IS-ZC444: ARTIFICIAL INTELLIGENCE

Lecture-07: Beyond Classical Search



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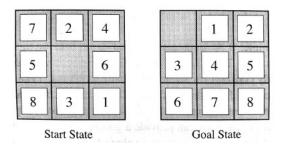
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Heuristic Function

Consider 8-puzzle problem



- Average solution cost is 22 steps
- Branching factor is \sim 3, so $3^{22} = 3.1 \times 10^{10}$ states in tree
- However, graph search would have 9!/2 = 181,440 states¹
- To find shortest solution using A*, we need to find appropriate heuristic functions

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Heuristic Function for 8-puzzle

h₁: number of misplaced tiles

h₂: sum of the distances of the tiles from their goal position^a

^aSince blocks cannot move along diagonal, distance would involve horizontal and vertical moves (city-block or Manhattan distance)

- For our example $h_1 = 8$, $h_2 = 18$
- Effective branching factor (b*): if the algorithm generates N
 number of nodes and the solution is found at depth d, then

$$N+1=1+(b^*)+(b^*)^2+(b^*)^3+...+(b^*)^d$$

you get N and d when you run the algorithm. Substitute N and d in above equation to get b^*

Effective branching factor can characterize the quality of heuristic

Heuristic Function for 8-puzzle

Experiment: A* algorithm is run on 1200 random problems with solution length 2 to 24 (100 for each even number).

100	Search Cost (nodes generated)			Effective Branching Factor		
d	IDS	$A^*(h_1)$	A*(h2)	IDS	$A^*(h_1)$	A*(h ₂)
2	10	6	6	2.45	1.79	1.79
4	112	13	12	2.87	1.48	1.45
6	680	20	18	2.73	1.34	1.30
8	6384	39	25	2.80	1.33	1.24
10	47127	93	39	2.79	1.38	1.22
12	3644035	227	73	2.78	1.42	1.24
14	-	539	113	Alexander	1.44	1.23
16		1301	211	HOUSE CHEET THE	1.45	1.25
18	, 1864 <u>- Jan</u> es	3056	363	" sweet, as	1.46	1.26
20	ages a = of	7276	676	. The state of	1.47	1.27
22	- -	18094	1219		1.48	1.28
24		39135	1641	_ /2/2	1.48	1.26

 $h_2(n) \ge h_1(n)$ so h_2 dominates h_1 .

From the table it is clear that h₂ is better than h₁

It is better to use heuristic function with larger values

Admissible Heuristics from relaxed problem

Consider following relaxations (relaxed 8-puzzle problem)

- Tile could move anywhere (not only one square towards adjacent empty block). h_1 would give exact number of steps
- Tile could move one square in any direction, even into an occupied square. h_2 would give exact number of steps

State space graph of relaxed problem is a supergraph of original state space because of removal of restrictions

- Solution in original problem space would work for relaxed setting
- However there could be a better solution in relaxed setting (having lower cost)

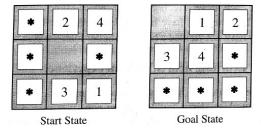
lower cost? means no overestimate on actual cost

A solution in relaxed problem space could be used as heuristic

$$h(n) = max\{h_1(n), h_2(n), ..., h_m(n)\}$$

Admissible Heuristics from subproblem

Following is a subproblem (only arrange 1-2-3-4)



- Pattern database stores each subproblem and solution cost
- Solution cost could be used as a heuristic
- we could also make database for 5-6-7-8
- But we cannot add the costs of 1-2-3-4 and 5-6-7-8
- Disjoint pattern database solution cost is taken as moves involving only 1,2,3,4 in the solution path. Now we can add cost of 1-2-3-4 and 5-6-7-8 database

Learning Heuristics from Experience

- Experience means solving lots of 8-puzzle problem
- Each optimal solution provides examples from which h(n) could be learned
- Using decision tree, neural nets or something else..
- Instead of state n we have to use feature say x₁(n) is number of misplaced tiles
- x₂(n) could be number of pairs of adjacent tiles that are not adjacent in the goal state

Heuristic would be obtained by combining $x_1(n)$, $x_2(n)$, ...

One such approach could be to use a liner combination

$$h(n) = c_1 x_1(n) + c_2 x_2(n) + ...$$



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- In many problems, path to the goal is irrelevant.
 - 8-queens order of adding queens is not important
 - Integrated circuit design
 - Factory floor layout
 - Job shop scheduling
 - Automatic programming
 - Telecommunication network optimization
 - Vehicle Routing
 - Portfolio management

Beyond Classical Search

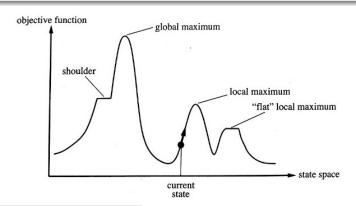
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- We go for local search algorithms in such scenarios
- These algorithms operate using a single current node (and move to its neighbor). Advantages include
 - Uses very little memory
 - 2 Can find reasonable solution in large or infinite (continuous) state space.

