separators
scanf	I/O function
(Separators
"%d"	Statement
,	Separators
&b	Statement
)	Separators
;	Separators
Int	Data Type
С	Identifier
[Separators
100	Operator
]	Separators
;	Separators
Int	Data Type
S	Identifier
=	Operator
0	Identifier
;	Separators

	1.,
If .	Keyword
(Separators
a	Identifier
>	Operator
b	Identifier
)	Separators
S	Identifier
а	Identifier
+	Operator
b	Identifier
;	Separators
Else	Keyword
S	Identifier
=	Operator
а	Identifier
-	Operator
b	Identifier
;	Separators
Int	Data Type
I	Identifier
;	Separators
For	Keyword
(Separators
I	Identifier
=	Operator
0	Identifier
;	Separators
I	Identifier
<	Operator
3	Identifier
;	Separators
S	Identifier
=	Operator
S	Identifier
+	Operator
i	Identifier
;	Separators
Printf	I/O function
(Separators
S	Identifier

=	Operator
,	Separators
S	Identifier
)	Separators
;	Separators
}	Separators

A1.2 Implementation in Lex

<u>Lex</u> is a computer program that generates lexical analyzers. Lex reads an input stream specifying the lexical analyzer and outputs source code implementing the lexer in the C programming language. It returns tokes which act as input to the yacc program which is responsible for the parsing

The code for the lex is as follows:

```
"%"
                         {return percent;}
"&"
                         {return and;}
"void"
                         {return VOID;}
\"[^\"\n]*\"
                                  {return DQ;}
\"[^\'\n]*\"
                                  {return SQ;}
"printf"
                         {return PRINT;}
"scanf"
                         {return INPUT;}
"+"
                         {return PLUS;}
"_"
                         {return MINUS;}
m \ge m
                         {return STAR;}
                         {return DIVIDE;}
"int"
                         { return INT;}
"char"
                         { return CHAR;}
                         { return semicolon;}
                         { return EQ;}
                         { return space;}
{digit}+
                                  {return num;}
{alpha}({alpha}|{digit})*
                                 {return var id;}
[!@#$%^&*<>,?a-zA-Z0-9] {return character;}
[\n]
                         { return NL;}
"if"
                         {return IF;}
"else"
                         {return ELSE;}
"for"
                         {return FOR;}
                         {return WHILE;}
"while"
"++"
                         {return INC;}
                         {return DNC;}
">="
                         {return GE;}
"<="
                         {return LE;}
"=="
                         {return EEQ;}
"&&"
                         {return AND;}
                         {return GT;}
">"
                         {return LT;}
                         {return OR;}
"!"
                         {return NOT;}
"!="
                         {return NE;}
"break"
                         {return BREAK;}
"["
                         {return SRB;}
"]"
                         {return SLB;}
[!@#$%^&*<>,?a-zA-Z0-9]* {return print_char;}
"return(0);"
                         {return RETURN;}
                         {ECHO; yyerror("unexpected character");}
%%
```

Fig 1.1: The above image shows the screenshot of a lex file with the extension -. I

A1.3 Design of Context Free Grammar for each function

The context free gramma for each function:

A parser generator is a program that takes as input a specification of a syntax, and produces as output a procedure for recognizing that language. Historically, they are also called compiler-compilers.

Minimum two data types:

The data types that are chosen is int and char. The grammar should be able declare variables of the abovementioned data types, it has to allow variable assignment and allow expressions involving the variables. declarations: declarations dec

```
I dec
dec:
      INT space var id semicolon NL
       | CHAR space var id semicolon NL
       INT space var id space EQ space num semicolon NL
       | CHAR space var id space SQ character SQ semicolon NL
```

Minimum two control statements:

```
The control statements that are accepted in the program are if-else, if and break statement.
control stmt: IF RB condition LB NL RCB NL exp NL LCB NL
```

```
| BREAK semicolon NL
condition: var id LE var id
         | var_id GE var_id
         | var id EEQ var id
         | var id AND var id
         | var id NE var id
         | var id OR var id
         | NOT var id
```

```
exp: var id space EQ space var id space PLUS space var id semicolon
       |var_id space EQ space var_id space MINUS space var_id semicolon
       Ivar id space EQ space var id space STAR space var id semicolon
       var id space EQ space var id space DIVIDE space var id semicolon
```

Minimum two looping statements:

| var id DNC semicolon NL

```
The loops used here are for loop and while loop.
loop stmt: FOR RB var id EQ num semicolon cond semicolon indc LB NL RCB NL exp NL LCB NL print stmt
         | WHILE RB cond LB RCB NL exp NL indc NL LCB NL print stmt
cond: var id LE var id
       | var id GE var id
       | var_id NE var_id
indc: var id INC
   | var id INC semicolon NL
   | var id DNC
```

```
exp: var_id space EQ space var_id space PLUS space var_id semicolon

|var_id space EQ space var_id space MINUS space var_id semicolon

|var_id space EQ space var_id space STAR space var_id semicolon

|var_id space EQ space var_id space DIVIDE space var_id semicolon

;
```

Input-output functions:

```
The standard input and output statements like printf and scanf are being implemented. input_stmt: INPUT RB DQ percent var_id DQ comma and var_id LB semicolon NL; print_stmt: PRINT RB DQ print_char DQ semicolon LB NL | PRINT RB DQ print_char DQ comma var_id LB semicolon NL:
```

Compound statements and two dimentional arrays:

Any statement inside curly braces are called compound statements. Here all the above statements are compound statements.

One dimentional array_declaration:

```
dec: INT space var_id SRB num SLB semicolon NL | CHAR space var_id SRB num SLB semicolon NL :
```

Two dimentional array declaration:

A1.4 Implementation in YACC

The implementation of yacc is as follows:

```
%{
void yyerror (char *s);
#include <stdio.h>
#include <stdlib.h>
%}
%start S
%token NL INT CHAR character comma RCB LCB space RB LB MAIN VOID semicolon var_id EQ PLUS MINUS STAR DIVIDE RETURN GT LT
%token SQ PRINT INPUT and percent DQ num IF ELSE SRB SLB INC DNC LE EEQ AND OR NOT RE NE GE FOR WHILE BREAK print_char
%right EQ
%left AND OR
%left LE GE EQ NE PLUS MINUS STAR DIVIDE GT LT
%%
S:
        prototype
prototype: VOID space MAIN RB LB NL RCB NL body RETURN NL LCB
           INT space MAIN RB LB NL RCB NL body RETURN NL LCB
body:
        declarations statements
declarations:
               declarations dec
                dec
dec:
        INT space var_id semicolon NL
        CHAR space var_id semicolon NL
        |INT space var_id space EQ space num semicolon NL
        CHAR space var_id space SQ character SQ semicolon NL
        INT space var_id SRB num SLB semicolon NL
        CHAR space var_id SRB num SLB semicolon NL
        | INT space var id SRB num SLB SRB num SLB semicolon NL
        CHAR space var id SRB num SLB SRB num SLB semicolon NL
```

Fig 1.1: Source code

```
statements: statements input stmt control stmt NL
input_stmt: INPUT RB DQ percent var_id DQ comma and var_id LB semicolon NL
control_stmt: IF RB condition LB NL RCB NL exp NL LCB NL loop_stmt
                | IF RB condition space logical op space condition LB NL RCB NL exp NL LCB NL loop stmt
                BREAK semicolon NL loop_stmt
condition: var id LE var id
           | var id GE var id
           | var_id EEQ var_id
           | var_id NE var_id
logical_op:
                AND
                OR
                NOT
exp: var_id space EQ space var_id space PLUS space var_id semicolon
        var_id space EQ space var_id space MINUS space var_id semicolon
        |var_id space EQ space var_id space STAR space var_id semicolon
        |var_id space EQ space var_id space DIVIDE space var_id semicolon
loop_stmt: FOR RB var_id EQ num semicolon cond semicolon indc LB NL RCB NL exp NL LCB NL print_stmt
           | WHILE RB cond LB RCB NL exp NL indc NL LCB NL print_stmt
cond: var_id LE var_id
        var_id GE var_id
        | var_id NE var_id
```

Fig 1.2: Source code

```
indc: var_id INC
     | var_id INC semicolon NL
     | var_id DNC
     | var_id DNC semicolon NL
print stmt: PRINT RB DQ print char DQ semicolon LB NL
             PRINT RB DQ print_char DQ comma var_id LB semicolon NL
             ;
%%
int main(){
printf("Enter the exp: ");
if (yyparse() == 0)
printf("accepted");
else{
printf("NOT accepted");
return(0);
}
}
int yywrap(){
return 1;
}
void yyerror (char *s) {fprintf (stderr, "%s\n", s);}
```

Fig 1.3: Source code

A1.5 Results and comments

Execution of the above program is as follows:

```
Enter the exp: int main () {
int a;
int b;
scanf("%d",&b);
int c[100];
int s=0;
if(a>b)
s=a+b;
else
s=a-b;
int i;
for(i=0;i<3;i++)
s=s+i;
printf("s = ",s);
}
Input accepted.</pre>
```

Fig : Results, the code typed in the above figure does not have any syntax error or it is according to the grammar defined, so the input is accepted

```
Enter the exp: void main ( ) {
int a;
int b[j];
}
NOT Accepted.
```

Fig : Results, the code typed in the above figure does have a syntax error in the third line, so the input is not accepted

Comments:

Yacc reads the grammar descriptions in yacc file and generates a syntax analyzer (parser), that includes function yyparse, in file y.tab.c. Included in file lab8.y are token declarations. The –d option causes yacc to generate definitions for tokens and place them in file y.tab.h. Lex reads the pattern descriptions in the lex file, includes file y.tab.h, and generates a lexical analyzer, that includes function yylex, in file lex.yy.c. Finally, the lexer and parser are compiled and linked together to create executable a.exe. From main we call yyparse to run the compiler. Function yyparse automatically calls yylex to obtain each token.

Limitations:

The above program works only for simple statements and does not work for nested statements. Nested if statements and nested looping statements are not supported by the program. The program should be in a particular order as per the grammar