Local Search Algorithms

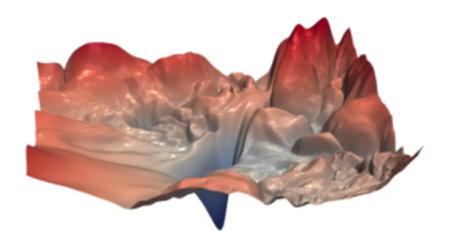
- Local search algorithms can solve pure optimization problems.
- They find optimum state based on some objective function ²

State space landscape can be used to proceed.



²Many optimization problems do not fit in standard search model. For example Darwinian evolution uses an objective function **reproductive fitness** but there is no goal test and path cost for this.

Real State Space Landscape



Hill Climbing Search

- Steepest ascent version
- Look around and choose the best
- Greedy-local search
- Makes quite rapid progress



Algorithm 1: Hill-Climbing(problem)

- Returned a state of problem that is local maximum.
- current ← Make-Node (problem.Initial-State)
- 3 while not terminate do
- 4 | neighbor ← a highest valued successor of current
- 5 | **if** *neighbor.Value* ≤ *current.Value* **then**
- 6 return current.State
- 7 | current ← neighbor

Hill Climbing Search

- Consider 8-queens problem
- In complete-state formulation has 8-queens in each state (one per column)
- Successor of a state are all possible states generated by moving a single queen to another square (each queen have 7 moves) so $8 \times 7 = 56$ successors.
- Heuristic cost function h is the number of pairs of queens that are attacking each other (h = 17 to h = 1)





Sometime algorithm reaches to a state where no further progress could be made.

Due to

- Local Maxima
- 2 Ridges
- Plateaux

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Put a limit on this (say 100) Now it solves 94% problems. But, takes 21 steps to solve and 64 for failure on an average.

Hill Climbing Variants

Stochastic hill climbing

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Random-restart hill climbing

if probability of success is p then run 1/p number of times.

- For our previous case ^a it needs 7 iterations.
- and 22 steps

^aWith random start 8-queens state, hill climbing gets stuck 86% of time.

Simulated Annealing (Metallurgy approach)

- Simulated annealing is hill climbing combined with random walk
- Step size is gradually reduced
- First Applied around 1980, for VLSI layout problem

```
Algorithm 2: Simulated-Annealing(problem, schedule)
```

```
1 current \leftarrow Make-Node (problem.Initial-State)

2 for t=1 to \infty do

3 | T \leftarrow schedule(t)

4 | if T=0 then return current

5 | next \leftarrow a randomly selected successor of current

6 | \Delta E \leftarrow next.Value - current.Value

7 | if \Delta E > 0 then current \leftarrow next

8 | else current \leftarrow next with probability e^{\Delta E/T}
```

Evolution

Recall Darwin's theory of evolution: "Survival of the fittest" 3

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Is it true? Let's formulate and try..



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Without guarantee, GA often finds an object of high fitness



Algorithm 3: GA (Fitness, F_{th} , p, r, m)

```
1 P ← generate p states at random
```

```
2 while max(Fitness(h_1), Fitness(h_2), ..., Fitness(h_p)) < F_{th} do
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Select: (1-r)p members of P
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4 Crossover: on (r \times p)/2 pairs to produce two offspring
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- 4 Crossover: on $(r \times p)/2$ pairs to produce two offspring
 - Mutation: randomly invert a bit of m percentage of population
- 6 **return** state h_i having maximum Fitness (h_i)
 - Fitness function is typically a heuristic
 - The fitness function is a criterion for ranking states to select states probabilistically for inclusion in the next generation population

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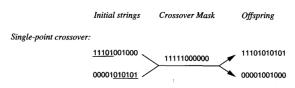
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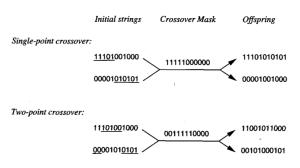
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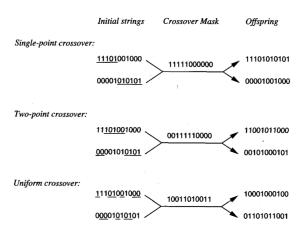
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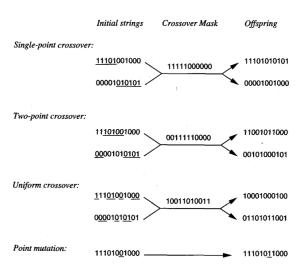
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- Elitist Model. select a small proportion of the fittest candidates in current population intact into the next generation









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• We know $Pr(h) = f(h)/(\sum f(h)) = f(h)/(n\overline{f}(t))$



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Thus, we can expect schema with above average fitness to be represented with increasing frequency on successive generations

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Full schema theorem thus provides a lower bound on the expected frequency of schema s, as follows

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Similar expression. More fit schemas will tend to grow in influence.

Thank You!

Thank you very much for your attention! Queries ?

(Reference⁴)

⁴1) Book - *AIMA*, ch-03/04, Russell and Norvig. 2) Book - Machine Learning, ch-09, *Tom Mitchell*